

# Lecture 1: Intro, Stacks & Queues

CSE 332: Data Structures & Parallelism

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# Data Structures?

**Clever** ways to organize information in order to enable *efficient* computation over that information

# Trade-Offs

- A data structure strives to provide many useful, efficient operations
- But trade-offs!
  - Time vs. Space
  - One operation more efficient if another less efficient
  - Generality vs. Simplicity vs. Performance
- That is why there are many data structures

# Terminologies

- Abstract Data Type (ADT)
  - Mathematical description of a "thing" with set of operations on that "thing"
- Data Structures
  - A specific organization of data and family of algorithms for implementing an ADT
- Implementation of a data structure
  - The actual code implementation in a specific language
- Algorithm
  - A high level, language-independent description of a step-by-step process

# Terminology Example: Stack

- The **Stack** ADT supports operations:
  - **push**: adds an item
  - **pop**: raises an error if isEmpty, else returns *most-recently pushed item* not yet returned by a pop
  - **isEmpty**: initially true, later true if there have been same number of pops as pushes
  - etc.
- A Stack data structure could use a linked-list or an array or something else, and associated algorithms for the operations
- One implementation is in the library `java.util.Stack`

# Why useful

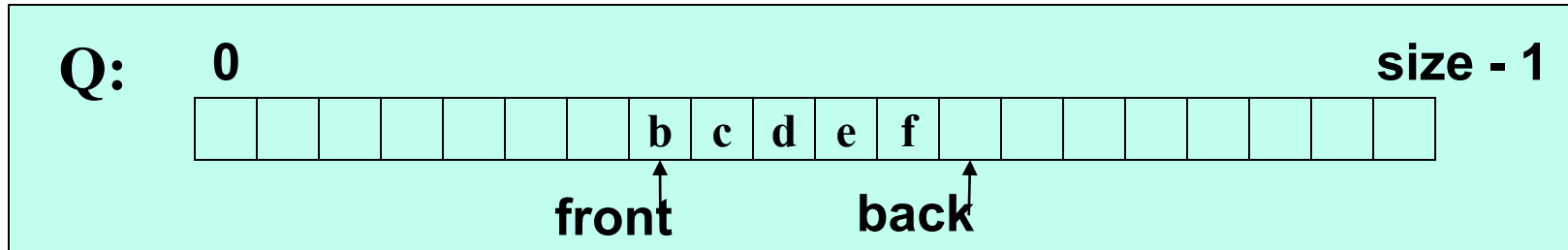
The ***Stack*** ADT is a useful abstraction because:

- It arises **all the time** in programming (see Weiss for more)
  - Recursive function calls
  - Balancing symbols (parentheses)
  - Evaluating postfix notation:  $3\ 4\ +\ 5\ *$
  - Clever: Infix  $((3+4) * 5)$  to postfix conversion (see Weiss)
- We can code up a **reusable library**
- We can **communicate** in high-level terms
  - “Use a stack and push numbers, popping for operators...”
  - Rather than, “create a linked list and add a node when...”

# Terminology Example: Queue

- The **Queue ADT** supports operations:
  - **enqueue**: adds an item at the end
  - **dequeue**: raises an error if isEmpty, else returns item at the start
  - **isEmpty**: initially true, later true if there have been same number of enqueue as dequeues
  - etc.
- A Queue data structure could use a linked-list or an array or something else, and associated algorithms for the operations
- One **implementation** is in the library **java.util.Queue**

# Circular Array Queue Data Structure



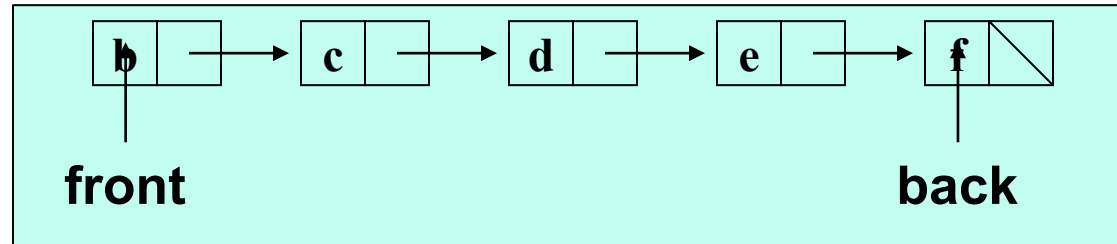
```
// Basic idea only!  
enqueue(x) {  
    Q[back] = x;  
    back = (back + 1) % size  
}
```

```
// Basic idea only!  
dequeue() {  
    x = Q[front];  
    front = (front + 1) % size;  
    return x;  
}
```

- What if **queue** is empty?
  - Enqueue?
  - Dequeue?
- What if **array** is full?
- How to *test* for empty?
- What is the *complexity* of the operations?



# Linked List Queue Data Structure



```
// Basic idea only!  
enqueue(x) {  
    back.next = new Node(x);  
    back = back.next;  
}
```

```
// Basic idea only!  
dequeue() {  
    x = front.item;  
    front = front.next;  
    return x;  
}
```

- What if **queue** is empty?
  - Enqueue?
  - Dequeue?
- Can **list** be full?
- How to *test* for empty?
- What is the *complexity* of the operations?