





LECTURE 34

Selection Sorting; Insertion Sorting; Shell sort; Java's sorting methods

Quiz (and Final) Study Guide

- Be able to execute the following algorithms
 - Counting sort
 - selection sort
 - bubble sort
 - insertion sort
 - Shell sort
 - merge sort
 - heapsort
 - quicksort
- Understand the differences in performance of these algorithms
- Read chapter 8 of the book (but don't take everything it says for granted (see slide 120)!
- For even more fun, watch
- 1. https://www.youtube.com/watch?v=OyQ6c-XHN8Y (Merge sort dance) and other dances by the same group
- https://www.youtube.com/watch?v=ZZuD6iUe3Pc (Comparison of different sorting algorithms)

Selection Sort

Selection Sort

- Find the minimum (or maximum)
- Put it in its place at the lowest (or the highest) index value
- Increase (or decrease) the index

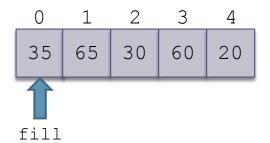
Trace of Selection Sort

- 1. for fill = 0 to n 2 do
- 2. Set posMin to the subscript of the smallest item in the subarray starting at subscript fill
- 3. Exchange the item at posMin with the one at

0	1	2	3	4
35	65	30	60	20

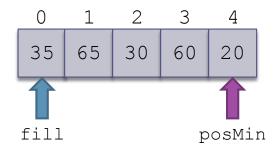
n	5
fill	
posMin	

- **▶1. for** fill = 0 **to** n 2 **do**
 - 2. Set posMin to the subscript of the smallest item in the subarray starting at subscript fill
 - 3. Exchange the item at posMin with the one at fill



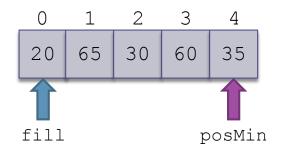
n	5
fill	0
posMin	

- 1. for fill = 0 to n 2 do
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 - 3. Exchange the item at posMin with the one at fill



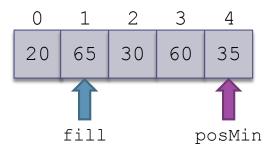
n	5
fill	0
posMin	4

- 1. for fill = 0 to n 2 do
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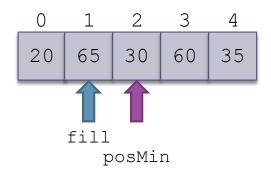
n	5
fill	0
posMin	4

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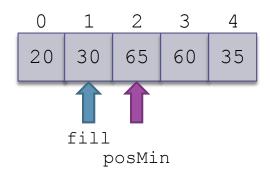
n	5
fill	1
posMin	4

- 1. for fill = 0 to n 2 do
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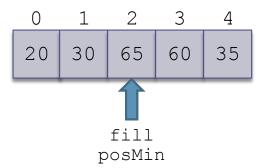
n	5
fill	1
posMin	2

- 1. for fill = 0 to n 2 do
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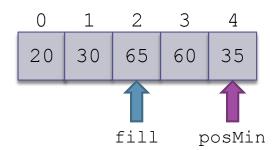
n	5
fill	1
posMin	2

- **▶ 1. for** fill = 0 to n 2 do
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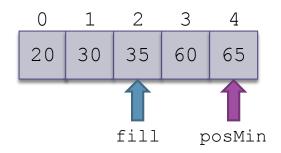
n	5
fill	2
posMin	2

- 1. for fill = 0 to n 2 do
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 - 3. Exchange the item at posMin with the one at fill



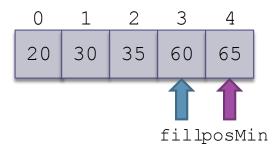
n	5
fill	2
posMin	4

- 1. for fill = 0 to n 2 do
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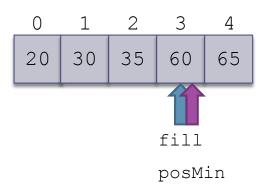
n	5
fill	2
posMin	4

- **▶1. for** fill = 0 **to** n 2 **do**
 - 2. Set posMin to the subscript of the smallest item in the subarray starting at subscript fill
 - 3. Exchange the item at posMin with the one at



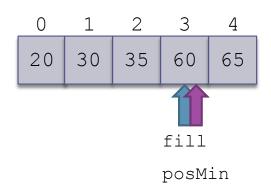
n	5
fill	3
posMin	4

- 1. for fill = 0 to n 2 do
- Set posMin to the subscript of the smallest item in the subarray starting at subscript fill
 - 3. Exchange the item at posMin with the one at fill



n	5
fill	3
posMin	3

- 1. for fill = 0 to n 2 do
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 - Exchange the item at posMin with the one at fill



n	5
fill	3
posMin	3

- 1. for fill = 0 to n 2 do
- 2. Set posMin to the subscript of the smallest item in the subarray starting at subscript fill
- 3. Exchange the item at posMin with the one at

0	1	2	3	4
20	30	35	60	65

n	5
fill	3
posMin	3

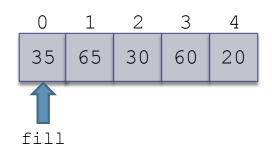
Trace of Selection Sort Refinement

n	5
fill	
posMin	
next	

0	1	2	3	4
35	65	30	60	20

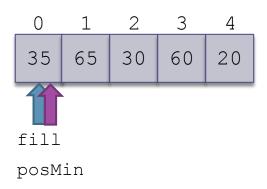
- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	
next	



- ▶ 1. for fill = 0 to n 2 do
 - 2. Initialize posMin to fill
 - 3. for next = fill + 1 to n 1 do
 - 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	0
next	

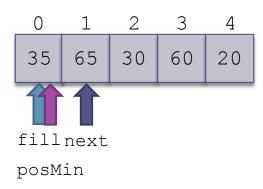


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1. for fill = 0 \text{ to } n - 2 \text{ do}
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2. Initialize posMin to fill

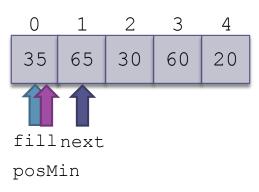
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	0
next	1



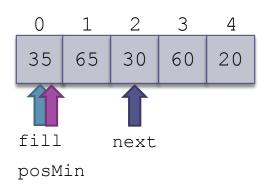
- 1. for fill = 0 to n 2 do
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 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	0
next	1



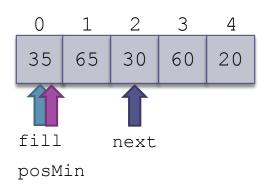
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 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	0
next	2



- 1. for fill = 0 to n 2 do
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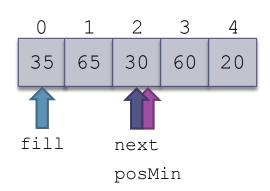
n	5
fill	0
posMin	0
next	2



- 1. for fill = 0 to n 2 do
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 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

5.

n	5
fill	0
posMin	2
next	2

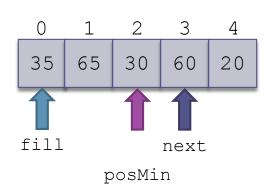


- 1. for fill = 0 to n 2 do
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Reset posMin to next

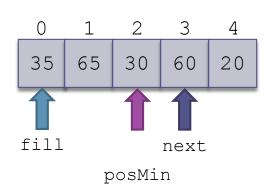
6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	2
next	3



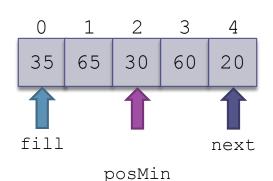
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n	5
fill	0
posMin	2
next	3



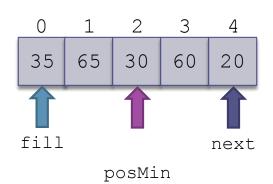
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 - 6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	2
next	4



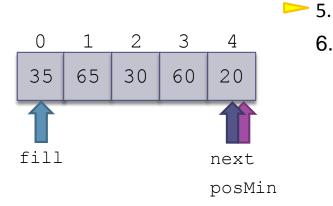
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n	5
fill	0
posMin	2
next	4



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n	5
fill	0
posMin	4
next	4

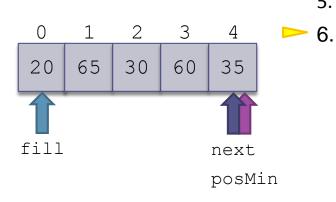


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Reset posMin to next

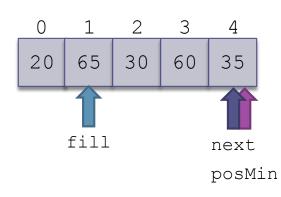
6. Exchange the item at posMin with the one at fill

n	5
fill	0
posMin	4
next	4



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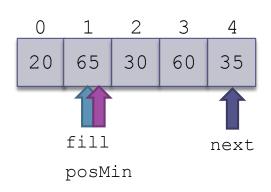
n	5
fill	1
posMin	4
next	4



```
▶ 1. for fill = 0 to n - 2 do
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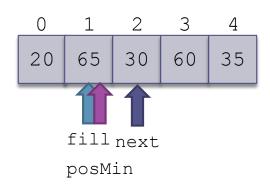
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
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- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

n	5
fill	1
posMin	1
next	4



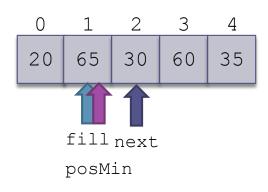
- 1. for fill = 0 to n 2 do
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 - 3. for next = fill + 1 to n 1 do
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n	5
fill	1
posMin	1
next	2



- 1. for fill = 0 to n 2 do
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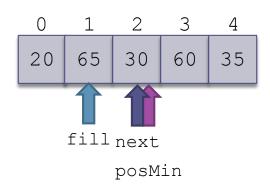
n	5
fill	1
posMin	1
next	2



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 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

5.

n	5
fill	1
posMin	2
next	2



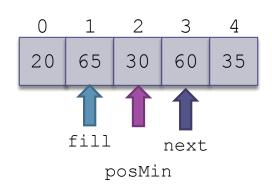
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Reset posMin to next

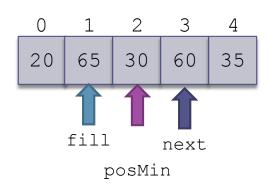
6. Exchange the item at posMin with the one at fill

n	5
fill	1
posMin	2
next	3



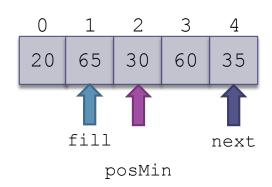
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n	5
fill	1
posMin	2
next	3



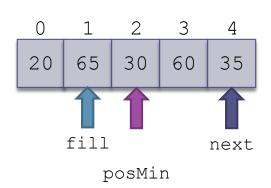
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n	5
fill	1
posMin	2
next	4



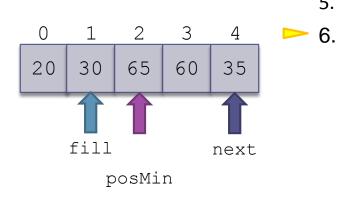
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fill	1
posMin	2
next	4



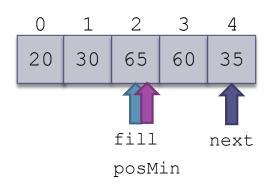
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n	5
fill	1
posMin	2
next	4



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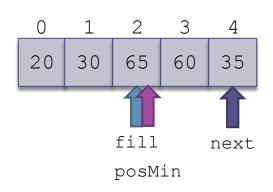
n	5
fill	2
posMin	2
next	4



```
► 1. for fill = 0 to n - 2 do
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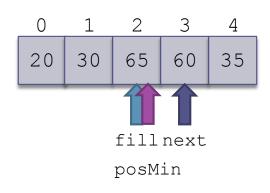
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

n	5
fill	2
posMin	2
next	4



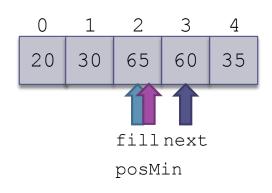
- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
 - 3. for next = fill + 1 to n 1 do
 - 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	2
posMin	2
next	3



- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- > 3. for next = fill + 1 to n 1 do
 - 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

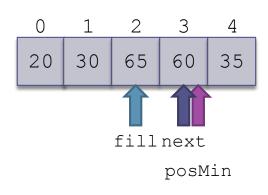
n	5
fill	2
posMin	2
next	3



- 1. for fill = 0 to n 2 do
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 - 5. Reset posMin to next
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5.

n	5
fill	2
posMin	3
next	3



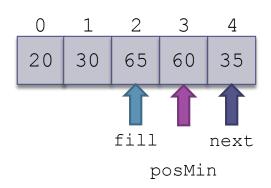
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1. for fill = 0 \text{ to } n - 2 \text{ do}
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- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin

Reset posMin to next

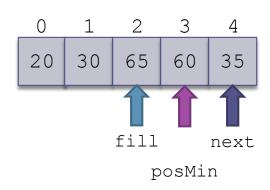
6. Exchange the item at posMin with the one at fill

n	5
fill	2
posMin	3
next	4



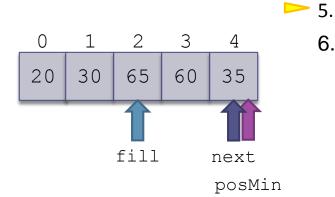
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 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	2
posMin	4
next	4

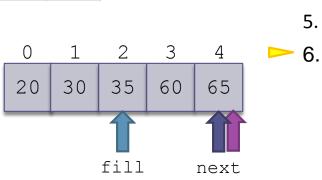


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- 3. for next = fill + 1 to n 1 do
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Reset posMin to next

6. Exchange the item at posMin with the one at fill

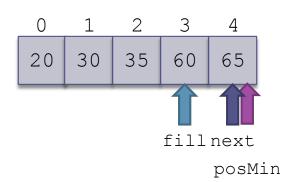
n	5
fill	2
posMin	4
next	4



posMin

- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
 - Exchange the item at posMin with the one at fill

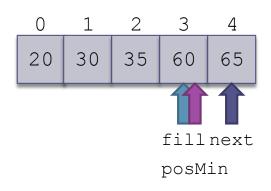
n	5
fill	3
posMin	4
next	4



```
▶ 1. for fill = 0 to n - 2 do
```

- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

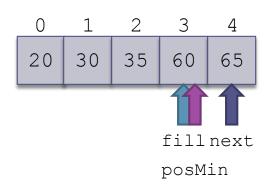
n	5
fill	3
posMin	3
next	4



```
1. for fill = 0 \text{ to } n - 2 \text{ do}
```

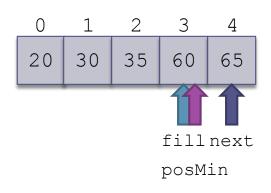
- 2. Initialize posMin to fill
 - 3. for next = fill + 1 to n 1 do
 - 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	3
posMin	3
next	4



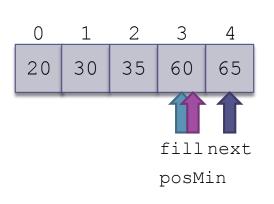
- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- > 3. for next = fill + 1 to n 1 do
 - 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	3
posMin	3
next	4



- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
 - 5. Reset posMin to next
 - 6. Exchange the item at posMin with the one at fill

n	5
fill	3
posMin	3
next	4



- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
 - Exchange the item at posMin with the one at fill

n	5
fill	3
posMin	3
next	4

0	1	2	3	4
20	30	35	60	65

- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

Analysis of Selection Sort

This loop is performed n-1 times

```
1. for fill = 0 to n - 2 do
```

- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

Analysis of Selection Sort (cont.)

```
2. Initialize posMin to fill

3. for next = fill + 1 to n - 1 do

4. if the item at next is less than the item at posMin

5. Reset posMin to next
exchanges

6. Exchange the item at posMin with the one at fill
```

1. for fill = 0 to n - 2 do

Analysis of Selection Sort (cont.)

This comparison is performed (n-1-fill) times for each value of fill = 0, 1, ... (n-2), resulting in (n-1) + (n-2) + ... 1 = n(n-1)/2 comparisons

```
1. for fill = 0 to n - 2 do
```

- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

Analysis of Selection Sort (cont.)

- 1. for fill = 0 to n 2 do
- 2. Initialize posMin to fill
- 3. for next = fill + 1 to n 1 do
- 4. if the item at next is less than the item at posMin
- 5. Reset posMin to next
- 6. Exchange the item at posMin with the one at fill

Code for Selection Sort (cont.)

□ Listing 8.1(SelectionSort.java, pages 426 - 427)

Insertion Sort

The simile

Arranging a hand of cards



1) 8 is dealt

The simile

Arranging a hand of cards





2) 7 is dealt

The simile

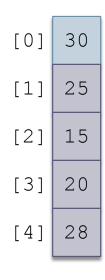
Arranging a hand of cards







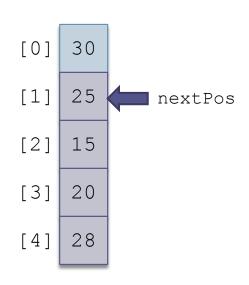
Trace of Insertion Sort



start with a sorted subarray consisting of only the first element

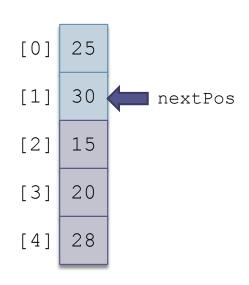
- 1. for each array element from the second
 (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





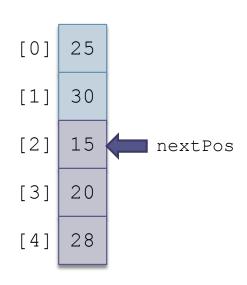
- 1. for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





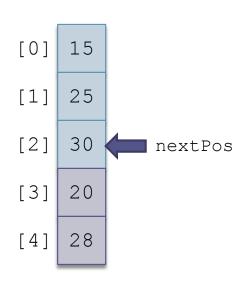
- for each array element from the second (nextPos = 1) to the last
- 2. Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





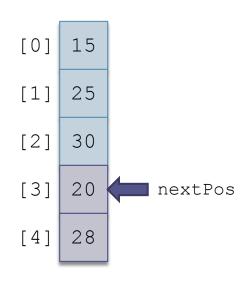
- 1. for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





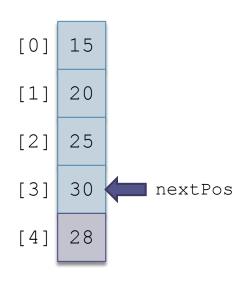
- 1. for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





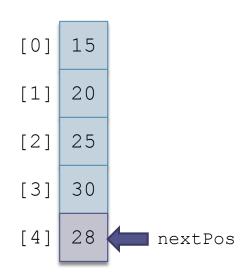
- 1. for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





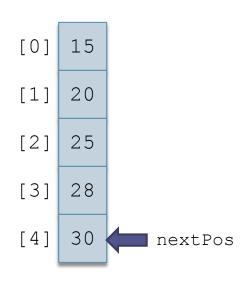
- for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





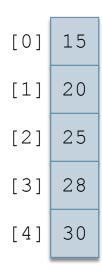
- for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





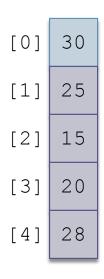
- for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element





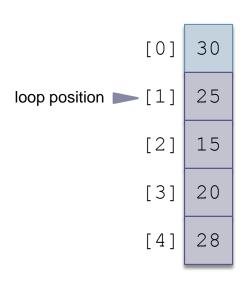
- for each array element from the second (nextPos = 1) to the last
- Insert the element at nextPos where it belongs in the array, increasing the length of the sorted subarray by 1 element

Trace of Insertion Sort Refinement



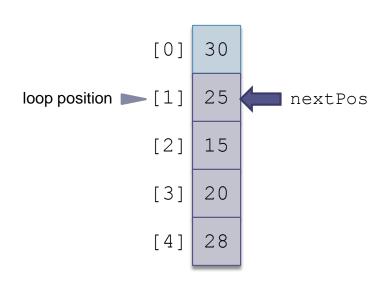
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- 6. **Decrement** nextPos by 1
- 7. Insert nextVal at nextPos

nextPos	1
nextVal	



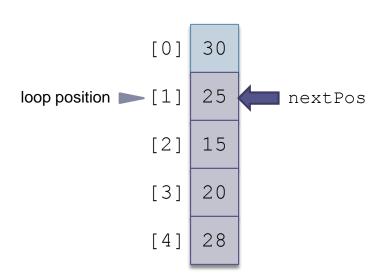
- 1. for each array element from the second
 (nextPos = 1) to the last
 - 2. nextPos is the position of the element to insert
 - Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	



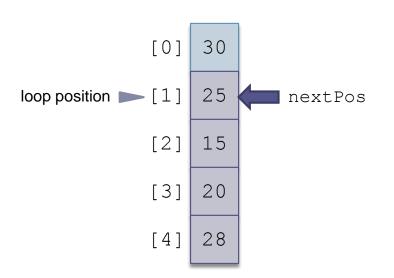
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
 - 3. Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	25



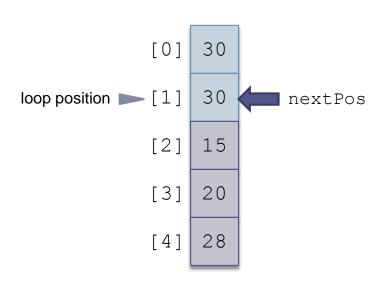
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	25



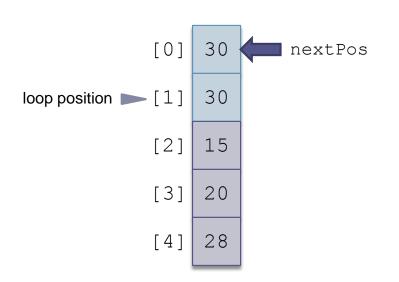
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	25



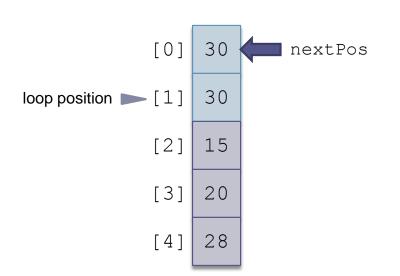
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	25



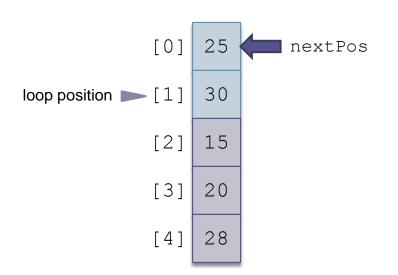
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	25



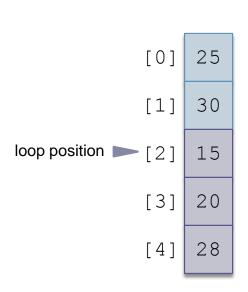
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- Save the value of the element to insert in next.Val
- ▶ 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	25



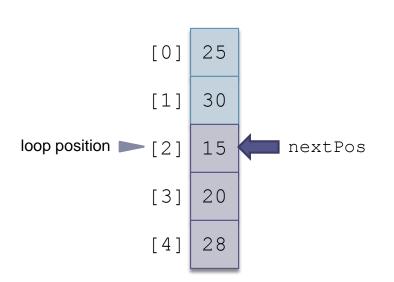
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- 6. **Decrement** nextPos by 1
- > 7. Insert nextVal at nextPos

nextPos	0
nextVal	25



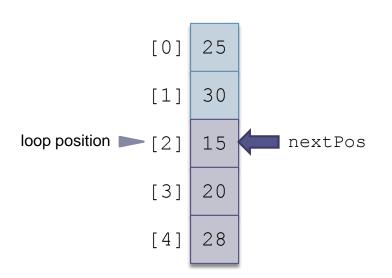
- 1. for each array element from the second
 (nextPos = 1) to the last
 - nextPos is the position of the element to insert
 - Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	25



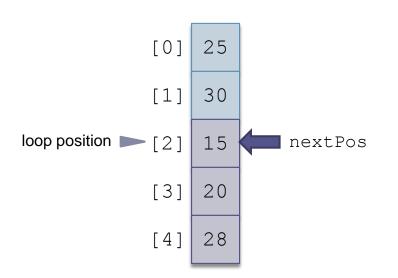
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
 - 3. Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	15



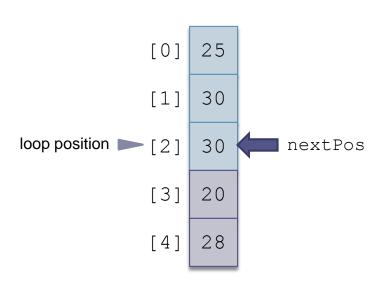
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	15



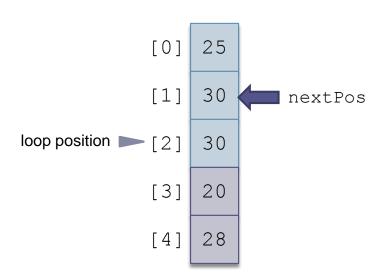
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- ▶ 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	15



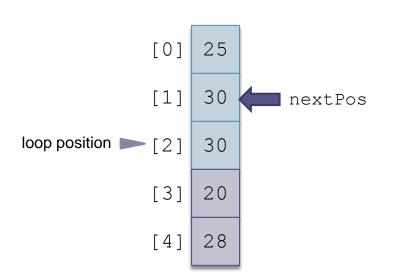
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	15



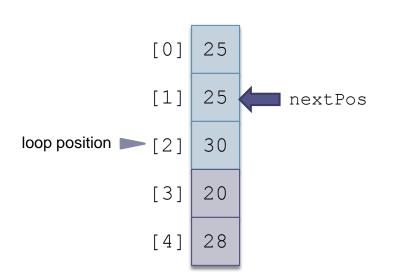
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	15



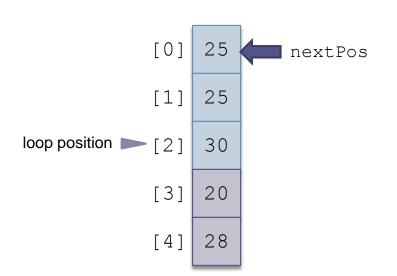
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos - 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	15



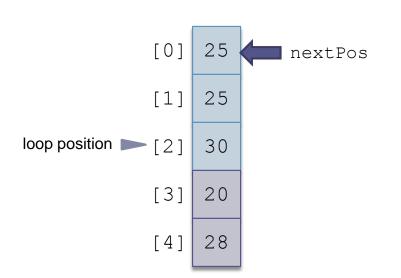
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	15



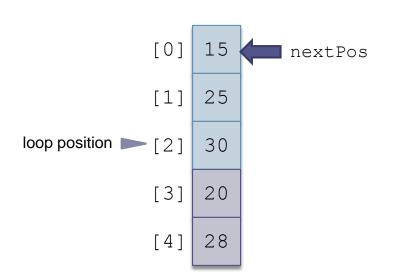
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	15



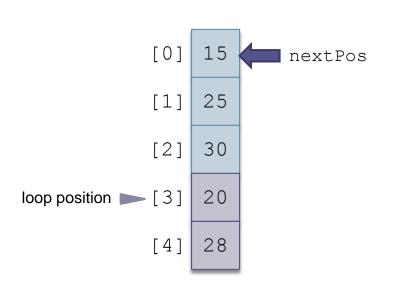
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- ▶ 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	0
nextVal	15



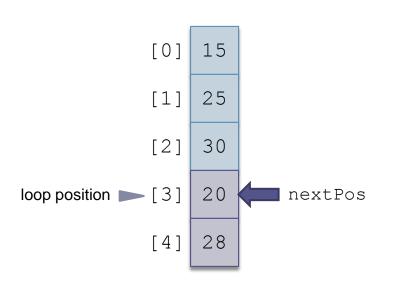
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- 6. **Decrement** nextPos by 1
- > 7. Insert nextVal at nextPos

nextPos	0
nextVal	15



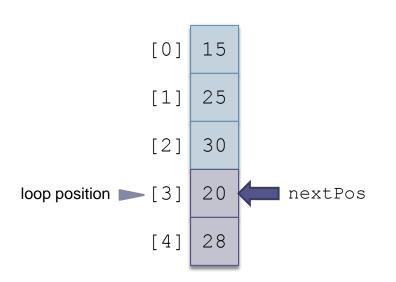
- 1. for each array element from the second
 (nextPos = 1) to the last
 - 2. nextPos is the position of the element to insert
 - 3. Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	15



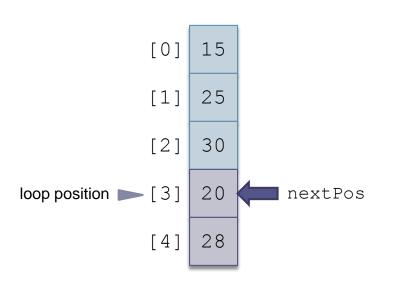
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
 - 3. Save the value of the element to insert in next Val
 - 4. while nextPos > 0 and the element at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	20



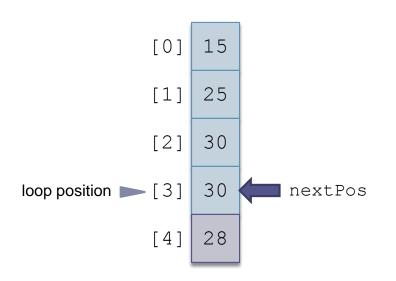
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	20



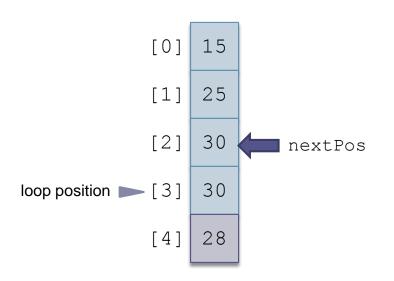
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- ▶ 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	20



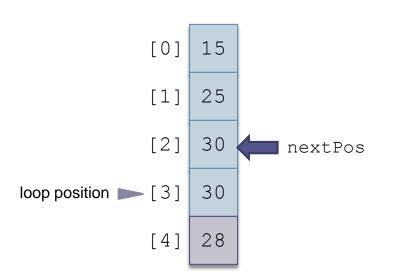
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos − 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	20



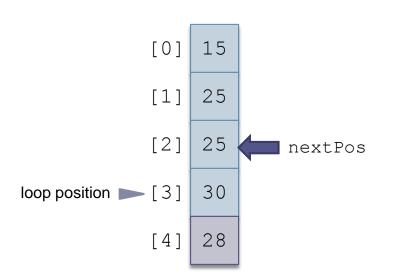
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	20



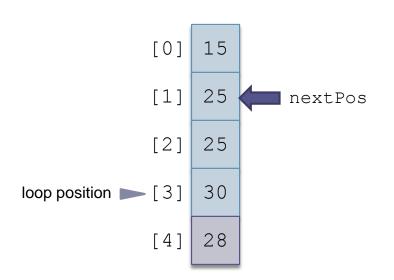
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- 3. Save the value of the element to insert in next Val
- → 4. while nextPos > 0 and the element at nextPos - 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	2
nextVal	20



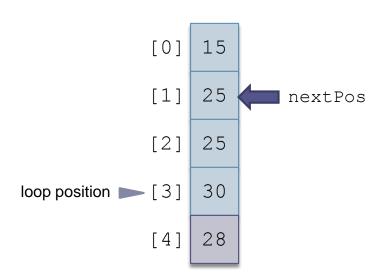
- 1. for each array element from the second
 (nextPos = 1) to the last
- nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	20



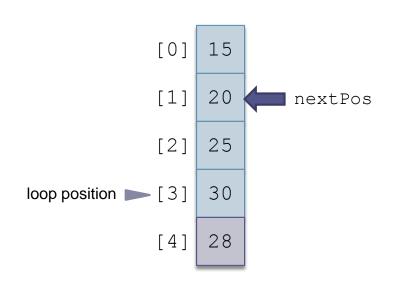
- 1. for each array element from the second
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- 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	20



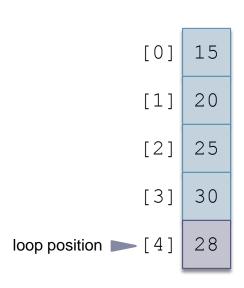
- 1. for each array element from the second
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 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	20



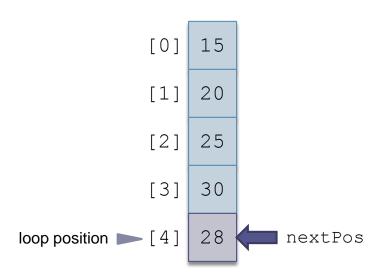
- 1. for each array element from the second
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- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- 6. **Decrement** nextPos by 1
- > 7. Insert nextVal at nextPos

nextPos	1
nextVal	20



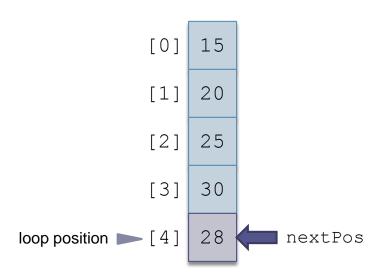
- 1. for each array element from the second
 (nextPos = 1) to the last
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 - 3. Save the value of the element to insert in next Val
 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	1
nextVal	20



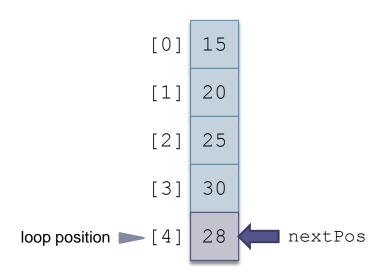
- 1. for each array element from the second
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 - 4. while nextPos > 0 and the element
 at nextPos 1 > nextVal
 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	4
nextVal	28



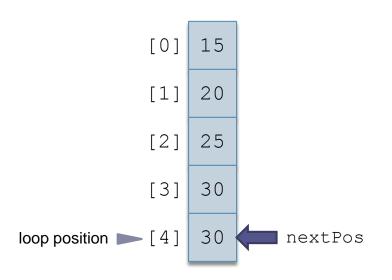
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 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	4
nextVal	28



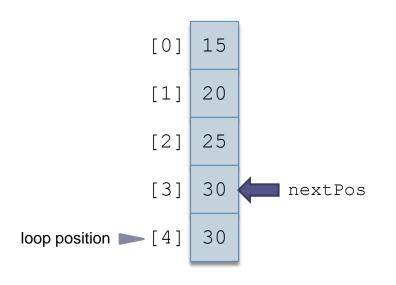
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- ▶ 4. while nextPos > 0 and the element
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 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	4
nextVal	28



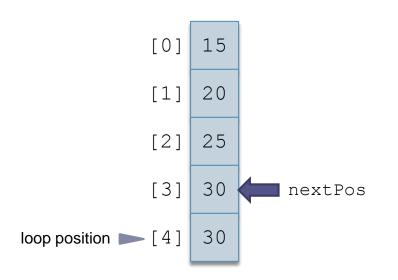
- 1. for each array element from the second
 (nextPos = 1) to the last
- 2. nextPos is the position of the element to insert
- Save the value of the element to insert in nextVal
- 4. while nextPos > 0 and the element at nextPos 1 > nextVal
- Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	28



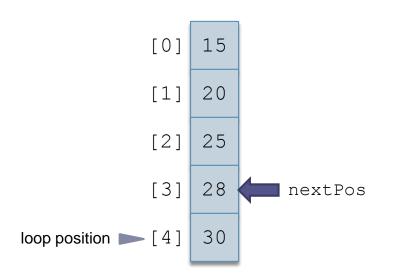
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- 5. Shift the element at nextPos 1 to position nextPos
- ▶ 6. Decrement nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	28



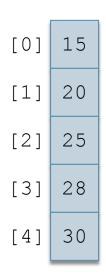
- 1. for each array element from the second
 (nextPos = 1) to the last
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 - 5. Shift the element at nextPos 1 to position nextPos
 - 6. **Decrement** nextPos by 1
 - 7. Insert nextVal at nextPos

nextPos	3
nextVal	28



- 1. for each array element from the second
 (nextPos = 1) to the last
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- 6. **Decrement** nextPos by 1
- > 7. Insert nextVal at nextPos

nextPos	3
nextVal	28



- 1. for each array element from the second
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 at nextPos 1 > nextVal
- 5. Shift the element at nextPos 1 to position nextPos
- 6. **Decrement** nextPos by 1
- 7. Insert nextVal at nextPos

Analysis of Insertion Sort

- \square The insertion step is performed n-1 times
- In the worst case, all elements in the sorted subarray are compared to nextVal for each insertion
- The maximum number of comparisons will then be:

$$1 + 2 + 3 + ... + (n-2) + (n-1) = n(n-1)/2$$

 And... we may do another O(n²) of assignments because of shifting



Code for Insertion Sort

Listing 8.3 (InsertionSort.java, page
434)

Comparison of Quadratic Sorts

Comparison of Quadratic Sorts (and a close look at the book...)

	Number of	f Comparisons	Number o	f assignments
	Best	Worst	Best	Worst
Selection sort	$O(n^2)$	$O(n^2)$	0	O(n)
Bubble sort	O(n)	$O(n^2)$	0	$O(n^2)$
Insertion sort	O(n)	$O(n^2)$	O(n)	$O(n^2)$
Sort by counting	$O(n^2)$	$O(n^2)$	0	0

Compare with the table in the book:

	Number of	f Comparisons	Number of Exchanges			
	Best	Worst	Best	Worst		
Selection sort	$O(n^2)$	$O(n^2)$	O(n)	O(n)		
Bubble sort O(n)		$O(n^2)$	O(1)	$O(n^2)$		
Insertion sort $O(n)$		$O(n^2)$	O(n)	$O(n^2)$		

Comparison of Quadratic Sorts (cont.)

Comparison of growth rates

n	m²	n log n
8	64	24
16	256	64
32	1,024	160
64	4,096	384
128	16,384	896
256	65,536	2,048
512	262,144	4,608

Comparison of Quadratic Sorts (cont.)

- Insertion sort
 - gives the best performance for most arrays
 - takes advantage of any partial sorting in the array and uses less costly shifts
- Bubble sort generally gives the worst performance—unless the array is nearly sorted
 - big-O analysis ignores constants and overhead
- None of the quadratic search algorithms are particularly good for large arrays (n > 1000)
- The best sorting algorithms provide n log n average case performance

Comparison of Quadratic Sorts (cont.)

- All quadratic sorts require storage for the array being sorted
- However, the array is sorted in place
- While there are also storage requirements for variables, for large n, the size of the array dominates and extra space usage is O(1)

Comparisons versus Exchanges

- In Java, an exchange requires a switch of two object references using a third object reference as an intermediary
- A comparison requires an execution of a compareTo method
- The cost of a comparison depends on its complexity, but is generally more costly than an exchange
- For some other languages, an exchange may involve physically moving information rather than swapping object references. In these cases, an exchange may be more costly than a comparison

Shell Sort: A Better Insertion Sort

Donald L. Shell (1924-2015)



Graduated with the Ph.D. degree in Mathematics from the University of Cincinnati in 1959 and immediately published the Shell Sort algorithm "A High-Speed Sorting Procedure" in the Communications of the ACM, Vol. 2, No. 7 [1959]

Has also made major breakthroughs in mathematics.

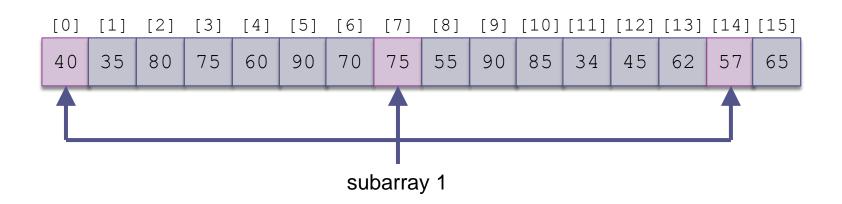
Served as a manager in GE and, for several years, owned a robotics company.

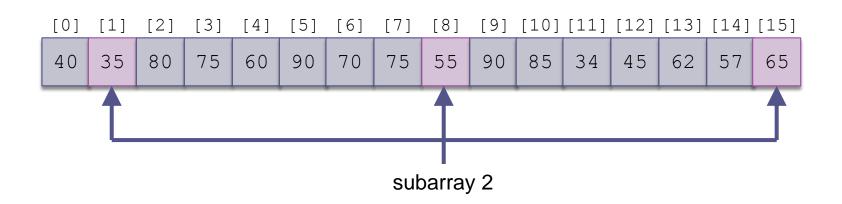
Shell Sort: A Better Insertion Sort

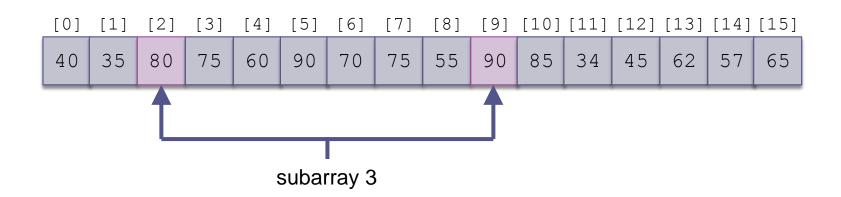
- □ A Shell sort is a type of insertion sort, but with $O(n^{3/2})$ performance
- A Shell sort can be thought of as a divide-andconquer approach to insertion sort
- Instead of sorting the entire array, Shell sort sorts many smaller subarrays using insertion sort before sorting the entire array

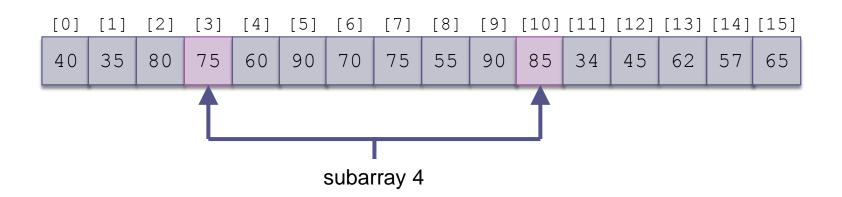
Trace of Shell Sort

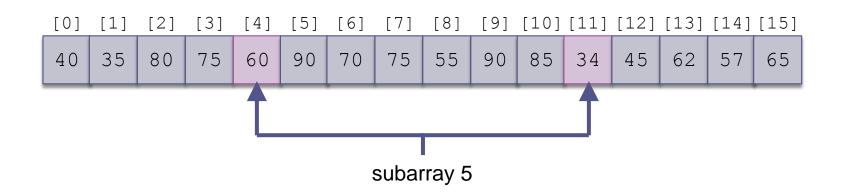
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40	35	80	75	60	90	70	75	55	90	85	34	45	62	57	65

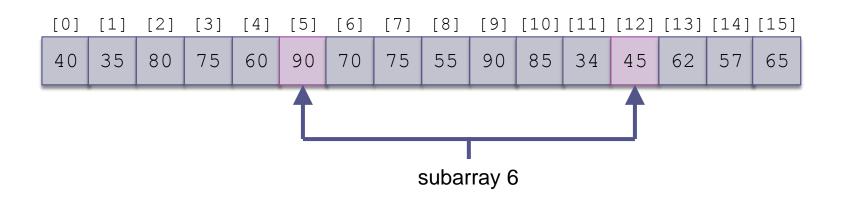


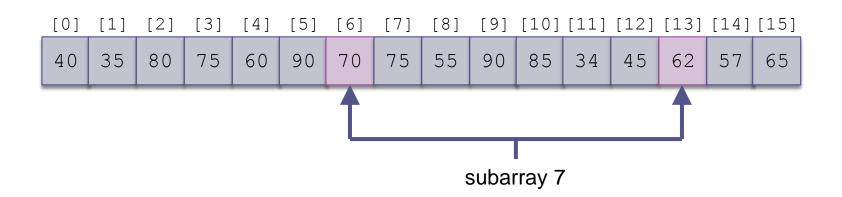


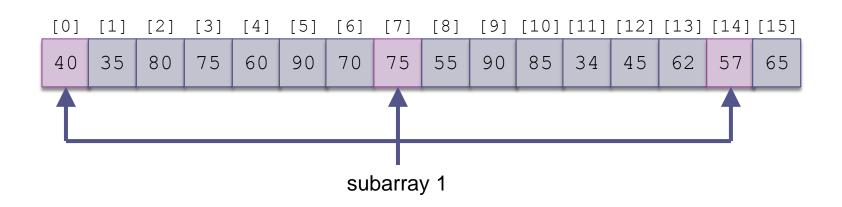


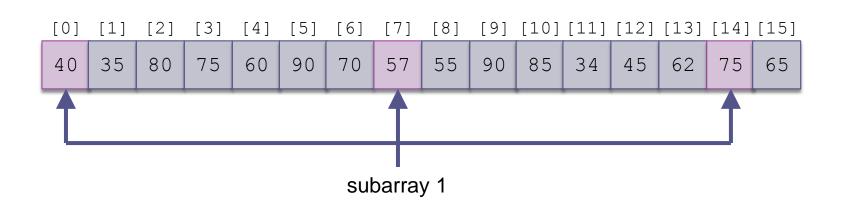


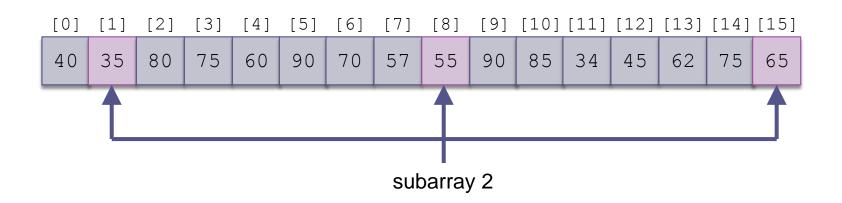


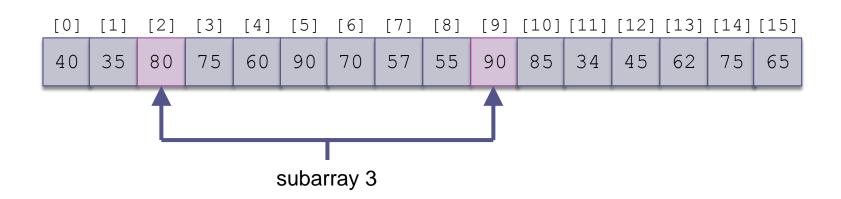


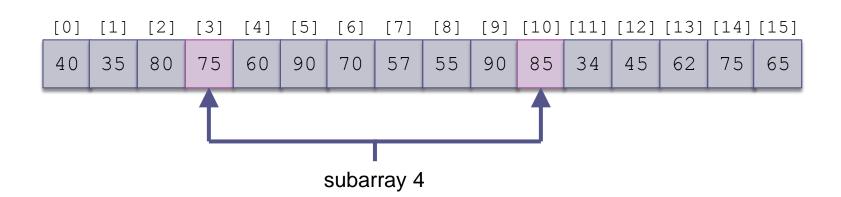


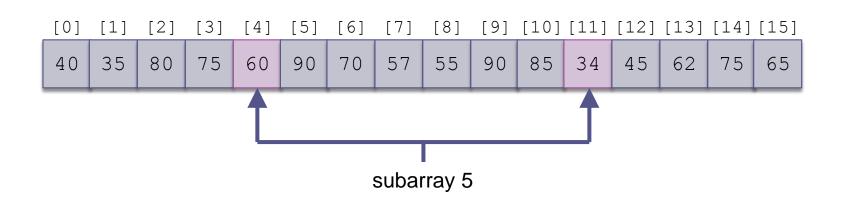


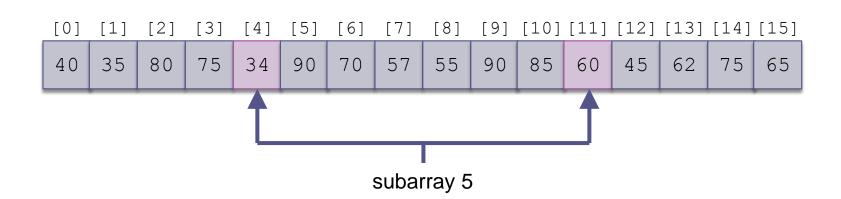


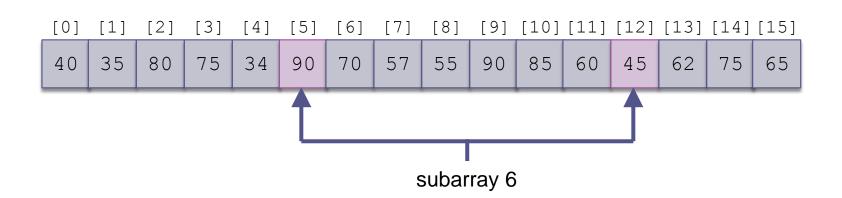


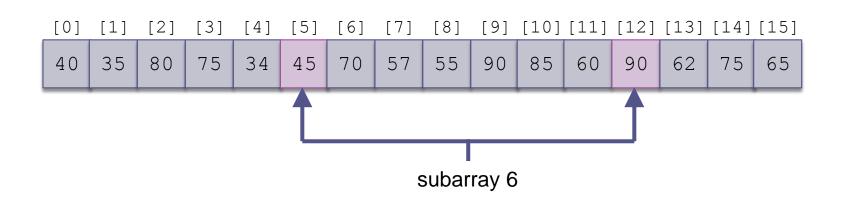


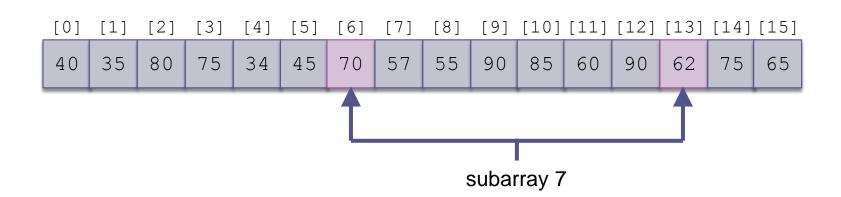


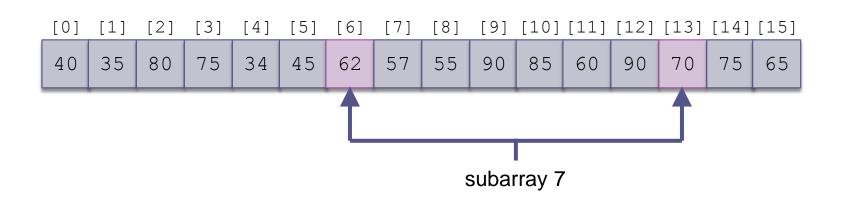












gap value 7

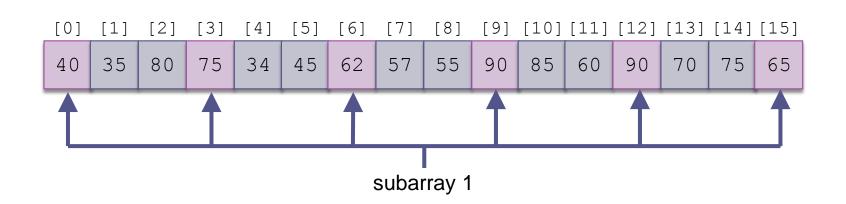
Sort on smaller gap value next

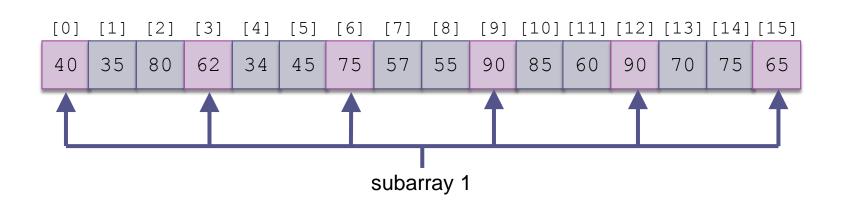
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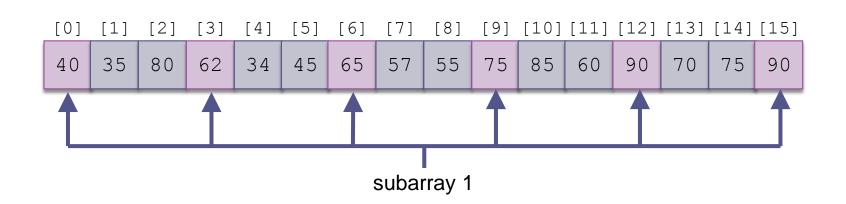
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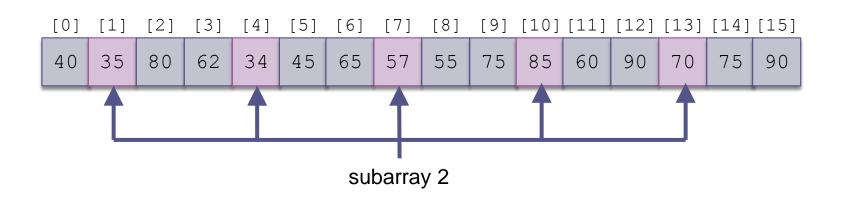
Sort on smaller gap value

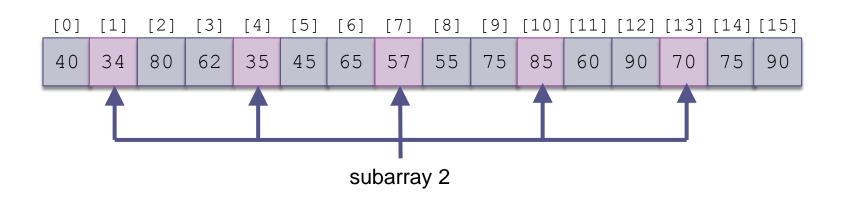
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40	35	80	75	34	45	62	57	55	90	85	60	90	70	75	65

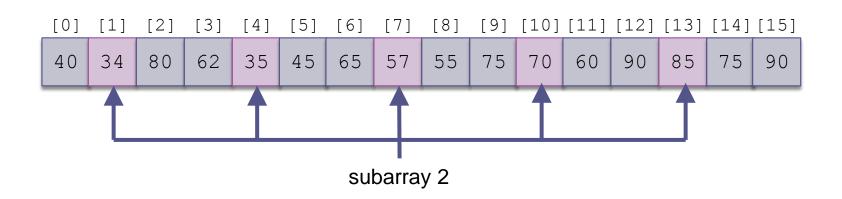


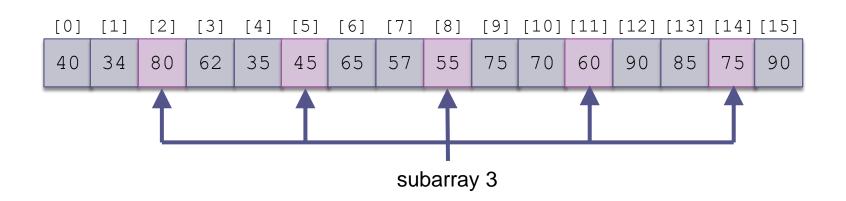


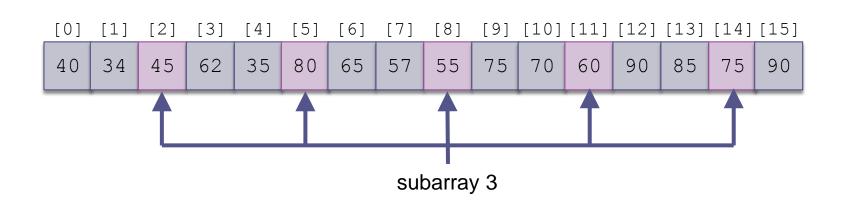


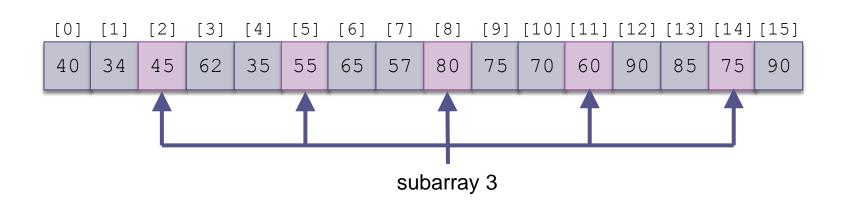


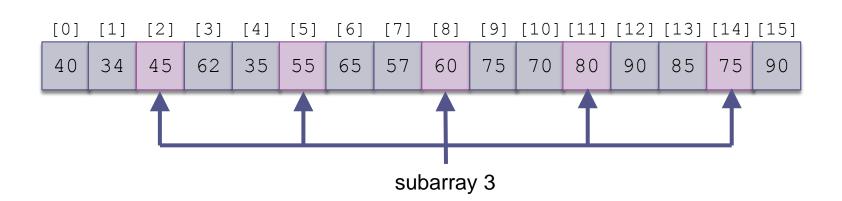


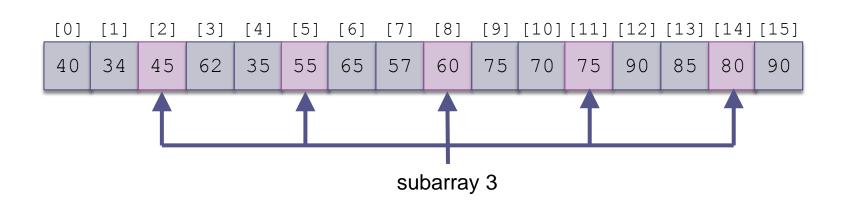












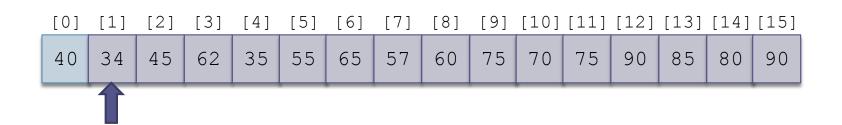
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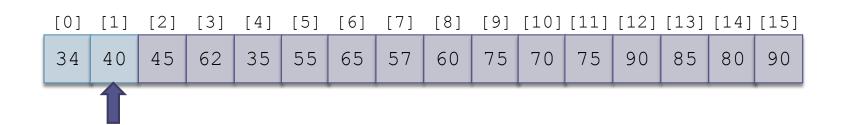
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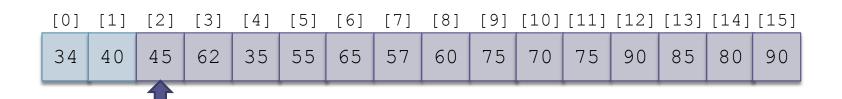
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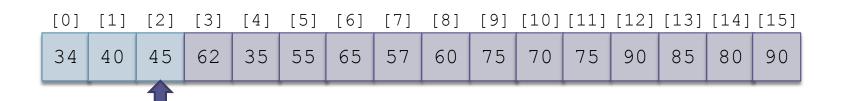
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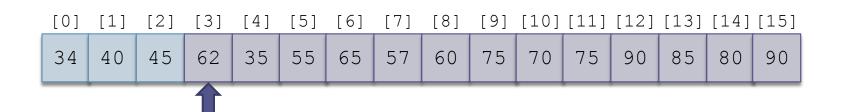
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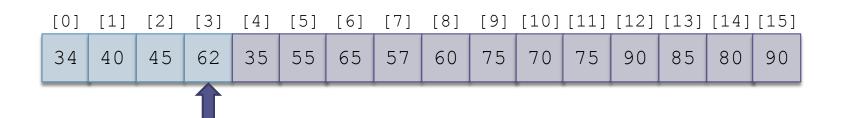
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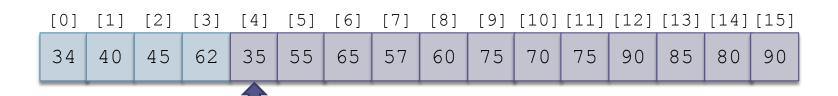
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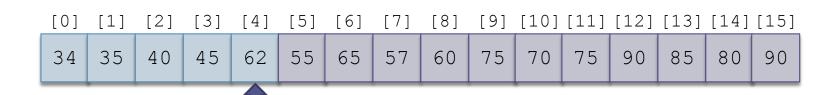
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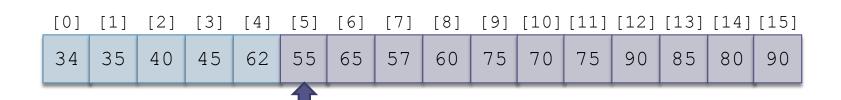
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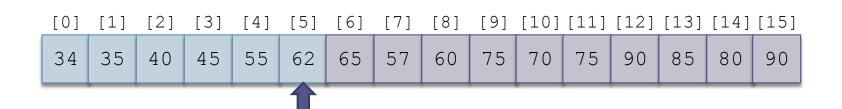
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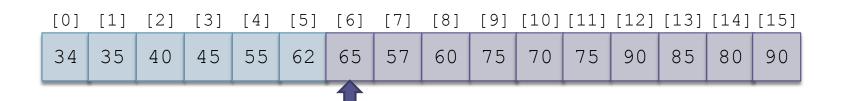
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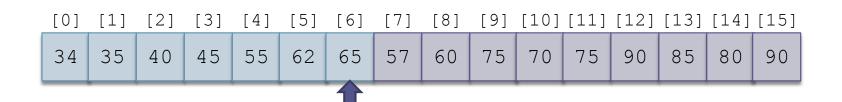
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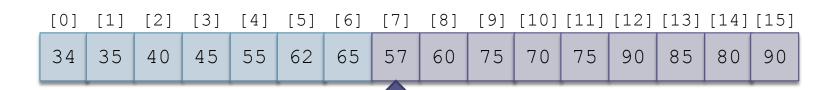
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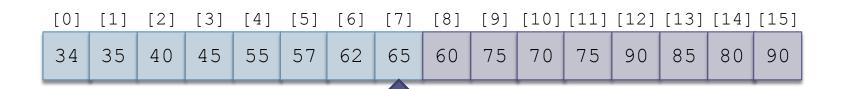
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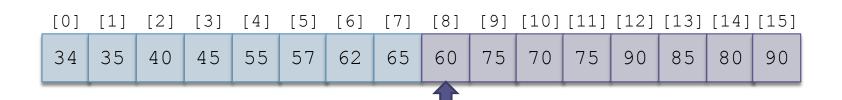
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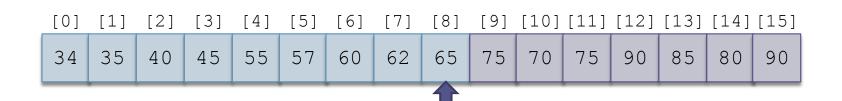
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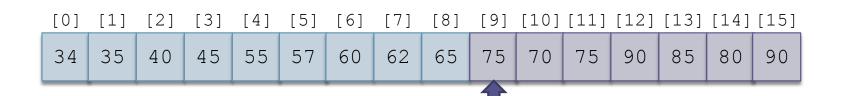
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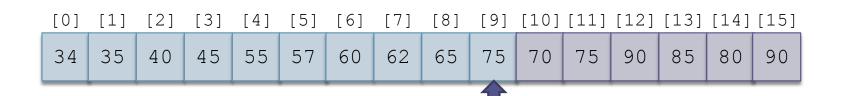
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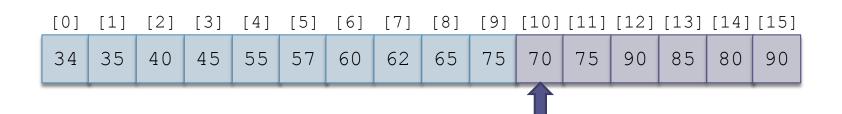
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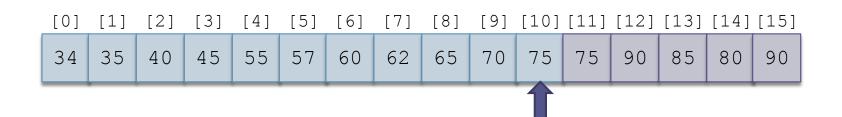
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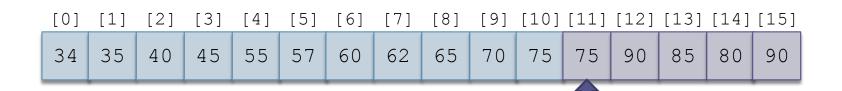
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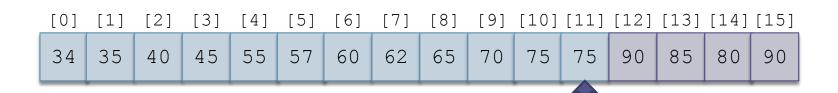
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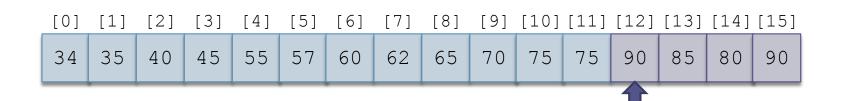
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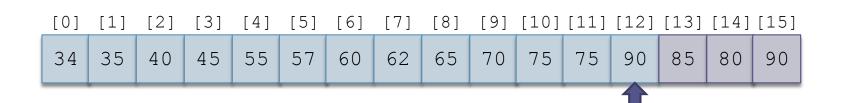
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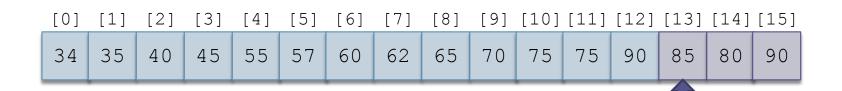
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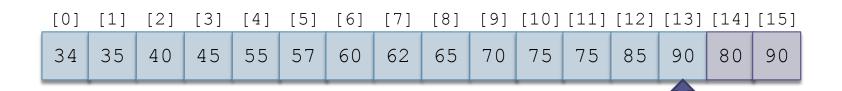
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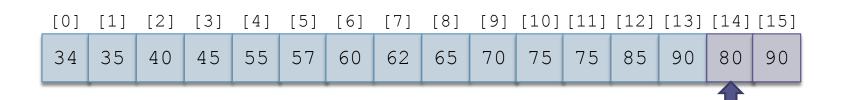
gap value 1



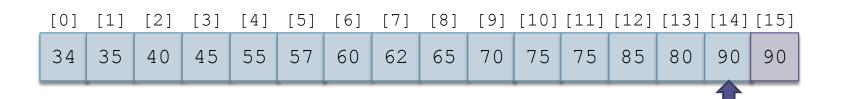
gap value 1



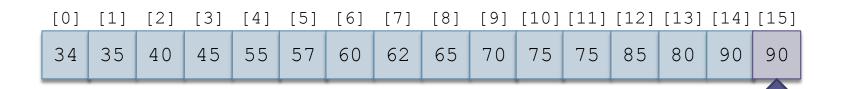
gap value 1



gap value 1



gap value 1



gap value 1

[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]	[12]	[13]	[14]	[15]
34	35	40	45	55	57	60	62	65	70	75	75	85	80	90	90

Shell Sort Algorithm

Shell Sort Algorithm

- Set the initial value of gap to n / 2
- 2. while gap > 0
- 3. for each array element from position gap to the last element
- 4. Insert this element where it belongs in its subarray.
- 5. if gap is 2, set it to 1
- 6. else gap = gap / 2.2. // chosen by experimentation

Shell Sort Algorithm (cont.)

Refinement of Step 4, the Insertion Step

- 4.1 nextPos is the position of the element to insert
- 4.2 Save the value of the element to insert in nextVal
- 4.3 while nextPos > gap and the element at nextPos gap > nextVal
- 4.4 Shift the element at nextPos gap to position nextPos
- 4.5 Decrement nextPos by gap
- 4.6 Insert nextVal at nextPos

Analysis of Shell Sort

- Because the behavior of insertion sort is closer to O(n) than O(n²) when an array is nearly sorted, presorting speeds up later sorting
- This is critical when sorting large arrays

Analysis of Shell Sort (cont.)

- A general analysis of Shell sort is an open research problem in computer science
- Performance depends on how the decreasing sequence of values for gap is chosen
- □ If successive powers of 2 are used for gap, performance is $O(n^2)$
- If successive values for gap are based on Hibbard's sequence,

```
Gap(k) = 2^k - 1, k = m, m-1, ... 1
```

it can be proven that the performance is $O(n^{3/2})$

Analysis of Shell Sort (cont.)

- Our algorithm selects the initial value of gap as $\frac{n}{2}$ and then divides by 2.2 and truncates the result
- Empirical studies of this approach show that the performance is $O(n^{5/4})$ or even $O(n^{7/6})$, but there is no theoretical basis for this result

Code for Shell Sort

Listing 8.4 (ShellSort.java, pages 439 440)

Testing the Sort Algorithms

Testing the Sort Algorithms

- Use a variety of test cases
 - small and large arrays
 - arrays in random order
 - arrays that are already sorted
 - arrays with duplicate values
- Compare performance on each type of array

Driver to Test Sort Algorithms

Listing 8.11(TestSort.java, page 461)

Comparison of Sort Algorithms

Summary

Using Java Sorting Methods

- The Java API provides a class Arrays with several overloaded sort methods for different array types
- The Collections class provides similar sorting methods for Lists
- Sorting methods for arrays of primitive types are based on the quicksort algorithm
- Sorting methods for arrays of objects and
 Lists are based on the merge sort algorithm
- \square Both algorithms are $O(n \log n)$

Using Java Sorting Methods (cont.)

Method sort in Class Arrays	Behavior
<pre>public static void sort(int[] items)</pre>	Sorts the array items in ascending order.
<pre>public static void sort(int[] items, int fromIndex, int toIndex)</pre>	Sorts array elements items[fromIndex] to items[toIndex] in ascending order.
<pre>public static void sort(Object[] items)</pre>	Sorts the objects in array items in ascending order using their natural ordering (defined by method compareTo). All objects in items must implement the Comparable interface and must be mutually comparable.
<pre>public static void sort(Object[] items, int fromIndex, int toIndex)</pre>	Sorts array elements items[fromIndex] to items[toIndex] in ascending order using their natural ordering (defined by method compareTo). All objects must implement the Comparable interface and must be mutually comparable.
<pre>public static <t> void sort(T[] items, Comparator<? super T> comp)</t></pre>	Sorts the objects in items in ascending order as defined by method comp.compare. All objects in items must be mutually comparable using method comp.compare.
<pre>public static <t> void sort(T[] items, int fromIndex, int toIndex, Comparator<? super T> comp)</t></pre>	Sorts the objects in items[fromIndex] to items[toIndex] in ascending order as defined by method comp.compare. All objects in items must be mutually comparable using method comp.compare.
Method sort in Class Collections	Behavior
<pre>public static <t comparable<t="" extends="">> void sort(List<t> list)</t></t></pre>	Sorts the objects in list in ascending order using their natural ordering (defined by method compareTo). All objects in list must implement the Comparable interface and must be mutually comparable.
<pre>public static <t> void sort (List<t> list, Comparator<? super T> comp)</t></t></pre>	Sorts the objects in list in ascending order as defined by method comp.compare. All objects must be mutually comparable.

Declaring a Generic Method



Declaring a Generic Method

FORM:

methodModifiers < genericParameters> returnType methodName(methodParameters)

EXAMPLE:

MEANING:

To declare a generic method, list the genericParameters inside the symbol pair \Leftrightarrow and between the methodModifiers (e.g., public static) and the return type. The genericParameters can then be used in the specification of the methodParameters.

Sample declarations:

Sample declarations:

Sample declarations:

```
public static <T> void sort(T[] items, Comparator<? super T> comp)
```

The second method parameter means that comp must be an object that implements the Comparator interface for type T or for a superclass of type T

Sample declarations:

```
public static <T> void sort(T[] items, Comparator<? super T> comp)
```

For example, we can define a class that implements

Comparator<Number> and use it to sort an array of Integer objects or an array of Double objects

Sample declarations:

Sample declarations:

```
public static <T extends Comparable<T>> void sort(List<T> list)
```

The method parameter list (the object being sorted) is of type List<T>

Example 8.2

```
public class Person implements Comparable<Person> {
     // Data Fields
     /* The last name */
     private String lastName;
     /* The first name */
     private String firstName;
     /* Birthday represented by an integer from 1 to 366 */
     private int birthDay;
     // Methods
     /** Compares two Person objects based on names. The result is based on the last names if they are
         different; otherwise, it is based on the first names.
         @param obj The other Person
         @return A negative integer if this person's name
                 precedes the other person's name;
                 0 if the names are the same:
                 a positive integer if this person's name follows the other person's name.
     * /
     @Override
     public int compareTo(Person obj) {
         Person other = obj;
         // Compare this Person to other using last names.
         int result = lastName.compareTo(other.lastName);
         // Compare first names if last names are the same.
         if (result == 0)
             return firstName.compareTo(other.firstName);
         else
             return result;
     // Other methods
```

Example 8.3

```
import java.util.Comparator;
public class ComparePerson implements Comparator<Person> {
     /** Compare two Person objects based on birth date.
        @param left The left-hand side of the comparison
        @param right The right-hand side of the comparison
        @return A negative integer if the left person's birthday
                precedes the right person's birthday;
                0 if the birthdays are the same;
                a positive integer if the left person's birthday
                follows the right person's birthday.
     * /
     @Override
     public int compare(Person left, Person right) {
         return left.getBirthDay() - right.getBirthDay();
```