

You have a function rand5() that generates a random integer from 1 to 5. Use it to write a function rand7() that generates a random integer from 1 to 7.

rand5() returns each integer with equal probability. rand7() must also return each integer with equal probability.

Gotchas

Simply running rand5() twice, adding the results, and taking a modulus won't give us an equal probability for each possible result.

Not convinced? Count the number of ways to get each possible result from 1..7.

Your function will have worst-case infinite runtime, because sometimes it will need to "try again."

However, at each "try" you only need to make two calls to rand5(). If you're making 3 calls, you can do better.

We can get away with worst-case O(1) space. **Does your answer have a non-constant space cost?** If you're using recursion (and your language doesn't have <u>tail-call optimization</u>,), you're potentially incurring a worst-case infinite space cost in the call stack.

Breakdown

Because we need a random integer between 1 and 7, we need at least 7 possible outcomes of our calls to rand5(). One call to rand5() only has 5 possible outcomes. So we must call rand5() at least twice.

Can we get away with calling rand5() exactly twice?

Our first thought might be to simply add two calls to rand5(), then take a modulus to convert it to an integer in the range 1..7:

However, this won't give us an equal probability of getting each integer in the range 1..7. Can you see why?

There are at least two ways to show that different results of rand7_mod() have different probabilities of occurring:

- Count the number of outcomes of our two rand5() calls which give each possible result of rand7_mod()
- 2. Notice something about the total number of outcomes of two calls to rand5()

If we count the number of ways to get each result of rand7_mod():

<pre>result of rand7_mod()</pre>	# pairs of rand5() results that give that result
1	4
2	3
3	3
4	3
5	3
6	4
7	5

So we see that, for example, there are five outcomes that give us 7 but only three outcomes that give us 5. We're almost twice as likely to get a 7 as we are to get a 5.

But even without counting the number of ways to get each possible result, we could have **noticed something about the** *total* **number of outcomes of two calls to rand5()**, which is 25 (5*5). If each of our 7 results of rand7_mod() were equally probable, we'd need to have the same number of outcomes for each of the 7 integers in the range 1..7. That means our total number of outcomes would have to be divisible by 7, and 25 is not.

Okay, so rand7_mod() won't work. **How do we get equal probabilities for each integer from 1** to 7?

Is there some number of calls we can make to rand5() to get a number of outcomes that is divisible by 7?

When we roll our die n times, we get 5^n possible outcomes. Is there an integer n that will give us a 5^n that's evenly divisible by 7?

No, there isn't.

That might not be obvious to you unless you've studied some number theory. It turns out every integer can be expressed as a product of prime numbers (its <u>prime factorization</u>.). It also turns out that every integer has only *one* prime factorization.

Since 5 is already prime, any number that can be expressed as 5^n (where n is a positive integer) will have a prime factorization that is all 5s. For example, here are the prime factorizations for 5^2 , 5^3 , 5^4 :

$$5^2 = 25 = 5 * 5$$

$$5^3 = 125 = 5 * 5 * 5$$

$$5^4 = 625 = 5 * 5 * 5 * 5$$

7 is also prime, so if any power of 5 were divisible by 7, 7 would be in its prime factorization. But 7 can't be in its prime factorization because its prime factorization is all 5s (and it has only *one* prime factorization). So no power of 5 is divisible by 7. BAM MATHPROOF.

So no matter how many times we run rand5() we won't get a number of outcomes that's evenly divisible by 7. What do we dooooo!?!?

Let's ignore for a second the fact that 25 isn't evenly divisible by 7. We can think of our 25 possible outcomes from 2 calls to rand5 as a set of 25 "slots" in a list:

Which we could then try to evenly distribute our 7 integers across:

```
results = [

1, 2, 3, 4, 5,

6, 7, 1, 2, 3,

4, 5, 6, 7, 1,

2, 3, 4, 5, 6,

7, 1, 2, 3, 4,

]
```

It *almost works*. We could randomly pick an integer from the list, and the chances of getting any integer in 1..7 are *pretty evenly* distributed. Only problem is that extra 1, 2, 3, 4 in the last row.

Any way we can sidestep this issue?

What if we just "throw out" those extraneous results in the last row?

21 is divisible by 7. So if we just "throw out" our last 4 possible outcomes, we have a number of outcomes that are evenly divisible by 7.

But what should we do if we get one of those 4 "throwaway" outcomes?

We can just try the whole process again!

Okay, this'll work. But how do we translate our two calls to rand5() into the right result from our list?

What if we made it a 2-dimensional list?

```
results = [
    [1, 2, 3, 4, 5],
    [6, 7, 1, 2, 3],
    [4, 5, 6, 7, 1],
    [2, 3, 4, 5, 6],
    [7, 1, 2, 3, 4],
]
```

Then we can simply treat our first roll as the row and our second roll as the column. We have an equal chance of choosing any column and any row, and there are never two ways to choose the same cell!

```
Python 3.6 ▼
def rand7_table():
    results = [
        [1, 2, 3, 4, 5],
        [6, 7, 1, 2, 3],
        [4, 5, 6, 7, 1],
        [2, 3, 4, 5, 6],
        [7, 0, 0, 0, 0],
    ]
    # Do our die rolls
    row = rand5() - 1
    column = rand5() - 1
    # Case: we hit an extraneous outcome
    # so we need to re-roll
    if row == 4 and column > 0:
        return rand7_table()
    # Our outcome was fine. return it!
    return results[row][column]
```

This'll work. But we can clean things up a bit.

By using recursion we're incurring a space cost in the call stack, and risking stack overflow. This is especially scary because our function *could* keep rerolling indefinitely (though it's unlikely).

How can we rewrite this iteratively?

```
def rand7_table():
    results = [
        [1, 2, 3, 4, 5],
        [6, 7, 1, 2, 3],
        [4, 5, 6, 7, 1],
        [2, 3, 4, 5, 6],
        [7, 0, 0, 0, 0],
    ]
    while True:
        # Do our die rolls
        row = rand5() - 1
        column = rand5() - 1
        # Case: we hit an extraneous outcome
        # so we need to re-roll
        if row == 4 and column > 0:
            continue
        # Our outcome was fine. return it!
        return results[row][column]
```

One more thing: we don't have to put our whole 2-d results list in memory. Can you replace it with some arithmetic?

We could start by coming up with a way to translate each possible outcome (of our two rand5() calls) into a different integer in the range 1..25. Then we simply mod the result by 7 (or throw it out and try again, if it's one of the last 4 "extraneous" outcomes).

How can we use math to turn our two calls to rand5() into a unique integer in the range 1..25?

What did we do when we went from a 1-dimensional list to a 2-dimensional one above? We cut our set of outcomes into sequential slices of 5.

How can we use math to make our first roll select which slice of 5 and our second roll select which item within that slice?

We could take something like:

```
outcome_number = roll1 * 5 + roll2
```

But since each roll gives us an integer in the range 1..5 our lowest possible outcome is two 1s, which gives us 5 + 1 = 6, and our highest possible outcome is two 5s, which gives us 25 + 5 = 30. So we need to do some adjusting to ensure our outcome numbers are in the range 1..25:

```
outcome_number = ((roll1-1) * 5 + (roll2-1)) + 1
```

(If you're a math-minded person, you might notice that we're essentially treating each result of rand5() as a digit in a two-digit base-5 integer. The first roll is the fives digit, and the second roll is the ones digit.)

Can you adapt our function to use this math-based approach instead of the results list?

Solution

Because rand5() has only 5 possible outcomes, and we need 7 possible results for rand7(), we need to call rand5() at least twice.

When we call rand5() twice, we have 5*5=25 possible outcomes. If each of our 7 possible results has an equal chance of occurring, we'll need each outcome to occur in our set of possible outcomes the same number of times. So our total number of possible outcomes must be divisible by 7.

25 isn't evenly divisible by 7, but 21 is. So when we get one of the last 4 possible outcomes, we throw it out and roll again.

So we roll twice and translate the result into a unique outcome_number in the range 1..25. If the outcome_number is greater than 21, we throw it out and re-roll. If not, we mod by 7 (and add 1).

```
def rand7():
    while True:
        # Do our die rolls
    roll1 = rand5()
    roll2 = rand5()
    outcome_number = (roll1-1) * 5 + (roll2-1) + 1

# If we hit an extraneous
    # outcome we just re-roll
    if outcome_number > 21:
        continue

# Our outcome was fine. return it!
    return outcome_number % 7 + 1
```

Complexity

Worst-case $O(\infty)$ time (we might keep re-rolling forever) and O(1) space.

What We Learned

As with the previous question about writing a rand5() function (/question/simulate-5-sided-die), the requirement to "return each integer with equal probability" is a real sticking point.

Lots of candidates come up with clever O(1)-time solutions that they are so sure about. But their solutions aren't actually uniform (in other words, they're not truly random).

In fact, it's impossible to have true randomness and non-infinite worst-case runtime.

If you don't understand why, go back over our proof using "prime factorizations," a little ways down in the breakdown section.

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