

30/7/24

# Matrix Computations

Frobenius norm =  $\|A_F\| = \left( \sum_{i=1}^n \sum_{j=1}^m |a_{ij}|^2 \right)^{1/2} = (\text{tr}(A^T A))^{1/2}$

$\|A_F\|$  is not induced by a vector norm but  $\|I\|_F = \sup_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2}$

but  $\|I\|_F = \sqrt{2}$

2/8/24

$M_i$  are called Gauss transforms

A lower triangular matrix with all ones in the diagonal is called a unit lower triangular matrix (like  $L$  in  $A=LU$ )

$$A^{(k)} = A^{(k-1)} - \begin{bmatrix} 0 & 0 \\ 0 & \hat{M}_k \end{bmatrix}$$

$$\hat{M}_k = \begin{bmatrix} m_{k+1,k} \\ \vdots \\ m_{n,k} \end{bmatrix} \begin{bmatrix} a_{kk}^{(k-1)} & \dots & a_{kn}^{(k-1)} \end{bmatrix}$$

$= (n-k) \times (n-k+1)$

$$\begin{bmatrix} k & n-k \\ n-k & n-k \end{bmatrix} \begin{bmatrix} A_{11}^{(k-1)} & A_{12}^{(k-1)} \\ A_{21}^{(k-1)} & A_{22}^{(k-1)} \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 \\ 0 & 0 \\ \vdots & \vdots \\ 0 & 0 \end{bmatrix}$$

$$\sum_{i=1}^k 2(n-i) = k+2(n-k) = n-k+1$$

obtained by shifting the column partition to the right

1st column of  $M$

$$\begin{bmatrix} a_{k+1,k}^{(k-1)} \\ \vdots \\ a_{n,k}^{(k-1)} \end{bmatrix}$$

$$Ax=b$$

$$LUx=b$$

$A=LU$  till now requires diagonal entries are non-zero

the diagonal entries are all 1 otherwise it won't be a unique decomposition

$$A=LU \Leftrightarrow \begin{bmatrix} k & n-k \\ n-k & n-k \end{bmatrix} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} k & n-k \\ n-k & n-k \end{bmatrix} \begin{bmatrix} L_{11} & \\ L_{21} & L_{22} \end{bmatrix} \begin{bmatrix} k & n-k \\ n-k & n-k \end{bmatrix} \begin{bmatrix} U_{11} & U_{12} \\ & U_{22} \end{bmatrix}$$

$$A_{11} =$$

4/5/24

Orthogonal  $P^T = P^{-1}$

5/5/24

$$A = LU$$

$$= L \begin{bmatrix} u_{11} & & 0 \\ & \ddots & \\ 0 & & u_{nn} \end{bmatrix} \begin{bmatrix} u_{11}^{-1} & & 0 \\ & \ddots & \\ 0 & & u_{nn}^{-1} \end{bmatrix} U$$

$\downarrow D$ 
 $\downarrow D^{-1}$

$$= L D V \quad [V = D^{-1} U]$$

16/5/24

For complex case (in  $\mathbb{R}^n$  for positive definite),  
 $x^* A x > 0 \Rightarrow x^* A x$  is a real (not complex)

\*  $A$  is positive definite  $\Rightarrow$  non singular (PD)

4: Suppose  $A$  is singular  $A = [A_1 \ A_2 \ \dots \ A_n]$

$\exists x_1, \dots, x_n$  (not all zero)

$$\sum x_i A_i = 0 \Rightarrow Ax = 0$$

Principal submatrix: Remove rows by corresponding columns (say row 3, 5, 6). All principal submatrices of a PD matrix are PD matrices.

4:  $\hat{A} = A[i_1, i_2, \dots, i_p]$

$$A^T = A \Rightarrow \hat{A}^T = \hat{A}$$

Let  $\hat{x} \in \mathbb{R}^p \setminus \{0\}$ . Let  $x \in \mathbb{R}^n \Rightarrow x_i = 0$  if  $i \notin \{i_1, \dots, i_p\}$   
 and  $x_i = \hat{x}_i$  if  $i \in \{i_1, \dots, i_p\}$

$$\text{Then } x \neq 0 \text{ and } 0 < x^T A x = \hat{x}^T \hat{A} \hat{x}$$

$\Rightarrow A$  is PD,  $x^T A x$  is PD (X n X)  
 $\rightarrow$  positive definite by  $S$  is non-singular s.t.  $S^T A S$  is defined.

$$\text{Let } B = S^T A S.$$

$$B^T = S^T A^T S = B$$

$$\text{Let } x \in \mathbb{R}^n \setminus \{0\} \quad x^T B x = x^T S^T A S x = y^T A y \text{ where } y \neq 0$$

$\left[ \begin{array}{l} S^{-1} \text{ exists} \\ y \neq 0 \end{array} \right]$

→  $A$  is PD  $\Rightarrow$  all the diagonal entries are  $+ve$ .

Prf:-  $e_i^T A e_i = a_{ii} > 0$

Inner product formulat<sup>n</sup> algo:-

Square root is 8 flops.

$\sum_{k=1}^{j-1} g_{kj}^2$  &  $\sum_{i=1}^{j-1} g_{ij} g_{ik}$  are inner products of rows/columns hence the name.

Flop count:-  $\frac{n^3}{3} + O(n^2)$  flops ( $\frac{1}{2}$  of LU decomposition!)

Outer product formulat<sup>n</sup> algo:-

( $i$ th col is found, then  $(i+1)$ th col)

Bordered form:-

$i$ th row,  $i$ th col &  $c_j$  then  $(i+1)$ th row,  $(i+1)$ th col

Sensitivity:-

The choice of norm (even for Frobenius norm) doesn't matter.

19/8/24 Sensitivity analysis:-

with  $\text{mag}(A) = 0$  when  $A$  is singular.

invhilb does not use Hilbert matrices to invert.

Dealing with ill conditioning when one column is larger than others:-

$$Ax = b$$

$$A = \begin{bmatrix} 1 & 2000 \\ 3 & -1000 \end{bmatrix}$$

$$A \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} x = b$$

see this matrix correctly in slides.

$$\begin{bmatrix} 1 \\ 1 \end{bmatrix}$$



Ill conditioning when row/cols are "nearly linearly dependent"

$$K(A) = \dots \quad B = A / \|A\|$$

$$K(B) = \frac{\max \text{mag}(A)}{\min \text{mag}(A)} = \frac{1}{\min \text{mag}(B)} \Rightarrow \min \text{mag}(B) \ll 1$$

$$\Rightarrow \exists x_0 \neq 0 \rightarrow \frac{\|Bx_0\|}{\|x_0\|} \ll 1$$

~~Let~~

→ However,  $\det A \approx 0$  is not always indicative of ill conditioning.

but we can multiply  $cA$  to make  $\det(cA)$  large or small.

$$\text{Let } \hat{x}_0 = \frac{x_0 \|A\|}{\|x_0\| \|A\|}$$

$$Bx_0 = B(\|A\| \|x_0\| \hat{x}_0) = A(\|x_0\| \hat{x}_0)$$

$$\frac{\|Bx_0\|}{\|x_0\|} = \frac{\|A(\|x_0\| \hat{x}_0)\|}{\|x_0\|}$$

$$= \|A \hat{x}_0\|$$

$$\therefore \|A \hat{x}_0\| \ll 1$$

$$\Rightarrow A \hat{x}_0 \approx 0$$

20/8/24 Finite precision systems

$$x_f = x_- + \beta^{-(p-1)} \times \beta^E$$

$$\frac{|f(x) - x|}{|x|} \leq \frac{|x_f - x_-|}{2|x|} = \frac{\beta^{1-p} \times \beta^E}{2 \times \beta^E} \leq \frac{\beta^{1-p}}{2}$$

★ Note that  $N_{\min}$  isn't the smallest no that can be represented in IEEE. In single precision,  $N_{\min} = 2^{-126}$  but smallest no represented is  $0.\underbrace{00\dots0}_{22}1 \times 2^{-126} = 2^{-149}$

In exercise, it says IEEE hidden bit = 0  
 $\& \ p=24$  (not 23)

$$\cancel{Ax=b}$$

$$(A+\delta A)x = b+\delta b$$

→ If  $A$  is non singular or  $\frac{\|\delta A\|}{\|A\|} < \frac{1}{K(A)}$  then  $A+\delta A$  is non-singular

Pf: Suppose  $A+\delta A$  is singular  
 →  $(A+\delta A)y = 0$  for some  $y$   
 $y = -A^{-1}\delta Ay = 0$

$$\|y\| = \|A^{-1}\delta Ay\| \leq \|A^{-1}\| \|\delta A\| \|y\|$$

$$1 \leq \|A^{-1}\| \|\delta A\|$$

→ Let  $A$  be non singular. Let  $\delta A \in \mathbb{F}^{n \times n}$   $\rightarrow \frac{\|\delta A\|}{\|A\|} < \frac{1}{K(A)}$

Consider  $Ax=b$  or  $(A+\delta A)(x+\delta x) = b+\delta b$

Then  $\frac{\|\delta x\|}{\|x\|} \leq K(A) \left( \frac{\|\delta A\|}{\|A\|} + \frac{\|\delta b\|}{\|b\|} \right) / \left( 1 - K(A) \cdot \frac{\|\delta A\|}{\|A\|} \right)$

Pf:

$$\delta Ax + A\delta x + \delta A\delta x = \delta b \quad = \|\delta A\| \|A^{-1}\| + \frac{\|\delta b\|}{\|b\|} \|A\| \|A^{-1}\|$$

$$\delta A\delta x = \delta b - A\delta x - \delta Ax$$

$$A\delta x = \delta b - \delta A\delta x - \delta Ax$$

$$\delta x = A^{-1}\delta b - A^{-1}\delta A\delta x - A^{-1}\delta Ax$$

$$\|\delta x\| \leq \|A^{-1}\| \|\delta b\| + \|A^{-1}\| \|\delta A\| \|\delta x\| + \|A^{-1}\| \|\delta A\| \|x\|$$

$$\|\delta x\| (1 - \|A^{-1}\| \|\delta A\|) \leq \|A^{-1}\| \|\delta b\| + \|A^{-1}\| \|\delta A\| \|x\|$$

$$\frac{\|\delta x\|}{\|x\|} \leq \frac{\|A^{-1}\| \|\delta b\| / \|x\| + \|A^{-1}\| \|\delta A\|}{1 - \|A^{-1}\| \|\delta A\|}$$

$$\leq \frac{\|A^{-1}\| \|\delta b\| \frac{\|A\|}{\|b\|} + \|A^{-1}\| \|\delta A\|}{1 - \|A^{-1}\| \|\delta A\|}$$

$$Ax=b$$

$$\|b\| \leq \|A\| \|x\|$$

23/0/24  
Rounding errors:-

$$f\left(\frac{x}{y}\right) = f\left(\frac{f(x)}{f(y)}\right) = \left(\frac{x(1+\epsilon_1)}{y(1+\epsilon_2)}\right)(1+\epsilon_3) = \frac{x}{y}(1+\epsilon_1)(1+\epsilon_3)(1+\epsilon_2)^{-1}$$

$|\epsilon_i| \leq u$

$$= \frac{x}{y} \left( 1 + \epsilon_1 + \epsilon_3 - \epsilon_2 + \epsilon_1\epsilon_3 - \epsilon_1\epsilon_2 - \epsilon_3\epsilon_2 - \epsilon_1\epsilon_2\epsilon_3 + \dots \right)$$

$$\therefore |\epsilon| \leq 3u + O(u^2)$$

$$\therefore f\left(\frac{x}{y}\right) = \left(\frac{x}{y}\right)(1+\epsilon)$$

$$f(x+y) = f(f(x) + f(y)) = (x(1+\epsilon_1) + y(1+\epsilon_2))(1+\epsilon_3)$$

$$= (x+y) \left( 1 + \frac{(\epsilon_1 + \epsilon_3 + \epsilon_1\epsilon_3)x}{x+y} + \frac{(\epsilon_2 + \epsilon_3 + \epsilon_2\epsilon_3)y}{x+y} \right)$$

$$|\epsilon| = \left| \frac{(\epsilon_1 + \epsilon_3 + \epsilon_1\epsilon_3)x}{x+y} + \frac{(\epsilon_2 + \epsilon_3 + \epsilon_2\epsilon_3)y}{x+y} \right|$$

$$\leq \left( \frac{|x|}{|x+y|} + \frac{|y|}{|x+y|} \right) [2u + O(u^2)]$$

Ex:-  $(10, 3, -2, 2)$  yrs    Doing  $(99.93 + 0.026) - (100 + 0.5)$

Swamping  $\rightarrow$  catastrophic cancellat<sup>n</sup>

$$f(99.93) = 9.99 \times 10 \quad f(0.026) = 0.026 \quad \therefore f(99.93 + 0.026) = 9.99 \times 10$$

$$f(100.5) = f(1.005 \times 10^2) = 100$$

Backward analysis:-

$$f(x+y) = f(f(x) + f(y)) = (x(1+\epsilon_1) + y(1+\epsilon_2))(1+\epsilon_3)$$

$$= x(1+\epsilon_1)(1+\epsilon_3) + y(1+\epsilon_2)(1+\epsilon_3)$$

$$= x(1+\hat{\epsilon}) + y(1+\hat{\epsilon})$$

$$= \hat{x} + \hat{y}$$

Why, All 4 ops are backward stable

where  $|\epsilon| \leq 2u + O(u^2)$   
 $|\hat{\epsilon}| \leq 2u + O(u^2)$



24/8/24

Thm:- Let  $w_i, i=1, \dots, n$  be floating pt nos. Then there exist  $\gamma_i$   
 $i=1, \dots, n$  satisfying  $|\gamma_i| \leq (n-1)u + O(u^2)$  &

$$f1\left(\sum_{i=1}^n w_i\right) = \sum_{i=1}^n w_i(1+\gamma_i)$$

Pf:-  $n=2$   $f1(w_1 + w_2) = \bar{w}_1 + \bar{w}_2$  where  $\frac{|\bar{w}_1 - w_1|}{|w_1|} \leq \frac{|\bar{w}_2 - w_2|}{|w_2|}$

Let  $\gamma_i = \frac{w_i - \bar{w}_i}{w_i}$  - Then  $\bar{w}_i = w_i(1+\gamma_i)$  where  $\gamma_i \leq u + O(u^2)$

$$\therefore f1(w_1 + w_2) = \sum_{i=1}^2 w_i(1+\gamma_i) \text{ where } |\gamma_i| \leq u + O(u^2) \forall i=1,2$$

Suppose  $1 \leq k \leq n-1$ .  $f1\left(\sum_{i=1}^n w_i\right) = f1\left(f1\left(\sum_{i=1}^k w_i\right) + f1\left(\sum_{i=k+1}^n w_i\right)\right) \rightarrow (3)$

$$f1\left(\sum_{i=1}^k w_i\right) = \sum_{i=1}^k w_i(1+\gamma_i), \quad |\gamma_i| \leq (k-1)u + O(u^2) \rightarrow (1)$$

$$f1\left(\sum_{i=k+1}^n w_i\right) = \sum_{i=k+1}^n w_i(1+\gamma_i'), \quad |\gamma_i'| \leq (n-k-1)u + O(u^2) \rightarrow (2)$$

$$\begin{aligned} \therefore f1\left(\sum_{i=1}^n w_i\right) &= f1\left(\sum_{i=1}^k w_i(1+\gamma_i) + \sum_{i=k+1}^n w_i(1+\gamma_i')\right) \\ &= \left(\sum_{i=1}^k w_i(1+\gamma_i) + \sum_{i=k+1}^n w_i(1+\gamma_i')\right)(1+\epsilon), \text{ where } |\epsilon| \leq u \\ &= \sum_{i=1}^n w_i(1+\gamma_i) \end{aligned}$$

$$\gamma_i = \gamma_i + \epsilon + \gamma_i \epsilon \text{ for } i=1, \dots, k$$

$$\gamma_i = \gamma_i' + \epsilon + \gamma_i' \epsilon \text{ for } i=k+1, \dots, n$$

$$\text{for } i=1:k, |\gamma_i| \leq |\gamma_i| + |\epsilon| + |\gamma_i||\epsilon| \leq (k-1)u + u + O(u^2) \leq (n-1)u + O(u^2)$$

Only for  $i=k+1:n$

Backward stability analysis for forward/backward substitution:-

$$\hat{y}_i = f1\left[b_i - \sum_{j=1}^{i-1} q_{ij} \hat{y}_j\right] \quad \forall i=1, \dots, n$$

$$f1\left(b_i - \sum_{j=1}^{i-1} q_{ij} \hat{y}_j\right) = f1\left[f1(b_i) - f1\left(\sum_{j=1}^{i-1} q_{ij} \hat{y}_j\right)\right]$$

$$= \sum_{j=1}^{i-1} q_{ij} \hat{y}_j(1+\alpha_{ij}) \text{ where } \alpha_{ij} \leq (i-1)u + O(u^2)$$

$$= b_i (1 + \delta_{ii}) - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j (1 + \alpha_{ij}) (1 + \gamma_{ij}) \quad \text{where } |\delta_{ij}| \leq \frac{(i-1)u}{1 + \alpha_{ij}}$$

$$\therefore \hat{y}_i = \left[ \frac{b_i (1 + \delta_{ii}) - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j (1 + \alpha_{ij}) (1 + \gamma_{ij})}{g_{ii}} \right] (1 + \epsilon_i) \quad \begin{matrix} \text{where } |\epsilon_i| \leq u \\ \text{being added/} \\ \text{subtracted} \end{matrix}$$

$$= b_i - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j \left( \frac{(1 + \alpha_{ij}) (1 + \gamma_{ij})}{1 + \delta_{ii}} \right) = 1 + \delta_{ij}$$

$$\frac{g_{ii}}{(1 + \delta_{ii}) (1 + \epsilon_i)} = 1 + \delta_{ii}$$

bound these using the prev bounds given

$$\therefore \hat{y}_n = b_i - \sum_{j=1}^{i-1} g_{ij} (1 + \delta_{ij}) \hat{y}_j$$

Alternate proof:-

$$\hat{y}_i = \text{fl} \left[ (b_i - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j) / g_{ii} \right] \quad \text{it is already a float - you don't need to write float again.}$$

$$= \text{fl} \left[ \text{fl} \left[ b_i - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j \right] / g_{ii} \right]$$

is an inner product  $b_i \cdot 1 - g_{ij} \hat{y}_1 - g_{ij} \hat{y}_2 - \dots$

$$= \text{fl} \left[ \frac{b_i (1 + \delta_{ii}) - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j (1 + \delta_{ij})}{g_{ii}} \right] \quad \text{where } |\delta_{ij}| \leq iu + \alpha_{ij} \quad \forall j=1, \dots, i-1$$

$$= \frac{b_i (1 + \delta_{ii}) - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j (1 + \gamma_{ij})}{g_{ii}} (1 + \epsilon_i) \quad \text{where } |\epsilon_i| \leq u$$

$$= b_i - \sum_{j=1}^{i-1} g_{ij} \hat{y}_j (1 + \delta_{ij}) (1 + \delta_{ii})^{-1}$$

$$\text{Let } 1 + \delta_{ij} = \begin{cases} (1 + \alpha_{ij}) (1 + \delta_{ii})^{-1} & \text{for } i \neq j, \dots, i-1 \\ (1 + \delta_{ii})^{-1} (1 + \epsilon_i) & \text{for } j = i \end{cases}$$

$$|\delta_{ij}| \leq (i-1)u + \alpha_{ij}$$



$$\delta G_{ij} = q_{ij} \delta y_j \quad j=1, \dots, n$$

$$\therefore \hat{y}_i = \left[ b_i - \sum_{j=1}^{i-1} (q_{ij} + \delta q_{ij}) \hat{y}_j \right] / (q_{ii} + \delta q_{ii})$$

Backward error analysis of Gaussian Elimination:-

$$a_{ij} = \sum_{k=1}^{\min\{i,j\}} l_{ik} u_{kj} = \sum_{k=1}^{j-1} l_{ik} u_{kj} + l_{ij} u_{jj}$$

$$l_{ij} = (a_{ij} - \sum_{k=1}^{j-1} l_{ik} u_{kj}) / u_{jj}$$

$$\hat{l}_{ij} = fl \left[ a_{ij} - \sum_{k=1}^{j-1} \hat{l}_{ik} u_{kj} \right]$$

← In book

(11) <sup>in slides</sup> →

$$A + \delta A = L_c U_c \rightarrow \textcircled{1} \quad |\delta A| \leq 2nu |L_c| |U_c| + O(u^2)$$

$$(L_c + \delta L_c) y_c = b \rightarrow \textcircled{2} \quad |\delta L_c| \leq 2nu |L_c| + O(u^2)$$

$$(U_c + \delta U_c) x_c = y_c \rightarrow \textcircled{3} \quad |\delta U_c| \leq 2nu |U_c| + O(u^2)$$

$$\textcircled{2} \text{ or } \textcircled{3} \Rightarrow (L_c + \delta L_c)(U_c + \delta U_c)x_c = b \Rightarrow \left[ L_c U_c + L_c \delta U_c + \delta L_c U_c + \delta L_c \delta U_c \right] x_c = b$$

$$\left[ A + \delta A + L_c \delta U_c + \delta L_c U_c + \delta L_c \delta U_c \right] x_c = b$$

$$|E| \leq |\delta A| + |L_c| |\delta U_c| + |\delta L_c| |U_c| + |\delta L_c| |\delta U_c|$$

$$\leq 2nu + 2u$$

2/9/24

$$\|\delta A\| \leq 2nu \|G_c^T\| \|G_c\| + O(u^2)$$

$$\Rightarrow \|\delta A\|_F \leq 2nu \|G_c^T\|_F \|G_c\|_F + O(u^2) \\ = 2nu \|G_c\|_F^2 + O(u^2)$$

$$\|G_c\|_F^2 = \text{tr}(G_c^T G_c) = \text{tr}(A + \delta A) = \text{tr}(A) + \text{tr}(\delta A)$$

$$\begin{aligned} \therefore \|\delta A\|_F &\leq 2nu \sqrt{n} (\|A\|_F + \|\delta A\|_F) + O(u^2) \\ &\leq \sum a_{ii} + \sum \delta a_{ii} \\ &\leq \sum |a_{ii}| + \sum |\delta a_{ii}| \\ &\leq \sqrt{n} (\sum |a_{ii}|^2)^{1/2} + \sqrt{n} (\sum |\delta a_{ii}|^2)^{1/2} \\ &\leq \sqrt{n} (\|A\|_F^2 + \|\delta A\|_F^2)^{1/2} \\ &\leq \sqrt{n} (\|A\|_F + \|\delta A\|_F) \end{aligned}$$

$$\Rightarrow \|\delta A\|_F (1 - 2n^{3/2}u) \leq 2n^{3/2}u \|A\|_F + O(u^2)$$

Least Squares Problem:-

$$[A_1 \dots A_m] \begin{bmatrix} x_1 \\ \vdots \\ x_m \end{bmatrix} = \sum x_i A_i \in \text{col}(A) = \text{Range}(A) \\ = \{Ay \mid y \in \mathbb{R}^m\}$$

Linear map def<sup>n</sup>:- A is linear map if

$$1) A(x+y) = Ax + Ay$$

$$2) A(\alpha x) = \alpha Ax$$

$$S^\perp = \{u \in \mathbb{R}^n, \langle u, s \rangle = 0 \forall s \in \mathbb{R}^n\}$$

3/4/24

Normal Eq<sup>n</sup>'s Method:- TFAE:-

$$\text{Rank } A = m \xrightarrow{\quad} Ax \neq 0 \forall x \Rightarrow$$

$$1) A^T A \text{ is non-singular}$$

~~2) A is non-singular~~

Condensed QR decomposition:-

$$A = QR$$

non-singular or Isometry, obtained by taking linearly independent columns of non-singular unitary matrix so

$\therefore R$  is also non-singular

$$\langle Qx, Qy \rangle = (Qy)^T (Qx) = y^T Q^T Q x = y^T x = \langle x, y \rangle$$

$$\|QB\|_2 = \max_{x \neq 0} \frac{\|QBx\|}{\|x\|} = \max_{x \neq 0} \frac{\|Bx\|}{\|x\|} = \|B\|_2$$

$$\|Q\|_2 = \max_{x \neq 0} \frac{\|Qx\|}{\|x\|} = \max_{x \neq 0} \frac{\|x\|}{\|x\|} = 1$$

$$\kappa_2(Q) = \frac{\max \text{mag}(Q)}{\min \text{mag}(Q)} = \frac{\max \frac{\|Qx\|_2}{\|x\|_2}}{\min \frac{\|Qx\|_2}{\|x\|_2}} \quad \text{but } \|Qx\|_2 = \|x\|_2$$

$$= 1$$

$\Rightarrow QQ^T$  is the orthogonal projector onto  $\text{span}\{q_1, \dots, q_m\}$

Pf:  $QQ^T q_i$

$$= [q_1 \ q_2 \ \dots \ q_m] \begin{bmatrix} q_1^T \\ q_2^T \\ \vdots \\ q_m^T \end{bmatrix} q_i = [q_1 q_1^T + q_2 q_2^T + \dots] q_i$$

$$= q_i q_i^T q_i + 0 + 0 + \dots$$

$$= q_i$$

$$q_i^T q_i = 1$$

$$q_i^T q_j = 0 \quad (i \neq j)$$

$$\text{Hence } QQ^T w = 0$$



6/9/24 Orthogonal

Project property:

If  $\mathbb{R}^n = U \oplus U^\perp$  and if  $Px = x$ , where  $x \in U$ , find  $P$ .

Ans:  $U = \{u_1, \dots, u_p\}$

$P = U(U^*U)^{-1}U^*$ . Then  $P^2 = P$  and  $P^* = P$ .

$Pu_i = U(U^*U)^{-1}U^*u_i = U(U^*U)^{-1}U^*Ue_i = Ue_i = u_i$

$v \in U$ . Then  $v = \sum_{i=1}^p \alpha_i u_i$  and  $Pv = P(\sum \alpha_i u_i) = \sum \alpha_i Pu_i = \sum \alpha_i u_i = v$

$v \in U^\perp$ .  $Pv = U(U^*U)^{-1}U^*v = 0$

2. In Least Squares Problem, If columns of  $A$  are l.i.,

$P = A(A^*A)^{-1}A^*$

$b_1 = Pb = A(A^*A)^{-1}A^*b$

rank  $A = m$

$A = Q \begin{bmatrix} R \\ 0 \end{bmatrix}$  is a condensed QR decomposition

$[Q^T Q = I_m, R \text{ is upper } \Delta^{lar} \text{ i.e. } m \times m]$

$[A_1 \dots A_m] = [q_1 \dots q_m] \begin{bmatrix} r_{11} & \dots & r_{1m} \\ & \ddots & \vdots \\ & & r_{mm} \end{bmatrix}$

$A_1 = q_1 r_{11}, A_2 = r_{12} q_1 + r_{22} q_2, \dots$

$\text{span}\{A_1, \dots, A_k\} = \text{span}\{r_{11} q_1, \sum_{i=1}^k r_{i2} q_i, \dots, \sum_{i=1}^k r_{ik} q_i\}$

Here now  $R^{-1}$  is upper  $\Delta^{lar}$  same as above how we proved  $\subseteq \text{span}\{q_1, \dots, q_k\}$

$Q = AR^{-1} \Rightarrow \text{span}\{q_1, \dots, q_k\} \subseteq \text{span}\{A_1, \dots, A_k\}$

$AR = Q$

full rank full rank  
 $\therefore R$  is full rank

$U = [u_1 \dots u_p]$  basis