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LAZY PROPAGATION

0 **2** [help](#)

i understood segmented tree...but nowher i can see clear explanations of "LAZY PROPAGATION"...i searched topcoder,geeksforggeeks and all possible blogs and forums....Is ther any book on advanced data structures??...i want to clearly understand this lazy propagation....pls help..thanks...

asked 20 Feb, 03:47

 chaseme
 546 ● 9 ● 23 ● 43
 accept rate: 0%

2 Answers:

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Ok, here I go.

21 Consider a standard interval tree (usually called segment tree, but I use a transliteration from the Slovak term), which supports operations "set the value of node i to v " and "return the maximum value from the interval $[u, v]$ " (semi-closed intervals are comfortable to use). This tree keeps the following information in a node: the interval $[u_i, v_i]$, the maximum value from this interval and at most two sons of this node.

Updates are done recursively: after the value in one of a node's sons is updated, the value in the node itself is updated. Queries are answered this way: you always query on $(node, interval)$, where $interval$ is the subinterval of the one $[v_i, u_i]$ in the $node$. If $interval$ is empty, the answer's trivial; if it's equal to $[v_i, u_i]$, it's just $value$ of the $node$ and otherwise, you compute it recursively.

Now, what does lazy propagation (lazy-loading is the Slovak transliteration) do? You don't update "one node to value v ", but "interval $[x, y]$ to all values equal to v ". You also add one more value to each node: the *modifier*, and a single function *modify(node)*. The *modifier* holds the following information: "all elements in this interval are actually equal to x " or "ignore this" (if there's no fitting x , for example during initialization). The *modify()* function updates all sons' *modifiers* to this one (we need to make sure that the *modifier* of an ancestor has higher priority than the one of a son), sets the *value* in the node equal to the *modifier* and changes *modifier* to "ignore", to avoid collisions with later updates. This basically processes all updates in the given node, but not in the underlying subtree.

Before any call of functions *update()* or *query()*, you call *modify()* on this node. That way, you'll be certain that earlier updates in this node have been processed and there can't be any collisions.

Notice that *query()* hardly changes. You still do the same, just with calling *modify()* beforehand on the node on which you call *query()*, and time complexity stays the same: $O(\log N)$.

Updating changes a bit, because it affects the modifier. Firstly, *update(node, interval)* is also called with *interval* being a subinterval of the one in the *node*. It's called in a similar way to *update()*: if *interval* is empty, just return from the function call; if it's the whole interval stored in the node, all elements are updated to the same value, which means you only change the *modifier* and return; the rest of the work would be done when any function is called on this node later on. And in the last other case, with *interval* being a non-trivial one, you call *update()* on the sons (if the *interval* is non-trivial, both exist), set *value* to their maximum and return. Just remember: don't forget to call *modify(node)* before *update(node, interval)* or *query(node, interval)*.

We can see that any function call only takes $O(1)$ time. The trick is in the width of recursion. It can be shown that if both recursive calls (whether of *update()* or *query()*) are non-trivial, then it won't happen again, which means no further recursive branching to depth larger than 1, and therefore $O(\text{tree depth}) = O(\log N)$ calls (or time).

There are many other variants of lazy-loaded interval trees, which depend on the types of updates and queries, and on the number of them - it's possible to have two types of updates or of queries, etc., but you need to handle specific commands in specific order to avoid collisions, and it's quite complicated.

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answered 20 Feb, 06:05

 xellos0
 3.7k ● 5 ● 28 ● 50
 accept rate: 10%

Can you please explain this sentence: all elements in this interval are actually equal to x or "ignore this" (if there's no fitting x , for example during initialization). From what I get, the x should be v instead.

[rajbsk](#) (25 Sep, 05:49)

@xello..reali thanks for the explanation...also i got a aother source to leaern this concept here

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answered 20 Feb, 11:54



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