# Find maximum value of Sum( i\*arr[i]) with only rotations on given array allowed

Given an array, only rotation operation is allowed on array. We can rotate the array as many times as we want. Return the maximum possible of summation of i\*arr[i].

Example:

```
Input: arr[] = {1, 20, 2, 10}
Output: 72
We can 72 by rotating array twice.
{2, 10, 1, 20}
20*3 + 1*2 + 10*1 + 2*0 = 72

Input: arr[] = {10, 1, 2, 3, 4, 5, 6, 7, 8, 9};
Output: 330
We can 330 by rotating array 9 times.
{1, 2, 3, 4, 5, 6, 7, 8, 9, 10};
0*1 + 1*2 + 2*3 ... 9*10 = 330
```

We strongly recommend you to minimize your browser and try this yourself first.

A **Simple Solution** is to find all rotations one by one, check sum of every rotation and return the maximum sum. Time complexity of this solution is  $O(n^2)$ .

We can solve this problem in O(n) time using an **Efficient Solution**.

Let  $R_j$  be value of i\*arr[i] with j rotations. The idea is to calculate next rotation value from previous rotation, i.e., calculate  $R_j$  from  $R_{j-1}$ . We can calculate initial value of result as  $R_0$ , then keep calculating next rotation values.

#### How to efficiently calculate R<sub>j</sub> from R<sub>j-1</sub>?

This can be done in O(1) time. Below are details.

```
Let us calculate initial value of i*arr[i] with no rotation
R_0 = 0*arr[0] + 1*arr[1] + ... + (n-1)*arr[n-1]
After 1 rotation arr[n-1], becomes first element of array,
arr[0] becomes second element, arr[1] becomes third element
and so on.
R_1 = 0*arr[n-1] + 1*arr[0] + ... + (n-1)*arr[n-2]
R_1 - R_0 = arr[0] + arr[1] + ... + arr[n-2] - (n-1)*arr[n-1]
After 2 rotations arr[n-2], becomes first element of array,
arr[n-1] becomes second element, arr[0] becomes third element
and so on.
R_2 = 0*arr[n-2] + 1*arr[n-1] + ... + (n?1)*arr[n-3]
R_2 - R_1 = arr[0] + arr[1] + ... + arr[n-3] - (n-1)*arr[n-2] + arr[n-1]
If we take a closer look at above values, we can observe
below pattern
R_j - R_{j-1} = arrSum - n * arr[n-j]
Where arrSum is sum of all array elements, i.e.,
arrSum = \sum arr[i]
        i<=0<=n-1
```

Below is complete algorithm:

Below are C++ and Python implementations of above idea.

# C++

```
// C++ program to find max value of i*arr[i]
#include <iostream>
using namespace std;
// Returns max possible value of i*arr[i]
int maxSum(int arr[], int n)
    // Find array sum and i*arr[i] with no rotation
   int arrSum = 0; // Stores sum of arr[i]
   int currVal = 0; // Stores sum of i*arr[i]
   for (int i=0; i<n; i++)
       arrSum = arrSum + arr[i];
       currVal = currVal+(i*arr[i]);
    // Initialize result as 0 rotation sum
   int maxVal = currVal;
   // Try all rotations one by one and find
   // the maximum rotation sum.
    for (int j=1; j<n; j++)
        currVal = currVal + arrSum-n*arr[n-j];
       if (currVal > maxVal)
           maxVal = currVal;
   }
    // Return result
    return maxVal;
}
// Driver program
int main(void)
{
   int arr[] = {10, 1, 2, 3, 4, 5, 6, 7, 8, 9};
   int n = sizeof(arr)/sizeof(arr[0]);
   cout << "\nMax sum is " << maxSum(arr, n);</pre>
   return 0;
}
```

# **Python**

```
'''Python program to find maximum value of Sum(i*arr[i])'''
# returns max possible value of Sum(i*arr[i])
def maxSum(arr):
# stores sum of arr[i]
 arrSum = 0
 # stores sum of i*arr[i]
 currVal = 0
 n = len(arr)
 for i in range(0, n):
 arrSum = arrSum + arr[i]
 currVal = currVal + (i*arr[i])
 # initialize result
 maxVal = currVal
# try all rotations one by one and find the maximum
 # rotation sum
 for j in range(1, n):
 currVal = currVal + arrSum-n*arr[n-j]
 if currVal > maxVal:
  maxVal = currVal
# return result
return maxVal
# test maxsum(arr) function
arr = [10, 1, 2, 3, 4, 5, 6, 7, 8, 9]
print "Max sum is: ", maxSum(arr)
```

#### Output:

```
Max sum is 330
```

Time Complexity: O(n)
Auxiliary Space: O(1)