



# Semi-Automatic Functional Conversion of microKanren

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One relation to solve many problems

Nondeterminism

Completeness of search

# Relational Conversion: Easy

Given a function

---

```
let rec add x y =  
  match x with  
  | 0 → y  
  | S x' → S (add x' y)
```

---

generate miniKanren relation

---

```
let rec addo x y z = conde [  
  (x ≡ 0 ∧ y ≡ z);  
  (fresh (x' z')  
    (x ≡ S x' ∧  
      addo x' y z' ∧  
      z ≡ S z')) ]
```

---

# Principal Directions of MINIKANREN Relations

Every argument of a relation can be either `in` or `out`

For addition relation `addo x y z` there are 8 directions:

- *Forward* direction: `addo in in out` — addition
- *Backward* direction: `addo out out in` — decomposition
- *Predicate*: `addo in in in`
- *Generator*: `addo out out out`
- `addo in out in` — subtraction
- `addo out in in` — subtraction
- `addo out in out`
- `addo in out out`

## Each Direction is a Function

# Each Direction is a Function (kinda)

Straightforward functions:

- *Forward* direction:  $\text{add}^o$  in in out — addition
- $\text{add}^o$  in out in — subtraction
- $\text{add}^o$  out in in — subtraction
- *Predicate*:  $\text{add}^o$  in in in

Relations:

- *Backward* direction:  $\text{add}^o$  out out in — decomposition
- *Generator*:  $\text{add}^o$  out out out
- $\text{add}^o$  out in out
- $\text{add}^o$  in out out

These relations are functions which return multiple answers (list monad)

# MINIKANREN Comes with an Overhead

Unifications

Occurs-check

Scheduling complexity

Given a relation and a principal direction, construct a functional program which generates the same answers as `MINIKANREN` would

Preserve completeness of the search

Both inputs and outputs are expected to be ground



## Example: Addition in Forward Direction

---

```
let rec addo x y z = conde [  
  (x ≡ 0 ∧ y ≡ z);  
  (fresh (x' z')  
    (x ≡ S x' ∧  
      addo x' y z' ∧  
      z ≡ S z')) ]
```

---

---

```
addIIIO :: Nat → Nat → Nat  
addIIIO x y =  
  case x of  
    0 → y  
    S x' → S (addIIIO x' y)
```

---

# Addition in Backwards Direction: Nondeterminism

---

```
let rec addo x y z = conde [  
  (x ≡ 0 ∧ y ≡ z);  
  (fresh (x' z')  
    (x ≡ S x' ∧  
      addo x' y z' ∧  
      z ≡ S z')) ]
```

---

---

```
add00I :: Nat → Stream (Nat, Nat)  
add00I z =  
  return (0, z) 'mplus'  
  case z of  
    0 → Empty  
    S z' → do  
      (x', y) ← add00I z'  
      return (S x', y)
```

---

# Free Variables in Answers: Generators

---

```
let rec addo x y z = conde [  
  (x ≡ 0 ∧ y ≡ z);  
  (fresh (x' z')  
    (x ≡ S x' ∧ z ≡ S z' ∧ addo x' y z') ) ]
```

---

---

```
addIOO :: Nat → Stream (Nat, Nat)  
addIOO x = case x of  
  0 → do  
    z ← genNat  
    return (z, z)  
  S x' → do  
    (y, z') ← addIOO x'  
    return (y, S z')
```

```
genNat :: Stream Nat  
genNat = Mature 0 (S <$> genNat)
```

---

# Predicates

---

```
let rec addo x y z = conde [  
  (x ≡ 0 ∧ y ≡ z);  
  (fresh (x' z')  
    (x ≡ S x' ∧  
      addo x' y z' ∧  
      z ≡ S z')) ]
```

---

---

```
addIII :: Nat → Nat → Nat → Stream ()  
addIII x y z =  
  case x of  
    0 | y == z → return ()  
      | otherwise → Empty  
    S x' →  
      case z of  
        0 → Empty  
        S z' → addIII x' y z'
```

---

# Order in Conjunctions

---

```
let rec multo x y z = conde [  
  ...  
  (fresh (x' r')  
    (x ≡ S x') ∧  
    (addo y r' z) ∧  
    (multo x' y r'))  
  ]
```

---

# Order in Conjunctions: Slow Version

---

```
multIIIO' :: Nat → Nat → Stream Nat
```

```
...
```

```
multIIIO' (S x') y    = do
  (r', r) ← addX y
  multIII x' y r'
  return r
```

```
multIII :: Nat → Nat → Nat → Stream ()
```

```
...
```

```
multIII (S x') y z = do
  z' ← multIIIO' x' y
  addIII y z' z
multIII _ _ _ = Empty
```

---

# Order in Conjunctions: Faster Version

---

```
multIIO :: Nat → Nat → Stream Nat
```

```
...
```

```
multIIO (S x') y = do  
  r' ← multIIO x' y  
  addXY y r'
```

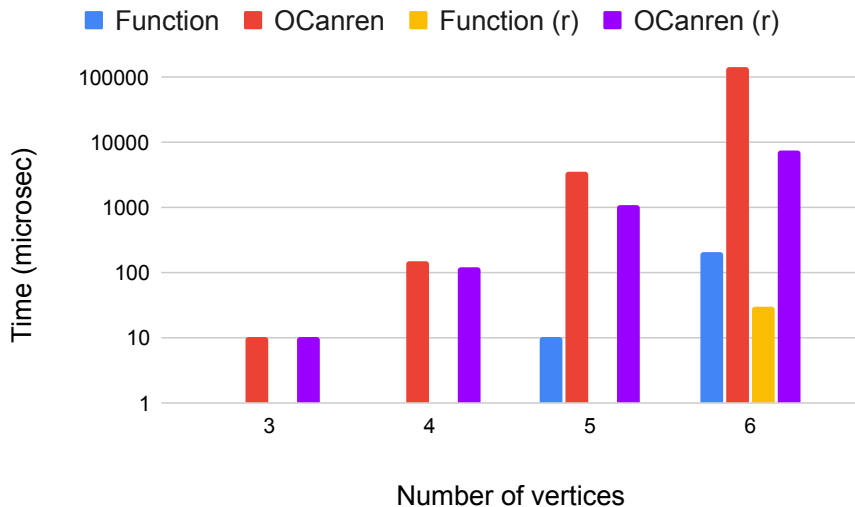
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We manually converted relational interpreters and measured execution time

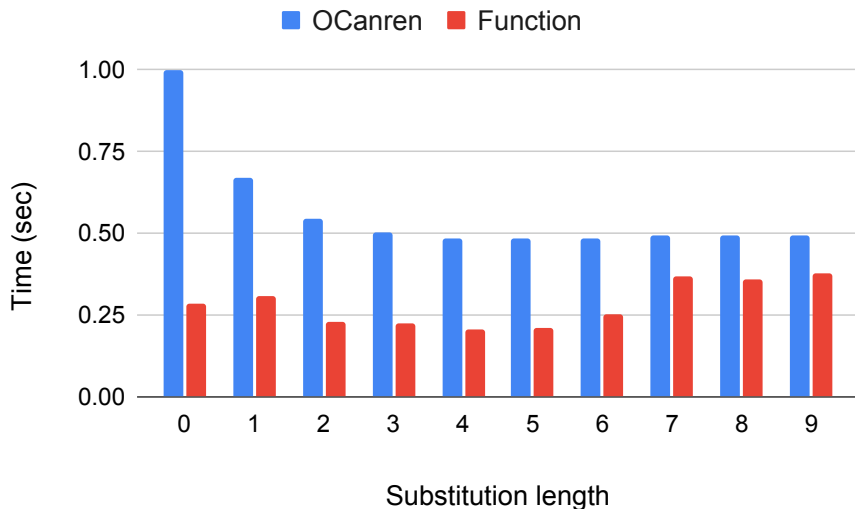
- Topologic sort
  - A verifier verifies that a vertex mapping sorts vertices topologically
  - Sort a DAG with an edge in between every pair of vertices
  - Two different representations: vertices sorted by their number, and with a reverse order
  - Sorting a graph with up to 6 vertices
- Logic formulas generation
  - Inverse computation of a logic formulas interpreter
  - Generate 10000 formulas which evaluate to true
  - Different substitution lengths



# Evaluation: Topologic Sort



# Evaluation: Logic Formulas Generation



## Conclusion

- We presented a functional conversion scheme as a series of examples
- The conversion speeds up implementations considerably

## Future work

- Implementation and formalization of the conversion scheme
- Finding a better way to order conjuncts
- Integration into a relational interpreters for solving framework