

GHS Implementation Report

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1 Introduction to the GHS Algorithm

The GHS Algorithm (named after it's authors Gallager, Humblet and Spira) was first published in a 1983 [paper](#). It gives an asynchronous distributed algorithm that runs identically on each node and in due time, each node

learns which of its incident edges belong in the minimum spanning tree.

The algorithm essentially works by creating, expanding and merging fragments of the graph. These fragments are guaranteed to be a part of the final MST of the complete graph.

Each node can be in one of 3 states : SLEEP, FIND and FOUND.

For a node, the edge between every neighbor can be in one of 3 states : BASIC, BRANCH and REJECT.

In basic state represents an edge that is yet to be tested for its presence in the MST. An edge in state branch represents an edge that is known to be in the final MST. And an edge in reject state has been rejected and cannot be a part of the MST.

At the start of the algorithm, all nodes are in state SLEEP, they are either woken up by a message, or they wake up spontaneously. For getting the best bound on messages, we shall assume that all nodes are woken up simultaneously at the start.

After that each node can either be in state FIND where it is looking for new edges for MST or it can be in state FOUND when it's passive.

Each node has a LEVEL and a NAME which is representative of its fragment. Two fragments combine when either the node with smaller requests a connection to the larger or when two nodes with same level combine and form a node with level incremented by one. The first case is called the LT rule of merging, and the second case is the EQ rule.

explain The nodes communicate through the following types of messages :

- CONNECT : Request for merging the two nodes.
- INITIATE : Communicates new fragment level and identity to the neighbors
- TEST : Tests to check whether an edge belongs to the MST
- ACCEPT : Communicates to the neighbor that the received edge tested is accepted into the MST
- REJECT : Communicates to the neighbor that the received edge cannot possibly be in the MST
- REPORT : Reports the best weight outgoing fragment for each node
- CHANGEROOT : The root node of the fragment is shifted to the node with the minimal outgoing edge

Under the assumption of initial awakening of all the nodes, GHS Algorithm guarantees the number of messages sent to be $< 2 \cdot E + 5 \cdot N \cdot \log_2(N)$. We shall see more about this in our experiments

We shall see the algorithm itself in next section :

2 Implementation Details

We have implemented a multi-threaded parallel implementation of the GHS Algorithm. Each node of the GHS is given a single thread in the implementation.

In this section, we shall outline the two parts of our implementation :

1. Inter-Thread Communication
2. GHS Algorithm at each node
3. Main Function that runs thread

2.1 Inter-Thread Communication

2.1.1 Messages

The messages sent by the threads(nodes) are dynamic arrays of strings.

We have created a message class to standardize all messages :

```
1  class Message
2  {
3  private:
4      std::vector<std::string> msg; //!< message content
5  public:
6      Message(std::vector<std::string> m) //!< Initiates a message
7      {
8          msg = m;
9      }
10     std::vector<std::string> getMessage() //!< returns the message
11     {
12         return msg;
13     }
14 };
15
```

Listing 1: Message Class

For the ease of sending messages, we have created the a helper function MSGCREATER() :

```

1
2 Message *GHSNode::msgCreator(std::vector<std::string> msg)
3 {
4     std::vector<std::string> m;
5     m.push_back(std::to_string(nodeid));
6
7     for(auto it : msg)
8     {
9         m.push_back(it);
10    }
11
12    Message *mg = new Message(m);
13    return mg;
14 }
15

```

Listing 2: msgCreator

This dynamic array has the structure : {sender_id, Message type, Parameter1, Parameter2, ...}

Here is an example of message sending (SENDMESSAGE() will be explained later) :

```

1 m.push_back("initiate"); //!< Initiate()
2 m.push_back(std::to_string(LN)); //!< LN
3 m.push_back(FN); //!< FN
4 m.push_back(SN); //!< SN
5
6 sendMessage(j, msgCreator(m)); //!< send Initiate(LN, FN, SN) on edge j
7

```

Listing 3: sendMessage Sample

2.1.2 Queues

To handle simultaneous messages, we have defined a thread safe data structure QUEUE built upon the std::queue provided by the Standard Template Library of C++ :

```

1 class Queue
2 {
3     private:
4         std::mutex mut; //!< mutex lock for thread safety
5         std::queue<Message *> q; //!< Message queue
6         int queueid; //!< Id of the node to which this queue belongs
7     public:
8         Queue()
9         {
10             queueid = -1; //!< Initialized as bad node
11         }

```

```

12 Queue(int qid)
13 {
14     queueid = qid; //!< ID of the queue
15 }
16 int getqueueid()
17 {
18     return queueid; //!< returns id for the node which owns this queue
19 }
20 void push(Message *m)
21 {
22     mut.lock(); //!< Puts the lock in place
23     q.push(m); //!< This push operation is atomic
24     mut.unlock(); //!< Unlocks the queue
25 }
26 Message *front()
27 {
28     mut.lock(); //!< Puts the lock in place
29     if(q.size() == 0)
30     {
31         mut.unlock(); //!< Unlocks the queue
32         return NULL; //!< Returns NULL if queue is empty
33     }
34     Message *temp = q.front(); //!< This front operation is atomic
35     mut.unlock(); //!< Unlocks the queue
36     return temp;
37 }
38 Message *pop()
39 {
40     mut.lock(); //!< Puts the lock in place
41     int temp = q.size(); //!< Temporary variable to store queue size
42     if(temp == 0)
43     {
44         mut.unlock(); //!< Unlocks the queue
45         return NULL; //!< Returns NULL if queue is empty
46     }
47     Message *tmp = q.front(); //!< Temporary variable to store the front of queue
48     q.pop(); //!< Pops the queue
49     if((int)q.size() != temp - 1) //!< Pop should change the size of queue by 1
50     {
51         std::cerr << "QUEUE Error" << std::endl;
52         exit(1);
53     }
54     mut.unlock(); //!< Unlocks the queue
55     return tmp;
56 }
57 bool empty()
58 {
59     mut.lock(); //!< Puts the lock in place

```

```

60     int temp = q.size(); //!< Temporary variable for queue size
61     mut.unlock(); //!< Unlocks the queue
62     return (temp == 0); //!< Returns if queue is empty
63 }
64 int getQueueSize()
65 {
66     mut.lock(); //!< Puts the lock in place
67     int temp = q.size(); //!< Size of queue
68     mut.unlock(); //!< Unlocks the queue
69     return temp; //!< returns size of queue
70 }
71 };
72

```

Listing 4: Queue

Each thread(node) gets it's own Queue instance, which is instantiated in the Network class. A thread which of the queue has all the operations available to it whereas a non-owner thread can only **push** to the queue.

Each node (i.e. each Instance of the GHSNode class) is provided with a pointer to it's queue at it's instantiation. The pointer to other node's queue can be taken from the Network class which we describe next.

2.1.3 Network

To send and receive messages from different nodes, we have a defined an internet of sorts for the threads. This is the NETWORK class :

```

1     class Network
2     {
3     private:
4         std::unordered_map<int, Queue *> msg_queues; //!< Network queues
5     public:
6         Network(std::vector<int> nodes) //!< Network Constructor
7         {
8             for(auto it : nodes)
9             {
10                 Queue *q = new Queue(it); //!< Initiates a new queue
11                 msg_queues[it] = q; //!< Initializes the hashmap
12             }
13         }
14         Queue *getQueue(int i) //!< Get's the ith queue
15         {
16             if(msg_queues.find(i) == msg_queues.end())
17             {
18                 std::cerr << "BAD network Request" << std::endl;
19             }
20             return msg_queues[i]; //!< Returns the pointer to ith queue
21         }
22     }

```

23 ;
24

Listing 5: Network

The NETWORK class is instantiated only once, and each node keeps a pointer to the single instance.

2.1.4 IsComplete

IsCOMPLETE is a class which we have defined to trigger the end of the algorithm. This is a thread join of sorts. Once this is complete we shutdown all nodes and ask them for the part of MST that they possess.

This class is instantiated only once and every node possesses a pointer to this instance so that any node can give out the halt signal.

```
1  struct IsComplete
2  {
3      bool complete; //!< If set true, the threads exit
4      IsComplete() //!< Constructor that sets the default value false
5      {
6          complete = false;
7      }
8  };
9
10
```

Listing 6: IsComplete

2.1.5 GHS Algorithm Constructor

To show what we pass to each node initially, we shall list out the constructor for the GHSNode :

```
1  GHSNode::GHSNode(int nid, std::unordered_map<int, int> &neighbors, Network *net,
2  IsComplete *iscom)
3  {
4      this->nodeid = nid;
5      this->nbd = neighbors;
6      this->network = net;
7      this->SN = "sleep";
8      this->isc = iscom;
9      this->nodequeue = net->getQueue(nid);
10 }
```

Listing 7: GHSNode constructor

As we can see, each node receives the pointer to its queue and the pointer to the network along with the IsComplete instance.

2.1.6 Sending and Receiving messages

Now to summarize the thread messaging, let us give the code snippets for the `SENDMESSAGE()` and `RECEIVEMESSAGE()` methods :

```

1  void GHSNode::sendMessage(int dest, Message *m)
2  {
3      Queue *q = network->getQueue(dest);
4      q->push(m);
5  }
6

```

Listing 8: sendMessage

```

1  bool GHSNode::recieveMessage()
2  {
3      msg = nodequeue->pop();
4      return (msg != NULL);
5  }
6

```

Listing 9: receiveMessage

This completes our summary of thread communication

2.2 GHS Algorithm implementation at each node

Each GHSNode has the following state variables :

Variable Name	Variable Title	Range of the variable
NodeId	Id of the node	\mathbb{N}
SN	State of the node	["sleep", "find", "found"]
FN	Name of the node	-
LN	Level of the node	\mathbb{N}
nbd	Adjacency list	$\{key : value\} \text{ key} \in \mathbb{N}, \text{ value} \in \mathbb{N}$
SE	State of edges	["baisc", "branch", "reject"]
best-edge	Currently the minimum weight outgoing edge	\mathbb{N}
best-weight	Current best outgoing weight	\mathbb{N}
test-edge	Edge currently being tested	\mathbb{N}
in-branch	Parent node of sorts	\mathbb{N}
find-count	number of MST Edges found	\mathbb{N}

The following types of messages can be sent :

Message Type	Description	Parameters
Connect	Request for merging the two nodes	Level
Initiate	Communicates new fragment level and identity to the neighbors	Level, Name, State
Test	Tests to check whether an edge belongs to the MST	Level, Name
Accept	Received test edge is accepted into the MST	-
Reject	Received test edge cannot possibly be in the MST	-
Report	Reports the best weight outgoing fragment for each node	best weight
ChangeRoot	Start Changing root of fragment	-

We will now list the 12 Algorithms as mentioned in the original GHS [paper](#) and our implementation for the same.

These are long pseudocodes and codes. If required one can skip to the [next](#) subsection on thread_runner.

2.2.1 Algorithm 1

Spontaneous awakening at each node at the starting instant. RUN() is the function that the thread starts.

```

1  void GHSNode::run()
2  {
3      wakeup();
4      runner();
5  }
6

```

Listing 10: run

2.2.2 Algorithm 2

Pseudocode :

```

wakeup()
begin
let  $m$  be adjacent edge of minimum weight
 $SE(m) \leftarrow Branch$ ;
 $LN \leftarrow 0$ ;
 $SN \leftarrow Found$ ;
 $Find - count \leftarrow 0$ ;
send  $Connect(0)$  on edge  $m$ 

```

Implementation :

```

1  void GHSNode::wakeup()
2  {
3      if (nbd.empty())
4      {
5          isc->complete = true;
6      }

```

```

7
8   for(auto it : nbd)
9   {
10      SE[it.first] = "basic";
11      basic.insert(std::make_pair(it.second, it.first));
12  }
13
14  int m = findMinEdge(); // let m be adjacent edge of minimum weight;
15  changestat(m, "branch"); //!< SE(m) <- Branch
16  LN = 0; //!< LN <- 0
17  SN = "found"; //!< SN <- Found
18  find_count = 0; //!< Find-count <- 0
19  std::vector<std::string> st;
20  st.push_back("connect");
21  st.push_back(std::to_string(LN));
22  sendMessage(m, msgCreator(st)); //!< send Connect(0) on edge m
23  }
24

```

Listing 11: wakeup

2.2.3 Algorithm 3

Pseudocode :

```

Response to Connect(L) on edge j
begin if SN = Sleeping then execute procedure wakeup;
    if L < LN
    then begin SE(j) ← Branch;
        send Initiate(LN, FN, SN) on edge j
        if SN = Find then
            find_count ← find_count + 1
        end
    else if SE(j) = Basic
    then place received message on end of queue
    else send Initiate(LN + 1, w(j), Find) on edge j
end

```

Implementation :

```

1   void GHSNode::handleConnect()
2   {
3       std::vector<std::string> mg = msg->getMessage();
4       int j = std::stoi(mg[0]); //!< edge j
5       int L = std::stoi(mg[2]); //!< L
6       if(SN == "sleep") //!< if SN = Sleeping
7       {

```

```

8      wakeup(); //!< execute procedure wakeup
9  }
10 if (L < LN) //!< If L < LN
11 {
12     changestat(j, "branch");
13
14     std::vector<std::string> m;
15     m.push_back("initiate"); //!< Initiate()
16     m.push_back(std::to_string(LN)); //!< LN
17     m.push_back(FN); //!< FN
18     m.push_back(SN); //!< SN
19
20     sendMessage(j, msgCreator(m)); //!< send Initiate(LN, FN, SN) on edge j
21
22     if (SN == "find") //!< if SN = Find then
23     {
24         find_count++; //!< find-count <- find-count +1
25     }
26
27     free(msg);
28     msg = NULL;
29 }
30 else if (SE[j] == "basic") //!< else if SE(j) = Basic
31 {
32     sendMessage(nodeid, msg); //!< place received message on end of queue
33 }
34 else
35 {
36     std::vector<std::string> st;
37     st.push_back("initiate"); //!< Initiate()
38     st.push_back(std::to_string(LN+1)); //!< LN +1
39     st.push_back(std::to_string(nbd[j])); //!< w(j)
40     st.push_back("find"); //!< Find
41     sendMessage(j, msgCreator(st)); //!< send Initiate(LN+1, w(j), Find) on edge
j
42     free(msg);
43     msg = NULL;
44 }
45 }
46

```

Listing 12: handleConnect()

2.2.4 Algorithm 4

Pseudocode :

Response to $Initiate(L, F, S)$ on edge j

```

begin  $LN \leftarrow L; FN \leftarrow F; SN \leftarrow S; in\_branch \leftarrow j;$ 
     $best\_edge \leftarrow nil; best\_wt \leftarrow \infty;$ 
    for all  $i \neq j$  such that  $SE(i) = Branch$ 
        do begin send  $Initiate(L, F, S)$  on edge  $i;$ 
            if  $S = Find$  then  $find\_count \leftarrow find\_count + 1;$ 
        end
    if  $S = Find$  then execute procedure  $test$ 
end

```

Implementation :

```

1  void GHSNode::handleInitiate()
2  {
3      std::vector<std::string> mg = msg->getMessage();
4      int j = stoi(mg[0]); //!< edge j
5      int L = stoi(mg[2]); //!< L
6      std::string F = mg[3]; //!< F
7      std::string S = mg[4]; //!< S
8      LN = L; //!< LN <- L
9      FN = F; //!< FN <- F
10     SN = S; //!< SN <- S
11     in_branch = j; //!< in-branch <- j
12     best_edge = -1; //!< best-edge <- nil
13     best_weight = INF; //!< best-wt <- infinity
14     for(auto it : branch) //!< forall
15     {
16         if(it.second == j)
17         {
18             continue;
19         } //!< i != j and SE(i) = branch
20         std::vector<std::string> st;
21         st.push_back("initiate"); //!< Initiate()
22         st.push_back(std::to_string(L)); //!< L
23         st.push_back(F); //!< F
24         st.push_back(S); //!< S
25         sendMessage(it.second, msgCreator(st)); //!< send Initiate(L, F, S) on edge i
26         if(S == "find") //!< if S = Find
27         {
28             find_count++; //!< find-count <- find-count + 1
29         }
30     }
31     if(S == "find") //!< if S = Find

```

```

32     {
33         test(); //!< execute procedure test
34     }
35     free(msg);
36     msg = NULL;
37 }
38

```

Listing 13: handleInitiate()

2.2.5 Algorithm 5

Pseudocode :

```

procedure test
    if there are adjacent edges in the state Basic
        then begin test_edge  $\leftarrow$  the minimum-weight adjacent edge in state Basic;
            send Test(LN, FN) on test_edge;
        end
    else begin test_edge  $\leftarrow$  nil; execute procedure report end

```

Implementation :

```

1  void GHSNode::test()
2  {
3      test_edge = findMinEdge(); //!< test-edge  $\leftarrow$  nil if no basic edges
4      //!< test-edge  $\leftarrow$  the minimum-weight adjacent edge in state Basic
5      if (test_edge != -1) //!< If there are adjacent edges in the state Basic
6      {
7          std::vector<std::string> st;
8          st.push_back("test"); //!< Test()
9          st.push_back(std::to_string(LN)); //!< LN
10         st.push_back(FN); //!< FN
11         sendMessage(test_edge, msgCreator(st)); //!< send Test(LN, FN) on test-edge
12     }
13     else //!< Else test-edge  $\leftarrow$  nil (by default)
14     {
15         report(); //!< execute procedure report()
16     }
17 }
18

```

Listing 14: test()

2.2.6 Algorithm 6

Pseudocode :

Response to $Test(L, F)$ on edge j

```
begin if  $SN = Sleeping$  then execute procedure wakeup;
    if  $L > LN$  then place received message on end of queue;
    else if  $F \neq FN$  then send Accept on edge  $j$  ;
        else begin if  $SE(j) = Basic$  then  $SE(j) \leftarrow Rejected$ ;
            if  $test\_edge \neq j$  then send Reject on edge  $j$ 
            else execute procedure test
        end
    end
end
```

Implementation :

```
1 void GHSNode::handleTest()
2 {
3     std::vector<std::string> mg = msg->getMessage();
4     int j = std::stoi(mg[0]); //!< edge j
5     if(SN == "sleep") //!< If SN = Sleeping
6     {
7         wakeup(); //!< execute procedure wakeup
8     }
9     int L = std::stoi(mg[2]); //!< L
10    std::string F = mg[3]; //!< F
11    if(L > LN) //!< L > LN
12    {
13        sendMessage(nodeid, msg); //!< place recieved message on end of queue
14    }
15    else if(F != FN) //!< F != FN
16    {
17        std::vector<std::string> st;
18        st.push_back("accept"); //!< Accept()
19        sendMessage(j, msgCreator(st)); //!< Send Accept on edge j
20        free(msg);
21        msg = NULL;
22    }
23    else
24    {
25        if(SE[j] == "basic") //!< if SE(j) = Basic
26        {
27            changestat(j, "reject");
28        }
29        if(test_edge != j) //!< test-edge != j
30        {
31            std::vector<std::string> st;
32            st.push_back("reject"); //!< Reject()
33            sendMessage(j, msgCreator(st)); //!< send Reject on edge j
34        }
35        else
```

```

36     {
37         test(); /*< execute procedure test
38     }
39     free(msg);
40     msg = NULL;
41 }
42 }
43

```

Listing 15: handleTest()

2.2.7 Algorithm 7

Pseudocode :

Response to *Accept* on edge j

```

begin test_edge  $\leftarrow$  nil;
    if  $w(j) < best\_wt$ 
        then begin best_edge  $\leftarrow j$ ; best_wt  $\leftarrow w(j)$  end
        execute procedure report
    end

```

Implementation :

```

1  void GHSNode::handleAccept()
2  {
3      std::vector<std::string> mg = msg->getMessage();
4      int j = std::stoi(mg[0]); /*< edge j
5      test_edge = -1; /*< test-edge <- nil
6      if(nbd[j] < best_weight) /*< w(j) < best-wt
7      {
8          best_edge = j; /*< best-edge <- j
9          best_weight = nbd[j]; /*< best-wt <- w(j)
10     }
11     free(msg);
12     msg = NULL;
13     report(); /*< execute procedure report()
14 }
15

```

Listing 16: handleAccept()

2.2.8 Algorithm 8

Pseudocode :

Response to *Reject* on edge j

```

begin if  $SE(j) = Basic$  then  $SE(j) \leftarrow Rejected$ ;
    execute procedure test
end

```

Implementation :

```

1  void GHSNode::handleReject()
2  {
3      std::vector<std::string> mg = msg->getMessage();
4      int j = std::stoi(mg[0]); //!< edge j
5      if(SE[j] == "basic") //!< if SE(j) = Basic
6      {
7          changestat(j, "reject");
8      }
9      free(msg);
10     msg = NULL;
11     test(); //!< execute procedure test()
12 }
13

```

Listing 17: handleReject()

2.2.9 Algorithm 9

Pseudocode :

```

procedure report
    if  $find\_count = 0$  and  $test\_edge = nil$ ;
    then begin  $SN \leftarrow Found$ ;
        send  $Report(best\_wt)$  on  $in\_branch$ ;
    end

```

Implementation :

```

1  void GHSNode::report()
2  {
3      if(find_count == 0 && test_edge == -1) //!< if find-count = 0 and test-edge
= nil
4      {
5          SN = "found"; //!< SN <- Found
6          std::vector<std::string> st;

```



```

7      st.push_back("report"); //!< Report()
8      st.push_back(std::to_string(best_weight)); //!< best-wt
9      sendMessage(in_branch, msgCreator(st)); //!< send Report(best-wt) on in-branch
10     }
11 }
12

```

Listing 18: report()

2.2.10 Algorithm 10

Pseudocode :

Response to $Report(w)$ on edge j

```

if  $j \neq in\_branch$ 
  then begin  $find\_count \leftarrow find\_count - 1$ 
    if  $w < best\_wt$  then begin  $best\_wt \leftarrow w; best\_edge \leftarrow j$  end
    execute procedure report
  end
else if  $SN = Find$  then place received message on end of queue
  else if  $w > best\_wt$ 
    then execute procedure  $change\_root$ 
  else if  $w = best\_wt = \infty$  then halt

```

Implementation :

```

1  void GHSNode::handleReport()
2  {
3      std::vector<std::string> mg = msg->getMessage();
4      int j = std::stoi(mg[0]); //!< edge j
5      int w = std::stoi(mg[2]); //!< w
6      if(j != in_branch) //!< if j != in-branch
7      {
8          find_count--; //!< find-count <- find-count - 1
9          if(w < best_weight) //!< w < best-wt
10         {
11             best_weight = w; //!< best-wt <- w
12             best_edge = j; //!< best-edge <- j
13         }
14         free(msg);
15         msg = NULL;
16         report(); //!< execute procedure report()
17     }
18     else if(SN == "find") //!< SN = Find
19     {
20         sendMessage(nodeid, msg); //!< place recieved message on end of queue
21     }

```

```

22     else if(w > best_weight) //!< w > best-wt
23     {
24         free(msg);
25         msg = NULL;
26         changeRoot(); //!< execute procedure change-root()
27     }
28     else if(w == best_weight && w == INF) //!< w = best-wt = infinity
29     {
30         isc->complete = true; //!< halt
31         free(msg);
32         msg = NULL;
33     }
34 }
35

```

Listing 19: handleReport()

Note that at halt step, we set the isc pointer to true, this tells the main thread that the nodes have found the MST.

2.2.11 Algorithm 11

Pseudocode :

```

procedure change_root
  if  $SE(best\_edge) = Branch$ 
    then send Change_root on best_edge
  else begin send Connect(LN) on best_edge;
     $SE(best\_edge) \leftarrow Branch$ 
  end

```

Implementation :

```

1  void GHSNode::changeRoot()
2  {
3      if(SE[best_edge] == "branch") //!< SE(best-edge) = Branch
4      {
5          std::vector<std::string> st;
6          st.push_back("changeroot"); //!< Change-root()
7          sendMessage(best_edge, msgCreator(st)); //!< send Change-root() on best-edge
8      }
9      else
10     {
11         std::vector<std::string> st;
12         st.push_back("connect"); //!< Connect()
13         st.push_back(std::to_string(LN)); //!< LN
14         sendMessage(best_edge, msgCreator(st)); //!< send Connect(LN) on best-edge
15         changestat(best_edge, "branch");

```

```

16     }
17 }
18

```

Listing 20: changeRoot()

2.2.12 Algorithm 12

Pseudocode :

Response to *Change_root*
execute procedure *change_root*

Implementation :

```

1  void GHSNode::handleChangeroot()
2  {
3      free(msg);
4      msg = NULL;
5      changeRoot(); //execute procedure change-root()
6  }
7

```

Listing 21: handleChangeroot()

2.3 Main Function/ Thread Runner

In this section we provide the snippet for our thread_runner funtion which initiates all the GHS Nodes and runs them.

```

1  Network *network = new Network(no); //!< Networks for the threads
2  IsComplete *isc = new IsComplete(); //!< Flag to check completion
3  TotMessage *tot = new TotMessage(); //!< Total message counter
4
5  std::vector<pthread_t> threads(n); //!< Vector of threads
6  std::vector<GHSNode *> nodes; //!< Vector of all GHSNodes
7
8  int i = 0;
9  for(auto it : adj_list)
10 {
11     GHSNode *temp = new GHSNode(it.first, it.second, network, isc, tot); //!<
Create new GHSNode
12     nodes.push_back(temp);
13
14     int errcode = pthread_create(&(threads[i]), NULL, run_thread, (void *)temp);
//!< Start the thread, if errcode != 0 then thread creation was not successful
15
16     if(errcode != 0)

```

```

17     {
18         std::cerr << "Thread Creation at index : " << i << " Failed.\n Exiting ....
" << std::endl;
19         exit(49);
20     }
21     i++;
22 }
23
24 while(!(isc->complete))
25 {
26     //threads still running
27     continue;
28 }
29 //GHS Complete
30
31

```

Listing 22: thread_runner Sample

Here each thread calls run_thread function which is a trivial function which fires up the node :

```

1 void* run_thread(void * node)
2 {
3     ((GHSNode *)node)->run();
4     pthread_exit(NULL);
5 }
6

```

Listing 23: run_thread()

For each node, we store classify edges into three sets : basic, branch and reject -

```

1 std::set<std::pair<int, int>> basic, branch, reject; //!< Set containing (
edge_weight, id) of edges in basic, branch and reject state respectively
2

```

Listing 24: branch classification

At the end of the algorithm, we collect the branch edges of all nodes via :

```

1 std::set<std::tuple<int, int, int>> ans; //!< MST Edges
2 for(auto it : nodes)
3 {
4     std::vector<int> v = it->getMSTEdges();
5     for(auto jt : v)
6     {
7         ans.insert(std::make_tuple(jt, mp[jt].first, mp[jt].second));
8     }
9 }
10 return ans;
11

```

Listing 25: MST Collection

Here jt is the weight of the edge, and $mp[jt]$ represents the nodes adjacent to it. And `getMSTEdges()` is as follows :

```

1  std::vector<int> GHSNode::getMSTEdges()
2  {
3      std::vector<int> v;
4      for(auto it : branch)
5      {
6          v.push_back(it.first);
7      }
8      return v;
9  }
10

```

Listing 26: `getMSTEdges()`

Further documentation for the code can be generated with doxygen as the code has been commented extensively with doxygen styled comments.

3 Complexity Analysis and Experiments

3.1 Experiment data

Firstly we present the results of the experiments we ran:

Sparse Graph Experiment(P=0.2)

Nodes	Edges	Number of Messages
20	39	198
50	252	1016
100	966	3381
150	2248	6646
200	3954	9954
250	6255	16049
300	8888	22054
350	12269	30942
400	15969	37643

Dense Graph Experiment(P=0.8)

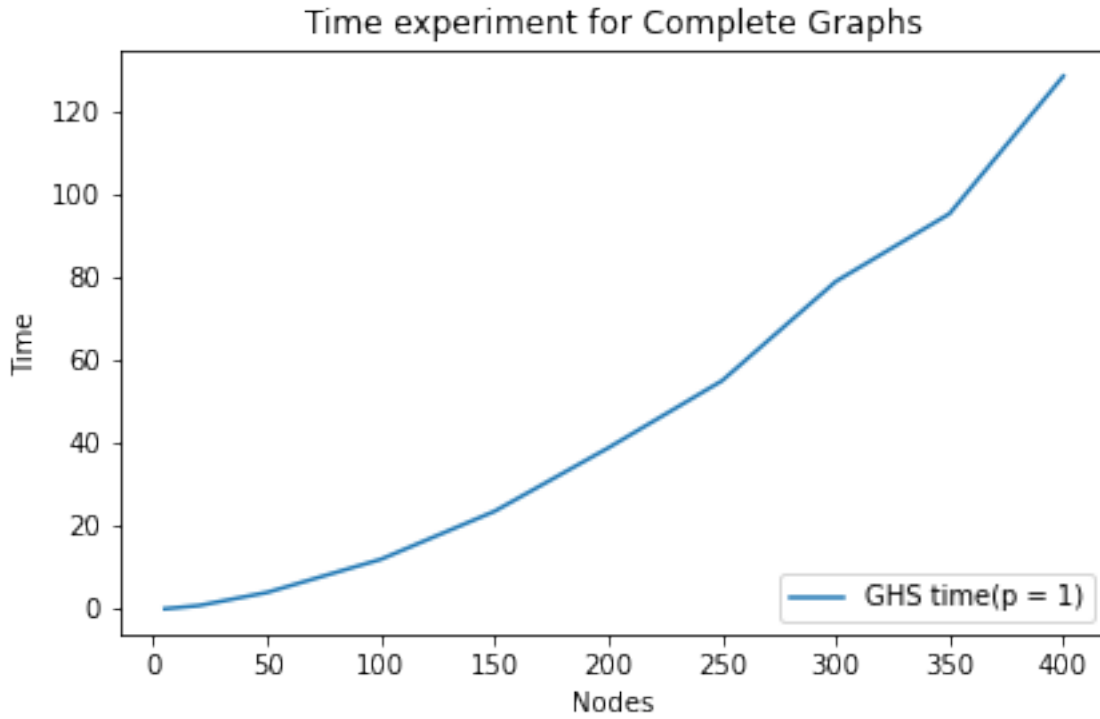
Nodes	Edges	Number of Messages
20	148	505
50	964	2436
100	4011	9049
150	9012	19525
200	15819	34471
250	24826	53218
300	35860	76051
350	48923	102848
400	63838	133375

We also took samples of running time for our code for complete graphs. However analysis of this is much more complicated as it involves major multi-threading overhead. Therefore we merely present the data here :

Runtime Experiment on Complete Graphs(P=1)

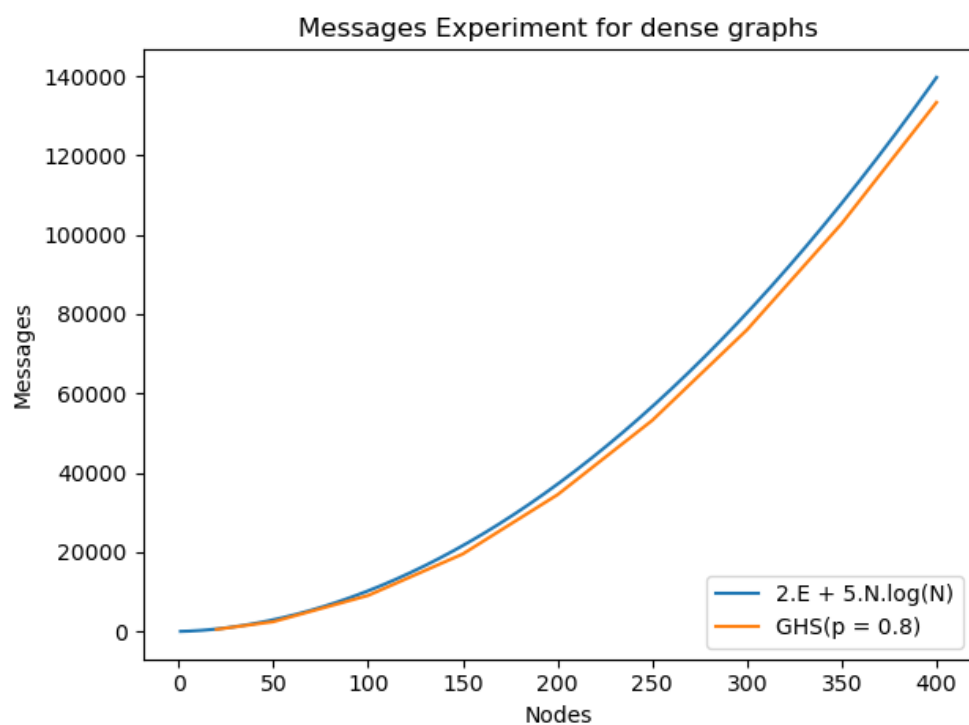
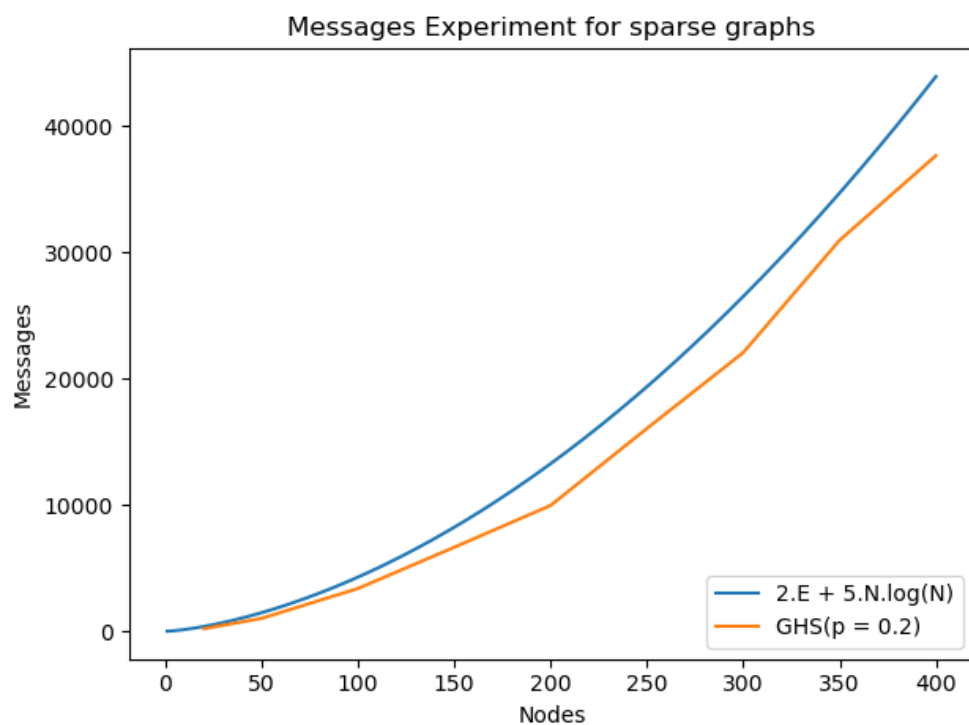
Nodes	Running time
5	0.000503
10	0.242961
20	0.770877
50	3.93151
100	11.9101
150	23.478
200	38.7509
250	54.9079
300	78.804
350	95.1269
400	128.236

Plot for the above data :



3.2 Complexity Analysis

Plots for the results on number of messages as given in the previous section are as follows :



In these plots

- The blue line represents the function : $2 \cdot E + 5 \cdot N \cdot \log_2(N)$ with N = number of nodes, E = number of Edges.
- Here roughly $E = \binom{N}{2} \cdot P$, where P is the probability that there is an edge between any two given nodes.
- The orange line represents the number of messages sent by all the nodes in our implementation.

Analysis :

- Clearly the orange line lies below the blue line for both, sparse and dense graphs.
- As probability increases, the gap between the two lines decreases.
- The blue line is therefore a tight upper bound on the total number of messages in GHS Algorithm.
- Therefore the total number of messages sent in the GHS Algorithm grows as $2 \cdot E + 5 \cdot N \cdot \log_2(N)$