Software Vulnerabilities: Exploitation and Mitigation Lab 6

Artem Kaliahin artem.kaliahin.001@student.uni.lu

April 2019

Question 1.1 Describe the source code.

main function - creates two arrays ia and ib with a length of 10, then calls sort2 function with 3 arguments (arrays ia, ab and number of arguments in command line argc) —>

sort2 function - takes 3 arguments (2 arrays and integer), then calls function sort twice. First time with array ia, integer argc and pointer to function lt. Second time - with array ib, integer argc and pointer to function gt —>

sort function - takes 3 arguments (during the first call - ia, argc and pointer to function lt, while second call takes ib, argc and pointer to gt), then it calls function gt or lt (depending on the argument) with arguments ia[argc] and ia[argc+1] (or ib[argc] and ib[argc+1]) —>

lt function - takes two arguments and returns 1 if first argument is less than the second, otherwise - returns 0 (in our case - returns 1 if ia[argc] < ia[argc+1], otherwise - returns 0).

gt function - takes two arguments and returns 1 if first argument is bigger than the second, otherwise - returns 0 (in our case - returns 1 if ib[argc] > ib[argc+1], otherwise - returns 0).

Question 1.2 Does gcc also support CFI?

Yes. CFI enforcement mechanism Vtable Verification (VTV) has been implemented in GCC version 4.9. [source]

Question 1.3 Describe all five options given to clang above.

 $-o\theta$ - sets optimization level to 0 (no optimization).

-fsanitize = cfi - enables ALL Clang's available CFI schemes (requires -flto flag and -fvisibility flag).

- -flto enables linking time optimization in 'full' mode.
- -fvisibility=hidden every declaration not explicitly marked with a visibility attribute has a hidden visibility (hides unnecessary symbols).
- -o cfi writes output to file "cfi".

Question 1.4 In the CFI version, there is the instruction ud2 which is not in the non-CFI version. What is this instruction doing?

Instruction ud2 - Raise invalid opcode exception, so called CFI-trap, which is generated by LLVM to crash the program (if CFI check fails, we execute a ud2 instruction, which will crash the program).

Question 1.5 What elements (instructions, functions, etc.) are in the CFI version and not in the non-CFI version.

There are 2 functions added by CFI - lt.cfi and gt.cfi. Instead of executing lt and gt directly, it jumps to lt.cfi and gt.cfi respectively (Picture 1) and executes lt.cfi and gt.cfi (when instructions for lt function in non-CFI version are the same as lt.cfi in CFI version, applicable to gt function as well).

0000000000400			
4005 d0:	e9 ab fe ff ff	jmpq	400480 <lt.cfi></lt.cfi>
4005 d5:	cc	int3	
4005 d6:	cc	int3	
4005 d7:	cc	int3	
0000000000400			
4005 d8:	e9 c3 fe ff ff	jmpq	4004a0 <gt.cfi></gt.cfi>
4005 dd:	cc	int3	ico ido agerer e
4005 de:	CC	int3	
4005 df:	cc	int3	

Picture 1

Instruction added to sort function are shown on Picture 2. They check for control flow integrity and crash the program if this check fails (instruction ud2 crashes the program if jbe doesn't jump on an instruction below, this jump requires flags CF = 1 or ZF = 1, which are updated by cmp instruction).

4004c8:	48 b8	d0	05	40	00	00	movabs \$0x4005d0,%rax
4004cf:	00 00	00					

4004 e1:	48 89 d7	mov	%rdx, %rdi
4004e4:	48 29 c7	sub	%rax,%rdi
4004e7:	48 89 f8	mov	%rdi,%rax
4004ea:	48 c1 e8 03	shr	\$0x3,%rax
4004ee:	48 c1 e7 3d	shl	\$0x3d,%rdi
4004f2:	48 09 f8	ог	%rdi,%rax
4004f5:	48 83 f8 01	CMP	\$0x1,%rax
4004f9:	48 89 55 e0	mov	%rdx,-0x20(%rbp)
4004fd:	76 02	jbe	400501 <sort+0x41></sort+0x41>
4004ff:	0f 0b	ud2	

Picture 2

While functions sort2 and main remain the same in CFI and non-CFI version.

Question 1.6 Explain how the added instructions check for CFI. Describre precisely under which conditions instruction ud2 is executed. Drow the CFG for the C program above.

4004 c8:	48	b8	d0	05	40	00	00	movabs \$0x40	005d0,%rax

4004 cf: 00 00 00

Copies \$0x4005d0 (address of *lt.cfi* function) to \$rax.

```
4004e1: 48 89 d7 mov %rdx, %rdi
```

Copies \$rdx to \$rdi (there is address of function that was given as an argument to function in \$rdx according to the convention - address of *lt.cfi* or *gt.cfi*, whereas the difference between addresses is 8 bytes).

```
4004e4: 48 29 c7 sub %rax,%rdi
```

Subtracts value in \$rax from value in \$rdi and puts difference in \$rdi. For now this value according to CFG can be either 0 or 8, because \$rdi can have either the same value as \$rax if \$lt.cfi\$ is called or 8 bytes more if \$gt.cfi\$ is called.

```
4004\,\mathrm{e}7: 48\,89\,\mathrm{f}8 mov \%\mathrm{rdi}, \%\mathrm{rax}
```

Copies the result of subtraction to \$rax.

4004ea:	48 c1 e8 03	shr	90x3,%rax
4004 ee :	48 c1 e7 3d	shl	\$0x3d.%rdi

Here we shift \$rax value 3 bits to the right and \$rdi value 62 bits to the left. In this case the only value that will eventually make zero in both registers after shifts is 0 or 8 (100_2) - this is valid behaviour. But it is also possible if the value is 24 (1100_2) , but in our case it is not a threat since it will jump to the program's _libc_csu_init library.

```
4004f2: 48 09 f8 or %rdi,%rax
```

This instruction puts to \$rax the result of OR operation on \$rdi and \$rax. It will be equal to 0 only if both registers have value 0.

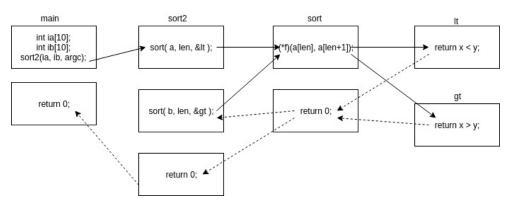
4004f5: 48 83 f8 01 cmp \$0x1,%rax

4004 fd: 76 02 jbe 400501 <sort+0x41>

Here we compare \$rax with 0x1 (subtract \$rax from 0x1). If control flow is not disrupted, \$rax will store 0-value as proved above and cmp will change flags needed to execute jbe (jump if below or equal) instruction. It will jump to address 400501 (if CF or ZF flags are equal to 1) and jump over instruction ud2 to the next one. Jump will also be executed if \$rax stores value 0x1.

4004 ff: 0f 0b ud2

Instruction ud2 crashes program. It will be executed ONLY when jump is not executed (above it is described precisely under which conditions jump is not executed).



Picture 3. Control Flow Graph