

Graph and Subgraph Isomorphism Using GNNs

An overview of the essential concepts in Stanford CS224W (Lectures 9 and 12)
with only oversimplified examples

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Graph Isomorphism

- Complexity of graph isomorphism is unknown
- Weisfeiler-Lehman graph kernel traditionally used to obtain graph-level embedding

WL color-refinement algorithm

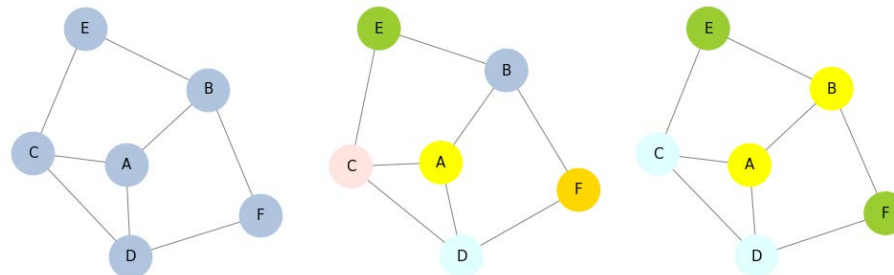
1. Assign initial color $c^{(0)}(v)$ to each node v
2. Iteratively refine node colors by

$$c^{(k+1)}(v) = \text{HASH} \left(c^{(k)}(v), \{c^{(k)}(u)\}_{u \in N(v)} \right)$$

- **HASH** function maps input to distinct values (colors)
- After K steps, $c^{(K)}(v)$ summarizes the structure of the K -hop neighborhood
- To run as GNN, need to implement **HASH** as AGG
 - Need notion of **distinguishing node embeddings**

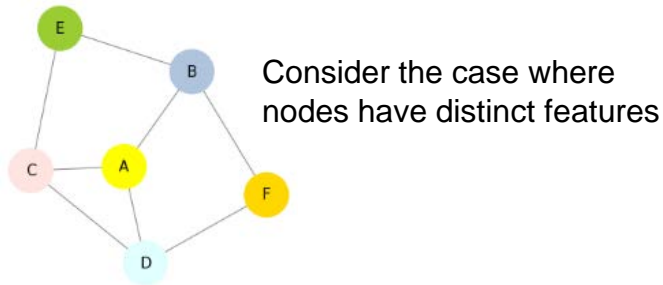
Distinguishing node embeddings

- Consider the task of keeping embeddings feature distinguishable
- Two factors affect embedding
 - (Initial) feature
 - Feature representation will affect whether features can be converted into each other
 - If $\begin{bmatrix} 1 & 0 \end{bmatrix}$, $\begin{bmatrix} 0 & 1 \end{bmatrix}$ and $\begin{bmatrix} 1 & 1 \end{bmatrix}$ represent three distinct features, then the sum of $\begin{bmatrix} 1 & 0 \end{bmatrix}$ and $\begin{bmatrix} 0 & 1 \end{bmatrix}$ would become $\begin{bmatrix} 1 & 1 \end{bmatrix}$, the third feature
 - To simplify this discussion assume that each distinct feature corresponds to one distinct dimension in the feature vector
 - Neighborhood structure
- Nodes with the same (initial) **feature** and **neighborhood structure** should be assigned the same embedding, and vice versa

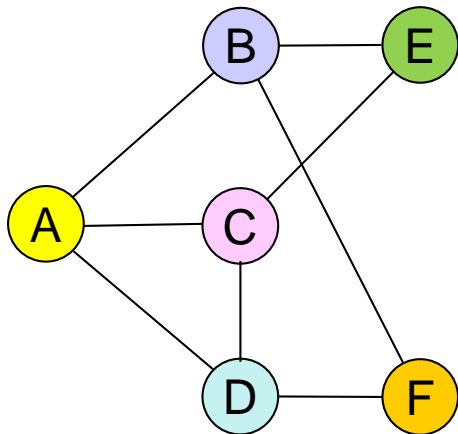


Computing node embeddings

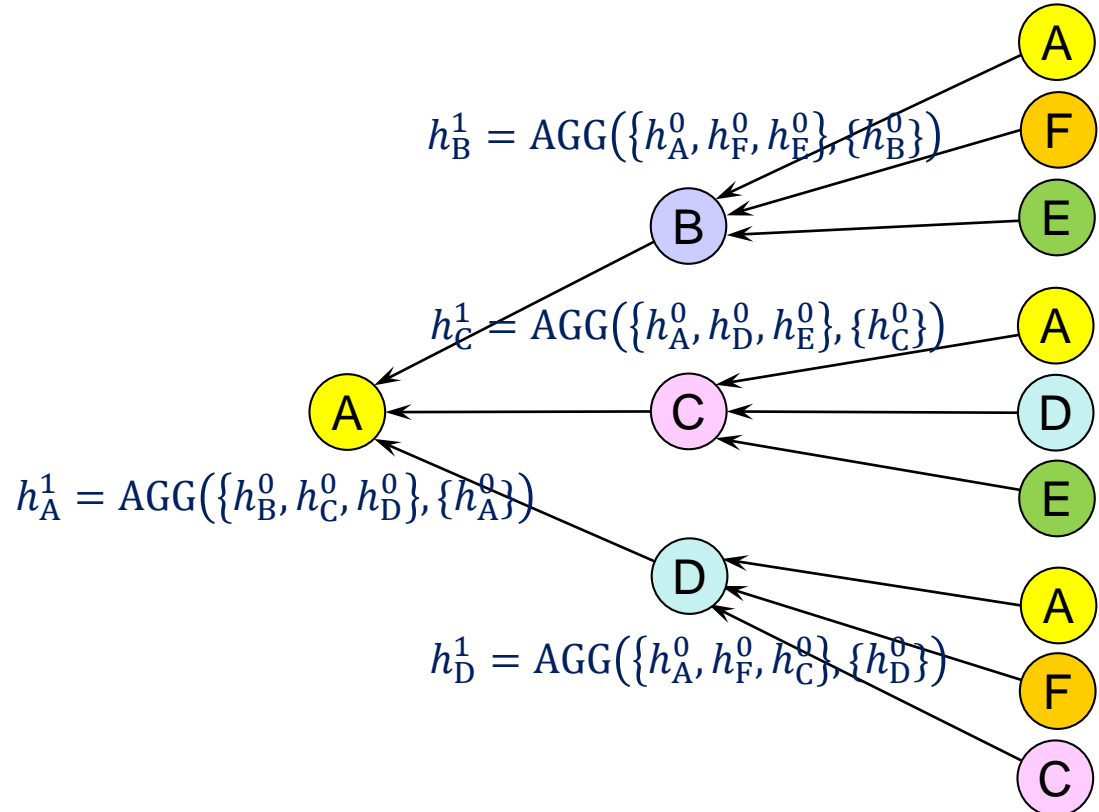
- Recall that given a GCN of 2 layers, the embedding of A is computed through the computation graph as follows



Rearranged w.r.t. distance from A



Computation graph of A's embedding



Distinguishability under sum

- Let $AGG = \text{sum}$, then for the GCN of 2 layers, the embeddings are as follows

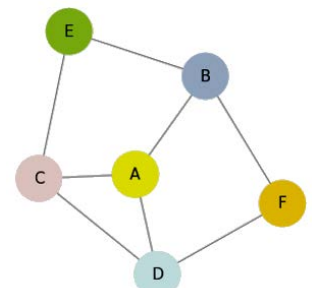
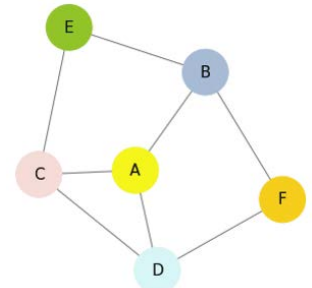
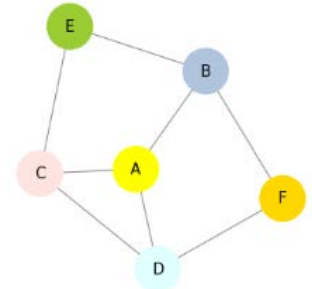
- Let $h_A^0 = [1, 0, 0, 0, 0, 0]$, $h_B^0 = [0, 1, 0, 0, 0, 0]$, $h_C^0 = [0, 0, 1, 0, 0, 0]$, $h_D^0 = [0, 0, 0, 1, 0, 0]$, $h_E^0 = [0, 0, 0, 0, 1, 0]$, $h_F^0 = [0, 0, 0, 0, 0, 1]$, and let AGG be **sum**. Then

- $h_B^1 = AGG(\{h_A^0, h_F^0, h_E^0\}, \{h_B^0\}) = [1, 1, 0, 0, 1, 1]$
- $h_C^1 = AGG(\{h_A^0, h_D^0, h_E^0\}, \{h_C^0\}) = [1, 0, 1, 1, 1, 0]$
- $h_D^1 = AGG(\{h_A^0, h_F^0, h_C^0\}, \{h_D^0\}) = [1, 0, 1, 1, 0, 1]$
- $h_A^1 = AGG(\{h_B^0, h_C^0, h_D^0\}, \{h_A^0\}) = [1, 1, 1, 1, 0, 0]$

- Finally the embedding of A is

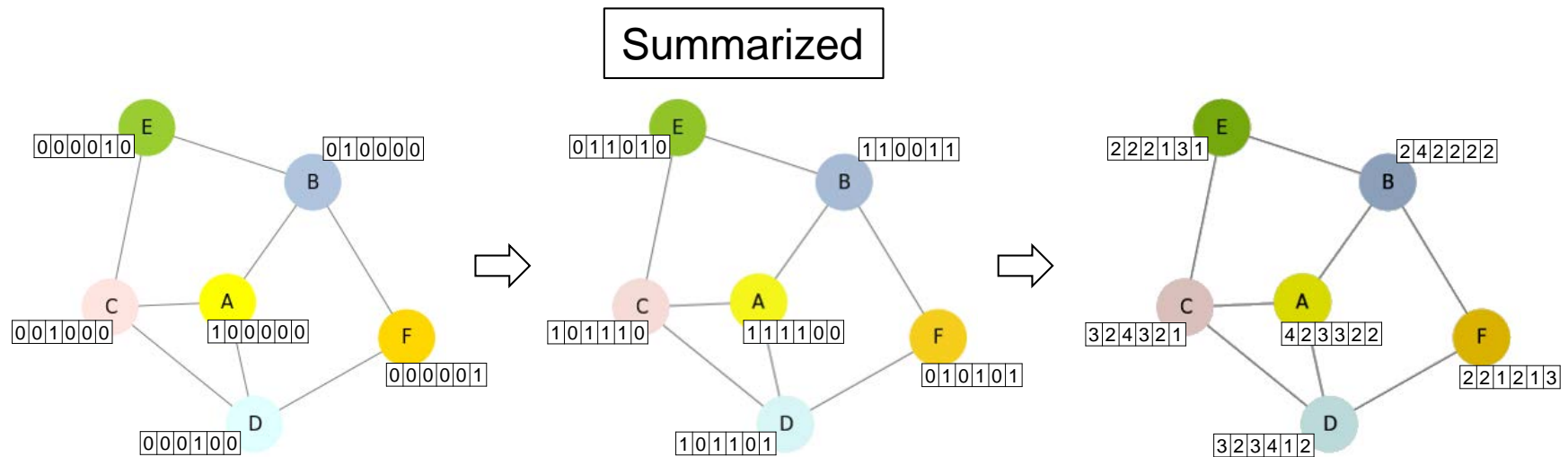
- $h_A^2 = AGG(\{h_B^1, h_C^1, h_D^1\}, \{h_A^1\}) = [4, 2, 3, 3, 2, 2]$

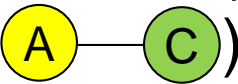
- Similarly, $h_B^2 = [2, 4, 2, 2, 2, 2]$
 $h_C^2 = [3, 2, 4, 3, 2, 1]$
 $h_D^2 = [3, 2, 3, 4, 1, 2]$
 $h_E^2 = [2, 2, 2, 1, 3, 1]$
 $h_F^2 = [2, 2, 1, 2, 1, 3]$
- } Distinct embeddings



Distinguishability under sum

- Let $AGG=\text{sum}$, then for the GCN of 2 layers, the embeddings are as follows

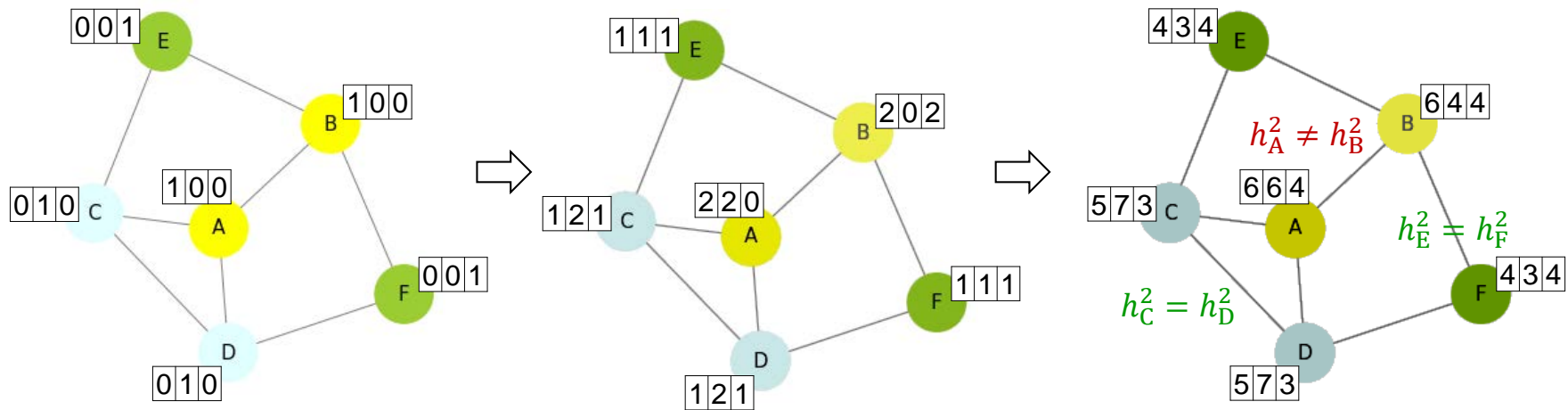


- By induction on the respective distinct feature dimension, the embeddings will be distinct for subsequent iterations
- If every node has **distinct feature**, then they have **distinct embeddings under sum** regardless of neighborhood structure or iterations (with the exception of the graph of only two nodes )

Distinguishability under **sum**

□ If some nodes have the same features?

- Let $h_A^0 = h_B^0 = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix}$, $h_C^0 = h_D^0 = \begin{bmatrix} 0 & 1 & 0 \end{bmatrix}$, $h_E^0 = h_F^0 = \begin{bmatrix} 0 & 0 & 1 \end{bmatrix}$, and let AGG be **sum**. Then, as can be seen from the following example



□ Two nodes with the **same feature** will have the same embedding under **sum** if they have the same **neighborhood structure**

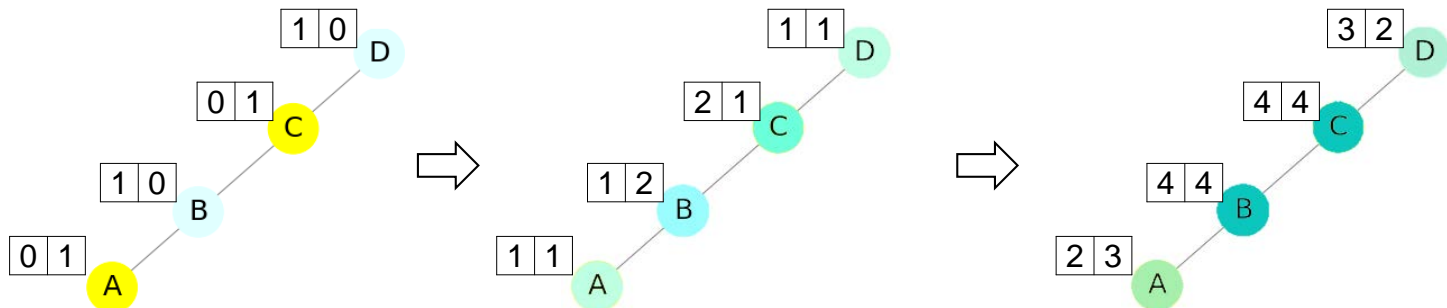
- However different features and neighborhood structure cannot guarantee distinct embeddings

Distinguishability under sum

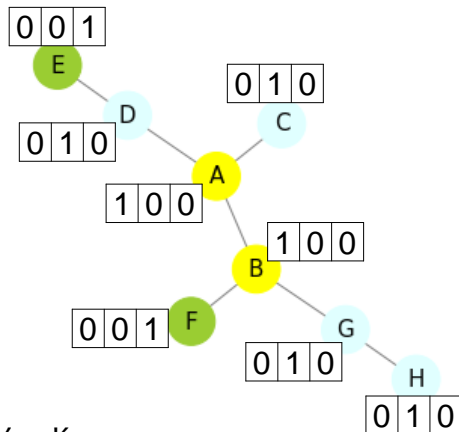
□ If some nodes have the same features?

- Different features and neighborhood structure cannot guarantee distinct embeddings for various reasons

□ Complementary neighborhood structure



□ Hard-to-predict cases



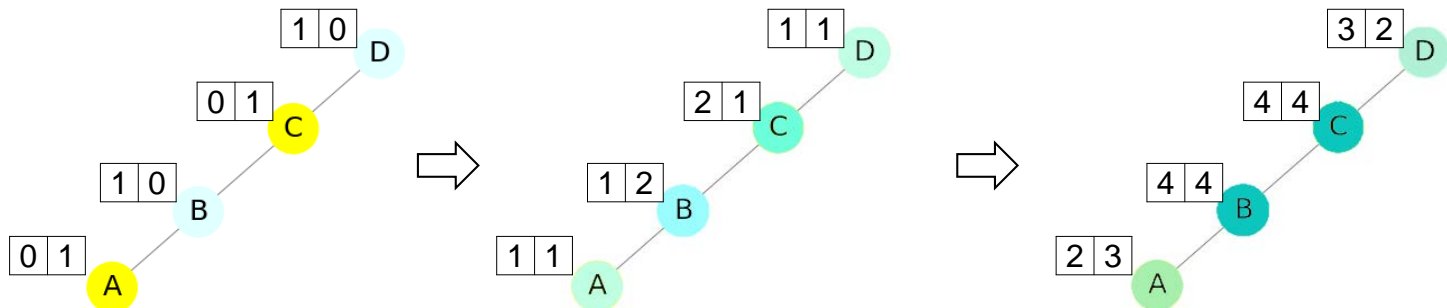
Will A and B ever become the same again after the first step?

Distinguishability under sum

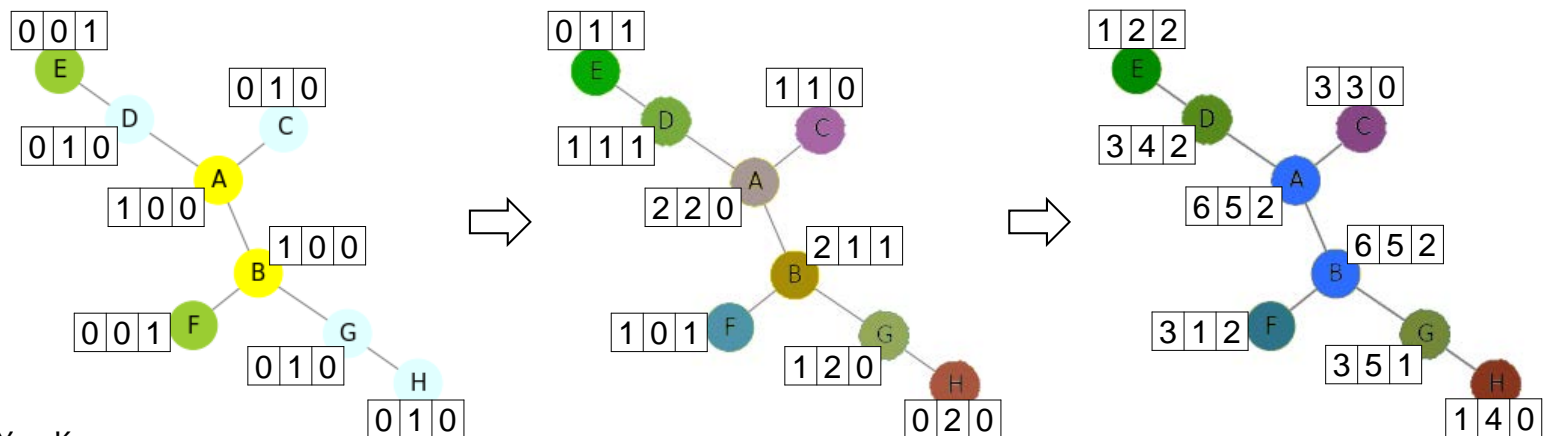
□ If some nodes have the same features?

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□ Complementary neighborhood structure

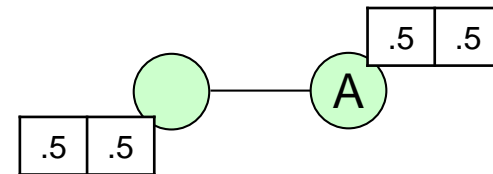
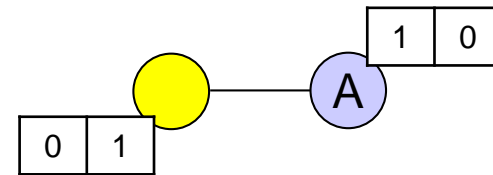
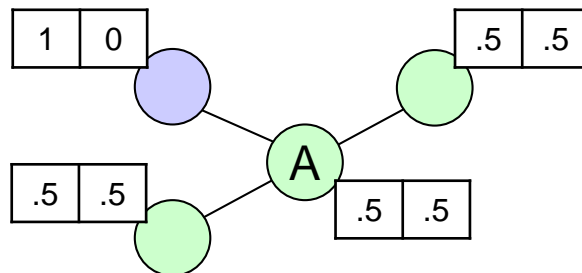
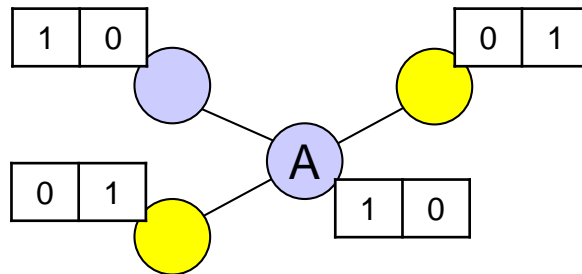


□ Hard-to-predict cases



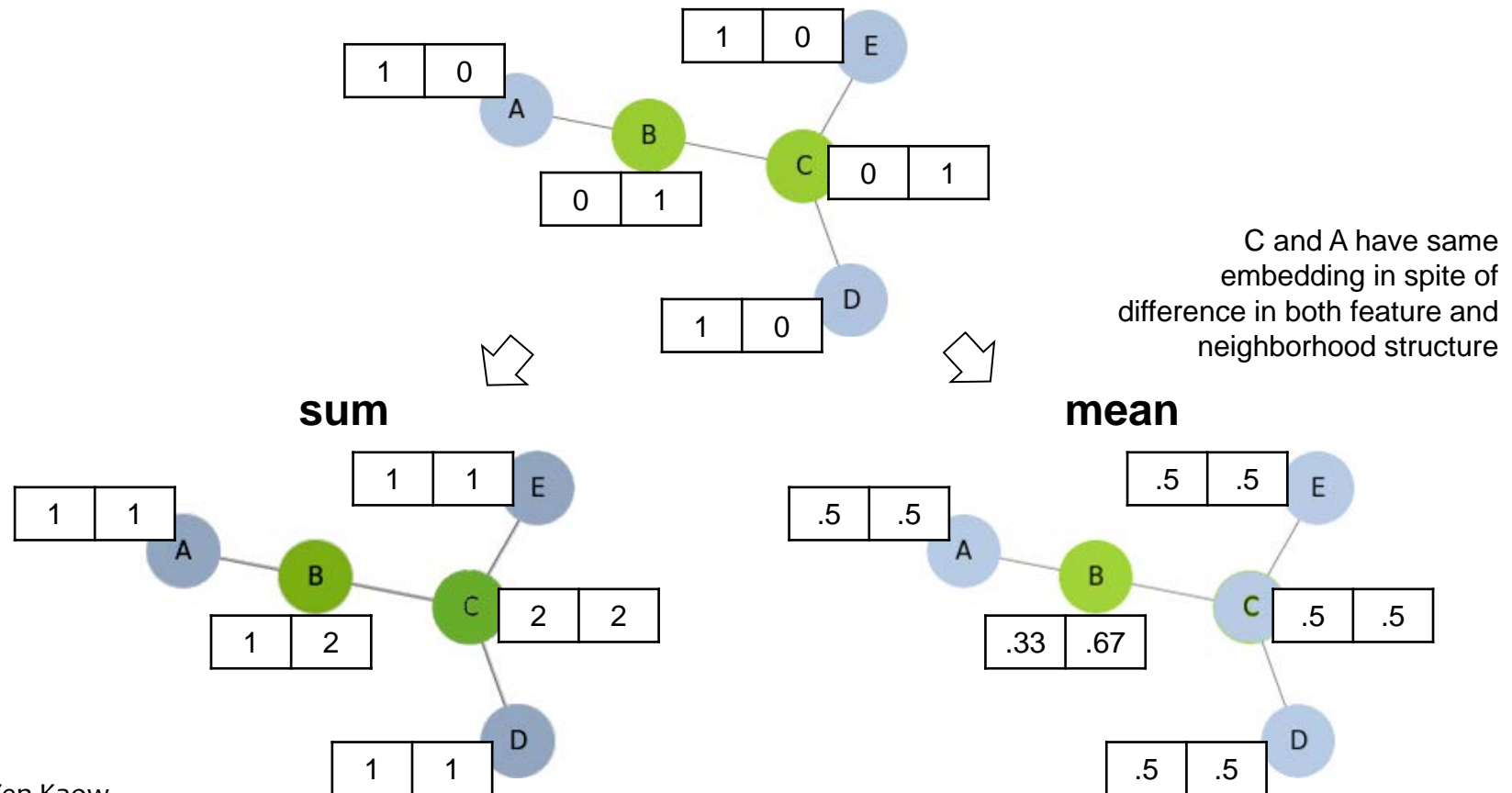
Distinguishability under **mean**

- Using **mean** as AGG results in even less desirable behavior
 - For instance, node A in both graphs below would give the same embedding under **mean** with one iteration



Distinguishability under **mean**

- Using **mean** as AGG results in even less desirable behavior
 - The following example shows **sum** to result in more consistent embeddings than **mean**



Injective function for isomorphism

- An **injective** function will output distinguishable embeddings for nodes of distinct feature and neighborhood structure
 - **sum, mean, and max are not injective**
- **Theorem** (Xu *et al.* 2019). Any **injective** AGG function can be expressed as $\Phi(\sum_{x \in S} f(x))$ for some non-linear Φ and linear f
- Since MLP is able to approximate any function, we can learn Φ and f with non-linear MLP_{Φ} and linear MLP_f

$$\text{AGG} = \text{MLP}_{\Phi} \left(\sum_{x \in S} \text{MLP}_f(x) \right)$$

⇒ **Graph Isomorphism Network (GIN)**

Subgraph Isomorphism

- Subgraph isomorphism is NP-complete
 - ⇒ Compare **neighborhood around each node**
- The **k -hop neighborhood around node $u \in Q$** is
 1. All the nodes within k hops from u , and
 2. All the edges in between those nodes
- Such a neighborhood is a subgraph of Q
 - However, not every subgraph of Q is a neighborhood of some node $u \in Q$
 - At most $k|V|$ neighborhoods for each k
- If P is a subgraph of Q , then every neighborhood of P is a subgraph of some neighborhood of Q

Order embedding space

- Idea: We want an d -dimensional embedding space z such that for every neighborhoods $p \in P$ and $q \in Q$

$$q \subseteq p \Leftrightarrow \forall_{i=1}^d z_q[i] \leq z_p[i]$$

- With embedding space, we can test subgraph isomorphism through the following

For each neighborhood $p \in P$ and $q \in Q$

If $(\exists i) [z_q[i] \leq z_p[i]]$, return false

Return true

Order embedding space

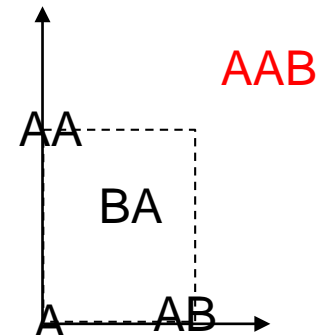
- Idea: We want an d -dimensional embedding space z such that for every neighborhoods $p \in P$ and $q \in Q$

$$q \subseteq p \Leftrightarrow \forall_{i=1}^d z_q[i] \leq z_p[i]$$

- Whether such an embedding space exist depends on D and the class of graphs
- Even for substring relation (\preceq) with only alphabet $\{A, B\}$, a 2D embedding space is insufficient for


$$q \preceq p \Leftrightarrow \forall_{i=1}^2 z_q[i] \leq z_p[i]$$




- AAB needs to cover the dotted box (since it includes both AA and AB)
- On the other hand, BA cannot be brought out of the dotted box (otherwise it would cover AA or AB)





Order embedding space

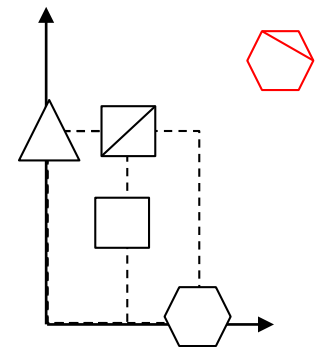
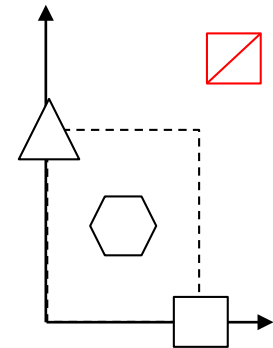
□ 2D embedding space example for graph

- □ needs to cover the dotted box (since it must cover both triangle and square) but that would cover 

- On the other hand,  cannot be brought out of the dotted box (otherwise it would cover  or )

- Toggling hexagon and square will allow  to be placed, but now  cannot be placed

- Assume that d is sufficiently large for reasonable embeddings



Training order embedding space

- Use a **node embedding space of k -hop**
- Denote the embedding of a node u as z_u
- Use the loss function

$$\text{loss}(u, v) = \sum_{i=1}^d \max(0, z_u[i] - z_v[i])^2$$

It is clear that

- $\text{loss}(u, v) = 0$ when $\forall_{i=1}^d (z_u[i] \leq z_v[i])$
- $\text{loss}(u, v) > 0$ otherwise
- Generate random pair of graphs P, Q and train GNN such that embeddings of $u_P \in P$ and $u_Q \in Q$ has
 - $\text{loss}(u_Q, u_P) = 0$ when $u_Q \subseteq u_P$, and
 - $\text{loss}(u_Q, u_P) > 0$ otherwise