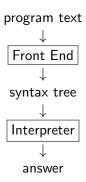
An Interpreter for a Simple Functional Language (LET)CS510

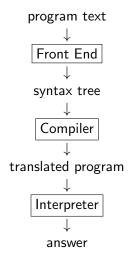
Expressions and Interpreters

- 1. Compiler vs Interpreter
- 2. A simple programming language: LET
- 3. Specification and Evaluation

Execution via Interpreter



Execution via Compiler

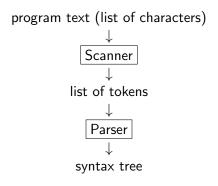


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```
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```

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\langle Expression \rangle ::= let \langle Identifier \rangle = \langle Expression \rangle in \langle Expression \rangle
```

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Non-examples

► (zero? -(55, -(x, 11)))

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- ► -(55, -(x, 11))
- zero? (-(55, -(x, 11)))
- let y = 23 in if zero?(y) then 4 else 6

- ► (zero? -(55, -(x, 11)))
- zero 4
- **►** +(1,2)

LET: Abstract Syntax

```
(define-datatype program program?
(a-program
          (exp1 expression?)))

(define-datatype expression expression?
;; ... continues in next slide ...
```

```
(define-datatype expression expression?
   (const-exp (num number?))
   (var-exp (var symbol?))
   (diff-exp
     (exp1 expression?)
5
     (exp2 expression?))
6
   (zero?-exp
7
     (exp1 expression?))
8
   (if-exp
9
     (exp1 expression?)
10
     (exp2 expression?)
11
     (exp3 expression?))
12
   (let-exp
13
     (var symbol?)
14
     (exp1 expression?)
15
     (body expression?)))
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Examples in Abstract Syntax

Let't revisit our earlier examples to translate them into the corresponding abstract syntax trees.

► Concrete syntax:

Abstract syntax:

Examples in Abstract Syntax

Concrete syntax:

► Abstract syntax

► Exercise: write the abstract syntax tree for this LET expression:

```
let y = 23 in if zero?(y) then 4 else 6
```

Interpreter for Expressions

▶ The next step is to define an interpreter for expressions

```
value-of: expression -> ???
```

- What should the return type of the interpreter be?
- ► Since we can write programs such as 2 and zero?(4) then either a number or a boolean
 - At least for now
- ▶ Let us define a new type for the return type of the interpreter that includes constructors for these two cases
- We'll call it the type of Expressed Values

Specification of Values (1/2)

Expressed Values: possible values of expressions.

$$ExpVal = Int + Bool$$

Here is a datatype declaration for ExpVal

Example expressions of this type are (num-val 2) and (bool-val #t)

Specification of Values (2/2)

Interface for the datatype of expressed values

Scheme Implementation

```
;; expval->bool : ExpVal -> Bool
2;; Page: 70
  (define expval->bool
    (lambda (v)
      (cases expval v
5
        (bool-val (bool) bool)
6
        (else (expval-extractor-error 'bool v)))))
7
8
  (define expval-extractor-error
    (lambda (variant value)
10
      (eopl:error 'expval-extractors
11
                   "Looking for a ~s, found ~s"
12
                    variant value)))
13
```

Interpreter for Expressions – The Need for Environments

- ► Now that we know the type of the interpreter for expressions value-of: expression -> expval we must move on to defining the interpreter itself
- ▶ Before doing so, however, one final observation
- ► The value of -(5,1) should clearly be (num-val 4)
- ▶ That of if zero?(-(4,4)) then 2 else 1 should clearly be 2
- ► What should the value of -(x,2) be?

Interpreter for Expressions – The Need for Environments

- ▶ What should the value of -(x,2) be?
- ▶ We need the value of x to be able to answer
- Hence we need environments
- ► The final type of the interpreter is therefore

```
value-of: { expression, environment } -> expval
```

Environments

```
(define-datatype environment environment?
(empty-env-record)
(extended-env-record
(syms (list-of symbol?))
(vals (list-of expval?))
(env environment?)))
```

- Two constructors
 - empty-env-record: constructs an empty environment
 - extended-env-record: extends a previous environment with a new association pair

Environment

- Function whose domain is a finite set of variables and whose range is the denoted values.
 - Denoted Values: values bound to variables.
 DenVal = Int + Bool
 - ► For now they coincide with Expressed Values
 - So we'll just use the latter
- ightharpoonup
 ho ranges over environments.
- [] denotes the empty environment.

Shorthands

- $[var = val]_{\rho}$ denotes (extend-env var val ρ).
- $[var_1 = val_1, var_2 = val_2]_{\rho} \text{ abbreviates}$ $[var_1 = val_1]([var_2 = val_2]_{\rho})$
- ▶ $[var_1 = val_1, var_2 = val_2, ...]$ denotes the environment in which the value of var_1 is val_1 , etc.

More Shorthands for Environments

```
[i = 1]

[v = 5]

[x = 10]
```

abbreviates

```
(extend-env 'i (num-val 1)
(extend-env 'v (num-val 5)
(extend-env 'x (num-val 10)
(empty-env))))))
```

▶ We'll call this environment init-env

Specifying the Behavior of the Interpreter for Expressions

► The value of a constant is the constant itself, as an expressed value

```
1 (value-of (const-exp n) \rho) = (num-val n)
```

▶ We must lookup the value of variables in the environment

```
1 (value-of (var-exp var) \rho) = (apply-env \rho var)
```

Note:

- ► This is not executable code (hence the shadow in the frame)
- ▶ It specifies the behavior of value-of in terms of equations
- On the next slide we show sample code for value-of that satisfies these equations

Implementing the first cases of value-of

```
(define value-of
      (lambda (exp env)
        (cases expression exp
3
          ;; corresponds to equation for constant
4
     expressions
          (const-exp (num) (num-val num))
          ;; corresponds to equations for
6
     variables
          (var-exp (var) (apply-env env var))
7
8
          ;; more cases still to come ....
9
10
   )))
11
```

Specifying the Behavior of the Interpreter for Expressions

► Difference is computed by first computing the values of the arguments and then performing the difference itself

```
(value-of (diff-exp exp_1 exp_2) \rho)

2 = (num-val)

3 (-

4 (expval->num (value-of exp_1 \rho))

5 (expval->num (value-of exp_2 \rho))))
```

Specifying the Behavior of the Interpreter for Programs

- ▶ A program is an expression that may contain free variables.
 - ▶ These represent the top-level declarations
- ► The value of the program is the value of the expression in a suitable environment.
- ▶ We assume the initial environment init-env defined above

```
1 (value-of-program exp) = (value-of exp init-env)
```

► This initial environment allows us to write examples involving variables *i*, *v* and *x* in our program without having to declare them

Specifying the Behavior of zero? and if-exp

```
(value-of (zero?-exp exp_1) \rho) =

(bool-val)

(zero? (expval->num (value-of exp_1 \rho))))

(value-of (if-exp exp_1 exp_2 exp_3) \rho) =

(if (expval->bool (value-of exp_1 \rho))

(value-of (exp_2 \rho))

(value-of (exp_3 \rho))
```

Specifying the Behavior of let

The right-hand side of the let is also an expression, so it can be arbitrarily complex. For example,

```
let x = 7
in let y = 2
in let y = let x = -(x,1)
in -(x,y)
in -(-(x,8), y)
```

- Here the x declared on the third line is bound to 6
- ► So the value of y is 4
- ▶ The value of the entire expression is ((-1) 4) = -5.
- ▶ We can write down the specification as an equation.

Behavior of the Interpreter for let-expressions

Note how the body is evaluated in an extended environment

```
(value-of (let-exp var exp_1 body) \rho) =

2 (let ((val (value-of exp_1 \rho)))

3 (value-of body [var = val]\rho))
```

 Important: static scoping is implemented by extending environments

The Interpreter for LET

- Code available from http://www.eopl3.com
- Directory chapter3/let-lang
- Open top.scm in Racket
- ▶ There are a number of tests in tests.scm
- You can run them with run-one. Eg.

```
1 > (run-one 'simple-let-1)
2 (num-val 3)
```