- Bufferless case
- 1. W > C · RTTmin + B
- 2. CRTTmin

Window increases by in on every ACK. It gradually increases towards BDP.

When it surpasses BDP by a small amount, packet loss will be detected and multiplicative decrease gets triggered. Therefore, we set CKITmin to be the time at which packet loss happens.

- 3. \frac{1}{2} CRTIMIN
- $\frac{4}{2}$  [CRTTmin ( $\frac{1}{2}$ )] =  $\frac{3}{4}$  CRTTmin
- $\frac{5}{1}$  throughput =  $\frac{3}{1}$  C
- $\frac{6}{}$  utilization = throughput / capacity =  $\frac{3}{4}$
- Non-zero Buffer
  - 7. Throughput =  $\begin{cases} \frac{W}{RTImin}, & W \leq C \cdot RTImin \\ C, & C \cdot RTImin \leq W \leq CRTImin + B \end{cases}$

When  $W \leq C$  RTTmin, the throughput is  $\frac{W}{RTTmin}$  Which is derived in class. In the other case, the quene starts to fill up, and throughput is the same as capacity C.

- 8 W, C, RTImin, W

M increases by  $\frac{1}{W}$  on every ACK. Pack loss happens when W surpasses CRT min +B.

11  $C \cdot RTT_{min} - \frac{1}{2}(CRTT_{min} + B) = (\frac{1}{2} C \cdot RTT_{min} - \frac{B}{2})$  units of time

12 We know that in the second phuse. W increases from CiRTImin to CiRTImin + B. W increases by 
$$\frac{1}{N}$$
 on every ACK.

It second is beastcally the time required for W to increase from B to BPP+1B. With large W, we can approximate this sum as integral

I trecond =  $\int_{0}^{B} RTImin + \frac{\pi^{2}}{C} dx$ 

$$= \int_{0}^{B} RTImin + \frac{\pi^{2}}{C} dx$$

13 thist + tsean = 
$$\frac{1-r}{2}$$
 CRTTmin + BRTTmin +  $\frac{B^2}{2C}$ 

average window size during thist =  $\frac{1}{2}$  [BDP]

=  $\frac{31r}{4}$  BDP

average throughput during 
$$t_{first} = \frac{3+r}{4} \frac{BDP}{RT_{min}}$$

$$= \frac{3+r}{2} \frac{BDP}{RT_{min}}$$

$$= \frac{3+r}{2} \frac{BP}{RT_{min}}$$

$$= \frac{3+r}{2} \frac{BP}{RT_{min$$

$$= \left( \frac{4 \, (^{2}R^{2}Y^{2} + 8 \, (^{2}R^{2}Y - (^{2}(^{2}+2x-3)R^{2}))^{2}}{4 [(^{2}R^{2}Y^{2} + 21 \, (^{2}R^{2} - (^{2}R^{2}(^{2}+2x-3)))]} \right)$$

$$= \left( \frac{3x^{2} + 67 + 43}{4x^{2} + 44 + 4} \right)$$

14 utilization = through put / 
$$C = \frac{3r^2 + 6r + 3}{4r^2 + 4r + 4}$$

Utilization U.S. 
$$\chi$$
 $\chi \in [0, \infty)$ 

$$\frac{16}{8} = \frac{12}{12} = \frac{12}{12} = \frac{100\%}{6}$$

$$= 7 \quad B = C \cdot RTImin$$

BDP+B = (1+r)BDP 
$$\frac{1}{M}$$
  $t_{first} = \frac{A}{BDP+B} = \frac{A}{A}$ 

$$= \frac{M-1-Y}{AM} = \frac{B}{A}$$

Time

 $t_{first} = \frac{B}{A}$ 
 $t_{first} = \frac{A}{A}$ 
 $t_{first} = \frac{A}{A}$ 

average window Size during 
$$t_{first} = \frac{1}{2} \begin{bmatrix} BDP + (tr)BDP \frac{1}{M} \end{bmatrix}$$
  
=  $\frac{M+1+r}{2M}BDP$ 

$$= \frac{M+l+r}{2M} \cdot \frac{M-l-r}{AM} \cdot \frac{M-l-r}{AM} \cdot \frac{RTTrain}{RTTrain}$$

$$= \frac{M-l-r}{AM} \cdot \frac{RTTrain}{RTTrain} + \left(\frac{RRTTrain}{RTTrain} + \frac{2B^2}{RTTrain}\right) + \left(\frac{RRTTrain}{RTTrain} + \frac{2B^2}{RTTrain}\right) + \left(\frac{RRTTrain}{RTTrain} + \frac{2B^2}{RTTrain}\right) + \left(\frac{RRTTrain}{RTTrain} + \frac{2B^2}{RTTrain}\right) + \left(\frac{RRTTrain}{RTTrain}\right) + \left(\frac{RRTTrain}{RTTrain}\right$$

vtilization = throughput /C

$$\frac{18}{2M(M-1)} = \frac{M^2 - 1}{2M(M-1)} = \frac{M+1}{2M} = \frac{1}{2} + \frac{1}{2M}$$

With B=0, utilization goes down as a function of M.

Since it's independent of A and U goes down as M incoenses

With large M, the rindow size decreases sharply after surpassing  $C \cdot RTTmin$ ,

And small window size means under utilization.

19. 
$$U(x) = 1$$
 =>  $M^{2}(2r+1) - (Y+1)^{2} = 2(M-1)M(Y+1)$   
=>  $Y = M-1$  =>  $B = (M-1)BDP$   
It goes up as a function of  $M$ .

What in this case would be C.RTInin + 13. Similar to the analysis above in <u>Q18</u>, we need to address the issue of having a very small window size after multiplicative decrease, in order to increase utilization. We can achieve that by increasing size of buffer.

Numerically,  $BDP + B = (M-1+1)BDP = MBDP \xrightarrow{decreuse} \xrightarrow{m} MBDP = BDP$ 



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