# CSCI-UA 480: Computer Networks Assignment 3

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## 1 Using Traceroute

#### 1.1

The traceroute works by continuously sending out probing packets with increasing the Time to Live field (initialized to 1), and listening for ICMP time exceed signal so that it knows which gateway the probing packet ends at. In each iteration, it increases TTL field, sends out 3 probing packets, and reports the gateways (not necessarily the same) and RTTs.

#### 1.2

The asterisk means we do not receive response (timeout with default value 5 seconds) for the probing packet we send out.

#### 1.3

My laptop had trouble with traceroute 139.130.4.5, so I connected to NYU VPN.

```
|Kai@MacBook-Pro-9 ~ % traceroute www.cs.nyu.edu
traceroute to cs.nyu.edu (128.122.49.30), 64 hops max, 52 byte packets

1 10.2.204.254 (10.2.204.254) 6.607 ms 4.434 ms 7.101 ms
2 10.2.05. (10.2.0.5) 3.672 ms 2.684 ms 3.051 ms
3 192.168.1.1 (192.168.1.1) 6.860 ms 2.529 ms 2.854 ms
4 1.140.35.58.broad.xw.sh.dynamic.163data.com.cn (58.35.140.1) 33.019 ms 5.771 ms 6.007 ms
5 101.95.91.101 (101.95.91.101) 10.931 ms
101.95.91.229 (101.95.91.229) 7.154 ms
61.152.50.5 (61.152.50.5) 5.660 ms
6 61.152.24.42 (61.152.24.42) 13.789 ms 11.588 ms
202.101.63.242 (202.101.63.242) 11.008 ms
7 ** 202.97.50.158 (202.97.50.158) 27.585 ms
8 202.97.90.53 (202.97.90.53) 20.847 ms
x 202.97.33.134 (202.97.33.134) 10.183 ms
202.97.33.134 (202.97.33.134) 11.260 ms
9 202.97.6.2 (202.97.6.2) 170.576 ms *
202.97.58.194 (202.97.58.194) 177.779 ms
10 202.97.50.58 (202.97.50.58) 174.788 ms 163.285 ms
202.97.50.74 (202.97.50.74) 147.030 ms
11 218.30.54.101 (218.30.54.101) 161.434 ms 650.056 ms 147.255 ms
289.149.128.166 (89.149.128.166) 199.267 ms 218.486 ms 264.514 ms
13 ip4.gtt.net (209.120.137.218) 206.456 ms 284.208 ms *
14 * **
15 128.122.254.108 (128.122.254.108) 371.668 ms
128.122.254.110 (128.122.254.110) 209.239 ms 206.372 ms
10 nyufw-outside-ngfw-vl3080.net.nyu.edu (128.122.254.116) 204.803 ms 322.676 ms 215.446 ms
17 **
18 128.122.11 (128.122.11) 206.958 ms 234.658 ms 315.289 ms
```

Figure 1: traceroute www.cs.nyu.edu

```
| KaigMacBook-Pro-9 ~ % traceroute gaia.cs.umass.edu | 128.119.245.120, 64 hops max, 52 byte packets | 1 lo.213.32.209 (10.213.19.213.19.102.11 lo.213.32.209) 20.990 ms 12.030 ms 5.405 ms | 1.0213.33.137 (10.213.33.137) 10.790 ms 8.179 ms 6.407 ms | 1 lo.22.24.245 (10.204.254) 3.484 ms 3.034 ms 5.316 ms | 2 lo.213.33.135 (10.213.33.135) 10.10.2213.7209) 9.550 ms 7.290 ms 7.018 ms | 5 lo.117.219.101 (10.117.219.101) 195.363 ms 199.001 ms 204.265 ms | 3 lo.213.33.150 (10.213.33.137) 10.10.117.219.101 (10.117.219.101) 195.363 ms 199.001 ms 204.265 ms | 3 lo.213.33.150 (10.213.33.135) 10.10.117.219.101 (10.117.219.101) 195.363 ms 199.001 ms 204.265 ms | 3 lo.213.33.150 (10.213.33.150) 10.117.219.101 (10.117.219.101) 195.363 ms 195.000 ms 195.000 ms | 19
```

Figure 2: traceroute gaia.cs.umass.edu and 139.130.4.5

#### 1.4

For www.cs.nyu.edu, we have 101.95.91.101 (Shanghai) -61.152.24.42 (Shanghai) -202.97.6.2 (Beijing) -89.149.128.166 (NYC) -128.122.1.1 (NYU)

For gaia.cs.umass.edu, we have 209.120.137.217 (NYC) -154.54.3.125 (NYC) -38.104.218.14 (Worcester, Massachusetts) -128.119.0.8 (UMass, Amherst) -128.119.3.32 (UMass, Amherst)

For 139.130.4.5, we have 101.95.91.225 (Shanghai) -202.97.50.146 (Liaoning, China) -213.248.92.129 (Illinois, Chicago) -213.248.92.107 (Illinois, Chicago) -203.50.6.96 (Sydney, Australia)

## 2 Interdomain vs. intradomain routing

#### 2.1

Interdomain routing is how packets can be transmitted from a router in one domain to a router in another domain.

#### 2.2

Intradomain routing is how packets can be trasmitted from a router to another router within the same domain.

#### 2.3

An interdomain algorithm aims to discover paths across different domains, and propagate that information to routers within the domain.

#### 2.4

An intradomain algorithm aims to compute the "shortest path" between any two routers within the same domain.

#### 2.5

Because they focus on different aspect of routing. Intradomain algorithm needs to handle routing among all routers within a domain (local), while interdomain algorithm needs to handle routing between different domains (global). They basically form a hierarchy.

### 2.6

The most important factor is the business need for the ISP. Local preference can be driven by business needs (such as the Kenyan ISP example discussed in lecture), and minimizing the length of path is not necessarily correlated with improved performance since it does not capture information about internal topology of each domain.

## 3 Transit and peer relationships

### 3.1

A (provider domain) provides Internet service for B (customer domain), or the other way around.

#### 3.2

A and B agree to carry each other's traffic.

#### 3.3

The South African ISP "does not advertise path to its servers to Kenyan ISP". [1]

# References

[1] Sivaraman Anirudh. Lecture 9: Interdomain routing. 2020. URL: https://cs.nyu.edu/~anirudh/CSCI-UA.0480-062/lectures/lec9.pdf.