

Lecture 10 : Confused Deputy

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2019

Previously...

Previously... in Lecture 1 (Introduction)

- ▶ Software Development Life-cycle
- ▶ Vulnerability Life-cycle
- ▶ Vulnerability Disclosure

Previously... in Lecture 2 (Buffer Overflow)

- ▶ A buffer on the stack
- ▶ Return address on the stack
- ▶ Overwrite return address
- ▶ Jump to shellcode on the stack

Previously... in Lecture 3 (ROP)

- ▶ NX bit (stack non-executable)
- ▶ Gadgets in already loaded code
- ▶ Chain gadgets (addresses of gadgets and data on the stack)
- ▶ Only data on the stack

Previously... in Lecture 4 (ASLR)

- ▶ Randomize code segment at program start
- ▶ Breaks gadget chains
- ▶ Bypass with information leak (e.g, vulnerability)

Previously... in Lecture 6 (CFI)

- ▶ Mechanism to allow only "intended" paths
- ▶ Binary instrumentation to add IDs
- ▶ Indirect jumps, call, returns check if ID of "destination" is correct
- ▶ Pure software implementation have 20% overhead

Previously... in Lecture 7 (Heap-Overflow)

- ▶ How a heap-overflow can be attacked depends on the heap management implementation
- ▶ The "unlink" attack present in early versions of glibc provides a "write anywhere" primitive to the attacker
- ▶ Recent implementations performs more check to prevent "unlink" based attacks

Previously... in Lecture 8 (Type Confusion)

- ▶ What is it? Manipulation of an object through another object.
- ▶ Consequences? Undefined behavior, hijack control flow.
- ▶ Why it works? No verification at runtime (otherwise runtime and/or memory overhead).

Previously... in Lecture 9 (SQL Injection)

- ▶ Code injection attacks enables bypass of authorization checks and/or execution of arbitrary code on the server
- ▶ Consequences: attacker gets access to privileged environment and/or can dump databases
- ▶ Protection include sanitization of the input and/or well defining what is code and what is data

Confused Deputy

- ▶ Confused: being unable to think with clarity or act with understanding and intelligence
- ▶ Deputy: An entity (person / software) empowered to act for another
- ▶ (specific type of privilege escalation vulnerability)

The vulnerability

- ▶ 1: Deputy software provides a service
 - ▶ 2: Deputy software is not well configured/programmed
 - ▶ 3: Attacker asks deputy software to use the service to its advantage
 - ▶ 4: Deputy software says, "Sure, no problem. Let's do it."
 - ▶ 5: Deputy software should have replied "No, sorry."
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- ▶ ex 1: return GPS coordinates
- ▶ ex 2: do not check that caller has GPS permission
- ▶ ex 3: give me GPS coordinates
- ▶ ex 4: here are the GPS coordinates, thank you for using my service
- ▶ ex 5: Sorry, I cannot give you the GPS coordinates since you do not have the GPS permission

Illustration 1: "The Confused Deputy"¹ (1988)

- ▶ Compiler FORT installed in directory SYSX
- ▶ User would call (SYSX)FORT to run the compiler on source code to generate a binary
- ▶ User could specify target file to receive debugging information
- ▶ Compiler engineers designed it to output statistics information in (SYSX)STATS
- ▶ This requires the OS to know that FORT has the right to write in its home directory (SYSX)
- ▶ Thus FORT has been given the right to its home directory (SYSX)

¹Hardy, Norm. "The Confused Deputy:(or why capabilities might have been invented)." ACM SIGOPS Operating Systems Review 22.4 (1988): 36-38.

Illustration 1: "The Confused Deputy" (1988)

- ▶ Billing information also saved to (SYSX)BILL...
- ▶ What could go wrong? → User can specify BILL as the output for debugging information...
- ▶ Problem: compiler runs with authority from 2 sources: the user and the compiler (can write to his home)
- ▶ Solution: new system call to "switch hats"
- ▶ The confused deputy was the compiler with the privilege to write to directories where normal users cannot

Illustration 2: Cross-Site Request Forgery (CSRF) ¹

- ▶ Tricks the victim into submitting a malicious request.
- ▶ Synonyms: XSRF, "Sea Surf", Session Riding, Cross-Site Reference Forgery and Hostile Linking.

¹[https://www.owasp.org/index.php/Cross-Site_Request_Forgery_\(CSRF\)](https://www.owasp.org/index.php/Cross-Site_Request_Forgery_(CSRF))

Illustration 2: Cross-Site Request Forgery (CSRF)

- ▶ User browses on his banking site
- ▶ Banking site accepts requests such as
"https://bank.lu/requests.html?amount=10&to=Dave"
- ▶ Without closing the banking session, user visits other websites
- ▶ One of them, "https://www.malicious.lu" sends an html page with the following link "https://bank.lu/requests.html?amount=10&to=Attacker"
- ▶ If the user clicks the request is sent.
- ▶ The confused deputy is the web-browser which has the privilege to send requests to the bank

Illustration 2: Cross-Site Request Forgery (CSRF)

- ▶ Secret value in cookie?
- ▶ Only accepting POST requests?
- ▶ HTTPS?
- ▶ Add a hash to all forms?

Illustration 3: Cross-Site Scripting (XSS) ¹

- ▶ Store code on web-server side, "www.compromised.lu"
- ▶ For instance, if servers allows users to add comments but do not filter comments properly
- ▶ Malory could send "< script > alert('hi there!'); < /script >" as a comment to page "www.compromised.lu/blog.html"
- ▶ This comment is stored in the database of web-server "www.compromised.lu"
- ▶ This comment is present in the html page generated by web-server when a user visits "www.compromised.lu/blog.html"
- ▶ Thus, all users visiting "www.compromised.lu/blog.html" will execute the javascript code of Malory: "< script > alert('hi there!'); < /script >"

¹[https://www.owasp.org/index.php/Cross-Site_Request_Forgery_\(CSRF\)](https://www.owasp.org/index.php/Cross-Site_Request_Forgery_(CSRF))

Illustration 3: Cross-Site Scripting (XSS)

- ▶ The confused deputy is the web-server which has the privilege to generate html pages to clients
- ▶ Solution?
- ▶ Sanitize input/output

Illustration 4: Android Application as Confused Deputies

- ▶ A developer can program multiple Android applications A1, A2, A3
- ▶ A1, A2, A3 all have GPS permission
- ▶ A1 has code to retrieve GPS coordinates and share this information to A2 and A3
- ▶ A1 is used to share GPS coordinates to A2 and A3 through a service component
- ▶ Malory could write application MA (with no GPS permission) and ask A1 for GPS permissions"

Illustration 4: Android Application as Confused Deputies

- ▶ The confused deputy is the Android application which has the privilege to retrieve GPS coordinates
- ▶ Solution?
- ▶ Restrict access to clients

- ▶ Confused Deputy Attacks are a type of privilege escalation vulnerability
- ▶ Attacker exploits the associated vulnerability (misconfiguration, logic error) to have more privilege
- ▶ Protection includes sanitization of the input and/or changing the configuration and/or patching the code logic

Question?

- ▶ Groups of 2
- ▶ Suggested topics:
 1. Heap exploitation on Debian 3.1
 2. Patch for CVE-2018-20343 (Ricardo, Alex)
 3. Complete exploit for CVE-2018-20343
 4. Study and PoC for CVE-2013-0912 (Chrome type confusion) (Adriano)
 5. Stable code injection through /proc/self/mem
 6. Explanation of a recent exploit targeting webbrowsers (Chrome, Firefox, etc.) (Yurii, Ervin)
 7. Exploitation of a PoC type confusion in C++ (Ihor, Artem)
 8. Break wordpress authentication mechanism.
- ▶ Deliverables: Presentation + Code (PoC)

Projet: Heap exploitation on Debian 3.1

- ▶ Explain the differences in the heap management code from debian 2.2 (lab 7) and debian 3.1
- ▶ Explain and develop a proof-of-concept to exploit a heap overflow on debian 3.1

Projet: Patch for CVE-2018-20343

- ▶ Understand CVE-2018-20343, a buffer overflow vulnerability
- ▶ You have to identify all instances of buffer overflow in the code (the code is not very big)
- ▶ You have to patch the vulnerable code

Projet: Complete exploit for CVE-2018-20343

- ▶ The current proof-of-concept only changes the value of EIP.
- ▶ You have to improve the PoC to enable an attacker to execute arbitrary code

Projet: Study and PoC for CVE-2013-0912 (Chrome type confusion)

- ▶ Reproduce the SVG code for the exploit based on information you find on the internet
- ▶ You should create a VM with a distribution from 2013 and have the vulnerable version of Chrome

Projet: Stable code injection through `/proc/self/mem`

- ▶ DosBox enables untrusted code to mount the host filesystem in the guest
- ▶ Thus untrusted code can write to `/proc/self/mem`
- ▶ You develop code to inject a shellcode into the virtual process of dosbox to execute arbitrary code
- ▶ You do this by writing to `/proc/self/mem`

Projet: Explanation of a recent exploit targeting webbrowsers (Chrome, Firefox, etc.)

- ▶ Contact me when you have found a CVE you want to explain.

Projet: Exploitation of a PoC type confusion in C++

- ▶ Write a proof-of-concept showing how to exploit a type confusion in C++ in a x86_64 architecture (latest debian)