CSCI235 Database Systems

Normalization in Practice

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A relational schema R = (A, B, C)

Functional dependencies: AB \rightarrow C

Keys?

If AB \rightarrow C is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

No other minimal keys

Normal form?

Left hand side of AB \rightarrow C is a minimal key K = (A, B)
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A relational schema R = (A, B, C)

Functional dependencies: AB \rightarrow C, C \rightarrow B

Keys?

If $AB \rightarrow C$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

If $C \rightarrow B$ then through augmentation rule $AC \rightarrow AB$

If $AC \rightarrow AB$ is valid in \mathbb{R} and it covers entire relational schema then its left hand side is a minimal key (\mathbb{A} , \mathbb{C})

Normal form?

Not BCNF because left hand side of $C \rightarrow B$ is not a minimal key

3NF because right hand side of $C \rightarrow B$ is a prime attribute

Decomposition into BCNF?

$$R1 = (C, B), R2 = (A, B)$$

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3/20

A relational schema R = (A, B, C)

Functional dependencies: AB \rightarrow C, C \rightarrow B, C \rightarrow A

Keys?

If $AB \rightarrow C$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, B)

If $C \rightarrow B$ and $C \rightarrow A$ then through union rule $C \rightarrow AB$

If $C \to AB$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

Normal form?

BCNF because left hand side of each functional dependency is a minimal key

A relational schema R = (A, B, C)

Functional dependencies: $A \rightarrow B$

Keys?

If $A \rightarrow B$ is valid in R then through augmentation rule $AC \rightarrow BC$

If $AC \to BC$ is valid in \mathbb{R} and it covers entire relational schema then its left hand side is a minimal key (\mathbb{A} , \mathbb{C})

Normal form?

not 2NF because a nonprime attribute B functionally depends (A \rightarrow B) on a subset of primary key (A, C)

Decomposition into BCNF?

$$R1 = (A, B), R2 = (A, C) \text{ or }$$

$$R1 = (A, B), R2 = (B, C)$$

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5/20

A relation al schema R = (A, B, C)

Functional dependencies: $A \rightarrow B$, $B \rightarrow A$

Keys?

If A \rightarrow B then through augmentation rule AC \rightarrow BC

If $AC \to BC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A, C)

If $B \rightarrow A$ then through augmentation rule $BC \rightarrow AC$

If $BC \to AC$ is valid in \mathbb{R} and it covers entire relational schema then its left hand side is a minimal key (\mathbb{B} , \mathbb{C})

Normal form?

Not BCNF because left hand side of $A \rightarrow B$ is not a minimal key

3NF because right hand side of A \rightarrow B is a prime attribute and right hand side of B \rightarrow A is a prime attribute

Decomposition into BCNF?

$$R1 = (A, B), R2 = (A, C) \text{ or }$$

$$R1 = (A, B), R2 = (B, C)$$

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A relational schema R = (A, B, C)

Functional dependencies: $A \rightarrow B$, $B \rightarrow C$

Keys?

If A \rightarrow B and B \rightarrow C then through transitivity rule A \rightarrow C

If $A \rightarrow B$ and $A \rightarrow C$ then through union rule $A \rightarrow BC$

If $A \to BC$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A)

Normal form?

Not 3NF because a non prime attribute c is transitively dependent on primary key A

2NF because no nonprime attribute depends on a part of primary key

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Decomposition into BCNF?

$$R1 = (A, B), R2 = (B, C) \text{ or }$$

$$R1 = (A, B), R2 = (A, C)$$

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A relational schema R = (A, B, C, D)
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Functional dependencies: $A \rightarrow B$, $A \rightarrow C$, $B \rightarrow D$

Keys?

If $A \rightarrow B$ and $A \rightarrow C$ then through union rule $A \rightarrow BC$

If $A \rightarrow B$ and $B \rightarrow D$ then through transitivity rule $A \rightarrow D$

If A \rightarrow BC and A \rightarrow D then through union rule A \rightarrow BCD

If $A \to BCD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (A)

Normal form?

Not 3NF because a non prime attribute D is transitively dependent on primary key A

2NF because no nonprime attribute depends on a part of primary key

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Decomposition into BCNF?

$$R1 = (A, B, C), R2 = (B, D)$$

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A relational schema R = (A, B, C, D)
   Functional dependencies: A \rightarrow B, B \rightarrow D, C \rightarrow B
   Keys?
  If A \rightarrow B and B \rightarrow D then through transitivity rule A \rightarrow D
  If A \rightarrow D and A \rightarrow B then through union rule A \rightarrow BD
  If A \rightarrow BD then through augmentation rule AC \rightarrow BCD
   If AC \rightarrow BCD is valid in R and it covers entire relational schema then its
   left hand side is a minimal key (A, C)
  If C \to B and B \to D then through transitivity rule C \to D
  If C \rightarrow D and C \rightarrow B then through union rule C \rightarrow BD
  If C \rightarrow BD then through augmentation rule AC \rightarrow ABD
   If AC \rightarrow BCD is valid in R and it covers entire relational schema then its
   left hand side is a minimal key (A, C)
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                                                                                              12/20
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Normal form?

Not 2NF because a nonprime attribute B depends on a part of a primary

key

(A, C)

Decomposition into BCNF?

$$R1 = (A, B), R2 = (B, C), R2 = (B, D)$$

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A relational schema R = (A, B, C, D)
Functional dependencies: A \rightarrow B, A \rightarrow C, B \rightarrow A, B \rightarrow C
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Keys?

If $A \rightarrow B$ and $A \rightarrow C$ then through union rule $A \rightarrow BC$

If A \rightarrow BC then through augmentation rule AD \rightarrow BCD

If $AD \to BCD$ is valid in \mathbb{R} and it covers entire relational schema then its left hand side is a minimal key (A, D)

If $B \rightarrow A$ and $B \rightarrow C$ then through union rule $B \rightarrow AC$

If B \rightarrow AC then through augmentation rule BD \rightarrow ACD

If $BD \to ACD$ is valid in \mathbb{R} and it covers entire relational schema then its left hand side is a minimal key (\mathbb{B} , \mathbb{D})

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Normal form?

Not 2NF because a nonprime attribute c depends on a part of a primary

key

(B, D)

Decomposition into BCNF?

$$R1 = (A, B), R2 = (B, C), R2 = (A, D)$$

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A relational schema R = (A, B, C, D)
Functional dependencies: AB \rightarrow C, C \rightarrow D, D \rightarrow A, D \rightarrow B
Keys?
If AB \rightarrow C and C \rightarrow D then through transitivity rule AB \rightarrow D
If AB \rightarrow D and AB \rightarrow C then through union rule AB \rightarrow CD
If AB \rightarrow CD is valid in R and it covers entire relational schema then its left
hand side is a minimal key (A, B)
If D \rightarrow A and D \rightarrow B then through union rule D \rightarrow AB
If D \rightarrow AB and AB \rightarrow C then through transitivity rule D \rightarrow C
If D \rightarrow C and D \rightarrow AB then D \rightarrow ABC
If D \rightarrow ABC is valid in R and it covers entire relational schema then its left
hand side is a minimal key (D)
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If $C \rightarrow D$ and $D \rightarrow AB$ then through transitivity rule $C \rightarrow AB$

If $C \rightarrow D$ and $C \rightarrow AB$ then through union rule $C \rightarrow ABD$

If $C \to ABD$ is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

Normal form?

BCNF because left hand side of each functional dependency is a superky

18/20

Example 11

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A relational schema R = (A, B, C, D)
   Functional dependencies: A \rightarrow B, B \rightarrow C, C \rightarrow D, D \rightarrow A
   Keys?
   If A \rightarrow B and B \rightarrow C then through transitivity rule A \rightarrow C
   If A \rightarrow C and C \rightarrow D then through transitivity rule A \rightarrow D
   If A \rightarrow B and A \rightarrow C and A \rightarrow D then through union rule A \rightarrow BCD
   If A \rightarrowBCD is valid in R and it covers entire relational schema then its left
   hand side is a minimal key (A)
   If B \rightarrow C and C \rightarrow D then through transitivity rule B \rightarrow D
   If B \rightarrow D and D \rightarrow A then through transitivity rule B \rightarrow A
   If B \rightarrow C and B \rightarrow D and B \rightarrow A then through union rule B \rightarrow ACD
   If B \rightarrowACD is valid in R and it covers entire relational schema then its left
   hand side is a minimal key (B)
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If C \to D and D \to A then through transitivity rule C \to A

If C \to A and A \to B then through transitivity rule C \to B

If C \to A and C \to B and C \to D then through union rule C \to ABD

If C \to ABD is valid in R and it covers entire relational schema then its left hand side is a minimal key (C)

If D \to A and A \to B then through transitivity rule D \to B

If D \to B and B \to C then through transitivity rule D \to C

If D \to A and D \to B and D \to C then through union rule D \to ABC

If D \to ABC is valid in R and it covers entire relational schema then its left hand side is a minimal key (D)
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Normal form?

BCNF because left hand side of each functional dependency is a superky

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References

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