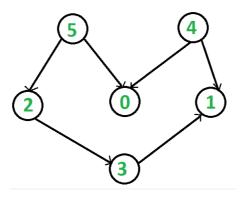
Kahn's algorithm for Topological Sorting

Topological sorting for **D**irected **A**cyclic **G**raph (DAG) is a linear ordering of vertices such that for every directed edge uv, vertex u comes before v in the ordering. Topological Sorting for a graph is not possible if the graph is not a DAG.

For example, a topological sorting of the following graph is "5 4 2 3 1 0?. There can be more than one topological sorting for a graph. For example, another topological sorting of the following graph is "4 5 2 0 3 1". The first vertex in topological sorting is always a vertex with indegree as 0 (a vertex with no in-coming edges).



A DFS based solution to find a topological sort has already been discussed.

In this article we will see another way to find the linear ordering of vertices in a directed acyclic graph (DAG). The approach is based on the below fact:

A DAG G has at least one vertex with in-degree 0 and one vertex with out-degree 0.

Proof: There's a simple proof to the above fact is that a DAG does not contain a cycle which means that all paths will be of finite length. Now let S be the longest path from u(source) to v(destination). Since S is the longest path there can be no incoming edge to u and no outgoing edge from v, if this situation had occurred then S would not have been the longest path => indegree(u) = 0 and outdegree(v) = 0

Algorithm:

Steps involved in finding the topological ordering of a DAG:

Step-1: Compute in-degree (number of incoming edges) for each of the vertex present in the DAG and initialize the count of visited nodes as 0.

Step-2: Pick all the vertices with in-degree as 0 and add them into a queue (Enqueue operation)

Step-3: Remove a vertex from the queue (Dequeue operation) and then.

- 1. Increment count of visited nodes by 1.
- 2. Decrease in-degree by 1 for all its neighboring nodes.
- 3. If in-degree of a neighboring nodes is reduced to zero, then add it to the queue.

Step 5: Repeat Step 3 until the queue is empty.

Step 5: If count of visited nodes is **not** equal to the number of nodes in the graph then the topological sort is not possible for the given graph.

How to find in-degree of each node?

There are 2 ways to calculate in-degree of every vertex:

Take an in-degree array which will keep track of

1) Traverse the array of edges and simply increase the counter of the destination node by 1.

```
for each node in Nodes
  indegree[node] = 0;
for each edge(src,dest) in Edges
  indegree[dest]++
```

Time Complexity: O(V+E)

2) Traverse the list for every node and then increment the in-degree of all the nodes connected to it by 1.

```
for each node in Nodes
If (list[node].size()!=0) then
for each dest in list
  indegree[dest]++;
```

Time Complexity: The outer for loop will be executed V number of times and the inner for loop will be executed E number of times, Thus overall time complexity is O(V+E).

The overall time complexity of the algorithm is O(V+E)

Below is C++ implementation of above algorithm. The implementation uses method 2 discussed above for finding indegrees.

C++

```
// A C++ program to print topological sorting of a graph
// using indegrees.
#include<bits/stdc++.h>
using namespace std;
// Class to represent a graph
class Graph
   int V; // No. of vertices'
    // Pointer to an array containing adjacency listsList
    list<int> *adj;
public:
   Graph(int V); // Constructor
   // function to add an edge to graph
   void addEdge(int u, int v);
    // prints a Topological Sort of the complete graph
    void topologicalSort();
};
Graph::Graph(int V)
{
    this->V = V;
    adj = new list<int>[V];
}
void Graph::addEdge(int u, int v)
{
    adj[u].push_back(v);
// The function to do Topological Sort.
void Graph::topologicalSort()
   // Create a vector to store indegrees of all
    // vertices. Initialize all indegrees as 0.
   vector<int> in_degree(V, 0);
    // Traverse adjacency lists to fill indegrees of
    // vertices. This step takes O(V+E) time
   for (int u=0; u<V; u++)
        list<int>::iterator itr;
        for (itr = adj[u].begin(); itr != adj[u].end(); itr++)
            in_degree[*itr]++;
    }
    // Create an queue and enqueue all vertices with
    // indegree 0
```

```
queue(III() q;
    for (int i = 0; i < V; i++)
        if (in_degree[i] == 0)
            q.push(i);
    // Initialize count of visited vertices
    int cnt = 0;
    // Create a vector to store result (A topological
    // ordering of the vertices)
    vector <int> top_order;
    // One by one dequeue vertices from queue and enqueue
    // adjacents if indegree of adjacent becomes \theta
    while (!q.empty())
        // Extract front of queue (or perform dequeue)
        \ensuremath{//} and add it to topological order
        int u = q.front();
        q.pop();
        top_order.push_back(u);
        // Iterate through all its neighbouring nodes
        // of dequeued node u and decrease their in-degree
        // by 1
        list<int>::iterator itr;
        for (itr = adj[u].begin(); itr != adj[u].end(); itr++)
            // If in-degree becomes zero, add it to queue
            if (--in_degree[*itr] == 0)
                q.push(*itr);
        cnt++;
   }
    // Check if there was a cycle
    if (cnt != V)
        cout << "There exists a cycle in the graph\n";
        return:
   }
    // Print topological order
    for (int i=0; i<top_order.size(); i++)</pre>
       cout << top_order[i] << " ";</pre>
    cout << endl;</pre>
}
// Driver program to test above functions
int main()
    \ensuremath{//} Create a graph given in the above diagram
    Graph g(6);
    g.addEdge(5, 2);
    g.addEdge(5, 0);
    g.addEdge(4, 0);
    g.addEdge(4, 1);
    g.addEdge(2, 3);
    g.addEdge(3, 1);
    cout << "Following is a Topological Sort of\n";</pre>
    g.topologicalSort();
    return 0;
}
```

Java

```
// A Java program to print topological sorting of a graph
// using indegrees
import java.util.*:
```

```
//Class to represent a graph
int V;// No. of vertices
//An Array of List which contains
//references to the Adjacency List of
//each vertex
List <Integer> adj[];
 public Graph(int V)// Constructor
 this.V = V;
 adj = new ArrayList[V];
 for(int i = 0; i < V; i++)
  adj[i]=new ArrayList<Integer>();
\ensuremath{//} function to add an edge to graph
 public void addEdge(int u,int v)
{
  adj[u].add(v);
 // prints a Topological Sort of the complete graph
public void topologicalSort()
 // Create a array to store indegrees of all
  // vertices. Initialize all indegrees as 0.
  int indegree[] = new int[V];
  // Traverse adjacency lists to fill indegrees of
  // vertices. This step takes O(V+E) time
  for(int i = 0; i < V; i++)
  ArrayList<Integer> temp = (ArrayList<Integer>) adj[i];
  for(int node : temp)
   indegree[node]++;
  }
  // Create a queue and enqueue all vertices with
  // indegree 0
  Queue<Integer> q = new LinkedList<Integer>();
  for(int i = 0;i < V; i++)
  if(indegree[i]==0)
   q.add(i);
  // Initialize count of visited vertices
  int cnt = 0:
  // Create a vector to store result (A topological
  // ordering of the vertices)
  Vector <Integer> topOrder=new Vector<Integer>();
  while(!q.isEmpty())
  // Extract front of queue (or perform dequeue)
  // and add it to topological order
  int u=q.poll();
  topOrder.add(u);
  // Iterate through all its neighbouring nodes
  // of dequeued node u and decrease their in-degree
  // by 1
  for(int node : adj[u])
   // If in-degree becomes zero, add it to queue
   if(--indegree[node] == 0)
    q.add(node);
```

```
cnt++;
  // Check if there was a cycle
 if(cnt != V)
  System.out.println("There exists a cycle in the graph");
 // Print topological order
 for(int i : topOrder)
  System.out.print(i+" ");
 }
}
}
// Driver program to test above functions
class Main
public static void main(String args[])
 // Create a graph given in the above diagram
 Graph g=new Graph(6);
 g.addEdge(5, 2);
    g.addEdge(5, 0);
    g.addEdge(4, 0);
    g.addEdge(4, 1);
    g.addEdge(2, 3);
    g.addEdge(3, 1);
    System.out.println("Following is a Topological Sort");
    g.topologicalSort();
}
```

Output:

```
Following is a Topological Sort
4 5 2 0 3 1
```