# Day 2 - 1/19/2024

### **Review**

## **Formal Languages**

A formal language is a set of strings. An alphabet is the set of characters that may appear in the strings.

A language over  $\{0, 1\}$ :

$$\{\varepsilon, 01, 0101, 010101, \ldots\}$$

## **Describing**

Two ways to describe a language:

- 1. Generator: Can produce all strings in a language
- 2. Recognizer: Can recognize all strings in a language

### **Grammars**

A grammar is a list of one or more productions (a string substitution rule).

Example:

1. 
$$S \rightarrow xSy$$

$$2. \quad S \to a$$

3. 
$$S \to \varepsilon$$

## Usage

- 1. Start with start symbol (S)
- 2. Apply production  $(\rightarrow)$
- 3. Go until no more *nonterminals* (uppercase letters)

Example: derive xxxyyy

$$S \to \overbrace{xSy}^1 \to \overbrace{xxSyy}^1 \to \overbrace{xxxSyyy}^1 \to \overbrace{xxxyyy}^3$$

Example: What language does this generate?

$$\{\varepsilon, a, xy, xxyy, ..., xay, xxayy, ...\}$$

Example: Based on Grammar D, write a derivation for xy.

#### Grammar D

- 1.  $S \rightarrow AB$
- $2. \quad A \to x$
- 3.  $B \rightarrow y$

$$S \to \overbrace{AB}^1 \to \overbrace{xB}^2 \to \overbrace{xy}^3$$

### **Exercises**

- 7. Write the grammar that generates the language:  $\{ab, abb, abbb, ...\}$ 
  - 1.  $S \rightarrow abX$
  - 2.  $X \to Xb$
  - 3.  $X \to \varepsilon$
- 8. How can you rewrite the grammar from 7. to also generate the string "a"?

Remove the b from the first production.

# The Chomsky Hierarchy

Includes 4 nested categories of languages, types 3, 2, 1, and 0.

# Regular Language

A regular language is one that can generated with a grammar that has one of:

- 1.  $A \to \varepsilon$
- $2. \quad A \rightarrow b$
- 3.  $A \rightarrow bC$

## **Context-Free Language**

A context-free language is one that can be generated from a grammar where each left-hand production consists of a single nonterminal.

$$n. \quad A \to [\text{anything}]$$

## **Context-Sensitive Language**

We don't care about context-sensitive languages.

## **Computably Enumerable Language**

A computably enumerable language is one that can be described by a grammar.

The recognizer is a **Turing machine**.

## Regular Languages

Regular languages have two important applications:

- 1. In most PLs, the set of all lexems of a particular kind forms a regular language.
- 2. He skipped the rest of the slide.

#### Grammar

As stated above:

1. 
$$A \to \varepsilon$$

$$2. \quad A \to b$$

3. 
$$A \rightarrow bC$$

Here's an example:

1. 
$$S \to \varepsilon$$

$$2. \quad S \to t$$

$$3. \quad SxB$$

4. 
$$ByS$$

What language does this grammar generate?

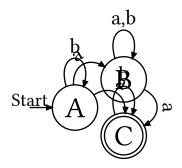
$$\{\varepsilon, t, xyt, \ldots\}$$

More examples occurred.

## Finite Automata

A deterministic finite automata (DFA) is a kind of recognizer for regular languages.

A DFA consists of a bunch of states and a translation between those states.



Incredible, I know.

Languages that are recognized by DFAs are regular languages.

Exercise: I won't draw the diagram but what language is recognized by it?

0 is an even number