Imprimitive groups:
· g a 2 françotively
Defn:
et $B = \{B_1,, B_n\}$ $B_i \subset \Omega$ ; $B_i \cap B_j = \emptyset$ for $i \neq j$
UBi = 2 (i.e. B is a partition of 2)
$\mathcal{Z}$ is a bloch system for $\mathcal{G} \cap \Omega$ when $\mathcal{B}_i^{\ \beta} \in \mathcal{B}  \forall g \in \mathcal{G}$
(i.e. Ris G-invariant).
Ex: frivial block systems: B = (li): ie 2)
$\mathcal{Z}_{\infty} = \{ -\alpha \}$
$\mathcal{E}_{\times}: \qquad \mathcal{G} = \langle (1, 2, 3, 4) \rangle$
R = { {1,3}, {2,4}}
Defn: Ne say klut & a Dingminitively
iff of transitively and admits a non- trivial block system.
(Otherwise we say that & a printively).

Lemma: Let  $B = \{B_1, \dots, B_n\}$  be a black system for  $g \cap \Omega$ . The action of g on blocks is framitive.

Proof: francitie =  $\forall 1 \leq i, j \leq h$   $\exists g \in g \leq f$ .  $B_i^s = B_j$ . Let  $\delta \in B_i$  and  $\gamma \in B_j$ . since  $g \in Q$  francitively.  $\Rightarrow \exists g \in g \leq f$ .  $\delta^s = \chi$ . The same g moves  $B_i \circ b \otimes B_j$ .

Corollany:

· 12 = B. B.1

· Bloch system is determined by a single block.

· If |G| = p;  $G \cap \Omega$  than tively, then  $G \cap primtively$ .

Lemma: Suppose & 1 2 framsitively and let

5 = 5 tolog(x) for some x & D.

Then there is a bijection between subgroups

{T: S&T&G} and blade systems B=1B1,...,Bu} for GQQ.

The bijection is given by

X,T ~~ ×T=:B, ~~ R=B,S.

The inverse is simply

B=(B1,..., Bn) ~ Stabe (B1).

Proof: let S ∈ T ∈ g and pich × €, Q × B 8 B2 Set B = x 2 - B 5. Claim: B is a block system for g Q \( \sigma \). Let BonBh # of i.e. So = yh for some S, y & B. Since  $B = X^T \Rightarrow \delta = X^a$ ,  $y = X^b$  for some a, b.  $\epsilon T$ .  $\Rightarrow \times^{a\cdot g} = \times^{bh}$  i.e  $\times^{a\cdot g \cdot h \cdot h \cdot l} = \times$ ⇒ aghib e Stabe(x) ≤ T ⇒ gh'eT. Bohin B = B i.e. Boh Jince T stabilizes B G-invariance is obvious by the definition of 2 Let B be a block system, x & B. & B. Hy fixes  $x_1 \Rightarrow x_1 \in \mathcal{B}_1 \cap \mathcal{B}_1^3 \neq \emptyset \Rightarrow \mathcal{B}_1^3 = \mathcal{B}_1$ Stab & (x) & Stab & (B.). Let se B, => x3=6 (by transitivity) => Se B,3, B, # ⇒ g ∈ Stabg (S). Therefore × stabg(B1) = B1. To make that the maps are inverse of each other we need to chedi: Stabe (x stabe (B1)) = Stabe (B1) =: T. If geg st (xT) > + ter Itast xtg=xt'

>> xtst"=x => tot eT => geT.

Definition:

\*\*Mogroup S<9 is called maximal

if S+9 and there is no subgroup T<9

s.f. S\*\*T.

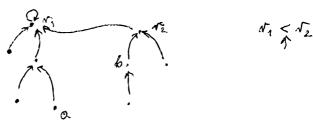
Corollary:

- a point stabilizer is a maximal subgroup.
- · Subgroup 3<9 is maximal iff 905/9 is primitive.

Finding blocks:

Let XC Q be a subset (a seed) we want to be CB,

- start with  $B_X = X$ ,  $(B_y = \{y\})$  for  $y \in \Omega \setminus X$
- Act via G= (5) on each of Bi and merge those which intersect.
- · Each Bi is represented by a unique element -> the representative
- for each  $\times \in \Omega$  we store the representative of B; which  $\times$  belongs to.



ALGORITHM: Union! Input:  $C_1 - a$  subset of  $\Omega$   $C_2 - a$  subset of  $\Omega$ Output: the union of Cy and Ce path haloing: begin while parent(x) \$ x x, parent = parent(parent(x)) 1, = representative (G) and = parent(x) [ = representative (G) 许与丰区 set-representative! (Te, Ti) return C, ALGORITHM: Block system INPUT: . S- a generaling set for G=<5> · D - a set with g-action  $\cdot$  B - a subset of  $\Omega$ OUTPUT: B - the finest Block system for & Q.D. s.t. B C B1 Begin for x e 12 set\_representative! (x, x) queue = [] I a queue of points that have changed infilize the witial for x & B set\_representative: (x, first(B)) partition push! (queue, x)

while ! "semply (queue) x = pap! (queue) y = representative(x) for se's \[
 \times = representative (x^s)
 \] B = representative (ys) if x +B Union! (x,B) push B to queue return r

The algorithm converges to a block system Proof: Since we've folling only winders the returned partition contains B in one of its Since every Union! moves up in the lattice of all partitions of 12 and

1.03 - the frivial black system is the maximal element the algorithm has to stop

We need to prove that the reduced patition B is G-invariant i.e. YBEB YxyEB Ygeg x3, y3 EBEB. Observe: . It's enough to check this for y = r(x). It's enough to check this for  $g \in S$ .  $\forall \mathcal{B} \in \mathcal{B} \ \forall x \in \mathcal{B} \ \forall s \in S \ \tau(x^s) = \tau(\mathcal{M}).$  $\forall \times \in \Omega \ \forall s \in S \ r(x^s) = r(r(x)^s).$ • It's enough to enforce the earthfion only for pts which changed label > it's the greve. Had we added all points whose representative is B we would be oh. But only B was added... Let  $\times \in \mathbb{R}$ : representative(x) =  $\beta \neq x$  $\Rightarrow$  x was on the greene  $\Rightarrow$  (\*) is satisfied for (x,B). If it happens that we change  $\tau(\beta) = \alpha$  then (\*) is enforced For  $(\beta, \alpha)$ , so  $\Gamma(x^3) = \Gamma(\beta^3) =$  $= r(r(\beta)^3) = r(\alpha^3) \square.$ Lemma: Let 16B, EB «bloch system for gr. 2. B, is union of orbits of Stabg (1).

Proof: If  $g \in \text{Moleg}(1)$ ,  $1 \neq \alpha \in B_1 \Rightarrow \{1, \alpha\}^3 = \{1, \alpha\}^3 \subset B_1^3 = B_1 \Rightarrow \alpha \in B_1$ 

Corollary: Any block that contains {1, \alpha}

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Practical tip: We closit need to know Stabog(1)

Exactly!

H's enough to find a few vandon elements

from stabog(1) to compute the orbit of \alpha.

(e.g. schreter generators from the frameworsal.)

Recall: Every infransitive group is a subdirect product of its transitive parts.

To there a universal description for imprimitive groups:

Defn: Let G, H be groups. let  $\psi: H \to Aut(g)$ be a homomorphism. Hhen  $g \times_{\psi} H = \langle (g,h) \in G \times H, \cdot_{\psi}, (1,1) \rangle$ is a group with

$$(g,h)\cdot\psi(a,b)=(g\cdot\psi(b)(a),hb)$$

$$(1,1) = (g,h) \cdot (a,b) \implies b = h^{-1}, \quad \alpha = \psi(h)(g^{-1})$$
Note:  $g \triangleleft g \bowtie_{\psi} H$ .