

# **SAut( $F_5$ ) HAS PROPERTY (T)**

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Kummamoto 2018

## Outline

**Group Laplacians & Property (T)**

**Positivity & SDP**

**The procedure**

**Concrete examples**

## Notation / Convention

- $G = \langle S \mid R \rangle$  is a finitely presented group generated by a fixed **symmetric generating set** (i.e.  $S^{-1} = S$ ).

## **GROUP LAPLACIANS**

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- ▶ spectrum of  $\Delta$  is real and non-negative;
- ▶ the second eigenvalue  $\lambda_1$  is called the **spectral gap**

$$0 = \lambda_0 \leq \lambda_1 \leq \dots$$

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- ▶ Provide a constructive (computable) proof for  $n = 5$ .

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### Theorem (Lubotzky & Pak, 2000)

*Let  $K$  be a finite group generated by  $k \leq n$  elements. If  $\text{SAut}(F_n)$  has property (T) with constant  $\kappa = \kappa(\text{SAut}(F_n), \{\text{transvections}\}) > 0$ , then PRA walk on  $\Gamma_n = \Gamma_n(K)$  has fast mixing rate, i.e.*

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$$\left\| Q_{(g)}^t - U \right\|_{\text{tv}} \leq \varepsilon \quad \text{for} \quad t \geq \frac{16}{\kappa^2} \log \frac{|\Gamma_n|}{\varepsilon} \sim O \left( \left( \frac{n}{\kappa} \right)^2 \log \frac{|K|}{\varepsilon} \right)$$

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### Note

We do observe fast mixing rate in practice for large  $n$ .

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### Corollary

Let  $G = \langle S | \dots \rangle$  be a finitely generated group. If there exists  $\lambda > 0$  such that  $\Delta^2 - \lambda\Delta \geq 0$ , then  $G$  has property (T) with

$$\sqrt{\frac{2\lambda}{|S|}} \leq \kappa(G, S).$$

**How to prove that  $\Delta^2 - \lambda\Delta \geq 0$  ?**

**How to prove that a polynomial  $f \geq 0$  ?**

## **POSITIVITY & SDP**

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## Hilbert's 17th problem

### Theorem (Hilbert's Positivstellensatz, 1888)

An everywhere non-negative polynomial  $p \in \Sigma^2 \mathbb{R}[x_1, \dots, x_n]$  (is a sum of squares) if and only if either

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$x^4y^2 + x^2y^2 - 3x^2y^2 + 1 \geq 0$  is SOS if You multiply it by  $x^2 + y^2 + 1$ .

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By evaluating

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For  $f$  to be a SOS we just need  $(p_{ij}) = P$  to be **semi-positive definite**, since then  $P = Q^T Q$  and (for  $(1, x, y)^T = \vec{x}$ )

$$f = \vec{x}^T P \vec{x} = \vec{x}^T Q^T \cdot Q \vec{x} = (Q \vec{x})^T \cdot (Q \vec{x}).$$

## Linear programming:

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## Semi-definite programming

- ▶ optimise linear functional
- ▶ on a polytope intersected with the cone of SPD matrices (spectrahedron)
- ▶ weak duality, non-unique solutions
- ▶ even feasibility is a hard problem!

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tries to maximise  $\lambda$  as long as  $(2x^2 + 4x + 1) - \lambda \geq 0$ .

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For any  $*$ -invariant element  $\xi \in \mathbb{R}[G]$

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This of no use for us: SOS decompositions  $\Delta^2 - \lambda \Delta + \varepsilon = \sum \xi_i^* \xi_i$  may be very different for different  $\varepsilon$ .

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for all  $\varepsilon$  simultaneously!

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3. Solve the problem (numerically):

maximize:  $\lambda$

subject to:  $P \succcurlyeq 0$ ,  $P \in \mathbb{M}_{\vec{x}}$

$$\lambda \geq 0$$

$$(\Delta^2 - \lambda \Delta)_t = \sum_{g^{-1}h=t} P_{g,h}, \quad \text{for all } t \in B_{2d}(e, S)$$

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3. Solve the problem (numerically):

maximize:  $\lambda$

subject to:  $P \succcurlyeq 0$ ,  $P \in \mathbb{M}_{\vec{x}}$

$$\lambda \geq 0$$

$$(\Delta^2 - \lambda \Delta)_t = \sum_{g^{-1}h=t} P_{g,h}, \quad \text{for all } t \in B_{2d}(e, S)$$

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5. Finally:  $\xi_g = \langle \vec{x}, \vec{q}_g \rangle$  and  $\Delta^2 - \lambda \Delta = \sum_{g \in \vec{x}} \xi_g^* \xi_g$ .

*How do we certify that the numerical result is sound?*

# Certifying correctness

## Lemma (Netzer&Thom)

Let  $r \in I[G] \subset \mathbb{R}[G]$  such that  $\text{supp}(r) \subset B_d(e)$ . Then

$$r + 2^{d-1} \|r\|_1 \cdot \Delta \in \Sigma^2 I[G].$$

## Corollary

If  $\Delta^2 - \lambda \Delta = \sum \xi_i^* \xi_i + r$ , then

$$\Delta^2 - (\lambda - 2^{d-1} \|r\|_1) \Delta = \sum \xi_i^* \xi_i + (r + 2^{d-1} \|r\|_1 \Delta) \in \Sigma^2 I[G],$$

i.e.  $\Delta$  has spectral gap of at least  $\lambda - 2^{d-1} \|r\|_1$ .

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$$P \rightarrow \sqrt{P} \rightarrow \sqrt{P}_{int} \rightarrow \sqrt{P}_{int}^{aug} \rightarrow Q \in \mathbb{M}_{\vec{x}}(\mathbb{R}IF)$$

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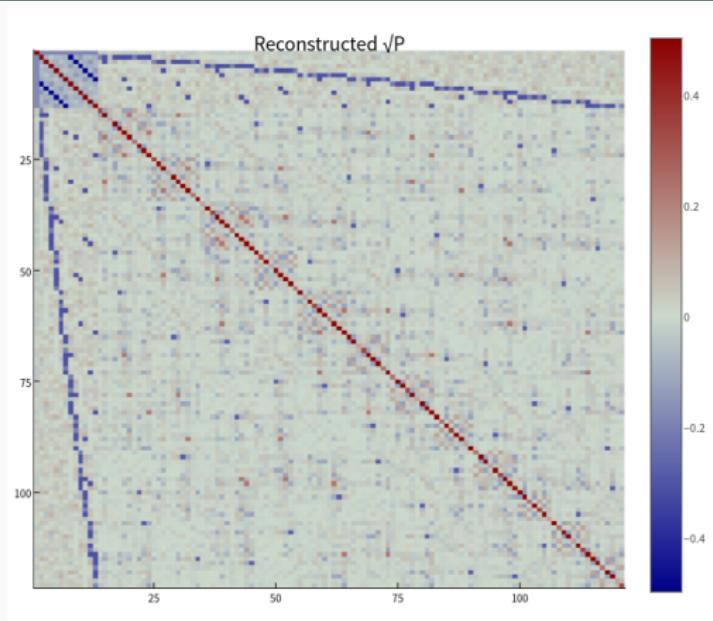
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## **CONCRETE EXAMPLES**

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$\sqrt{P} = Q \in \mathbb{M}_{\vec{x}}$ , where  $\vec{x} = B_2(e, E(3))$ , i.e. rows and columns are indexed by elements in  $(SL(3, \mathbb{Z}), E(3))$  of word length  $\leq 2$ . In this case

$$\Delta^2 - 0.28\Delta = \sum_{i=1}^{121} \xi_i^* \xi_i + r, \quad \|r\|_1 \in [3.8508, 3.8511] \cdot 10^{-7}$$

$G$	$n$	$m$	$\lambda$	$\ r\ _1 <$	$lb_\kappa$	$< \kappa$	$ub_\kappa$
SL(3, $\mathbb{Z}$ )	390,287	935,021	0.5405	$5.2 \cdot 10^{-7}$	0.19	0.30014	0.81650
SL(4, $\mathbb{Z}$ )	93,962	263,122	1.3150	$5.2 \cdot 10^{-8}$	0.00106	0.33103	0.70711
SL(5, $\mathbb{Z}$ )	628,882	1,757,466	2.6500	$2.0 \cdot 10^{-4}$	0.00105	0.36400	0.63246

$\text{SAut}(F_4)$

$G$	$n$	$m$	$\lambda$	$\ r\ _1 <$
$\text{SAut}(F_4)$	3,157,730	1,777,542	0.0100	7.4

(after weeks of computation)

$\mathrm{SAut}(F_5)$

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$G$	$n$	$m$
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- ▶ Find a finite group  $K < \mathrm{Aut}(\mathrm{SAut}(F_n))$  which keeps the generating set  $S$  and (thus)  $\Delta^2 - \lambda\Delta$  invariant ( $K = \mathbb{Z}_2 \wr S_5$ );

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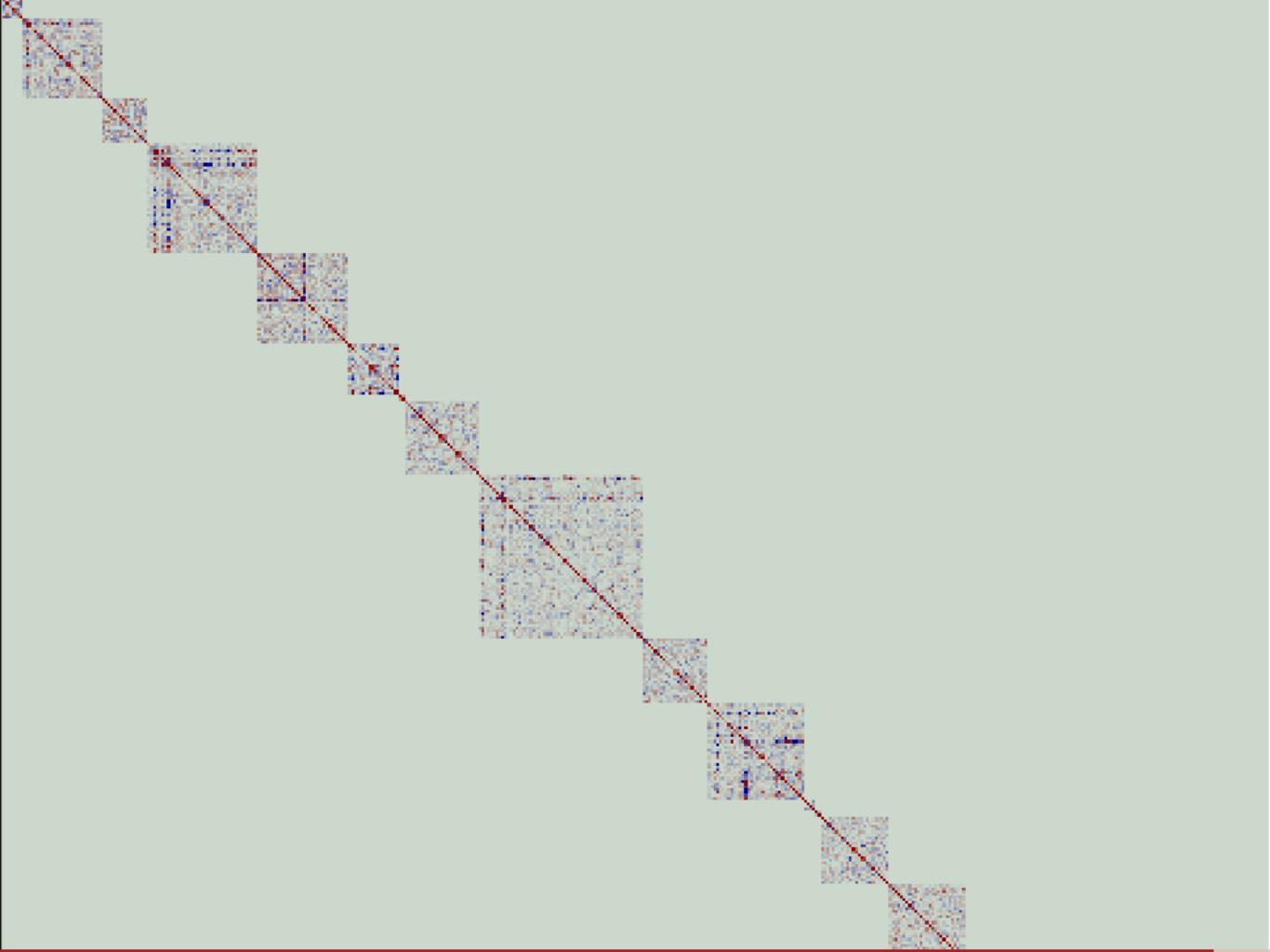
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- ▶ Solve the smaller problem and reconstruct the solution  $P$  of the larger one.



$\text{SAut}(F_5)$

$G$	$n$	$m$	$\lambda$	$\ r\ _1 <$	$< \kappa$
$\text{SAut}(F_5)$	13,233	7,230	1.3000	$2.1 \cdot 10^{-6}$	0.18028

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*(The Method of Mechanical Theorems, Archimedes)*