

Adding Procedures to LET (PROC)

CS510

PROC: A Language with Procedures

Extending our language LET with procedures:

1. Extending the concrete and abstract syntax.
2. Extending the set of Expressed Values.
3. Specification and Implementation of the interpreter.

The PROC-Language

An Interpreter for PROC

Procedural Representation

PROC: Concrete Syntax

Extending the concrete syntax of LET

$\langle Program \rangle$	$::=$	$\langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$\langle Number \rangle$
$\langle Expression \rangle$	$::=$	$\langle Identifier \rangle$
$\langle Expression \rangle$	$::=$	$- (\langle Expression \rangle , \langle Expression \rangle)$
$\langle Expression \rangle$	$::=$	$zero? (\langle Expression \rangle)$
$\langle Expression \rangle$	$::=$	$if \langle Expression \rangle$ $then \langle Expression \rangle \text{ else } \langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$let \langle Identifier \rangle = \langle Expression \rangle \text{ in } \langle Expression \rangle$

PROC: Concrete Syntax

Extending the concrete syntax of LET

$\langle Program \rangle$	$::=$	$\langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$\langle Number \rangle$
$\langle Expression \rangle$	$::=$	$\langle Identifier \rangle$
$\langle Expression \rangle$	$::=$	$- (\langle Expression \rangle, \langle Expression \rangle)$
$\langle Expression \rangle$	$::=$	$zero? (\langle Expression \rangle)$
$\langle Expression \rangle$	$::=$	$if \langle Expression \rangle$ $then \langle Expression \rangle \text{ else } \langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$let \langle Identifier \rangle = \langle Expression \rangle \text{ in } \langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$proc (\langle Identifier \rangle) \langle Expression \rangle$
$\langle Expression \rangle$	$::=$	$(\langle Expression \rangle \langle Expression \rangle)$

Examples of Expressions in PROC

```
1 let f = proc (x) -(x,11)
2 in (f (f 77))
3
4 (proc (f) (f (f 77))    ;; rator
5   proc (x) -(x,11))    ;; rand
6
7 let x = 200
8 in let f = proc (z) -(z,x)
9     in (f 1)
10
11 let x = 200
12 in let f = proc (z) -(z,x)
13     in let x = 100
14         in let g = proc (z) -(z,x)
15             in -((f 1), (g 1))
```

PROC: Abstract Syntax (1/3)

```
1 (define-datatype program program?  
2   (a-program  
3     (exp1 expression?)))
```

PROC: Abstract Syntax (2/3)

```
1 (define-datatype expression expression?
2   (const-exp
3     (num number?))
4   (diff-exp
5     (exp1 expression?)
6     (exp2 expression?))
7   (zero?-exp
8     (exp1 expression?))
9   (if-exp
10    (exp1 expression?)
11    (exp2 expression?)
12    (exp3 expression?))
13   (var-exp
14     (var symbol?))
15   (let-exp
16     (var symbol?)
17     (exp1 expression?)
18     (body expression?))
```


The new variants: proc-exp and call-exp (3/3)

```
1  (proc-exp  
2    (var symbol?)  
3    (body expression?))  
4  (call-exp  
5    (rator expression?)  
6    (rand expression?)))
```

Concrete Syntax vs Abstract Syntax

► Concrete

```
1 let f = proc (x) -(x,11) in (f (f 77))
```

► Abstract

```
1 (a-program
2   (let-exp
3     'f
4     (proc-exp 'x (diff-exp (var-exp 'x)
5                             (const-exp 11)))
6     (call-exp (var-exp 'f)
7               (call-exp (var-exp 'f)
8                           (const-exp 77))))))
```

Concrete Syntax vs Abstract Syntax

► Concrete

```
1 let f = proc (x) -(x,11) in (f (f 77))
```

► Abstract

```
1 (a-program  
2   (let-exp  
3     'f  
4     (proc-exp 'x (diff-exp (var-exp 'x)  
5       (const-exp 11)))  
6     (call-exp (var-exp 'f)  
7       (call-exp (var-exp 'f)  
         (const-exp 77)))))
```

Concrete Syntax vs Abstract Syntax

► Concrete

```
1 let f = proc (x) -(x,11) in (f (f 77))
```

► Abstract

```
1 (a-program
2   (let-exp
3     'f
4     (proc-exp 'x (diff-exp (var-exp 'x)
5                             (const-exp 11)))
6     (call-exp (var-exp 'f)
7               (call-exp (var-exp 'f)
8                           (const-exp 77))))))
```

The PROC-Language

An Interpreter for PROC

Procedural Representation

The Interpreter for PROC

- ▶ What is the result of evaluating the expression 2?

The Interpreter for PROC

- ▶ What is the result of evaluating the expression 2?

`num-val 2`

- ▶ What is the result of evaluating the expression `zero?(2)`?

The Interpreter for PROC

- ▶ What is the result of evaluating the expression 2?

`num-val 2`

- ▶ What is the result of evaluating the expression `zero?(2)`?

`bool-val #f`

- ▶ What is the result of evaluating the expression `proc(x) -(7,x)`?

The Interpreter for PROC

- ▶ What is the result of evaluating the expression 2?

num-val 2

- ▶ What is the result of evaluating the expression zero?(2)?

bool-val #f

- ▶ What is the result of evaluating the expression proc(x) -(7,x)?

- ▶ Its not a number
- ▶ Its not a boolean
- ▶ Its a **closure**
 - ▶ Record that stores the formal parameter and body
 - ▶ More details to follow

The Interpreter for PROC

- ▶ Before (for the LET-language)

value-of: { expression, **environment** } -> expval

- ▶ Now

value-of: { expression, **environment** } -> expval

- ▶ What's the difference?

- ▶ The definition of expressed values
- ▶ Our syntax now supports procedures

- ▶ Before

ExpVal = Int + Bool

- ▶ Now

ExpVal = Int + Bool + **Clos**

The Interpreter for PROC

value-of: { expression, environment } -> expval

► Expressed values before (LET)

```
1 (define-datatype expval expval?
2   (num-val
3     (num number?))
4   (bool-val
5     (bool boolean?))
```

► Now (PROC)

```
1 (define-datatype expval expval?
2   (num-val
3     (num number?))
4   (bool-val
5     (bool boolean?))
6   (proc-val
7     (clos clos?)))
```

Defining Closures

What should the value of the following expression be?

`proc (x) -(x,11)`

- ▶ A datatype that records the parameter and its body
- ▶ However, is that enough?
- ▶ What about the procedures `f` and `g` below?

```
1 let x = 200
2 in let f = proc (z) -(z,x)
3     in let x = 100
4         in let g = proc (z) -(z,x)
5             in -((f 1), (g 1))
```

- ▶ The value of `f` and `g` depends on the value of `x`.

Defining Closures

What should the value of the following expression be?

`proc (x) -(x,11)`

- ▶ A datatype that records the parameter and its body
- ▶ However, is that enough?
- ▶ What about the procedures `f` and `g` below?

```
1 let x = 200
2 in let f = proc (z) -(z,x)
3     in let x = 100
4         in let g = proc (z) -(z,x)
5             in -((f 1), (g 1))
```

- ▶ The value of `f` and `g` depends on the value of `x`.

Defining Closures

```
1 let x = 200
2 in let f = proc (z) -(z,x)
3     in let x = 100
4         in let g = proc (z) -(z,x)
5             in -((f 1), (g 1))
```

- ▶ The value of a procedure depends on the environment in which it is evaluated.
 - ▶ f and g only differ in the value of x .
- ▶ Summary of what a closure should contain:
 1. The formal parameter of the procedure
 2. The body of the procedure
 3. The environment extant at the point where the procedure was evaluated

Representing Closures

Datatype Representation

```
1 (define-datatype clos clos?  
2   (closure  
3     (var symbol?)  
4     (body expression?)  
5     (saved-env environment?)))
```

- ▶ An alternative representation using Scheme procedures is also possible

Specifying the Behavior of the Interpreter – `proc-exp`

```
1 (value-of (proc-exp var body)  $\rho$ ) =  
2   (proc-val (closure var body  $\rho$ ))
```

- ▶ Recall from above: `closure` is a constructor with arguments:
 - ▶ a formal parameter `var`,
 - ▶ an expression `body`,
 - ▶ and an environment ρ .
- ▶ It constructs elements of type `clos?`
- ▶ **Note:** In the EOPL book this constructor is called `procedure`
- ▶ `proc-val` tags the procedure so that it can be considered of type `expval?`
 - ▶ Remember that `value-of` returns an expressed value.

Implementation

New clauses for value-of

```
1 (define value-of
2   (lambda (exp env)
3     (cases expression exp
4       (proc-exp (var body)
5         (proc-val (closure var body env)))
6
7       (call-exp (rator rand)
8         ... to be defined ...
9
10    )))
```

Specifying the Behavior of the Interpreter – `call-exp`

- ▶ A procedure call is represented as

`(call-exp rator rand)`

in the abstract syntax

- ▶ We must therefore give meaning to

`(value-of (call-exp rator rand) ρ)`

1. Evaluate rator to a procedure (check it is `proc-val`)
 2. Evaluate rand to an argument
 3. Pass argument to procedure
- ▶ Lets specify its behavior using equations

Evaluating Procedure Calls

- ▶ We must give meaning to $(\text{value-of } (\text{call-exp rator rand}) \rho)$
 1. Evaluate rator to a procedure (check it is `proc-val`)
 2. Evaluate rand to an argument
 3. Pass argument to procedure

```
1 (value-of (call-exp rator rand)  $\rho$ )  
2   = (let ((proc (expval->proc (value-of rator  $\rho$ )))  
3         (arg (value-of rand  $\rho$ )))  
4     (apply-procedure proc arg))
```

- ▶ *expval*->*proc* checks argument is actually a procedure (i.e. it was constructed with `proc-val`)
- ▶ What does *apply-procedure* do?

Evaluating Procedure Calls

Consider the example:

```
1 (apply-procedure
2   (closure 'x
3           (diff-exp
4             (var-exp 'x)
5             (num-val 11))
6            $\rho$ )
7   (num-val 20))
```

- The value of this expression is the value of the body $-(x, 11)$ where x 's value is 20.

In general:

```
1 (apply-procedure (closure var body  $\rho$ ) val)
2   = (value-of body [var=val] $\rho$ )
```

Evaluating Procedure Calls

```
1 (apply-procedure (closure var body  $\rho$ ) val)
2   = (value-of body [var=val] $\rho$ )
```

- The code corresponding to this specification is:

```
1 (define apply-procedure
2   (lambda (clos1 val)
3     (cases clos clos1
4       (closure (var body saved-env)
5         (value-of body
6           (extend-env var val saved-env))))))
```

Implementation

New clauses for value-of

```
1 (define value-of
2   (lambda (exp env)
3     (cases expression exp
4       (proc-exp (var body)
5         (proc-val (closure var body env)))
6
7       (call-exp (rator rand)
8         (let ((proc (expval->proc (value-of rator
9           env))))
10           (arg (value-of rand env)))
11           (apply-procedure proc arg)))
12   )))
```

The PROC-Language

An Interpreter for PROC

Procedural Representation

Representing Procedures using Procedures

- ▶ We have defined `clos?`, `closure`, and `apply-procedure` for the procedures represented as inductive types
- ▶ We'll see one other implementation: using the procedures of Scheme

```
1 (define clos?  
2   (lambda (val)  
3     (procedure? val)))
```

- ▶ `procedure?` is Scheme's built-in predicate

Representing Procedures

procedural representation

$\text{procedure} : \text{Var} \times \text{Exp} \times \text{Env} \rightarrow \text{Proc}$

```
1 (define closure
2   (lambda (var body env)
3     (lambda (val)
4       (value-of body (extend-env var val env))))))
```

$\text{apply-procedure} : \text{Proc} \times \text{ExpVal} \rightarrow \text{ExpVal}$

```
1 (define apply-procedure
2   (lambda (clos1 val)
3     (clos1 val)))
```

The Interpreter for PROC

- ▶ Code available from <http://www.eopl3.com>
- ▶ Directory `chapter3/proc-lang/ds-rep`
- ▶ Open `top.scm` in Racket
- ▶ There are a number of tests in `tests.scm`
- ▶ You can run them with `run-one`. Eg.

```
1 > (run-one 'nested-procs)
2 (num-val -1)
```