

[illegible]

2) Let w, z be strings over the alphabet $\{0, 1, \$\}$. Define a pushdown automaton which accepts any string $w\$z$, where w^{-1} is a subsequence of z , and both w and z do not contain “\$”.

Ex: “110\$001001”

Ex: “0101\$01001110101”

Input	0				1				\$				ϵ			
Stack	0	1	\$	ϵ	0	1	\$	ϵ	0	1	\$	ϵ	0	1	\$	ϵ
st																$(q_1, \$)$
q_1				$(q_1, 0)$				$(q_1, 1)$				(q_2, ϵ)				
q_2	(q_2, ϵ)			(q_2, ϵ)		(q_2, ϵ)		(q_2, ϵ)							(fin, ϵ)	
fin																

Pumping Lemma

1) Let $\Sigma = \{0, 1\}$. Prove that the set $\{w^2 \mid w \in \Sigma^*\}$ is not context-free.

$$w = 0^p 1 0^p 1$$

$$00\dots 00 \quad 0 \quad 1 \quad 0 \quad 00\dots 001$$

$$\backslash \quad u \quad / \quad \backslash v^p / \quad \backslash x / \quad \backslash y^p / \quad \backslash \quad z \quad /$$

$$|vyx| \leq p$$

$$|vyx| = 2p + 1$$

Therefore, $\{w^2 \mid w \in \Sigma^*\}$ is not context free.