



## LECTURE 23

Using binary search trees

# Assignment

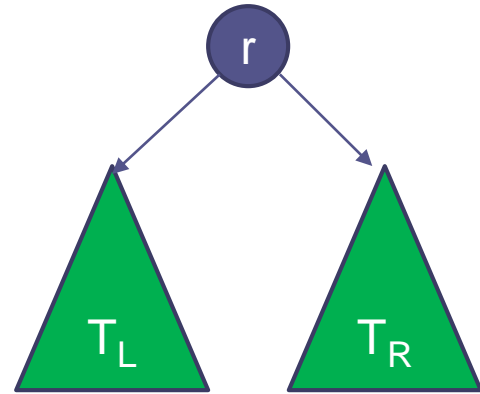
- Read section 6.4
- Do all self-check exercises (use the code developed at recitations)
- Make sure that you understand every detail of the code presented in the book (and these slides)
- Make sure that you can prove that the number of nodes in a *perfect binary tree* of height  $n$  is  $2^n - 1$ .

# Binary Search Tree

- Binary search trees
  - ▣ All elements in the left subtree precede those in the right subtree
- A formal definition:

A set of nodes  $T$  is a binary search tree if either of the following is true:

- $T$  is a single node; or
- The root of  $T$  has two subtrees,  $T_L$  and  $T_R$ , such that  $T_L$  and  $T_R$  are binary search trees and the values in the root node of  $T$  is not less than any value in  $T_L$  and not greater than any value in  $T_R$

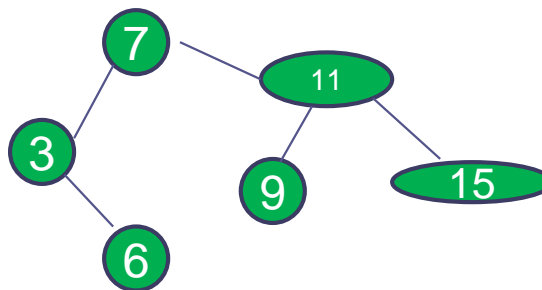


$$\forall x \in T_L, \forall y \in T_R: x \leq r \leq y$$

Note: The definition on p. 300 is inaccurate!

# Using binary search trees

- A binary search tree presents a sorted sequence
- New elements must be inserted (or removed) so that the order is maintained



1) insert 5

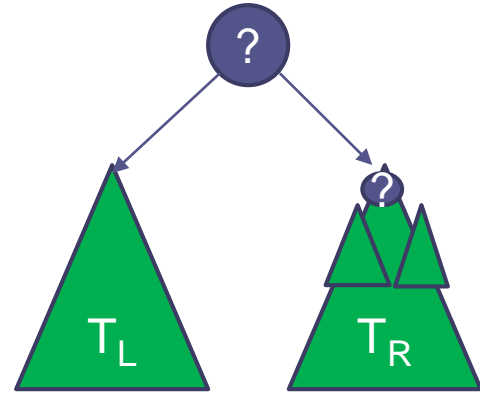
2) remove 7

# Binary Search Tree (BST)

- Searching a **balanced** BST can be done in  $O(\log n)$  time
- The same algorithm used on an unbalanced tree may take  $O(n)$  time (Why?)

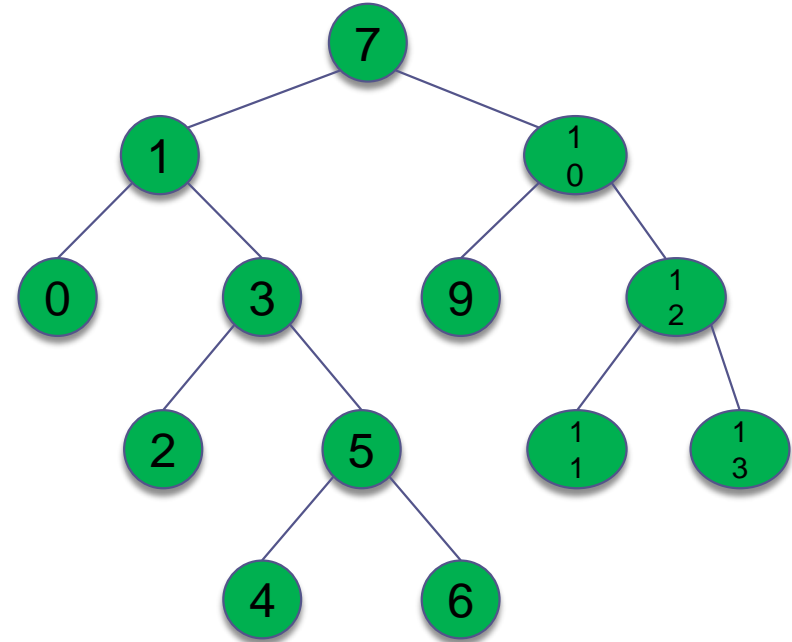
# Recursive Algorithm for Searching a Binary Tree

1. **if** the tree is empty
2.     return null (*target is not found*)
- else if** the target matches the root node's data  
3.     return the data stored at the root node
- else if** the target is less than the root node's data  
4.     return the result of searching the left subtree of the root
- else**  
5.     return the result of searching the right subtree of the root



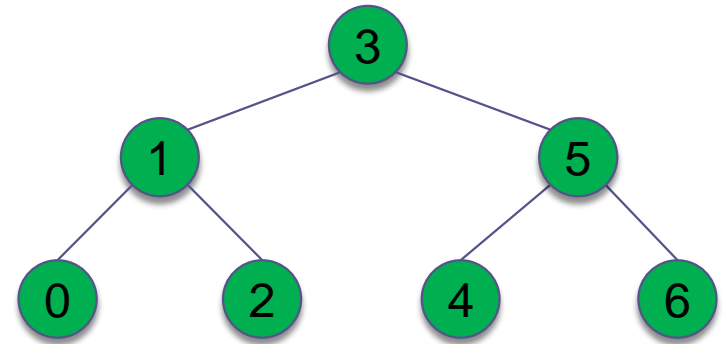
# Full binaryTrees

- A full binary tree is a binary tree where each node nodes has either two children or or no children



# Perfect Binary Trees

- A *perfect binary tree* is a full binary tree of height  $n$  with exactly  $2^n - 1$  nodes
- $1 + 2 + \dots + 2^{n-1} = 2^n - 1$

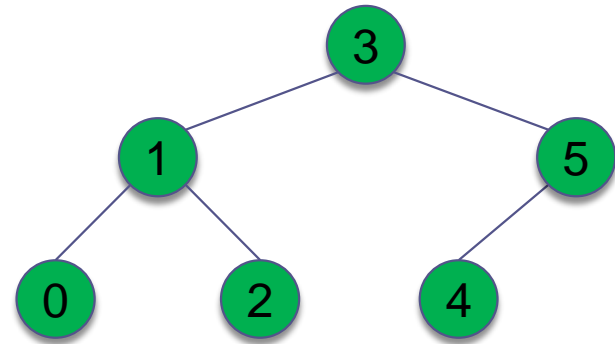


(cf. *Towers of Hanoi*)

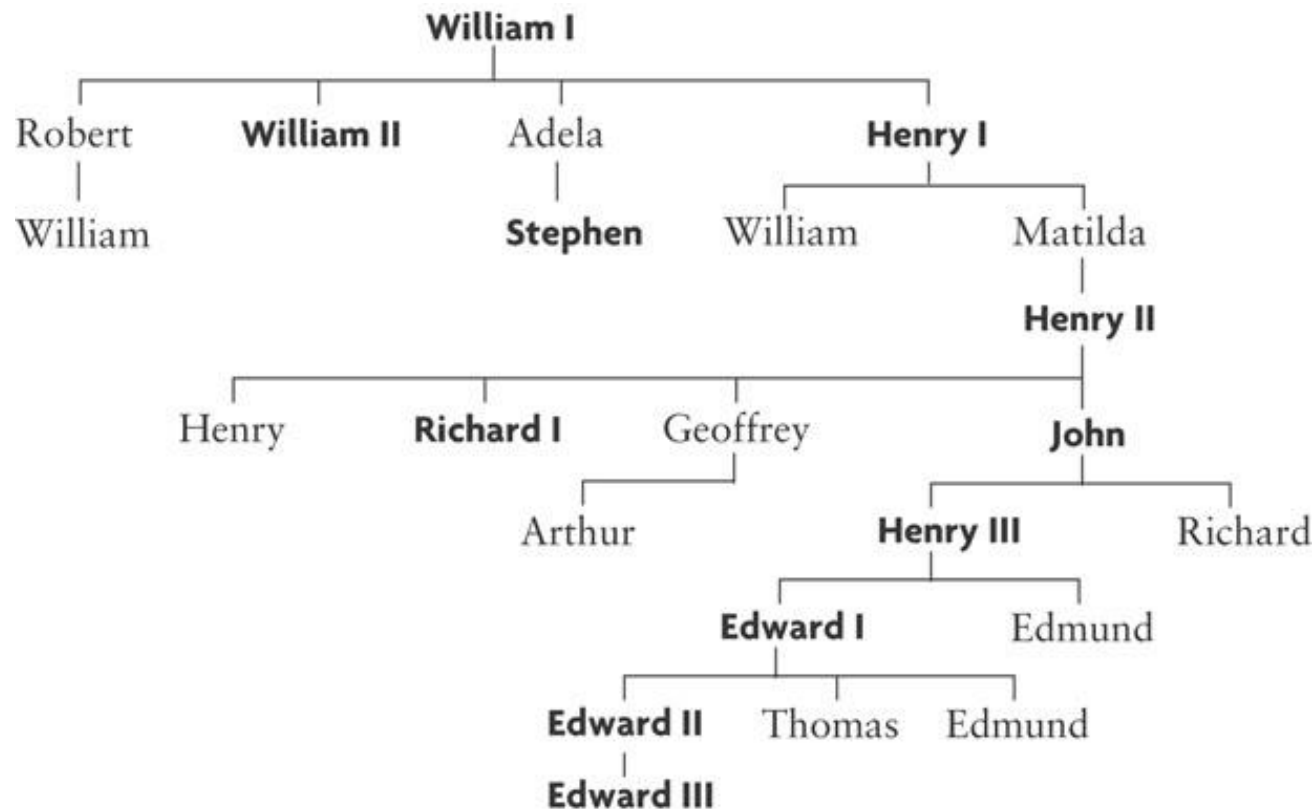


# Complete Binary Trees

- A *complete binary tree* is a perfect binary tree through level  $n - 1$  with some extra leaf nodes at level  $n$  (the tree height), all toward the left

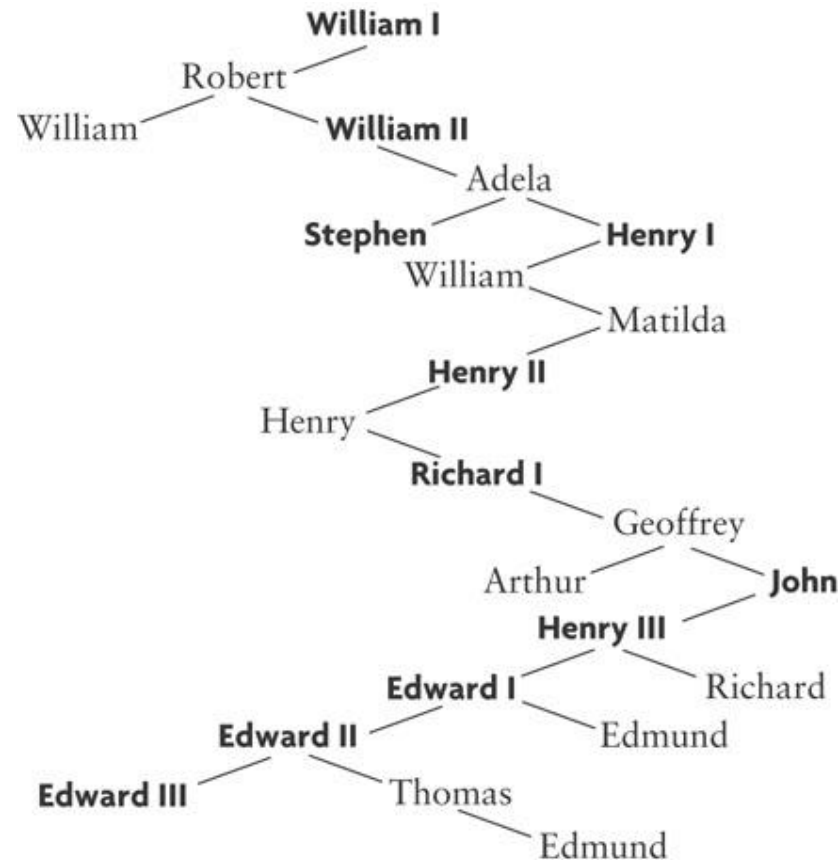


# General Trees



# General Trees

- A general tree can be converted into a binary tree
- The left branch of a node is the oldest child, and each right branch is connected to the next younger sibling (if any)



# Tree Traversals

- How can we list all nodes in a tree?
  - ▣ We can do this by walking through the tree in a prescribed order and visiting the nodes as they are encountered
  - ▣ This process is called *tree traversal*
- Three common kinds of tree traversal are
  - ▣ *Inorder*
  - ▣ *Preorder*
  - ▣ *Postorder*

# Tree Traversals

- *Preorder*: visit root node, traverse  $T_L$ , traverse  $T_R$
- *Inorder*: traverse  $T_L$ , visit root node, traverse  $T_R$
- *Postorder*: traverse  $T_L$ , traverse  $T_R$ , visit root node

## Algorithm for Preorder Traversal

1. if the tree is empty
2.     Return.
- else
3.     Visit the root.
4.     Preorder traverse the left subtree.
5.     Preorder traverse the right subtree.

## Algorithm for Inorder Traversal

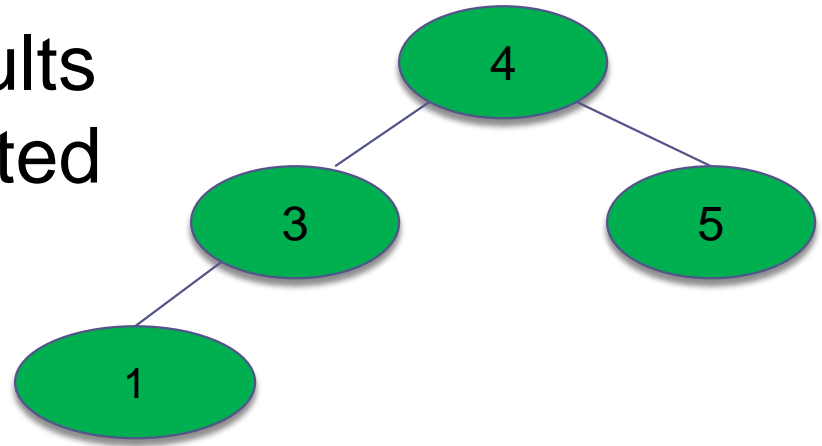
1. if the tree is empty
2.     Return.
- else
3.     Inorder traverse the left subtree.
4.     Visit the root.
5.     Inorder traverse the right subtree.

## Algorithm for Postorder Traversal

1. if the tree is empty
2.     Return.
- else
3.     Postorder traverse the left subtree.
4.     Postorder traverse the right subtree.
5.     Visit the root.

# Traversals of Binary Search Trees and Expression Trees

- An inorder traversal of a binary search tree results in the nodes being visited in sequence by increasing data value



*1, 3, 4, 5*

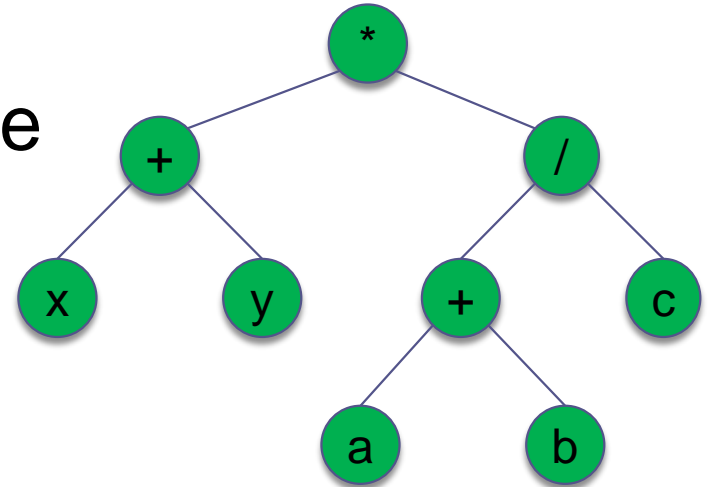
# Traversals of Binary Search Trees and Expression Trees (cont.)

- An inorder traversal of an expression tree results in the sequence

$x + y * a + b / c$

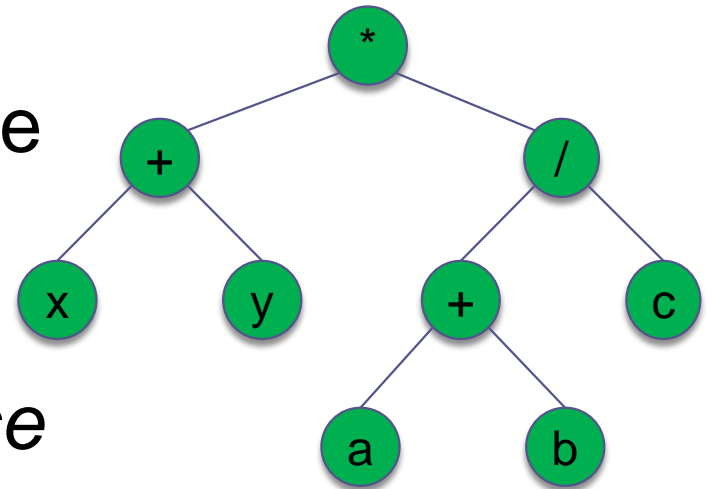
- If we insert parentheses where they belong, we get the infix form:

$(x + y) * ((a + b) / c)$



# Traversals of Binary Search Trees and Expression Trees (cont.)

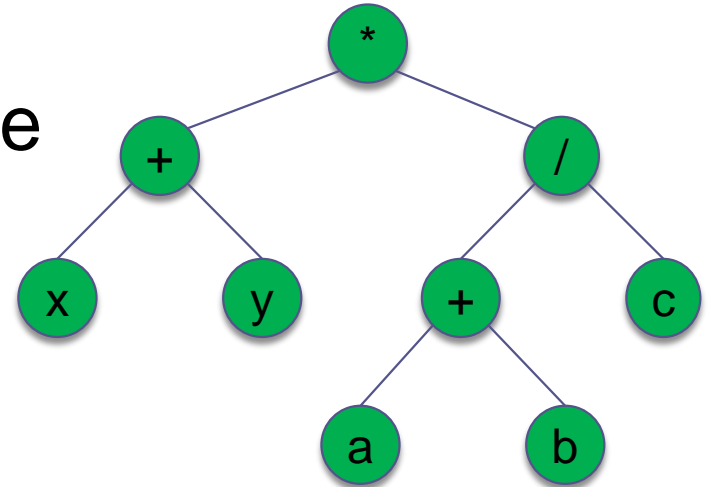
- A postorder traversal of an expression tree results in the sequence  
 $x y + a b + c / *$
- This is the *postfix* or *reverse polish* form of the expression
- Operators follow operands





# Traversals of Binary Search Trees and Expression Trees (cont.)

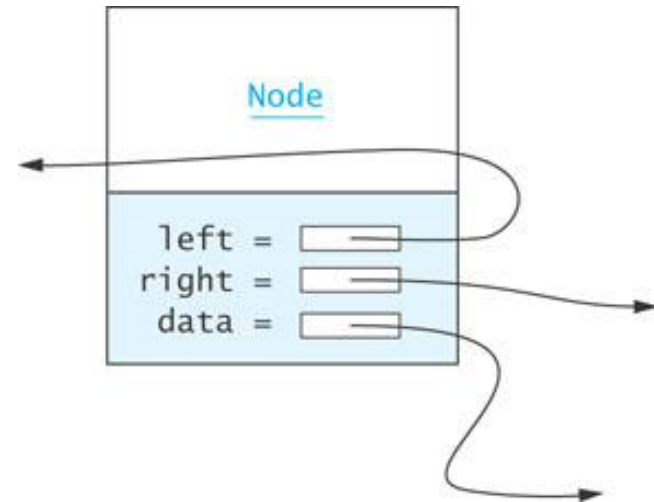
- A preorder traversal of an expression tree results in the sequence  
 $* + x y / + a b c$
- This is the *prefix* or *forward polish* form of the expression
- Operators precede operands



# Implementing a `BinaryTree` Class

# Node<E> Class

- Just as in a linked list, a node consists of 1) **a data part** and 2) **the links to successor nodes**
- The data part is a reference to type E
- A binary tree node must have links to both its left and right subtrees

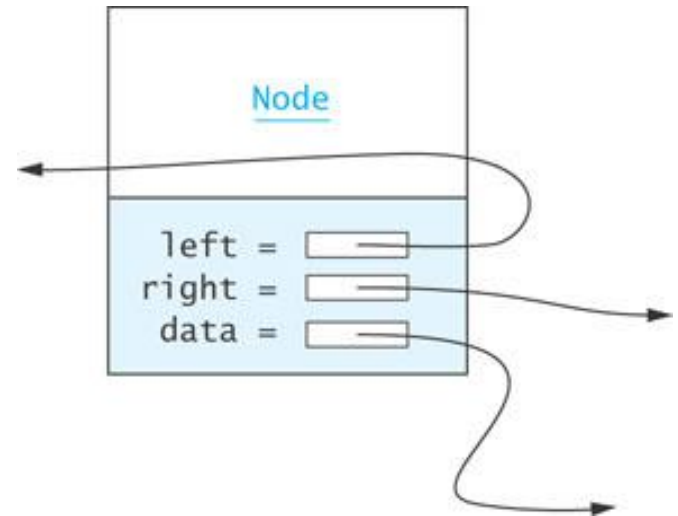


# Node<E> Class (cont.)

```
protected static class Node<E>
    implements Serializable {
    protected E data;
    protected Node<E> left;
    protected Node<E> right;

    public Node(E data) {
        this.data = data;
        left = null;
        right = null;
    }

    public String toString() {
        return data.toString();
    }
}
```



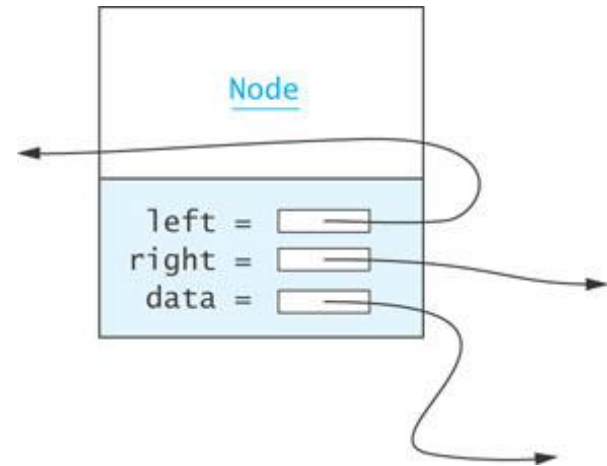
`Node<E>` is declared as an  
inner class within  
`BinaryTree<E>`

# Node<E> Class (cont.)

```
protected static class Node<E>
    implements Serializable {
    protected E data;
    protected Node<E> left;
    protected Node<E> right;

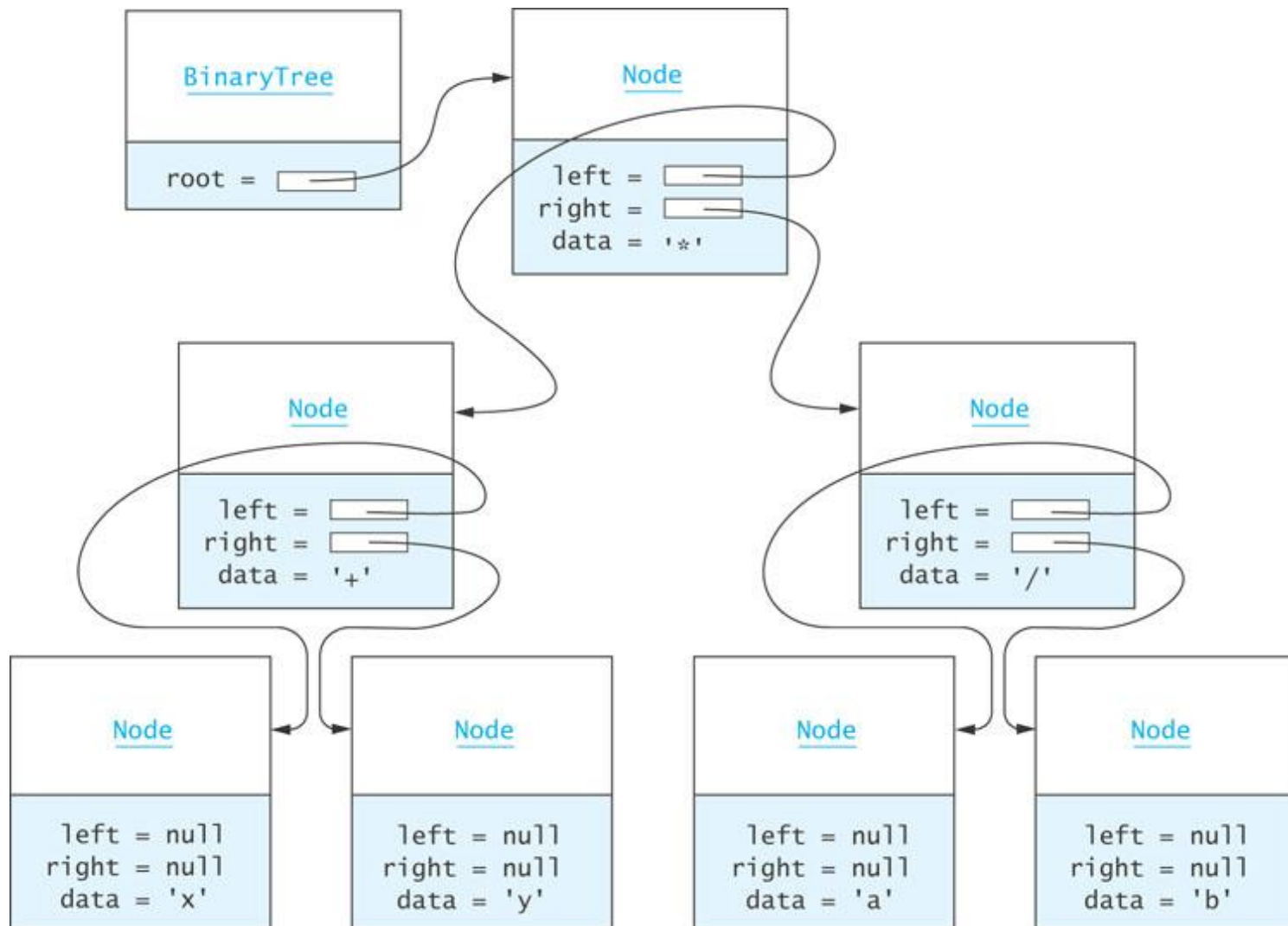
    public Node(E data) {
        this.data = data;
        left = null;
        right = null;
    }

    public String toString() {
        return data.toString();
    }
}
```



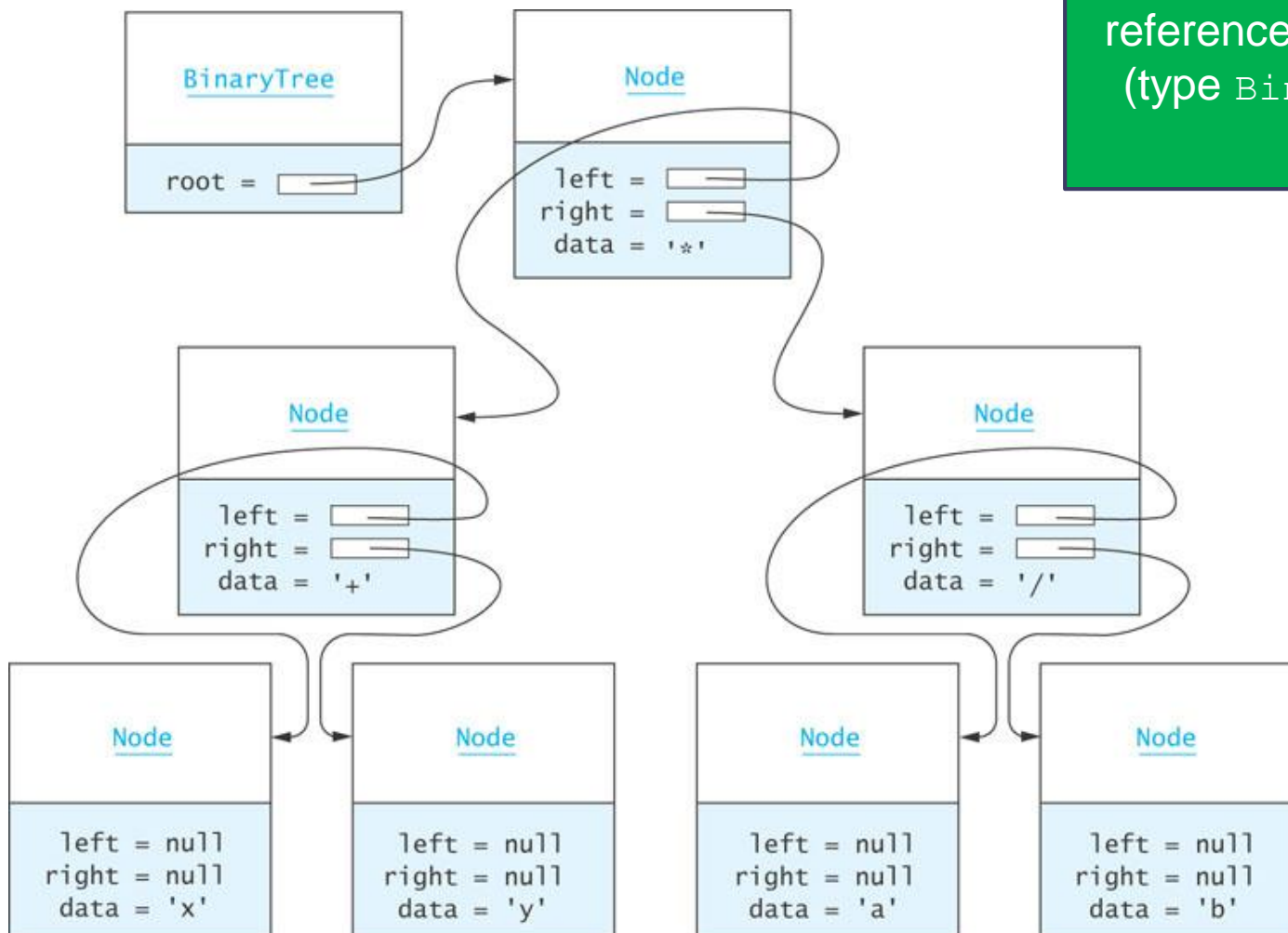
`Node<E>` is declared protected. This way we can use it as a superclass.

# BinaryTree<E> Class (cont.)



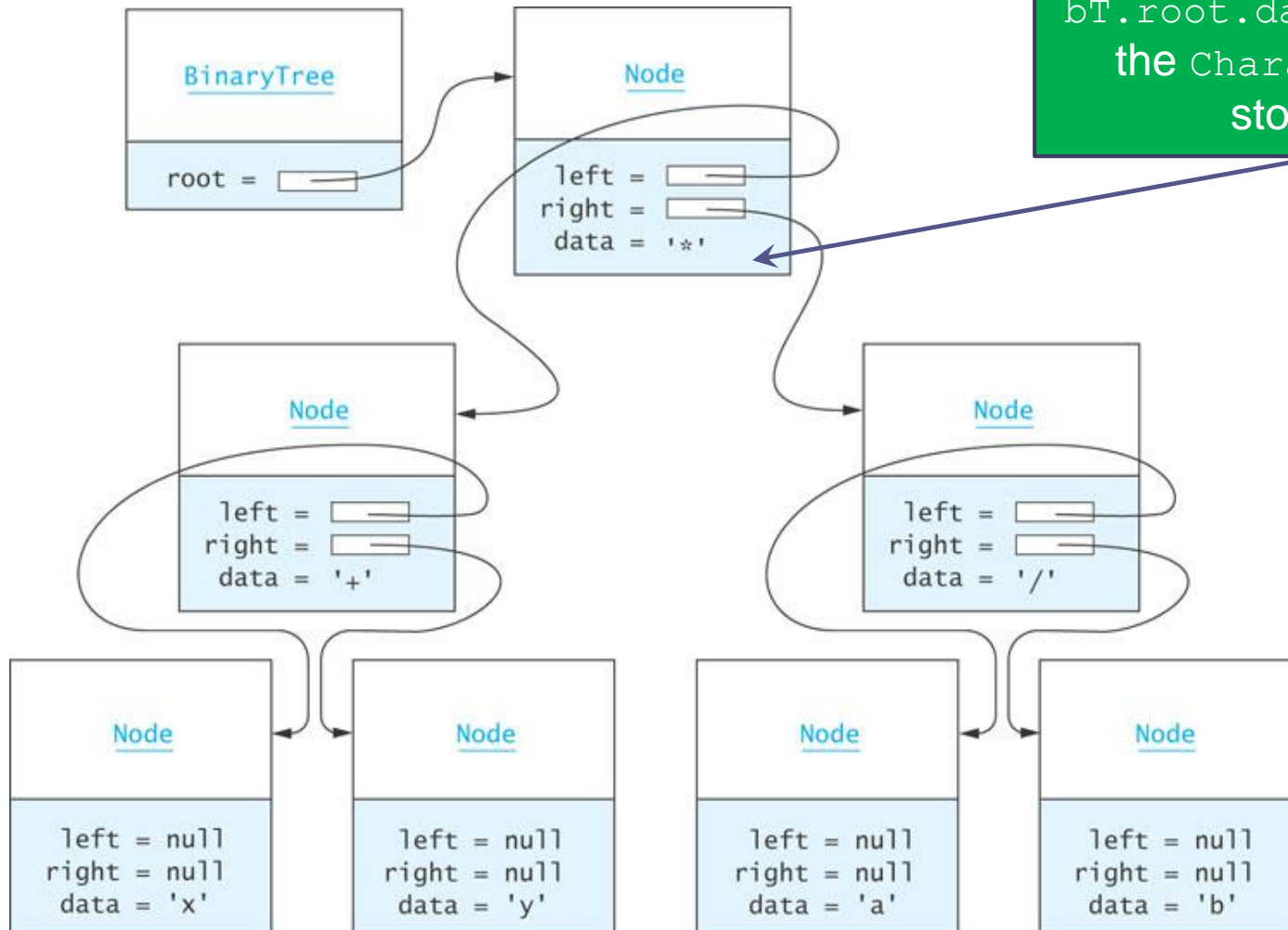
# BinaryTree<E> Class (cont.)

Assuming the tree is referenced by variable `bT` (type `BinaryTree`) then



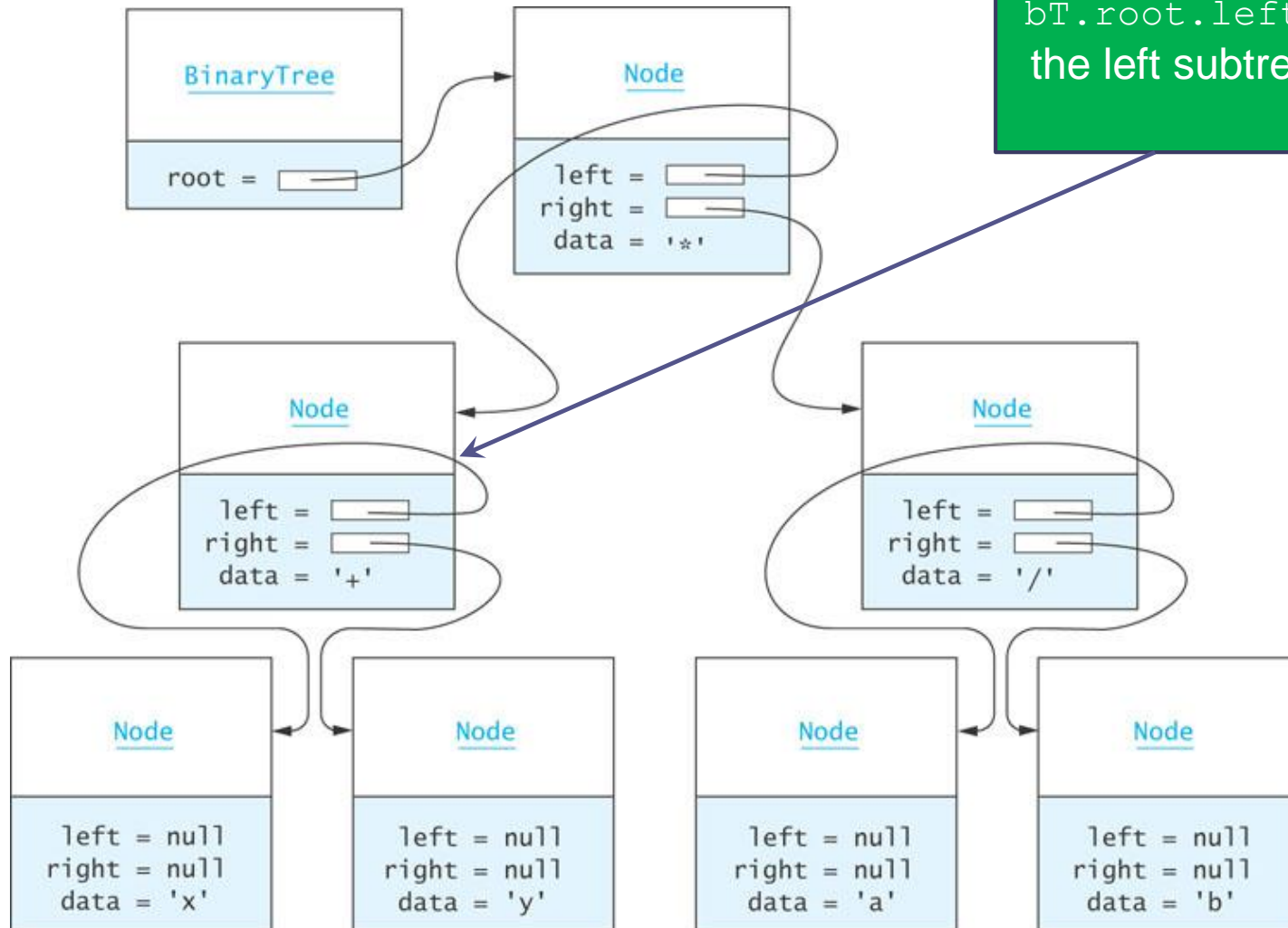
# BinaryTree<E> Class (cont.)

bT.root.data references  
the Character object  
storing '\*'



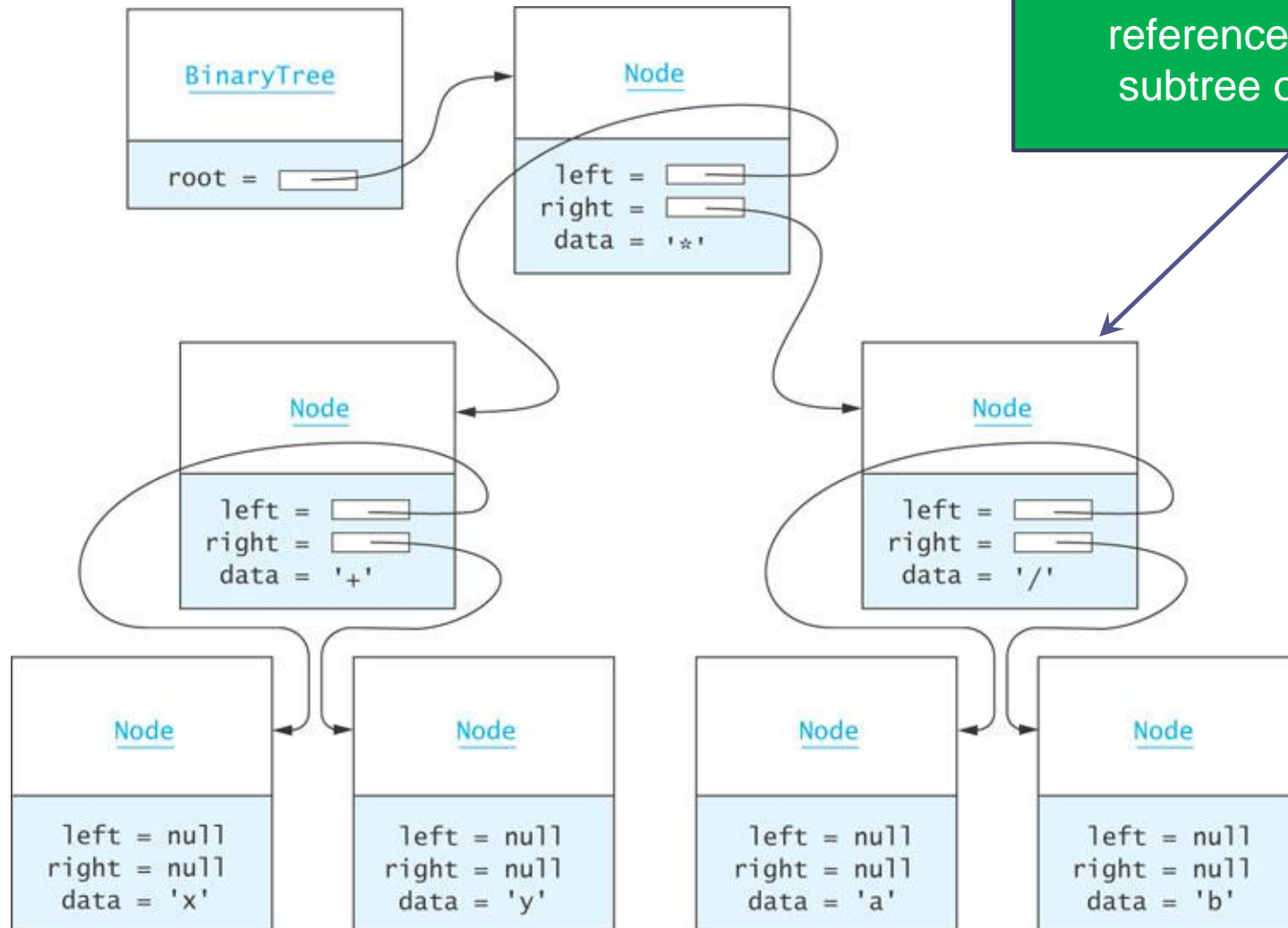


# BinaryTree<E> Class (cont.)



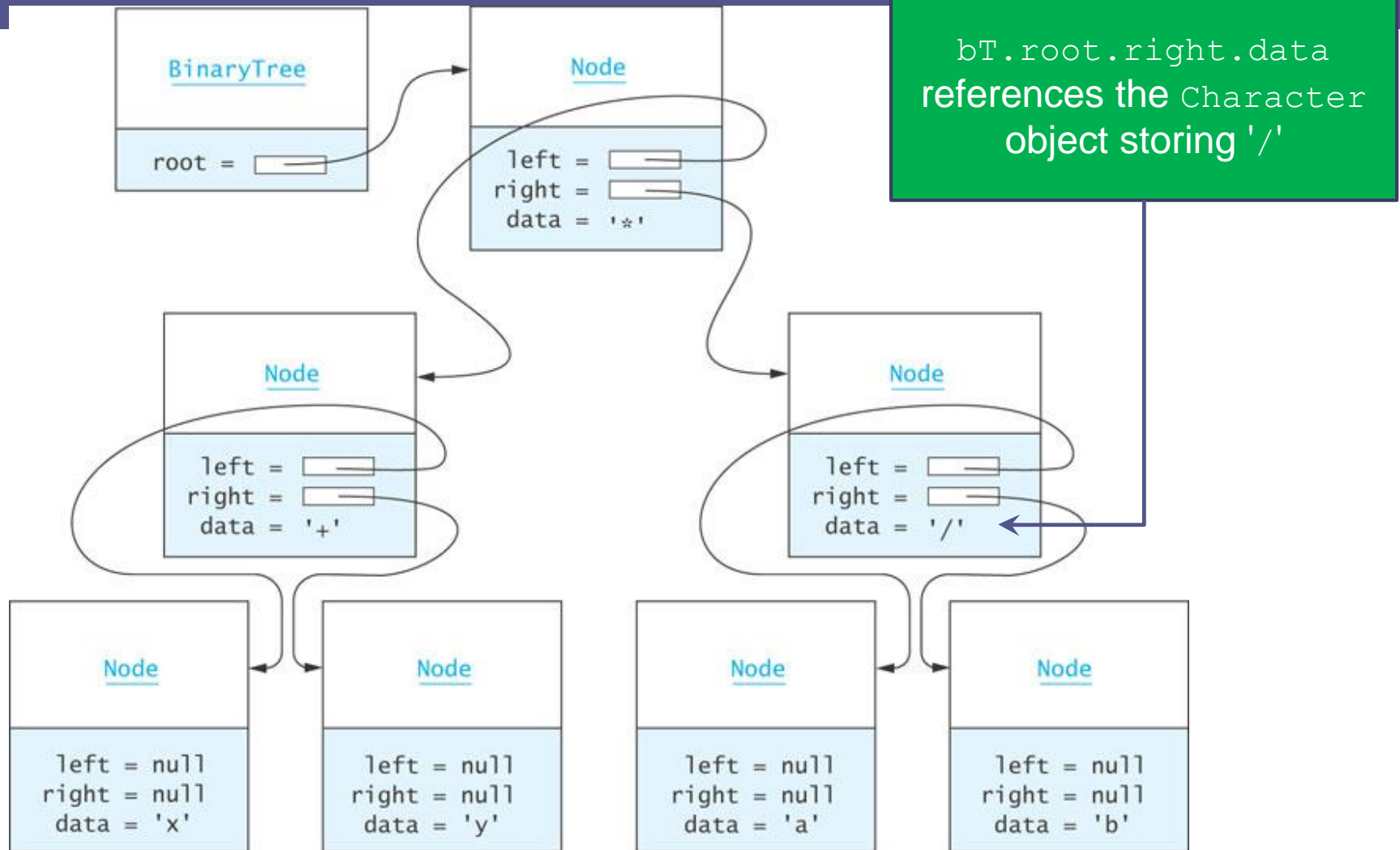
`bT.root.left` references  
the left subtree of the root

# BinaryTree<E> Class (cont.)



`bT.root.right`  
references the right  
subtree of the root

# BinaryTree<E> Class (cont.)



# BinaryTree<E> Class (cont.)

Data Field	Attribute
protected Node<E> root	Reference to the root of the tree.
Constructor	Behavior
public BinaryTree()	Constructs an empty binary tree.
protected BinaryTree(Node<E> root)	Constructs a binary tree with the given node as the root.
public BinaryTree(E data, BinaryTree<E> leftTree, BinaryTree<E> rightTree)	Constructs a binary tree with the given data at the root and the two given subtrees.
Method	Behavior
public BinaryTree<E> getLeftSubtree()	Returns the left subtree.
public BinaryTree<E> getRightSubtree()	Returns the right subtree.
public E getData()	Returns the data in the root.
public boolean isLeaf()	Returns <b>true</b> if this tree is a leaf, <b>false</b> otherwise.
public String toString()	Returns a String representation of the tree.
private void preOrderTraverse(Node<E> node, int depth, StringBuilder sb)	Performs a preorder traversal of the subtree whose root is node. Appends the representation to the StringBuilder. Increments the value of depth (the current tree level).
public static BinaryTree<E> readBinaryTree(Scanner scan)	Constructs a binary tree by reading its data using Scanner scan.

# BinaryTree<E> **Class** (cont.)

## □ Class heading and data field declarations:

```
import java.io.*;

public class BinaryTree<E> implements Serializable {
    // Insert inner class Node<E> here

    protected Node<E> root;

    . . .
}
```

# BinaryTree<E> **Class** (cont.)

- The `Serializable` interface defines no methods
- It provides a marker for classes that can be written to a binary file using the `ObjectOutputStream` and read using the `ObjectInputStream`