

Lecture 9 : SQL Injections

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`https://dropit.uni.lu/invitations?share=38fb341606a357367b7a&dl=0`

Previously...

Previously... in Lecture 1 (Introduction)

- ▶ Software Development Life-cycle
- ▶ Vulnerability Life-cycle
- ▶ Vulnerability Disclosure

Previously... in Lecture 2 (Buffer Overflow)

- ▶ A buffer on the stack
- ▶ Return address on the stack
- ▶ Overwrite return address
- ▶ Jump to shellcode on the stack

Previously... in Lecture 3 (ROP)

- ▶ NX bit (stack non-executable)
- ▶ Gadgets in already loaded code
- ▶ Chain gadgets (addresses of gadgets and data on the stack)
- ▶ Only data on the stack

Previously... in Lecture 4 (ASLR)

- ▶ Randomize code segment at program start
- ▶ Breaks gadget chains
- ▶ Bypass with information leak (e.g, vulnerability)

Previously... in Lecture 6 (CFI)

- ▶ Mechanism to allow only "intended" paths
- ▶ Binary instrumentation to add IDs
- ▶ Indirect jumps, call, returns check if ID of "destination" is correct
- ▶ Pure software implementation have 20% overhead

Previously... in Lecture 7 (Heap-Overflow)

- ▶ How a heap-overflow can be attacked depends on the heap management implementation
- ▶ The "unlink" attack present in early versions of glibc provides a "write anywhere" primitive to the attacker
- ▶ Recent implementations performs more check to prevent "unlink" based attacks

Previously... in Lecture 8 (Type Confusion)

- ▶ What is it? Manipulation of an object through another object.
- ▶ Consequences? Undefined behavior, hijack control flow.
- ▶ Why it works? No verification at runtime (otherwise runtime and/or memory overhead).

SQL Injection

- ▶ SQL: Structured Query **Language**; language to query relational databases
- ▶ Injection: introduction of code in the existing code of the target

SQL: Structured Query Language

- ▶ `SELECT * from users;`
- ▶ `SELECT * from users WHERE SUBSTRING(username, 0, 1) = "a";`
- ▶ `SELECT COUNT(*) from users WHERE username = "root";`

Injection (ex in C)

```
// example from OWASP  
// https://www.owasp.org/index.php/Command_injection  
  
#include <stdio.h>  
#include <unistd.h>  
  
int main(int argc, char **argv) {  
    char cat[] = "cat ";  
    char *command;  
    size_t commandLength;  
  
    commandLength = strlen(cat) + strlen(argv[1]) + 1;  
    command = (char *) malloc(commandLength);  
    strncpy(command, cat, commandLength);  
    strncat(command, argv[1], (commandLength - strlen(cat)) );  
  
    system(command);  
    return (0);  
}
```

Injection (ex in SQL)

```
<? php
$conn->query("SELECT * from users where username == \"\$username\" and password ==
↳ \"\$password\"");
?>
```

SQL Injection Vulnerability Exploitation

```
<? php
$conn->query("SELECT * from users where username == \"\$username\" and password ==
↳ \"\$password\"");
?>
```

- ▶ Comments in SQL: --
- ▶ What if username is: " OR 1 == 1 --
- ▶ Results in authorisation check bypass!

"Blind" SQL Injection Vulnerability Exploitation

```
<? php
$conn->query("SELECT * from users where username == \"\$username\" and password ==
↳ \"\$password\"");
?>
```

- ▶ How to dump a database?
- ▶ 1. Send queries to the database server through the web page.
- ▶ 2. The webserver might generate a different webpage if the query is successful or not
- ▶ 3. For every query the attacker can deduce a tiny bit of information about the database (ex: one character of the first element of row 3)

Preventing SQL Injection Vulnerabilities

```
<? php

$stmt = $conn->prepare("SELECT * from users where username == \"?\"/>
    ↪ \"?\");
$stmt->bind_param("ss", $un, $pw); /* types and variable bindings */

$un = ... ;
$pw = ... ;
$stmt->execute();
?>
```

- ▶ Use "prepared" statements

- ▶ Code injection attacks enables bypass of authorization checks and/or execution of arbitrary code on the server
- ▶ Consequences: attacker gets access to privileged environment and/or can dump databases
- ▶ Protection include sanitization of the input and/or well defining what is code and what is data

Question?

- ▶ Groups of 2
- ▶ Suggested topics:
 1. Heap exploitation on Debian 3.1
 2. Patch for CVE-2018-20343 (Ricardo, Alex)
 3. Complete exploit for CVE-2018-20343
 4. Study and PoC for CVE-2013-0912 (Chrome type confusion) (Adriano)
 5. Stable code injection through /proc/self/mem
 6. Explanation of a recent exploit targeting webbrowsers (Chrome, Firefox, etc.) (Yurii, Ervin)
 7. Exploitation of a PoC type confusion in C++ (Ihor, Artem)
 8. Break wordpress authentication mechanism.
- ▶ Deliverables: Presentation + Code (PoC)

Projet: Heap exploitation on Debian 3.1

- ▶ Explain the differences in the heap management code from debian 2.2 (lab 7) and debian 3.1
- ▶ Explain and develop a proof-of-concept to exploit a heap overflow on debian 3.1

Projet: Patch for CVE-2018-20343

- ▶ Understand CVE-2018-20343, a buffer overflow vulnerability
- ▶ You have to identify all instances of buffer overflow in the code (the code is not very big)
- ▶ You have to patch the vulnerable code

Projet: Complete exploit for CVE-2018-20343

- ▶ The current proof-of-concept only changes the value of EIP.
- ▶ You have to improve the PoC to enable an attacker to execute arbitrary code

Projet: Study and PoC for CVE-2013-0912 (Chrome type confusion)

- ▶ Reproduce the SVG code for the exploit based on information you find on the internet
- ▶ You should create a VM with a distribution from 2013 and have the vulnerable version of Chrome

Projet: Stable code injection through `/proc/self/mem`

- ▶ DosBox enables untrusted code to mount the host filesystem in the guest
- ▶ Thus untrusted code can write to `/proc/self/mem`
- ▶ You develop code to inject a shellcode into the virtual process of dosbox to execute arbitrary code
- ▶ You do this by writing to `/proc/self/mem`

Projet: Explanation of a recent exploit targeting webbrowsers (Chrome, Firefox, etc.)

- ▶ Contact me when you have found a CVE you want to explain.

Projet: Exploitation of a PoC type confusion in C++

- ▶ Write a proof-of-concept showing how to exploit a type confusion in C++ in a x86_64 architecture (latest debian)