Maximum size rectangle binary sub-matrix with all 1s

Given a binary matrix, find the maximum size rectangle binary-sub-matrix with all 1's.

We strongly recommend you to minimize your browser and try this yourself first.

We have discussed a dynamic programming based solution for finding largest square with 1s.

In this post an interesting method is discussed that uses largest rectangle under histogram as a subroutine. Below are steps. The idea is to update each column of a given row with corresponding column of previous row and find largest histogram area for for that row.

Illustration:

```
step 1: 0 1 1 0 maximum area = 2
step 2:
    row 1 1 2 2 1 area = 4, maximum area becomes 4
    row 2 2 3 3 2 area = 8, maximum area becomes 8
    row 3 3 4 0 0 area = 6, maximum area remains 8
```

Below is C++ implementation. It is strongly recommended to refer this post first as most of the code taken from there.

```
// C++ program to find largest rectangle with all 1s
// in a binary matrix
#include<bits/stdc++.h>
using namespace std;
// Rows and columns in input matrix
#define R 4
#define C 4
// Finds the maximum area under the histogram represented
// by histogram. See below article for details.
// http://www.geeksforgeeks.org/largest-rectangle-under-histogram/
int maxHist(int row[])
    // Create an empty stack. The stack holds indexes of
    // hist[] array/ The bars stored in stack are always
    // in increasing order of their heights.
    stack<int> result;
    int top_val; // Top of stack
    int max_area = 0; // Initialize max area in current
                     // row (or histogram)
```

```
int area = υ; // initialize area with current top
   // Run through all bars of given histogram (or row)
   int i = 0;
   while (i < C)
       // If this bar is higher than the bar on top stack,
       // push it to stack
       if (result.empty() || row[result.top()] <= row[i])</pre>
            result.push(i++);
       else
           // If this bar is lower than top of stack, then
           // calculate area of rectangle with stack top as
            // the smallest (or minimum height) bar. 'i' is
           // 'right index' for the top and element before
            // top in stack is 'left index'
           top_val = row[result.top()];
            result.pop();
           area = top_val * i;
            if (!result.empty())
               area = top_val * (i - result.top() - 1 );
            max_area = max(area, max_area);
       }
   }
   // Now pop the remaining bars from stack and calculate area
   // with every popped bar as the smallest bar
   while (!result.empty())
       top_val = row[result.top()];
       result.pop();
       area = top_val * i;
       if (!result.empty())
            area = top_val * (i - result.top() - 1 );
       max_area = max(area, max_area);
   return max_area;
}
// Returns area of the largest rectangle with all 1s in A[][]
int maxRectangle(int A[][C])
   // Calculate area for first row and initialize it as
   // result
   int result = maxHist(A[0]);
   // iterate over row to find maximum rectangular area
   // considering each row as histogram
   for (int i = 1; i < R; i++)
       for (int j = 0; j < C; j++)
            // if A[i][j] is 1 then add A[i -1][j]
            if (A[i][j]) A[i][j] += A[i - 1][j];
       // Update result if area with current row (as last row)
       // of rectangle) is more
       result = max(result, maxHist(A[i]));
   }
   return result;
}
// Driver code
int main()
   int A[][C] = \{ \{0, 1, 1, 0\},
```

Output:

```
Area of maximum rectangle is 8
```

Time Complexity : O(R x X)