Solving Triangular Systems

A Lower Triangular System of Equations

Consider

$$Lx = b$$

where

$$L = [I_{ii}] \leftarrow n \times n$$
 lower triangular matrix

$$b \leftarrow n \times 1$$
 vector.

In expanded form,

$$\begin{aligned}
 I_{11}x_1 &= b_1 \\
 I_{21}x_1 + I_{22}x_2 &= b_2 \\
 \vdots &\vdots &\vdots \\
 I_{k1}x_1 + I_{k2}x_2 + \dots + I_{kk}x_k &= b_k \\
 \vdots &\vdots &\vdots \\
 I_{n1}x_1 + I_{n2}x_2 + \dots + I_{nk}x_k + \dots + I_{nn}x_n &= b_n
 \end{aligned}$$

Row Oriented Forward Substitution

$$x_1 = b_1/I_{11}; \quad x_k = \left(b_k - \sum_{j=1}^{k-1} I_{kj} x_j\right)/I_{kk}, \ k=2,\ldots,n;$$
 for $k=1:n$ for $j=1:k-1$
$$b_k = b_k - I_{kj} b_j;$$
 end if $I_{kk} \neq 0$
$$b_k = b_k/I_{kk};$$
 else exit {'matrix is singular'} end end

Row Oriented Forward Substitution

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end

Total flops:

$$\overline{1 + \sum_{k=2}^{n} \left\{ \left(\sum_{j=1}^{k-1} 2 \right) + 1 \right\}} = \sum_{k=1}^{n} 2(k-1) + n = n^2.$$



Column Oriented Forward Substitution

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\begin{array}{l} \text{for } j=1:n \\ \qquad \qquad \text{if } I_{jj} \neq 0 \\ \qquad \qquad \qquad b_j = b_j/I_{jj}; \qquad \text{(1 flop)} \\ \qquad \qquad \text{end} \\ \qquad \qquad \text{for } i=j+1:n \qquad \leftarrow \text{entries of $L$ are being accessed column wise} \\ \qquad \qquad b_i = b_i - I_{ij}b_j; \qquad \text{(2 flops)} \\ \qquad \qquad \text{end} \\ \end{array}
```

Total flops: n^2 .

Likewise, $n \times n$ upper triangular systems may be solved in n^2 flops each by performing row and column oriented backward substitutions.

Blocked forward subtitution:

The $n \times n$ lower triangular system system Lx = b may be rewritten as

$$\left[\begin{array}{cc} L_{11} \\ L_{21} \end{array}\right] \left[\begin{array}{c} x_1 \\ x_2 \end{array}\right] = \left[\begin{array}{c} b_1 \\ b_2 \end{array}\right],$$

where the blocks of L are

$$L_{11} \rightarrow k \times k$$
 lower triangular; $L_{21} \rightarrow (n-k) \times k$; $L_{22} \rightarrow (n-k) \times (n-k)$ lower triangular

and the partitions x_1, x_2 of x and b_1, b_2 of b are such that

$$x_1, b_1 \rightarrow k \times 1$$
 and $x_2, b_2 \rightarrow (n-k) \times 1$.

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Pseudocode for blocked forward substitution:

- 1. Solve $L_{11}x_1 = b_1$ for x_1 .
- 2. Compute $b_2 L_{21}x_1$.
- 3. Solve $L_{22}x_2 = b_2 L_{21}x_1$ for x_2 .

