

# Performance

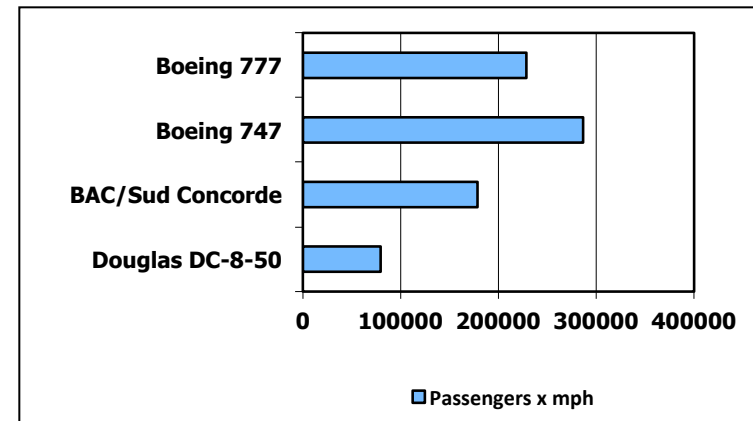
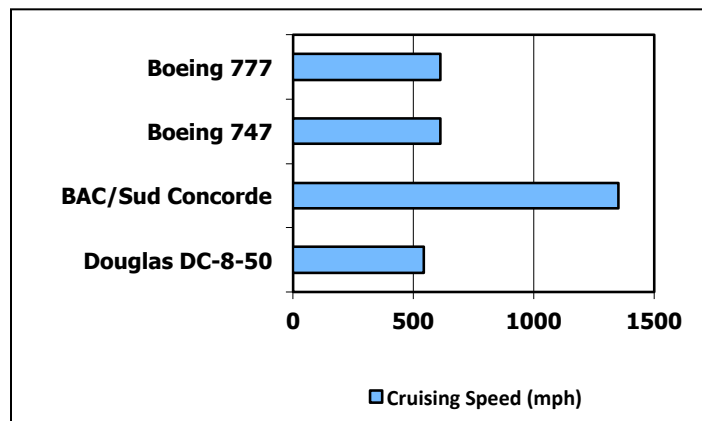
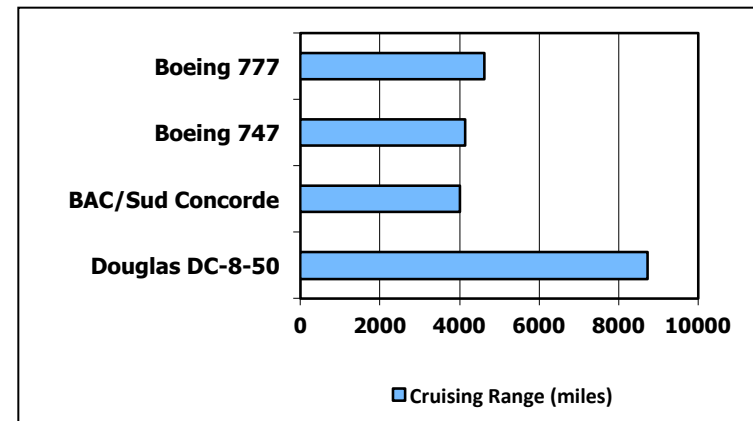
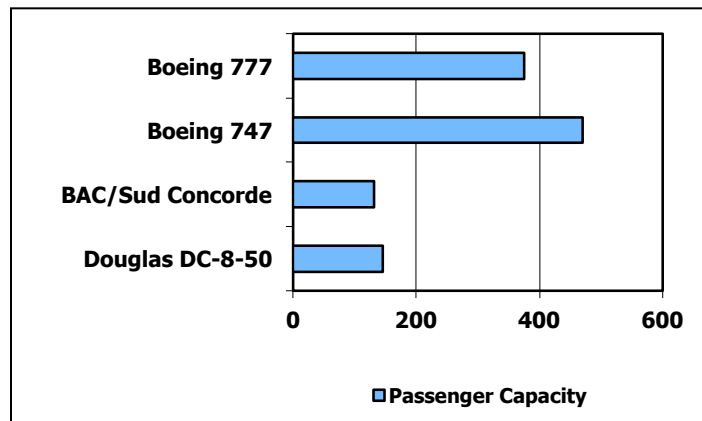
Lecture 19  
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Jae W. Lee ([jaewlee@snu.ac.kr](mailto:jaewlee@snu.ac.kr))  
Computer Science and Engineering  
Seoul National University

***Slide credits:*** [CS:APP3e] slides from CMU; [COD5e] slides from Elsevier Inc.

# Performance Example

■ Question: Which aircraft performs the best??



# Today

**Textbook: [P&H] 1.6**

- **Performance Metrics: Time and Rate**
- **Summarizing Performance**

# Performance Metrics #1: Time

## ■ Wall-clock time, response time, or elapsed time

- Actual time from start to completion
- Includes everything: CPU time for other programs as well as for itself, I/O, operating system overheads, etc

## ■ CPU (execution) time

- CPU time spent for a given program
- user CPU time + system CPU time
- e.g., results of UNIX `time` command

```
90.7u  12.9s  2:39  65%
```

# Performance Metrics #1: Time

## ■ Decomposition of CPU (Execution) Time

$$\begin{aligned}\text{CPU time} &= \frac{\text{Seconds}}{\text{Program}} \\ &= \frac{\text{Cycles}}{\text{Program}} \times \frac{\text{Seconds}}{\text{Cycle}} \\ &= \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}\end{aligned}$$

*This equation is called "**Iron Law of CPU Performance.**"*

# Performance Metrics #1: Time

## ■ More on CPI (Clocks or Cycles Per Instruction)

$$\text{CPI} = \frac{\sum_{i=1}^n (\text{CPI}_i \times I_i)}{\text{Instruction Count}}$$

### ■ CPI Example

Instruction Class	Frequency	CPI <sub>i</sub>
ALU operations	43%	1
Loads	21%	2
Stores	12%	2
Branches	24%	2

$$\text{CPI} = 0.43 \times 1 + 0.21 \times 2 + 0.12 \times 2 + 0.24 \times 2$$

# Performance Metrics #1: Time

## ■ Comparing CPIs of two CPUs

- Example question: What is the CPI of CPU<sub>S</sub> and CPU<sub>Q</sub>?

Instruction Type	Instr. count (millions)	Cycles per Instr. (CPI)	
		CPU <sub>S</sub>	CPU <sub>Q</sub>
Arithmetic & Logic	10	1	1
Load & Store	5	4	2
Branch	4	2	3
Miscellaneous (7 E )	1	4	4

$$\text{CPI}_S = (10 \times 1 + 5 \times 4 + 4 \times 2 + 1 \times 4) / (10 + 5 + 4 + 1) = 2.1$$

$$\text{CPI}_Q = (10 \times 1 + 5 \times 2 + 4 \times 3 + 1 \times 4) / (10 + 5 + 4 + 1) = 1.8$$

**Question:** So, CPU<sub>Q</sub> always performs better?

# Performance Metrics #1: Time

## ■ Factors involved in the CPU Time

$$\text{CPU time} = \frac{\text{Seconds}}{\text{Program}} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}$$

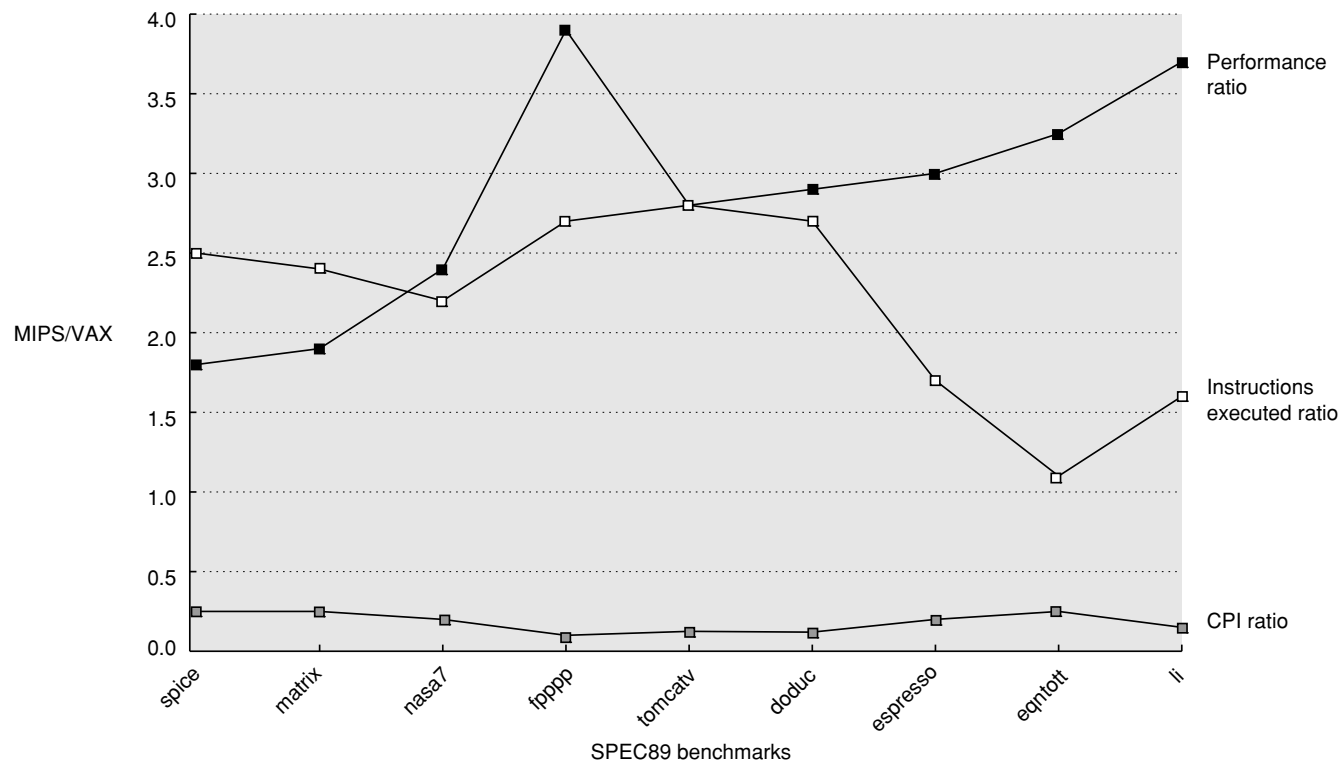
	$\frac{\text{Instructions}}{\text{Program}}$	$\frac{\text{Cycles}}{\text{Instruction}}$	$\frac{\text{Seconds}}{\text{Cycle}}$
Program	✓		
Compiler	✓		
ISA	✓	✓	
Organization		✓	✓
Technology			✓



# Performance Metrics #1: Time

## ■ RISC vs. CISC arguments

### ■ MIPS (typical RISC) vs. VAX8700 (typical CISC)



Source : Hennessy & Patterson *Computer Architecture: A Quantitative Approach, 5th Ed.(Appencix L)*, Morgan Kaufmann, 2012

# Performance Metrics #2: Rate

## ■ MIPS (million instructions per second)

- $$\text{MIPS} = \frac{\text{Instruction count}}{\text{Execution time} \times 10^6}$$
- Specifies performance (roughly) inversely to execution time
- Easy to understand; faster machines means bigger MIPS
- Problems
  - It does not take into account the capabilities of the instructions.
  - It varies between programs on the same computer.
  - It can even vary inversely with performance!!

## ■ MFLOPS (million floating-point operations per second)

# Performance Metrics: Ratio

- “X is n times faster than Y” means:

$$\frac{\text{Execution Time}_Y}{\text{Execution Time}_X} = n$$

- “X is n% faster than Y” means:

$$\frac{\text{Execution Time}_Y}{\text{Execution Time}_X} = 1 + \frac{n}{100}$$

- “X is n order of magnitude faster than Y” means:

$$\frac{\text{Execution Time}_Y}{\text{Execution Time}_X} = 10^n$$

# Summarizing Performance

- Arithmetic mean  
(Time)

$$\frac{1}{n} \sum_{i=1}^n T_i$$

- Harmonic mean  
(Rate)

$$\frac{n}{\sum_{i=1}^n \frac{1}{R_i}}$$

- Geometric mean  
(Ratio)

$$\sqrt[n]{\prod_{i=1}^n \text{Ratio}_i}$$

# Summarizing Performance: Arithmetic Mean

## ■ Used to summarize performance given in times

- Average Execution Time =  $( \sum_{i=1}^n \text{Execution Times} ) / n$
- Assumes each benchmark is run an equal no. of times

## ■ Weighted Arithmetic Mean

- Weighted Average Execution Time =  $\sum_{i=1}^n ( W_i \times \text{Execution Times} ) / \sum_{i=1}^n W_i$
- One possible weight assignment: equal execution time on some machine

# Summarizing Performance: Harmonic Mean

## ■ Used to summarize performance in rates (e.g., MIPS, FLOPS):

- Harmonic Mean =  $n / \sum_{i=1}^n (1 / R_i)$

- Example

- Four programs execute at 10, 100, 50 and 20 MFLOPS, respectively
- Harmonic mean is  $4 / (1/10 + 1/100 + 1/50 + 1/20) = 22.2$  MFLOPS

## ■ Weighted Harmonic Mean

- Weighted Harmonic Mean =  $\sum_{i=1}^n W_i / \sum_{i=1}^n (W_i / R_i)$

# Summary

$$\text{CPU time} = \frac{\text{Seconds}}{\text{Program}} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}$$

- **“Execution time is the only and unimpeachable measure of performance”**
  - CPU time equation can predict performance by estimating the effects of changing features.
- **Measuring performance requires good care**
  - Good ways to summarize performance
  - Good workloads (benchmarks)