

[[A]]: $V^* \rightarrow S \rightarrow (S, V_\perp)_\perp$
[[I]]: $S \rightarrow S_\perp$
[[C]]: $S \rightarrow B_\perp$
[[E]]: $S \rightarrow V_\perp$

Algorithm

[[a]](v^*)(s) = pop_al_context(**[[i]]**(s'))
 where
 $a = (e^*, i)$
 $s = (w^*, \quad \quad \quad c'^*, h)$
 $s' = (w^*, \quad c :: c'^*, h)$
 $c = (\rho, \perp, 0)$
 $\rho = \text{bind}(\mathcal{E}, e^*, v^*)$

 where
 pop_al_context(
 $w^*, (_, v_\perp, _) :: c'^*, h$
) = $(w^*, c'^*, h), v_\perp$

Instruction

[[i1; i2]](s) = if $v_\perp = \perp$ then **[[i2]]**(s') else s'
 where
 $s' = \text{[[i1]]}(s)$
 = $(w^*, c :: c'^*, h)$
 $c = (_, v_\perp, _)$

[[if c then i1 else i2]](s) = if **[[c]]**(s) then **[[i1]]**(s) else **[[i2]]**(s)
[[either i1 or i2]](s) = **[[i1]]**(s)
[[enter e1: e2 after i]](s) = cleanup(**[[i]]**(s'))
 where
 $v = \text{[[e1]]}(s)$
 $v^* = \text{[[e2]]}(s)$
 $s = (\quad \quad \quad w^*, c :: c^*, h)$
 $s' = (w' :: w^*, c' :: c^*, h)$
 $c = (\rho, v_\perp, n)$
 $c' = (\rho, v_\perp, n + 1)$
 $w' = (v, \mathcal{E}, v^* ++ [\text{ConstructorV}(\text{"end"}, [])])$

where
 cleanup(s) = if $n > 0$ then cleanup(s_{res}) else s_{res}
 where
 $s = (w :: w^*, c :: c^*, h)$
 $s' = (w' :: w^*, c :: c^*, h)$
 $w = (v_{\text{ctx}}, v_{\text{val}}^*, v :: v_{\text{instr}}^*)$
 $w' = (v_{\text{ctx}}, v_{\text{val}}^*, \quad \quad \quad v_{\text{instr}}^*)$
 $v = \text{ConstructV}(\text{str}, v'^*)$
 $s_{\text{res}} = \text{[[lookup(p, str)]]}(v'^*)(s').1$
 = $(w_{\text{res}}^*, \quad \quad \quad c :: c^*, h)$
 $c = (\rho, v_\perp, n)$

[[assert c]](s) = if **[[c]]**(s) then s else \perp

[[push e]](s) = s'
 where
 $v = \text{[[e]]}(s)$
 $s = (w :: w^*, c^*, h)$
 $s' = (w' :: w^*, c^*, h)$
 $w = (v1, \quad \quad \quad v2^*, v3^*)$
 $w' = (v1, v :: v2^*, v3^*)$

[[pop e]](s) = s'
 where
 $s = (w :: w^*, c :: c^*, h)$
 $s' = (w' :: w^*, c' :: c^*, h)$
 $w = (v1, v :: v2^*, v3^*)$
 $w' = (v1, \quad \quad \quad v2^*, v3^*)$
 $c = (\rho, v_\perp, n)$
 $c' = (\rho', v_\perp, n)$
 $\rho' = \text{bind}(\rho, e, v)$

[[pop all e]](s) = s'
 where
 $s = (w :: w^*, c :: c^*, h)$
 $s' = (w' :: w^*, c' :: c^*, h)$
 $w = (v1, v2^*, v3^*)$
 $w' = (v1, \quad \mathcal{E}, v3^*)$
 $c = (\rho, v_\perp, n)$
 $c' = (\rho', v_\perp, n)$
 $\rho' = \text{bind}(\rho, e, v2^*)$

[[let e1 e2]](s) = s'
 where
 $v = \text{[[e2]]}(s)$
 $s = (w^*, c :: c^*, h)$
 $s' = (w^*, c' :: c^*, h)$
 $c = (\rho, v_\perp, n)$
 $c' = (\rho', v_\perp, n)$
 $\rho' = \text{bind}(\rho, e1, v)$

[[trap]](s) = \perp

[[nop]](s) = s

[[return e]](s) = s'
 where
 $v = \text{[[e]]}(s)$
 $s = (w^*, c :: c^*, h)$
 $s' = (w^*, c' :: c^*, h)$
 $c = (\rho, _, n)$
 $c' = (\rho, \quad v, n)$

[[execute e]](s) = **[[lookup(p, str)]]**(v^*)(s).1
 where
 $v = \text{[[e]]}(s)$
 = ConstructV (str, v^*)

[[execute all e]](s) = s_n
 where
 $v = \text{[[e]]}(s)$
 = $v_0 \dots v_{n-1}$
 $v_i = \text{ConstructV}(\text{str}_i, v_i^*)$
 $s_0 = s$
 $s_{i+1} = \text{[[lookup(p, str}_i\text{)]]}(v_i^*)(s_i).1$

[[exit]](s) = s'
 where
 $s = (w :: w'^*, c^*, \quad \quad \quad h)$
 $s' = (\quad \quad \quad w'^*, c'^*, \quad \quad \quad h)$
 $c'^* = \text{decrease_first_al_context_whose_number_is_nonzero}(c^*)$

[[ref x e]](s) = s'
 where
 $v = \text{[[e2]]}(s)$
 $s = (w^*, c :: c^*, h)$
 $s' = (w^*, c' :: c^*, h')$
 $c = (\rho, v_\perp, n)$
 $c' = (\rho', v_\perp, n)$
 $h', a = \text{alloc}(h, v)$
 $\rho' = \rho + (x \rightarrow a)$

[[replace e1 [p*] with e2]](s) = s'
 where
 $s = (w^*, c :: c^*, h)$
 $s' = (w^*, c :: c^*, h')$
 $c = (\rho, _, _)$
 $a = \text{[[e1]]}(s)$
 $v = \text{[[e2]]}(s)$
 $v' = \text{replace}(h(a), p^*, v)$
 $h' = h + (a \rightarrow v')$

where
 replace($v1, [], v2$) = $v2$
 replace($v1, p :: p'^*, v2$) = $v1 + (p \rightarrow \text{replace}(v1(p), p'^*, v2))$

Condition

[[C]]: $S \rightarrow B_\perp$
[[not c]](s) = $\neg \text{[[c]]}(s)$
[[c1 \oplus c2]](s) = **[[c1]]**(s) \oplus **[[c2]]**(s)
[[e1 \otimes e2]](s) = **[[e1]]**(s) \otimes **[[e2]]**(s)
[[e is of case t]](s) = **[[e]]**(s) == ConstructV($t, _$)

Expression

[[E]]: $S \rightarrow V_\perp$
[[n]](s) = n
[[t]](s) = t
[[e1 \oplus e2]](s) = **[[e1]]**(s) \oplus **[[e2]]**(s)
[[e*]](s) = 1
 where
 $n = |e^*|$
 $l(i) = \text{[[e*[i]]]}(s)$ (for $0 \leq i < n$)
[[e1^e2]](s) = 1
 where
 $v = \text{[[e1]]}(s)$
 $n = \text{[[e2]]}(s)$
 $l(i) = v$ (for $0 \leq i < n$)
[[e1 ++ e2]](s) = 1
 where
 $l1 = \text{[[e1]]}(s)$
 $l2 = \text{[[e2]]}(s)$
 $n1 = |l1|$
 $n2 = |l2|$
 $l[i] = \text{if } i < n1 \text{ then } l1[i] \text{ else } l2[i-n2]$ (for $0 \leq i < n1 +$
 $n2$)

[[|e|]](s) = $| \text{[[e]]}(s) |$
[[{(t \rightarrow e)*}]](s) = r
 where
 $n = |t^*| = |e^*|$
 $r[t^*[i]] = \text{[[e*[i]]]}(s)$

[[e[p]]](s) = **[[e]]**(s)[p]
[[e1[p*] <+ e2]](s) = ...
[[e1[p*] +> e2]](s) = ...
[[e1[p*] := e2]](s) = ...
[[t(e*)]](s) = ConstructV(t, v^*)
 where
 $v^*[i] = \text{[[e*[i]]]}(s)$ (for $0 \leq i < |e^*|$)

[[e1, e2]](s) = ($v1, v2$)
 where
 $v1 = \text{[[e1]]}(s)$
 $v2 = \text{[[e2]]}(s)$

[[f(e*)]](s) = v_\perp
 where
 $v^* = \text{[[e*]]}(s)$
 $v_\perp = \text{[[lookup(p, f)]]}(v^*)(s).2$

[[ref e]](s) = ...

[[current context]](s) = ...

[[x]](s) = $h(\rho(x))$
 where
 $s' = (w^*, c :: c'^*, h)$
 $c = (\rho, _, _)$

[[e^{x*, iter}]](s) = ...