Secure Serverless Computing using Dynamic Information Flow Control

Kalev Alpernas, Cormac Flanagan, Sadjad Fouladi, Leonid Ryzhyk, Mooly Sagiv, Thomas Schmitz, and Keith Winstein















Forbe

Strong Cloud Growth Continues to Drive Microsoft Higher

Growth was strong across all segments, but Azure was still the highlight.

Danny Vena (TMFLifelsGood) Oct 28, 2018 at 10:37AM

Over the past several years, **Microsoft** (NASDAQ:MSFT) has accomplished an amazing feat, transforming from a shrink-wrapped business software provider to one of the biggest names in cloud computing. The company has been successful challenging the cloud leader, **Amazon** Web Services, and making meaningful headway, currently ranking in a strong second place.

As the global digital transformation gains steam, investors were watching closely for signs that Microsoft's cloud momentum would continue. When the company reported the financial results of its fiscal 2019 first quarter, which ended Sept. 30, 2018, Microsoft showed that it has much more growth in the tank, producing a record first quarter.



AUTHOR



Danny Vena (TMFLifeIsGood)

Daniel W. Vena, CPA, CGMA is long-term investor searching for intangibles that provide explosive growth opportunities in his investments. He served on active duty with the US Army and has a Bachelors degree in accounting.



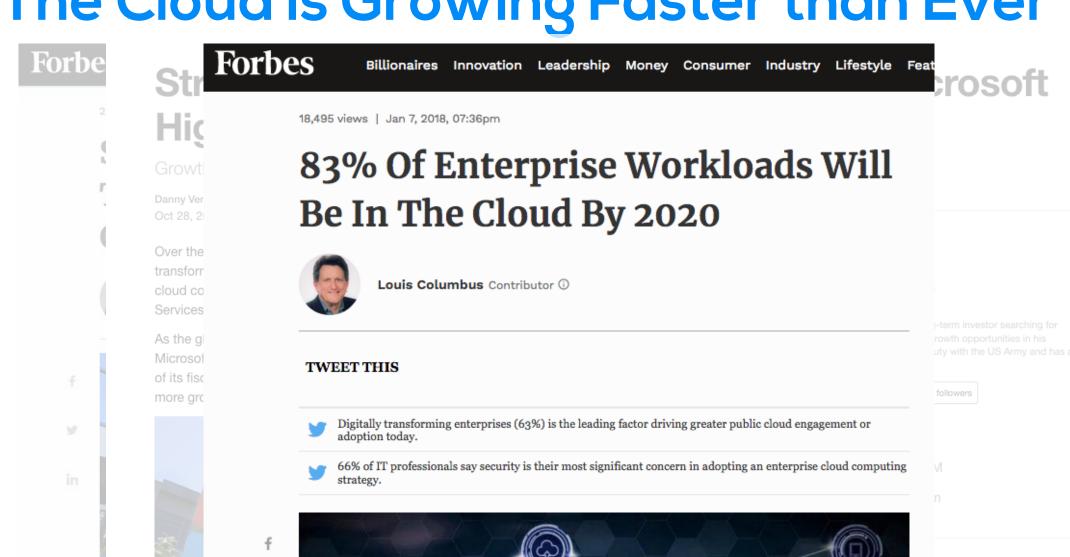
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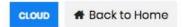
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Is security now a driver of cloud adoption?

JAMES NUNNS EDITOR





Markets Tech Media Success Perspectives Video

U.S. Edition + =



Facebook just had its worst hack ever — and it could get worse

By Donie O'Sullivan, CNN Business

Updated 9:22 AM ET, Thu October 4, 2018









New York (CNN) – On Sunday, September 16, engineers at Facebook detected some unusual activity on the social media platform's networks. It was an attack, the <u>biggest security breach</u> in Facebook's history. And it would take the company 11 more days to stop it.

Now, almost a week since the public was first told of the attack, we still barely know anything about what happened.

We don't know who the hackers were, or what they were looking for. We don't know whether they were targeting particular people in certain countries. We don't know how long they had access to users' information. And we don't know what, if anything, they took.





Markets Tech Media Success Perspectives Video

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Fac

Cathay Pacific got hacked, compromising the data of millions of passengers

By Donie (

Updated 9

By Jethro Mullen, CNN Business

Updated 12:22 AM ET, Thu October 25, 2018









New York

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Now, almabout wh

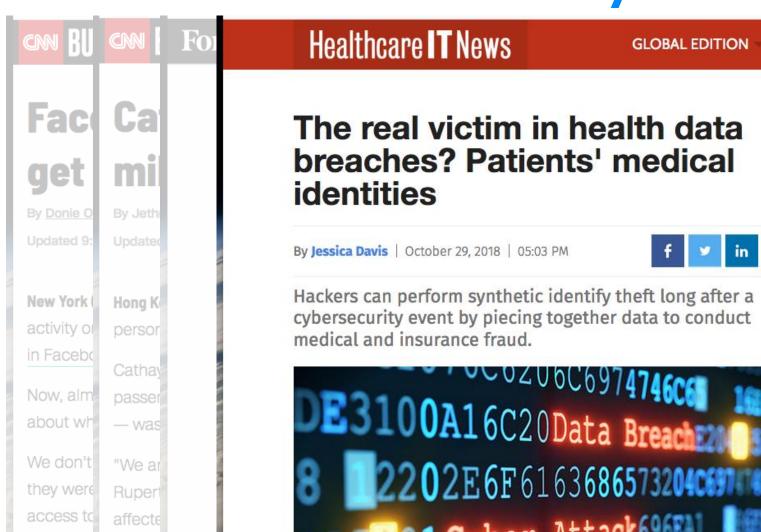
We don't

Hong Kong (CNN Business) – One of Asia's top airlines has discovered a data breach in which the personal information of more than 9 million passengers may have been stolen.

Cathay Pacific (CPCAY) said late Wednesday that a wide range of data — including passengers' names, dates of birth, phone numbers, email addresses and passport numbers — was exposed in a hack of its information systems earlier this year.

"We are very sorry for any concern this data security event may cause our passengers," CEO Rupert Hogg said in a statement. The Hong Kong-based carrier is in the process of contacting affected people, he added.





The real victim in health data breaches? Patients' medical





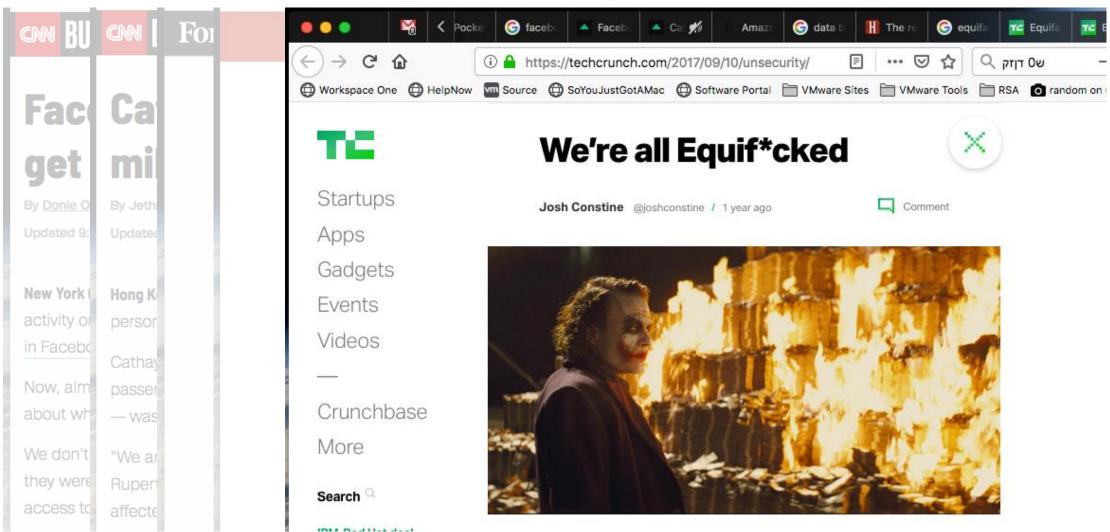
GLOBAL EDITION TOPICS



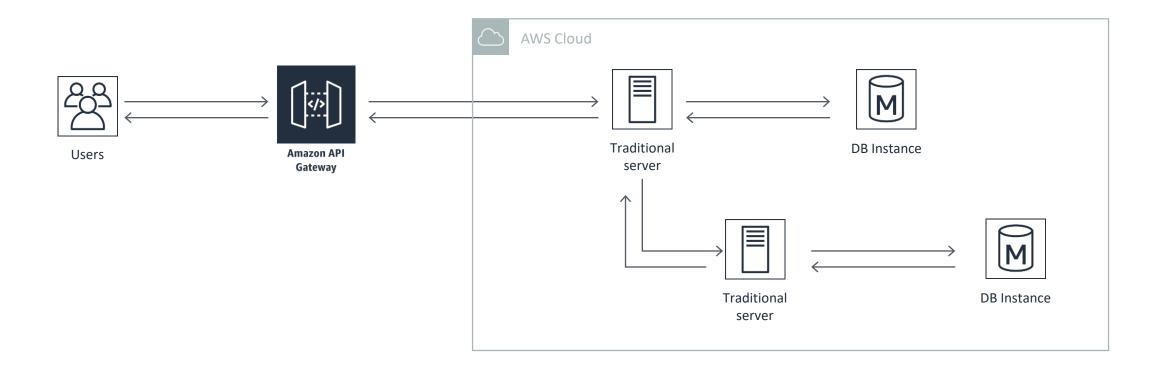


SIGN UP

MAIN MENU

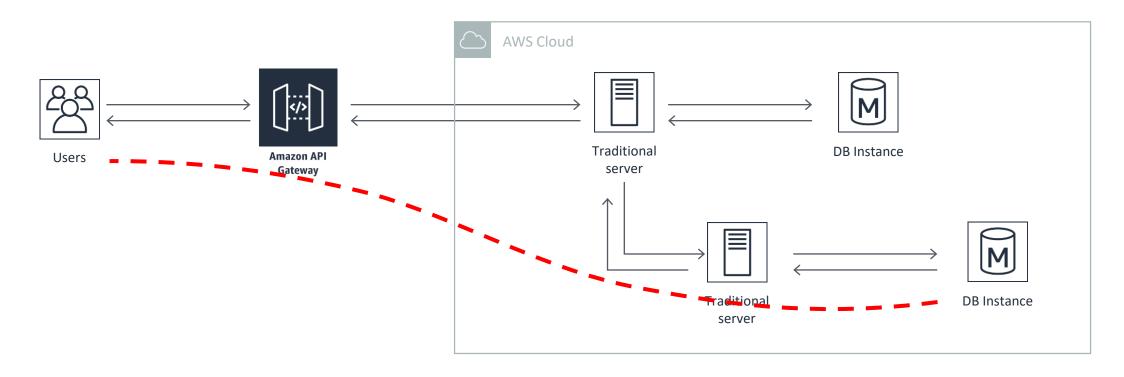


Traditional (monolithic) three tiered application

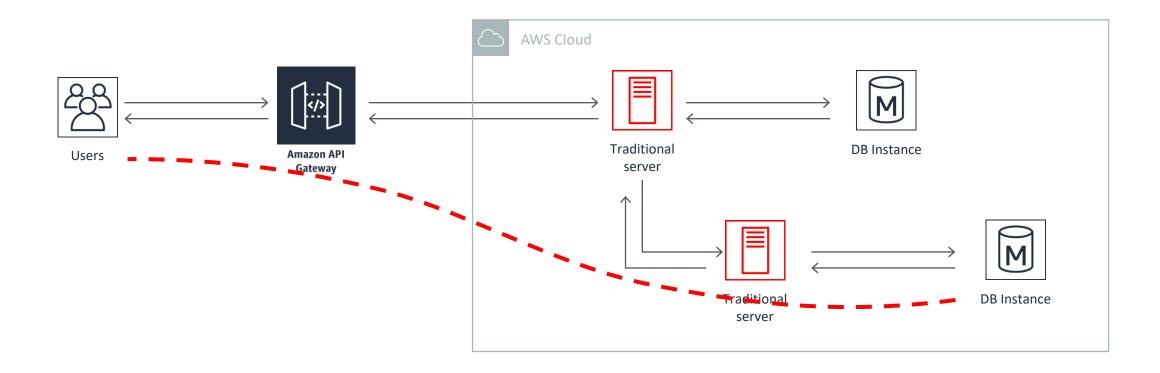


Goal: keep data secure – Information Flow Control (IFC)

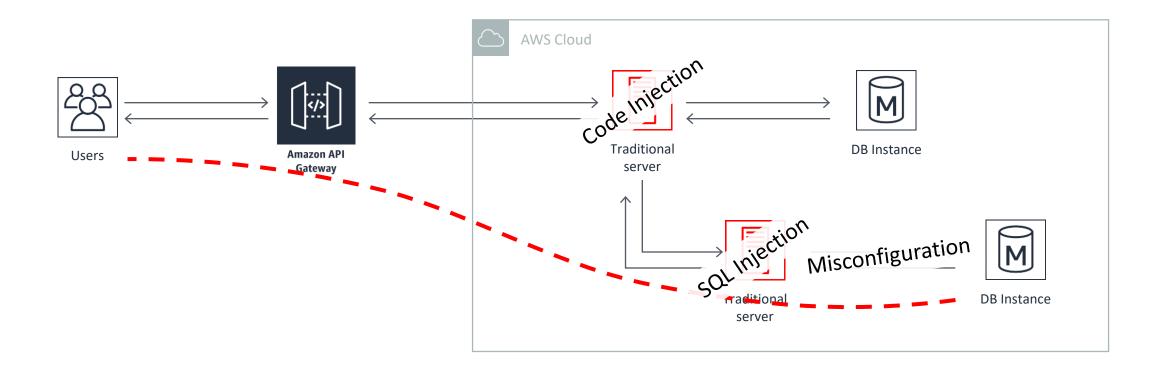
Non-Interference



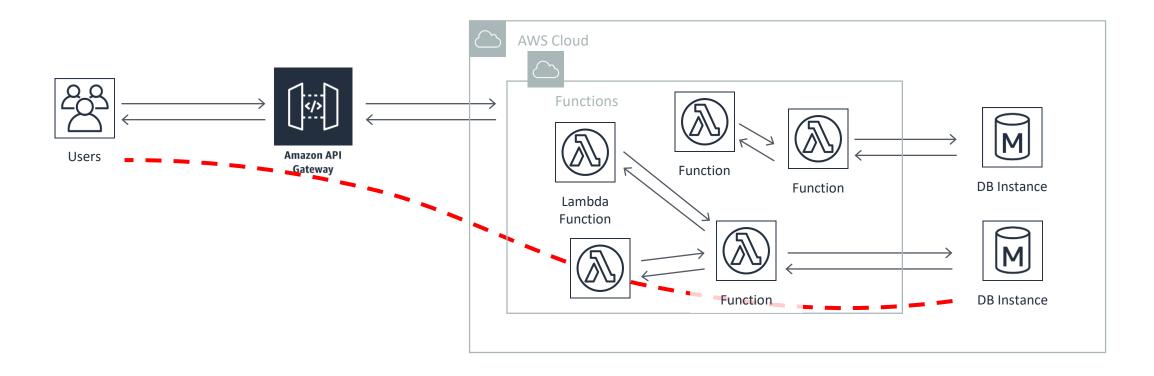
Servers are vulnerable



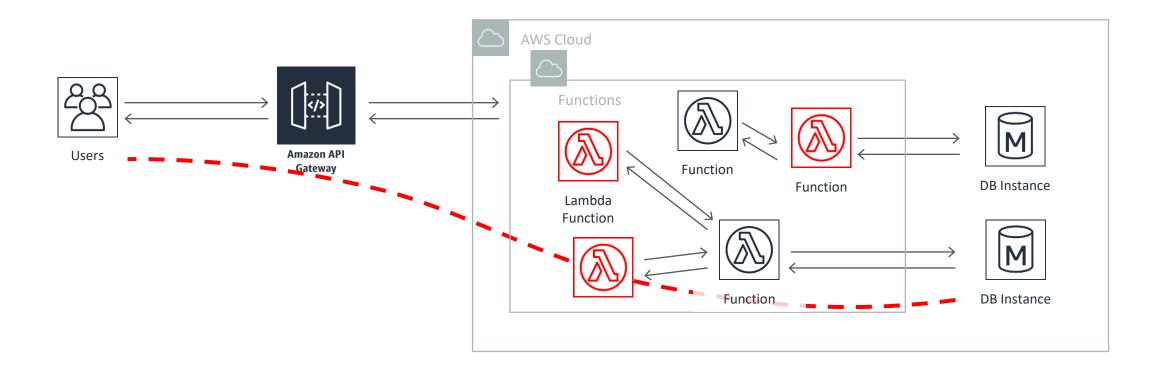
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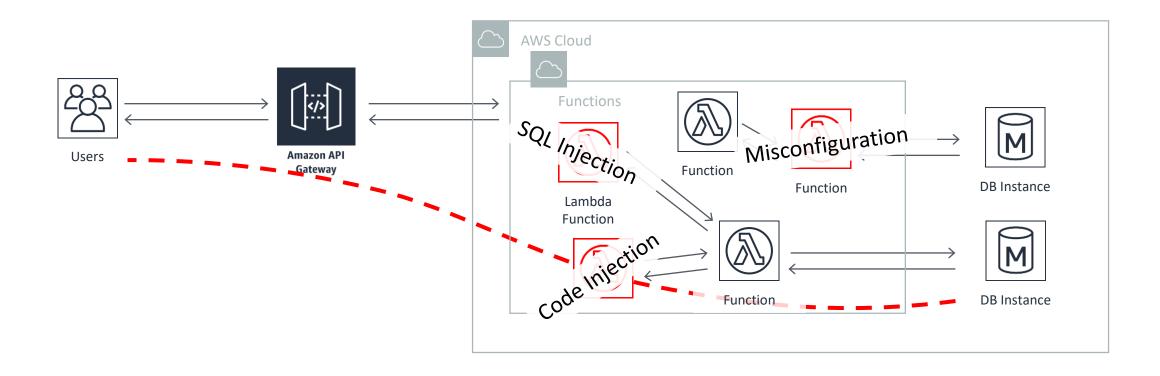
What happens when we replace them with serverless functions?



Functions are vulnerable

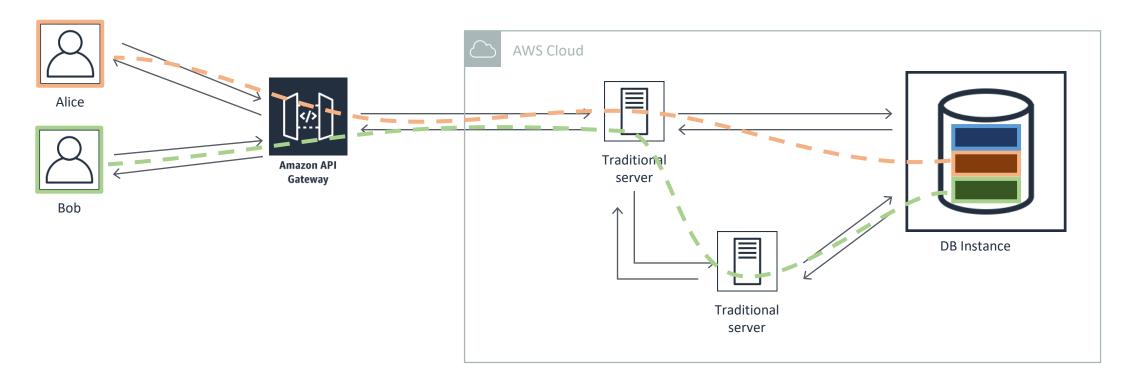


Polyglot, distributed environment makes things even harder



Monolithic Applications

Long running processes, serving multiple users, managing complex local state



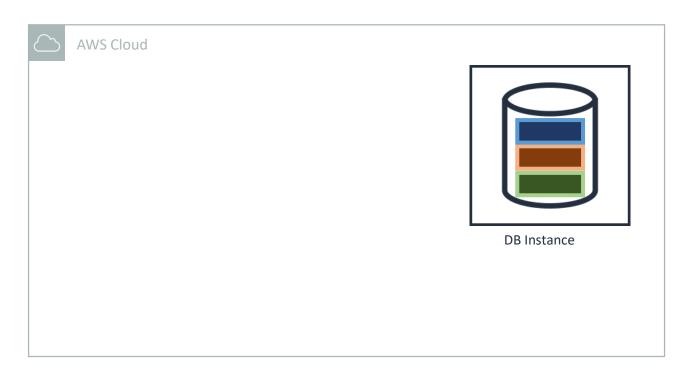


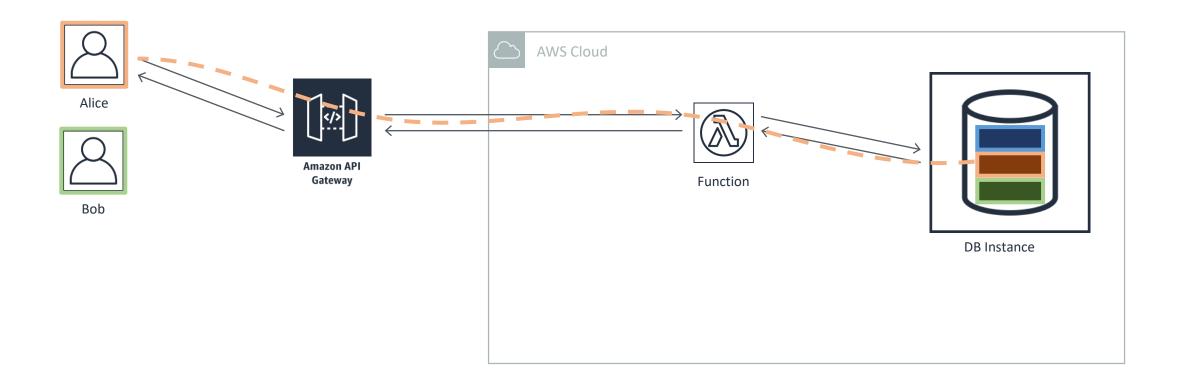
Alice



Bob







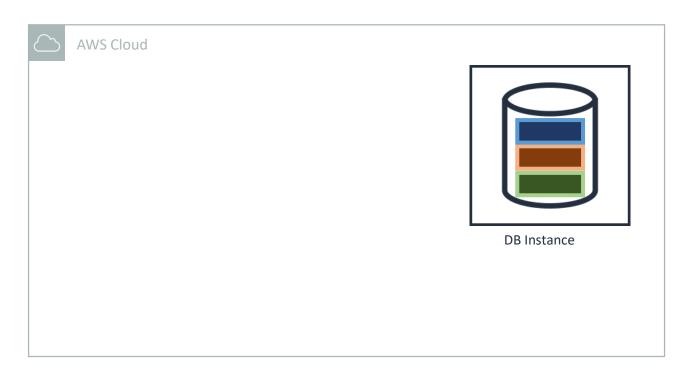


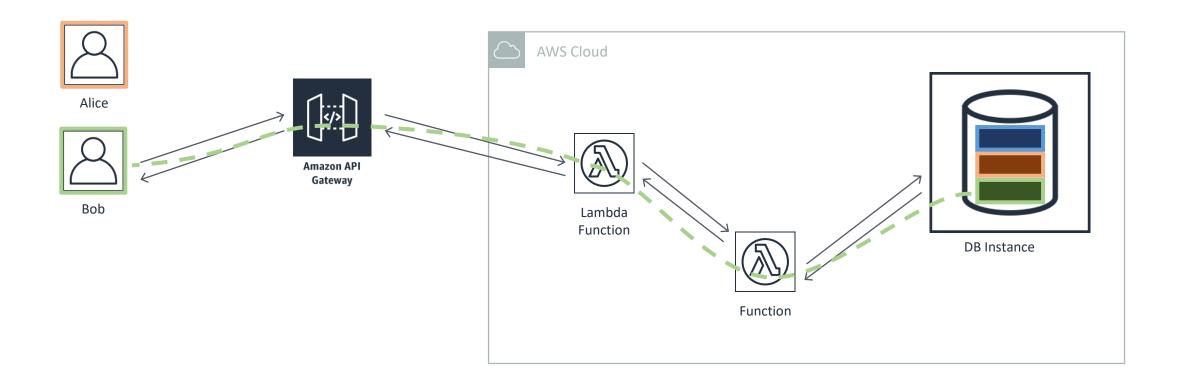
Alice



Bob







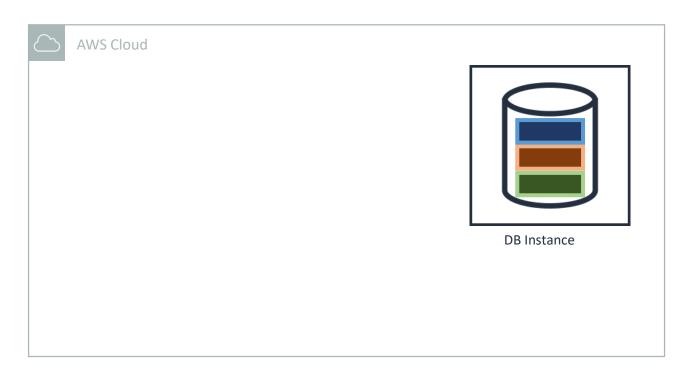


Alice



Bob





Underneath Everything, there is Hope

Serverless computing is our opportunity to implement efficient IFC

Perform IFC tracking at **function boundaries**



How do we exploit Serverless to design an efficient IFC system?

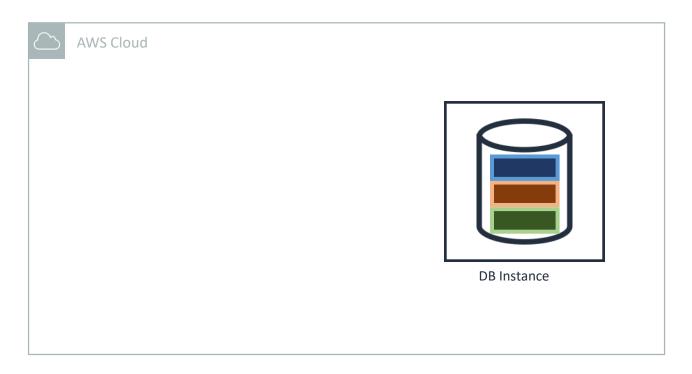


Alice

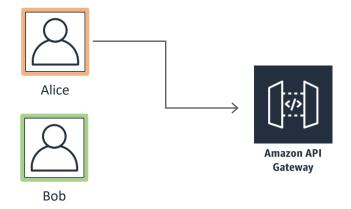


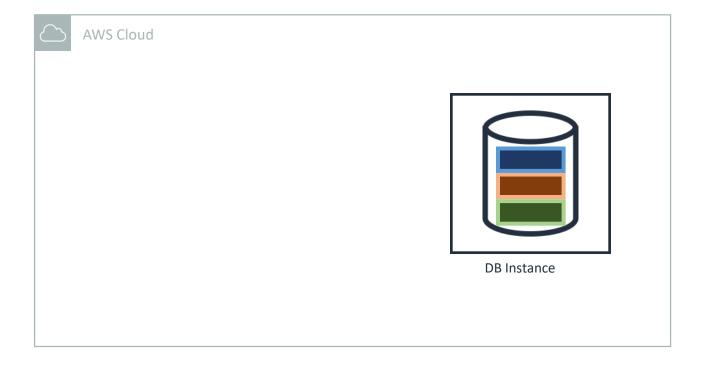
Bob



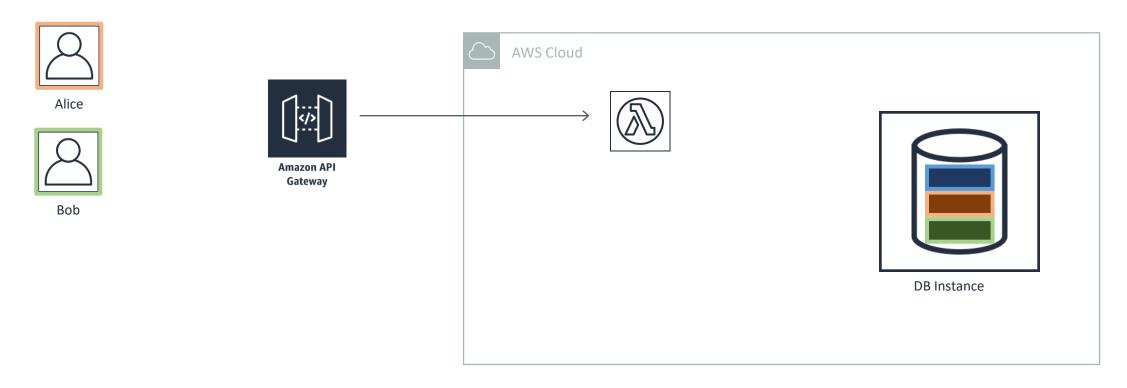


Alice makes a request





The Serverless framework spawns a function to serve the request



Functions may spawn additional functions

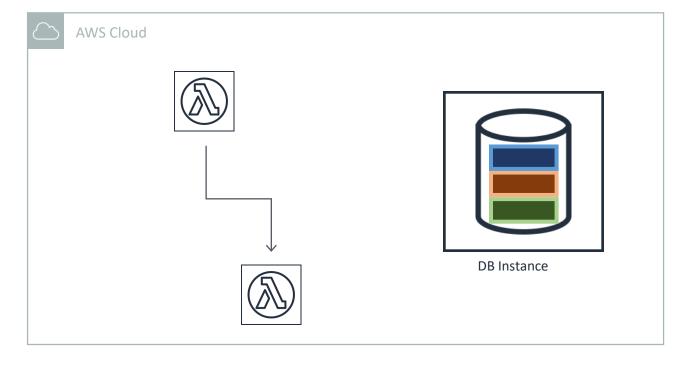


Alice



Bob





Functions interact with DB

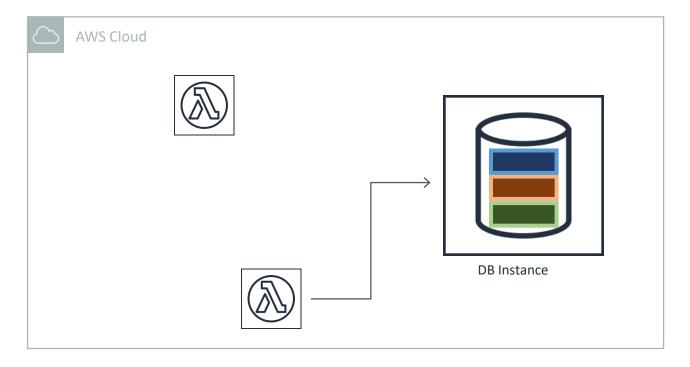


Alice



Bob





Function reads labeled data and is marked with the label

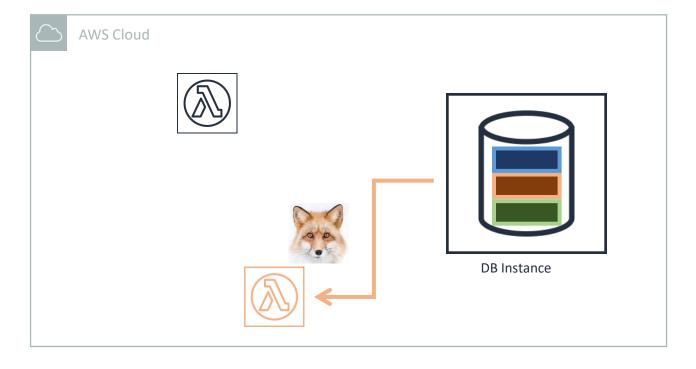


Alice



Bob





Label is propagated across function calls

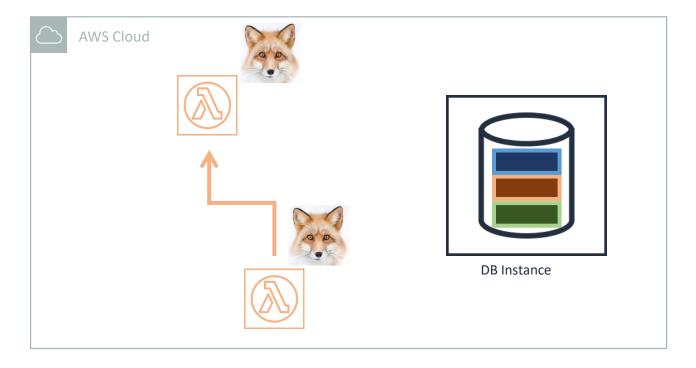


Alice

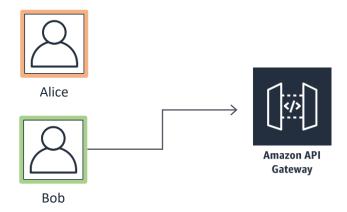


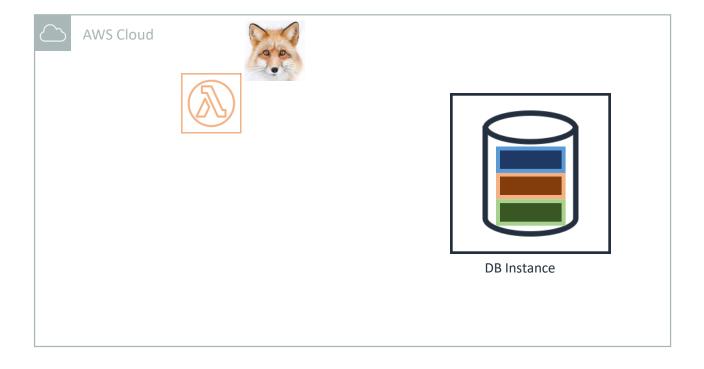
Bob



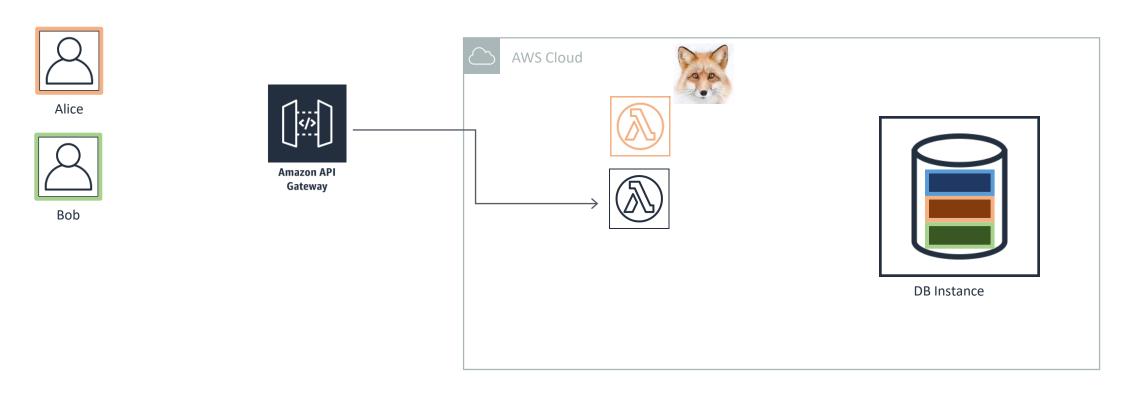


Other requests may occur in parallel





Parallel requests spawn new functions, that are independent from one another



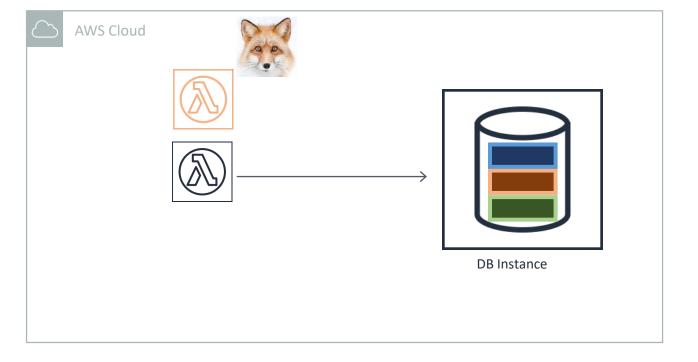


Alice



Bob





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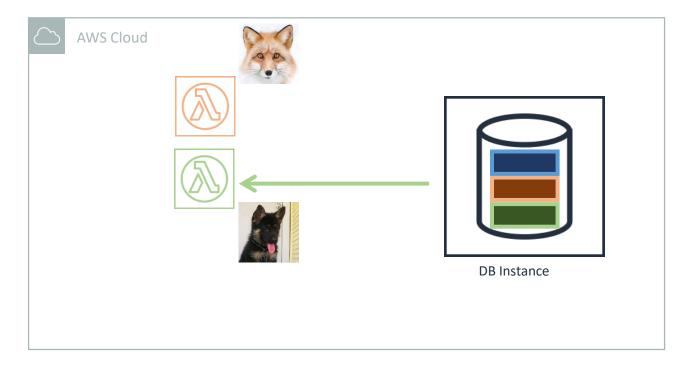


Alice

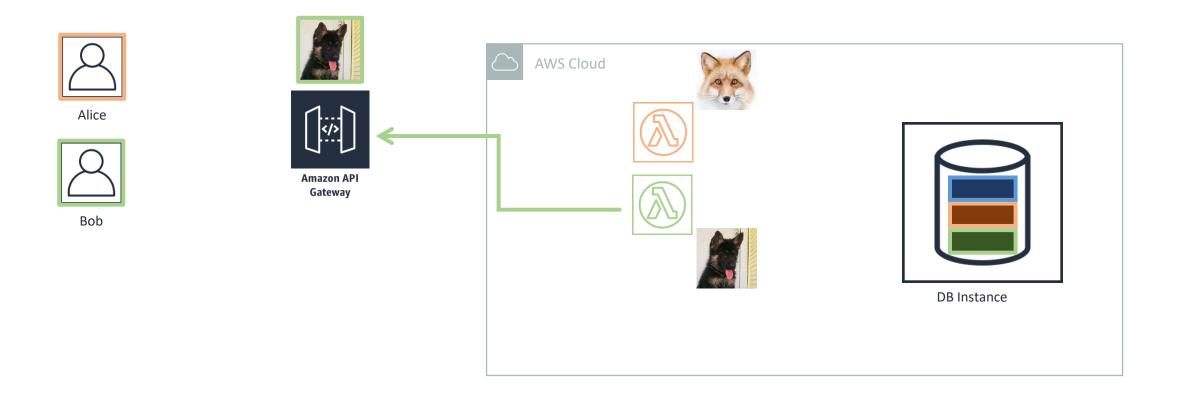


Bob





Function returns result



Before the result is sent to the user, we can check that the security policy is not violated

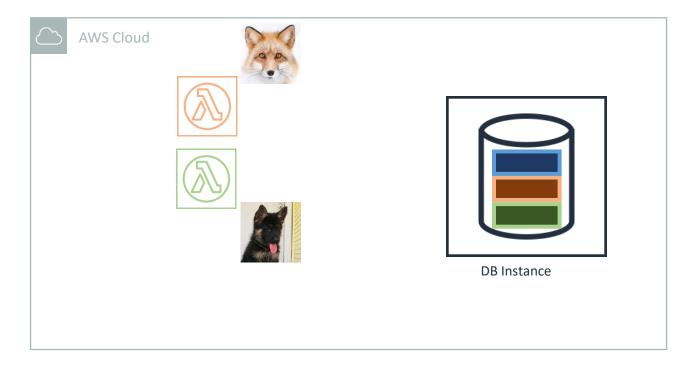


Alice

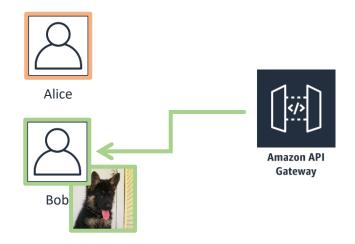


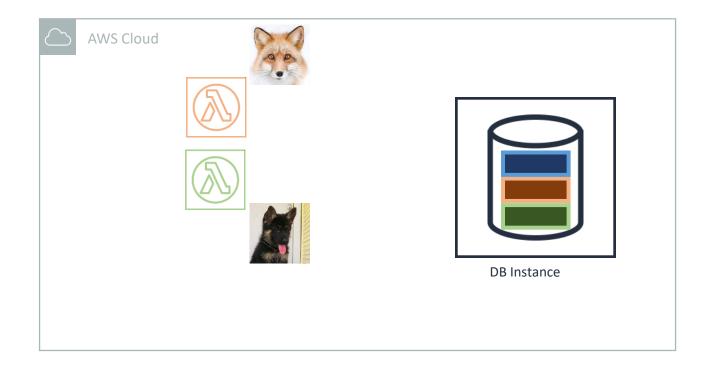
Bob





And send the response if it does not violate the security policy

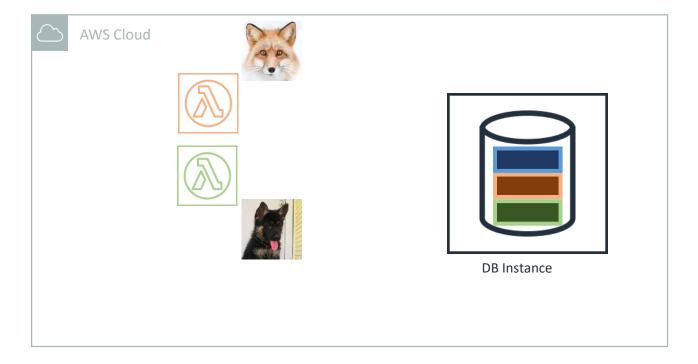




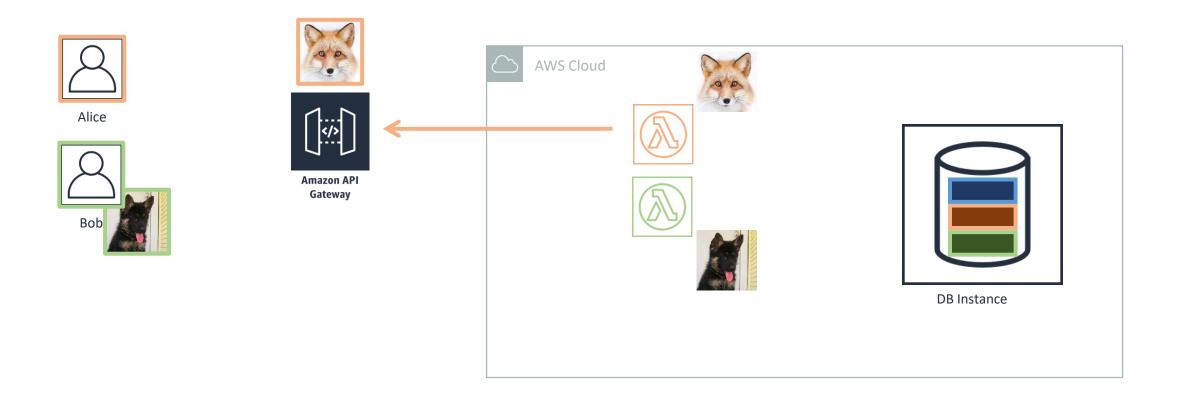








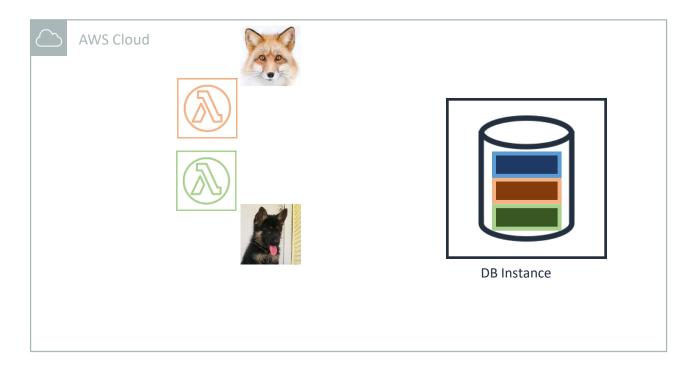
Function returns result



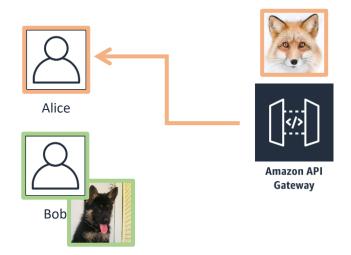
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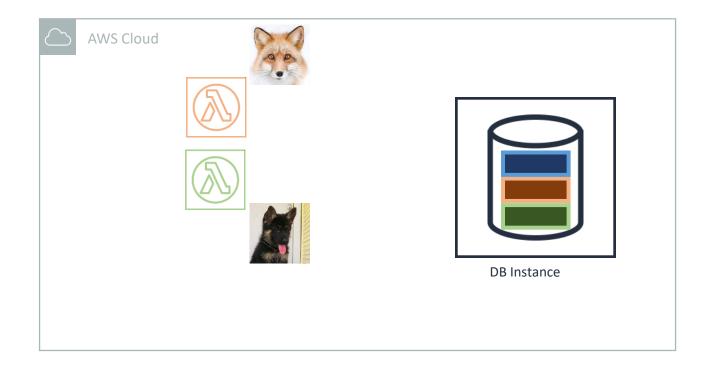






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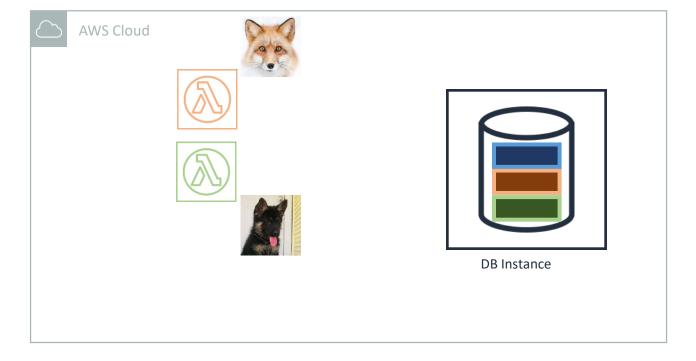




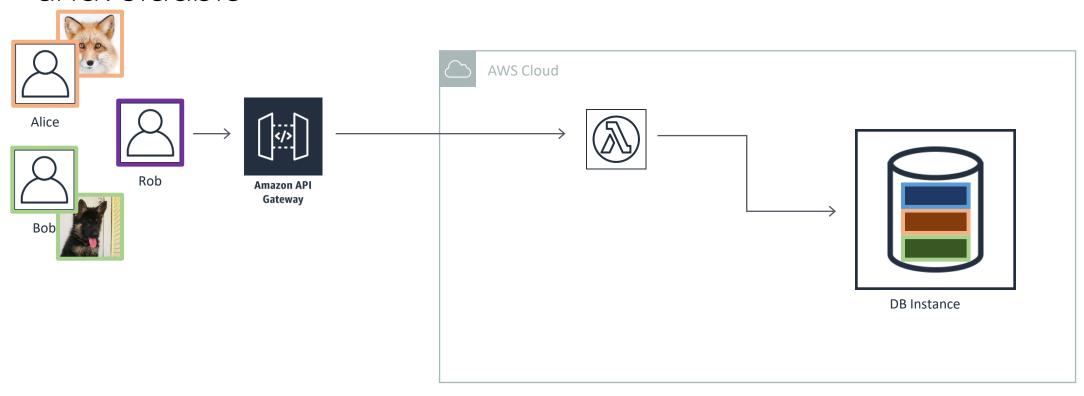




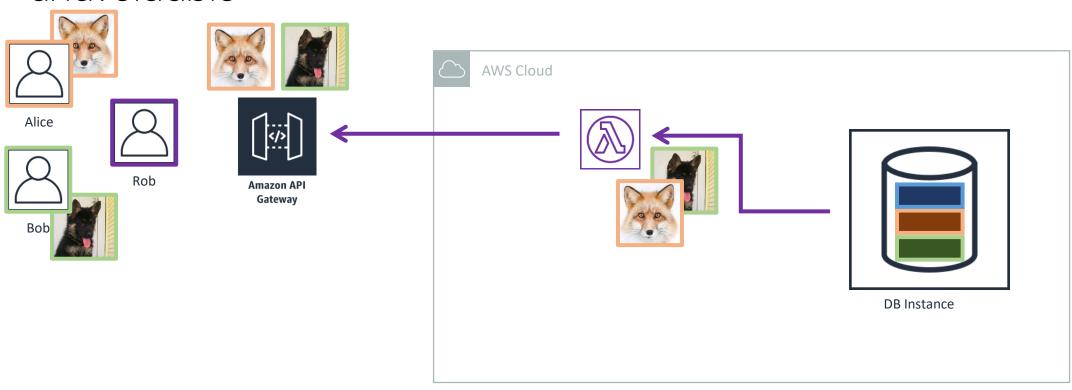




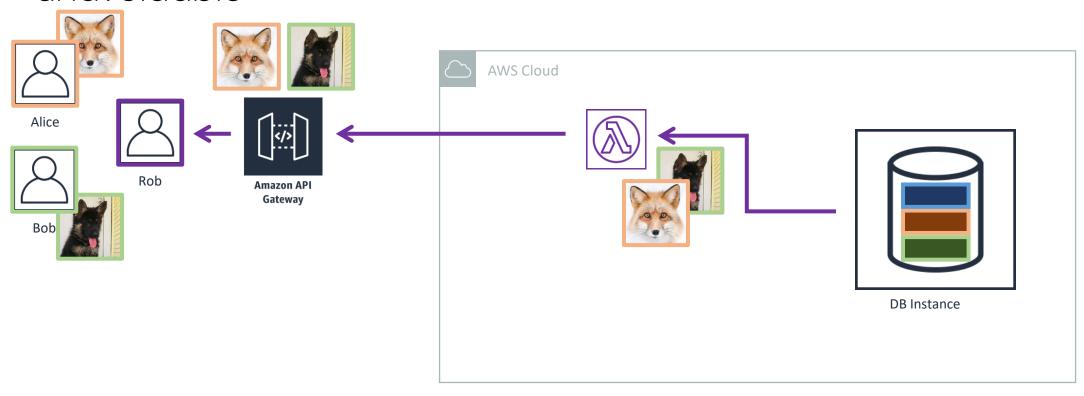
Security policy might have a hierarchy – making shared dbs unavoidable



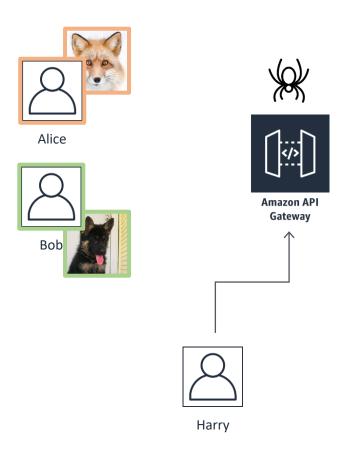
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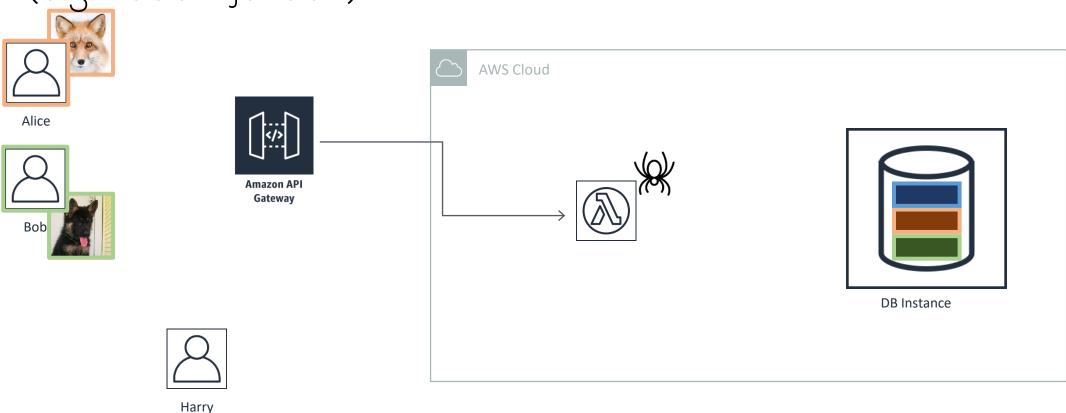


What happens when a hacker tries to steal data?





The hacker might exploit some vulnerability in the application (e.g. code injection)



Trapeze guarantees security even if the underlying function code is completely compromised



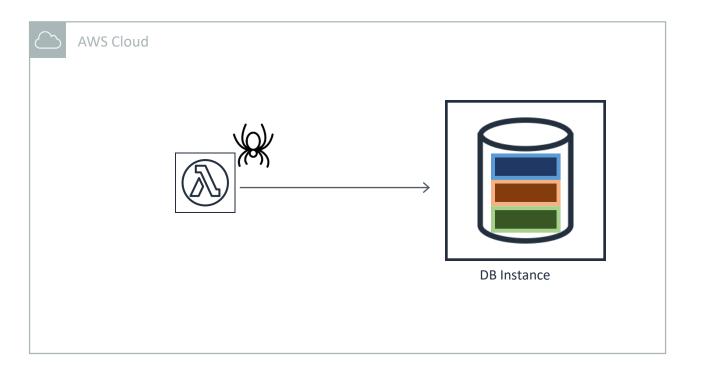
Alice







Harry



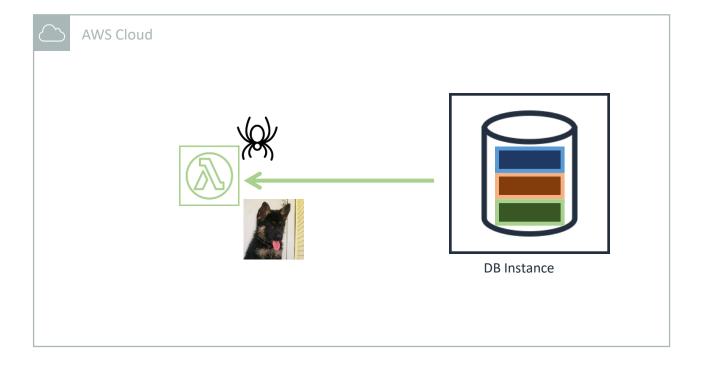
The malicious code may read a secret, and is marked with a label



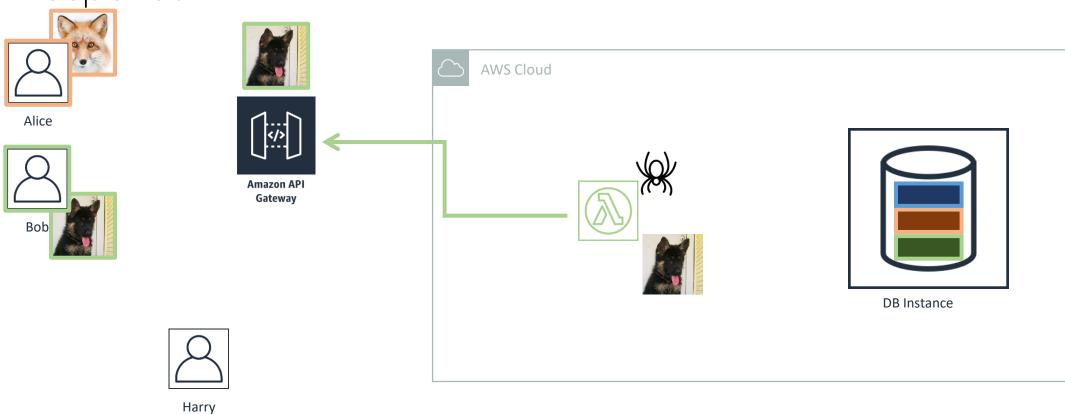








The malicious function then tries to send the secret in a response



At which point we can check if the hacker is allowed to see the data

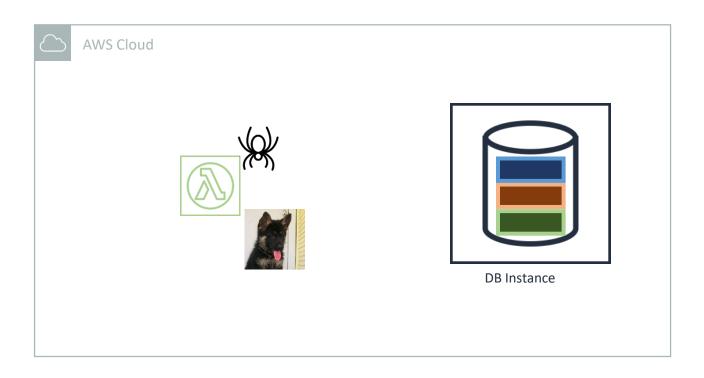












And block the request from being sent



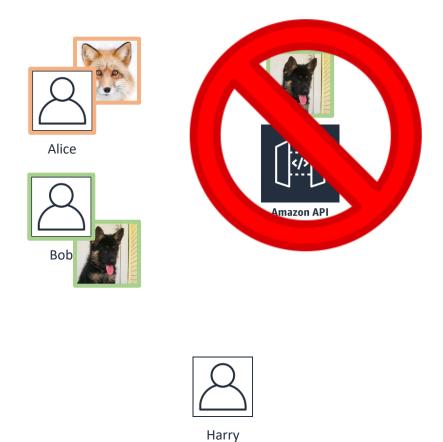


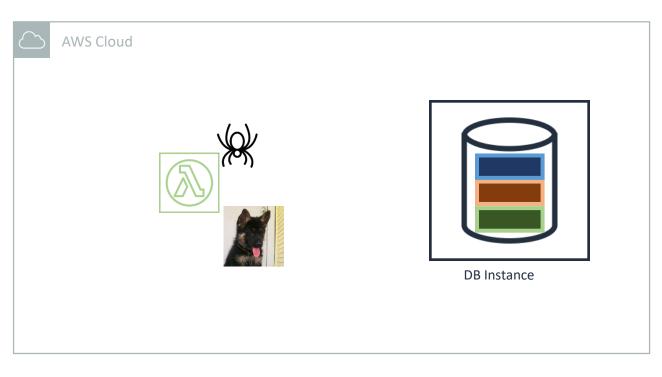




Harry

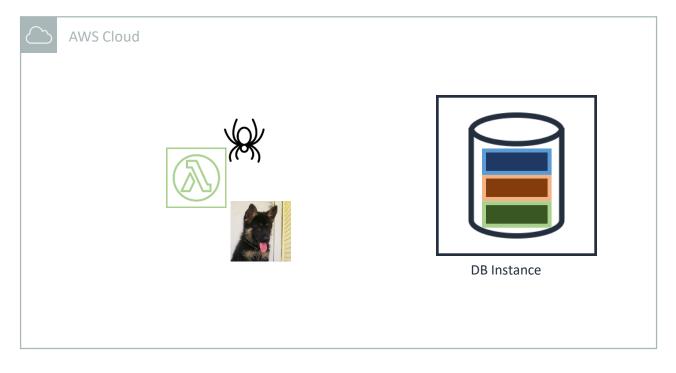
But - the attacker learned of the existence of secret data





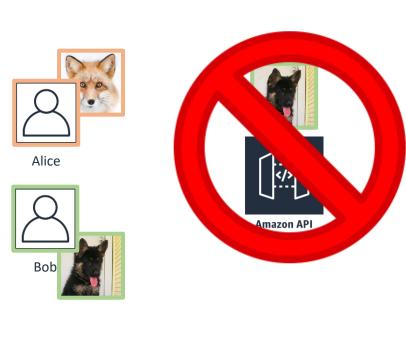
Which is bad if the existence of a secret is itself a secret







Termination Insensitive Non-Interference is insufficient!

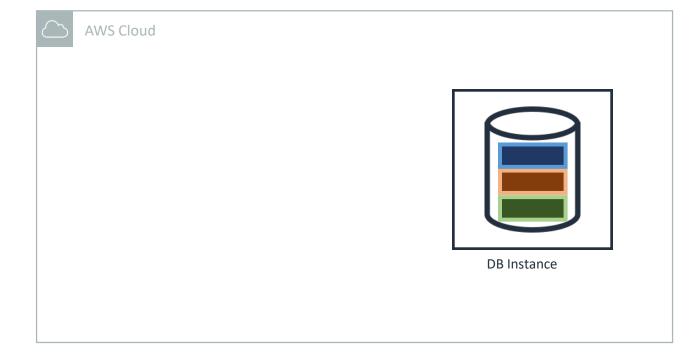






Attacker can encode a secret value by making certain values 'inaccessible'





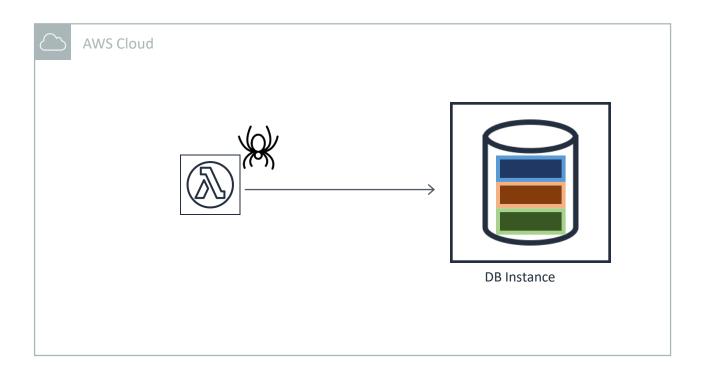


Harry

Attacker reads a secret value



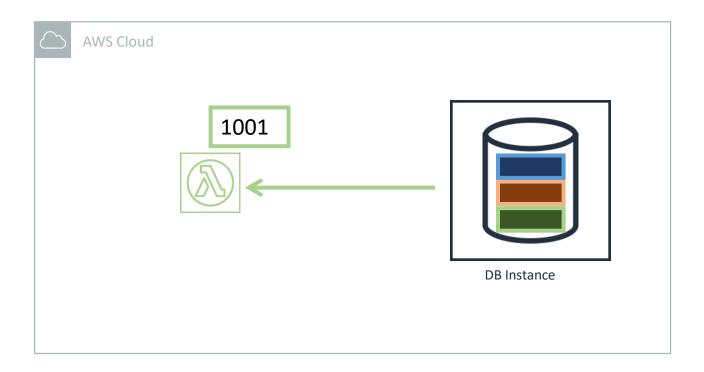




Attacker reads a secret value (1001)



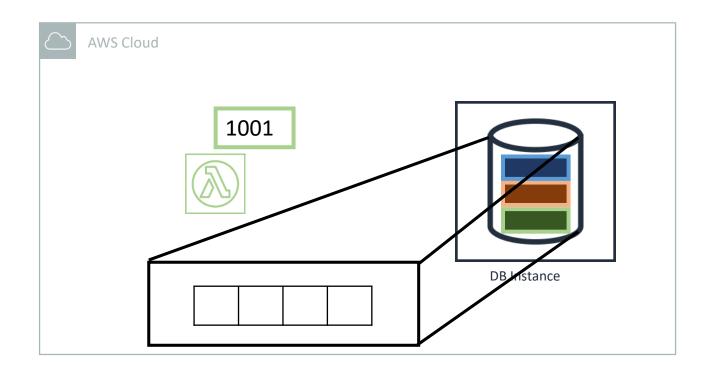




Write something to predefined locations



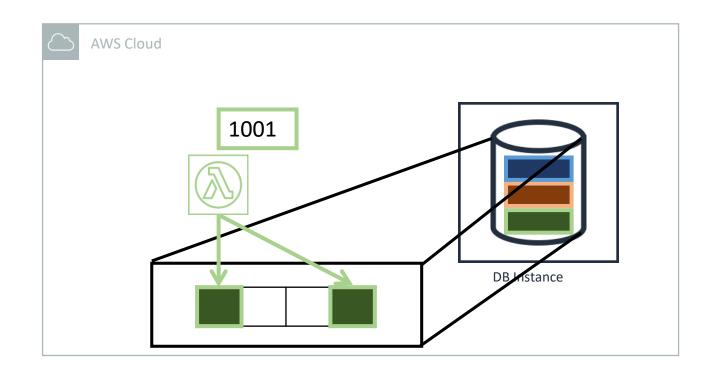




Encode secret as a memory storage pattern



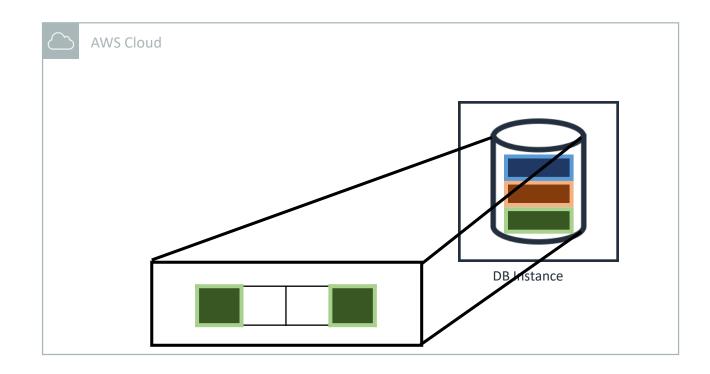




Attacker can leak every bit by accessing each cell separately



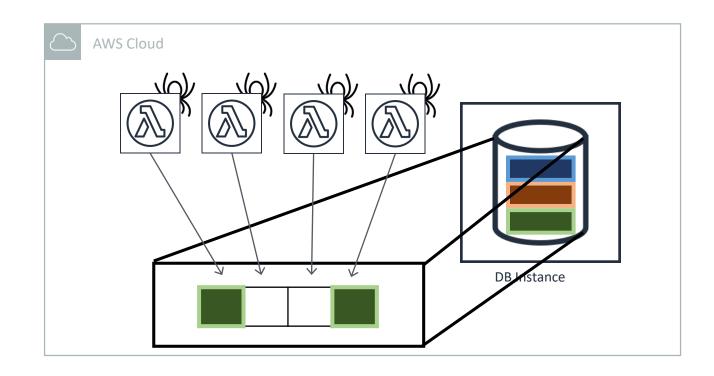




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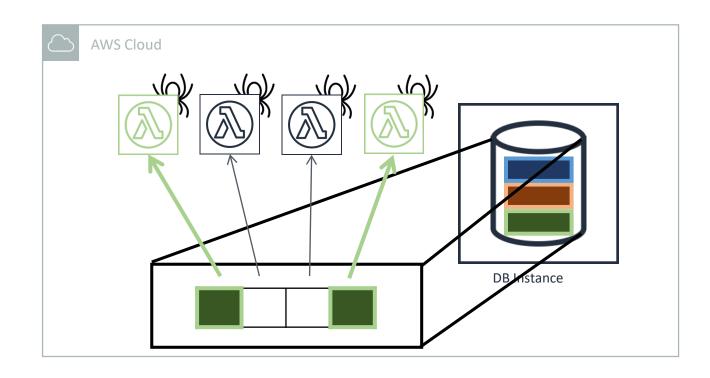




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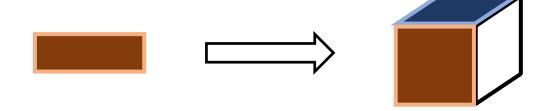




How do we get Termination **Sensitive** Non-Interference?

- Static labeling instead of floating labels
 - Label assigned at function invocation
 - Output behavior depends on invocation, not secret

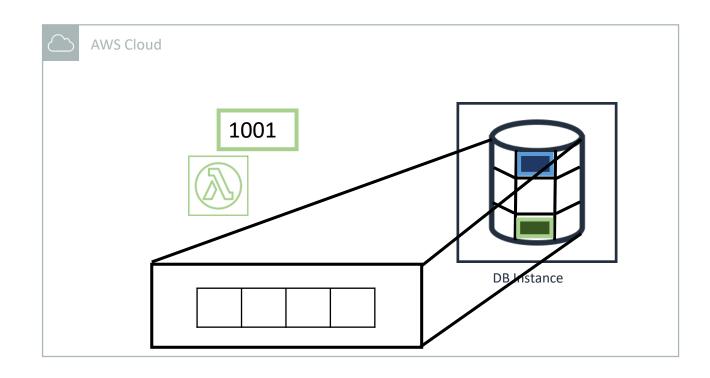
Faceted data store



Modified read and write semantics



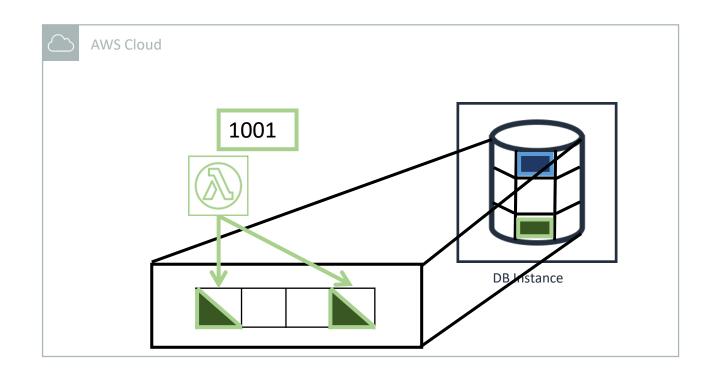




Writes add a facet to the location, instead of overwriting



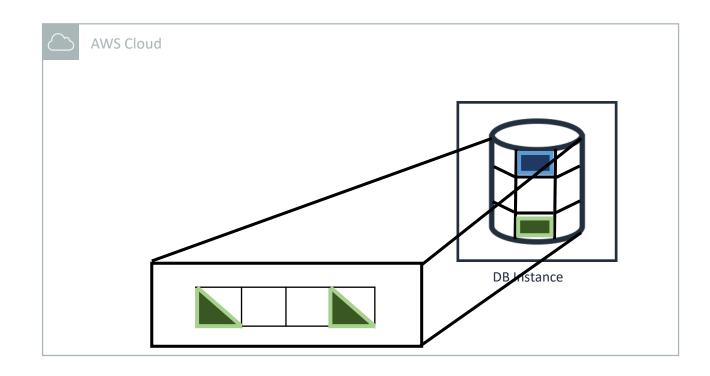




An execution only sees the facets that is allowed to see



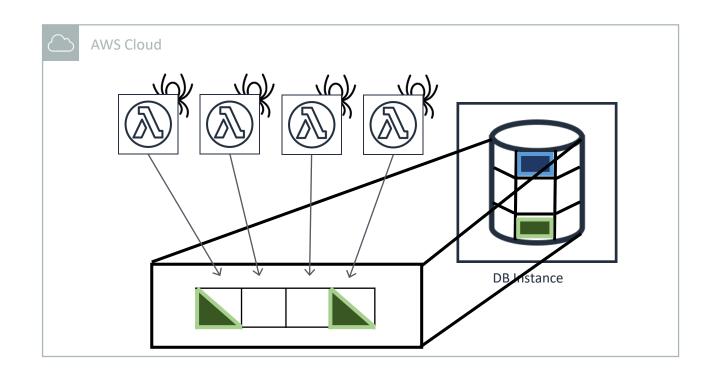




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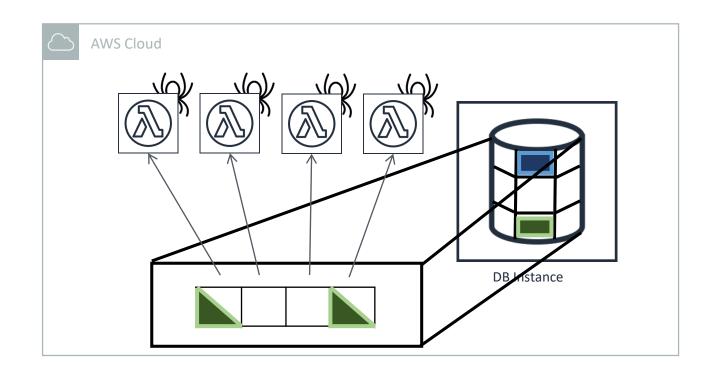




An execution without the correct label can't distinguish between an empty cell, and a faceted cell







Trapeze

- Runtime enforcement of *Termination Sensitive Non-Interference*
 - Attacker cannot leak information by affecting execution termination
- Function executions get a security label when invoked
 - Label matches the security privilege of the invoker
- Faceted data store semantics
 - Items in the DB may have multiple values
 - Depending on the security label of the function that wrote the data

Trapeze - Goals

Practical IFC system for serverless applications

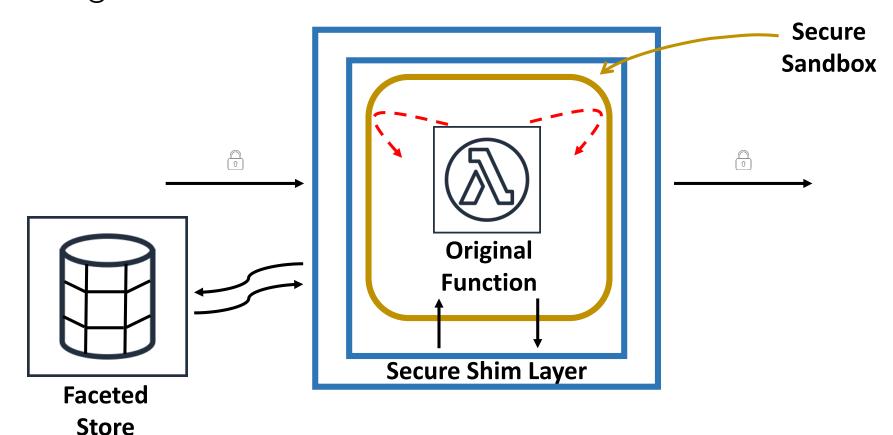
Low overhead in runtime

Allow running existing application securely with minimal modification

 Transparency – secure applications have the same semantics as when running without Trapeze

Trapeze Architecture

Shim layer wraps functions, controlling communication and enforcing faceted store semantics



Evaluation

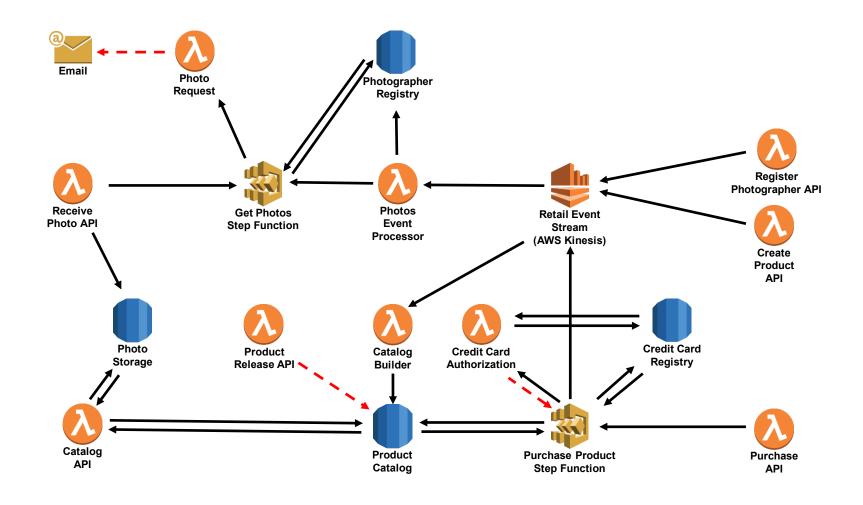
Test cases:

- Nordstrom Hello, Retail!
- 99
- Image feature extraction

Questions:

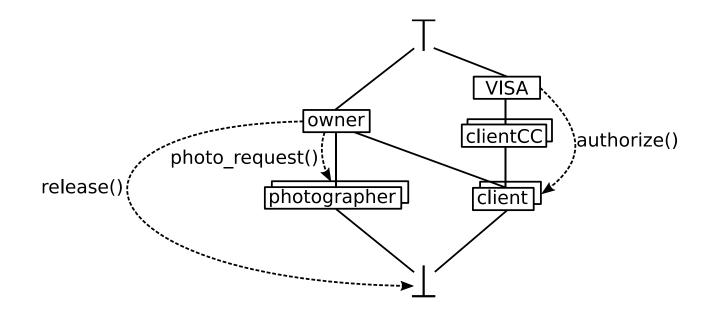
- Performance
- Ease of deployment
- Transparency

Evaluation - Hello, Retail!



Evaluation - Hello, Retail!

Security policy:



Evaluation Results

Scenario	# λ	λ runtime (ms)		Δ(%)	App sada	Modified		
		baseline	Trapeze	Δ(/0)	App. code	Modified		
Hello, Retail!								
Update & Purchase	435	51,246	98,563	92.33	1300	3		
Build & Browse	955	257,799	391,900	52.03				
99								
mosh	111	681,173	654,448	-3.92	- 8800	2		
git	917	2,602,500	2,660,721	2.24				
vim	234	1,242,873	1,338,128	7.66				
openssh	654	1,626,233	1,649,139	1.41				
Image Feature Extraction								
Upload	1	475	525	10.4	95	0		
Extract	1	1,882	2,114	12.3				

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Extract	1	1,882	2,114	12.3				

Summary

- Serverless computing
 - Security Risks
 - Security opportunities
- Trapeze
 - A practical IFC system for serverless applications
 - Guarantees TSNI
 - Low overhead (mostly) thanks to function granularity
 - High transparency

GitHub https://github.com/kalevalp/trapeze