

Lecture 13:

Non-Comparison Sorting

CSE 332: Data Structures & Parallelism

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Take Handouts!

(Raise your hand if you need one)

Today

- Special Sorting Algorithm 1: Bucket Sort
- Special Sorting Algorithm 2: Radix Sort
- Sorting Summary

- Introduction to Graphs!

Special Sorting Algorithm 1: Bucket Sort

- Also called Bin Sort
- Intuition: Small range of integers, get a tally
- Algorithm:
 1. Find the `min` and `max` value
 2. Make an aux array to represent the range between `min` and `max`
 3. Go through the original array and start tallying each number
 4. Copy the aux into the original array

Bucket Sort: Visualization

size = n

4	3	1	2	1	1	2	3	4	2
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min = 1

max = 4

1 **2** **3** **4**

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size = k

size = n

1	1	1	2	2	2	3	3	4	4
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Any Questions?

Bucket Sort: Analysis

1. Stable?

- Yes!

2. In-Place?

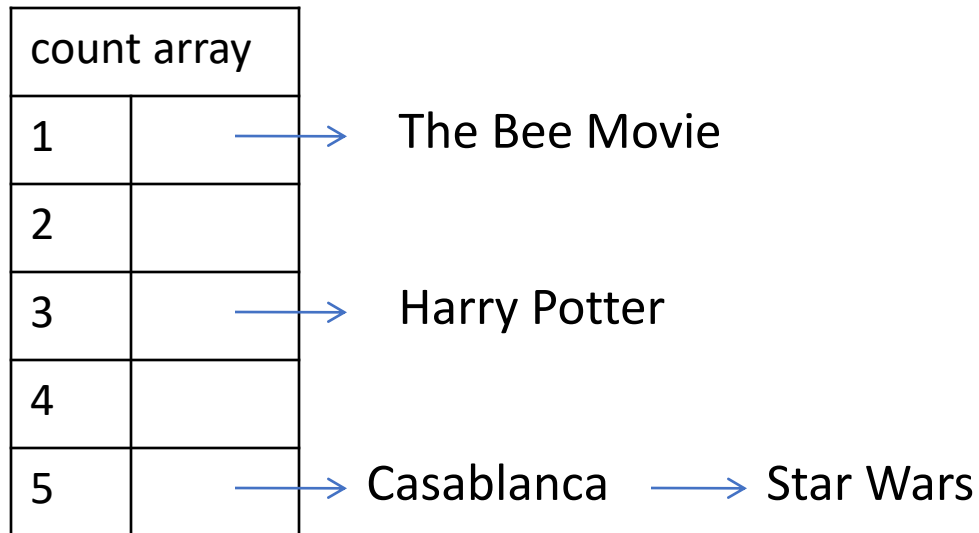
- No

3. Fast?

- Yes! (in terms of asymptotics)
 - Let k = range between `min` and `max` value
 - Worst Case: $\mathcal{O}(n + k)$
 - Note: k can be bigger than n
- **Good constant factors!**

Bucket Sort: Non Integers

- Most real lists aren't just #'s; we have data
- Each bucket is a list (say, linked list)
- To add to a bucket, place at end $\mathcal{O}(1)$ (keep pointer to last element)



- Example: Movie ratings:
1=bad,... 5=excellent
- Input=
5: Casablanca
3: Harry Potter movies
1: The Bee Movie
5: Star Wars

Result: 1: Rocky V, 3: Harry Potter, 5: Casablanca, 5: Star Wars
This result is **stable**; Casablanca still before Star Wars

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Special Sorting Algorithm 2: Radix Sort

- Intuition: Exploit integer digits
- Algorithm:
 1. For each digit,
Start from the **Least Significant Digit (LSD)** to **Most Significant Digit (MSD)**
 - Run bucket sort on just that digit (e.g., just the 1s place)

Radix Sort: Visualization (1s Place)

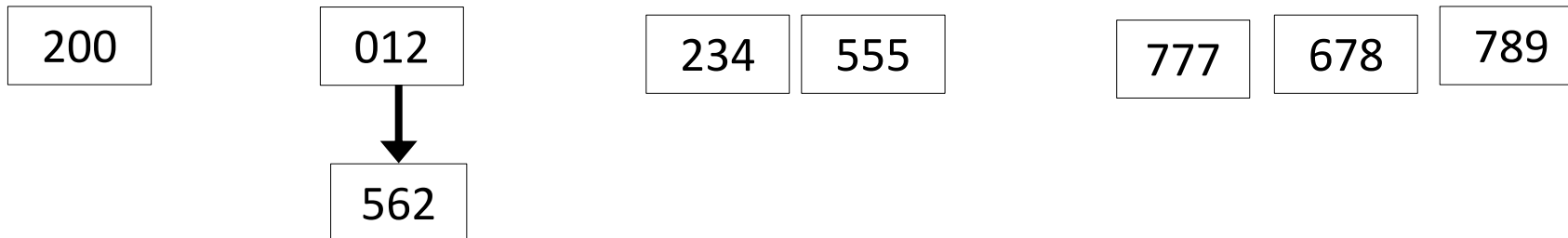
012	234	789	555	678	200	777	562
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0	1	2	3	4	5	6	7	8	9

Radix Sort: Visualization (1s Place)

012	234	789	555	678	200	777	562
------------	------------	------------	------------	------------	------------	------------	------------

0	1	2	3	4	5	6	7	8	9



200	012	562	234	555	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

Radix Sort: Visualization (10s Place)

200	012	562	234	555	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

0	1	2	3	4	5	6	7	8	9

Radix Sort: Visualization (10s Place)

200	012	562	234	555	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

0	1	2	3	4	5	6	7	8	9



200	012	234	555	562	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

Radix Sort: Visualization (100s Place)

200	012	234	555	562	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

0	1	2	3	4	5	6	7	8	9

Radix Sort: Visualization (100s Place)

200	012	234	555	562	777	678	789
------------	------------	------------	------------	------------	------------	------------	------------

0	1	2	3	4	5	6	7	8	9



012	200	234	555	562	678	777	789
------------	------------	------------	------------	------------	------------	------------	------------

Radix Sort: Analysis

1. Stable?

- Yes!

2. In-Place?

- No

3. Fast?

- Yes! (in terms of asymptotics)
 - Let n = input size
 - d = number of digits (e.g., previous example is 3 digits because only until 100s place)
 - k = range of values for each digit (e.g., 10 for base-10 integers)
 - Worst Case: $\mathcal{O}(d \cdot (n + k))$
 - Note: k can be bigger than n
- **Good constant factors!**

Radix Sort: Runtime

- Let n = input size
- d = number of digits (e.g., previous example is 3 digits because only until 100s place)
- k = range of values for each digit (e.g., 10 for base-10 integers)
- Worst Case: $\mathcal{O}(d \cdot (n + k))$
 - How many passes? d number of passes
 - Work per pass = 1 Bucket Sort = $\mathcal{O}(n + k)$

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