

Statement-Level Control Structures

CHAPTER 8 TOPICS

- Introduction
- Selection Statements
- Iterative Statements
- Unconditional Branching
- Guarded Commands
- Conclusions



LEVELS OF CONTROL FLOW

- Within expressions
 - low level, discussed in chapter 7
- Among program units
 - high level, will be discussed in chapter 9
- Among program statements



CONTROL STATEMENTS: EVOLUTION

- FORTRAN I control statements were based directly on IBM 704 hardware
- Much research and argument in the 1960s about the issue
 - One important result:

It was proven that all algorithms represented by flowcharts can be coded with only two-way selection and pretest logical loops



CONTROL STRUCTURE

- Control statements
 - Too many less readable
 - Too few require the use of low level statement, which is also less readable
- A *control structure* is a control statement and the statements whose execution it controls
- Design question
 - Should a control structure have multiple entries?



SELECTION STATEMENTS

- A *selection statement* provides the means of choosing between two or more paths of execution
- Two general categories:
 - Two-way selectors
 - Multiple-way selectors



TWO-WAY SELECTION STATEMENTS

General form:

```
if control_expression
  then clause
  else clause
```

- o Design Issues:
 - What is the form and type of the control expression?
 - How are the **then** and **else** clauses specified?
 - How should the meaning of nested selectors be specified?



TWO-WAY SELECTION: EXAMPLES

- FORTRAN: IF (boolean_expr) statement
- Problem: can select only a single statement; to select more, a **GOTO** must be used, as in the following example

IF (.NOT. condition) GOTO 20

. . .

20 CONTINUE

- Negative logic is bad for readability
- This problem was solved in FORTRAN 77
- Most later languages allow compounds for the selectable segment of their single-way selectors



TWO-WAY SELECTION: EXAMPLES

o ALGOL 60:

```
if (boolean_expr)
```

then statement (then clause)

else statement (else clause)

• The statements could be single or compound



NESTING SELECTORS

o Java example

```
if (sum == 0)
    if (count == 0)
        result = 0;
else result = 1;
```

- Which if gets the else?
- Java's static semantics rule: else matches with the nearest if
- Indentation has no effect on semantics



NESTING SELECTORS (CONTINUED)

• To force an alternative semantics, compound statements may be used:

```
if (sum == 0) {
   if (count == 0)
     result = 0;
}
else result = 1;
```

- The above solution is used in C, C++, and C#
- Perl requires that all then and else clauses to be compound



MULTIPLE-WAY SELECTION STATEMENTS

- Allow the selection of one of any number of statements or statement groups
- Obesign Issues:
 - 1. What is the form and type of the control expression?
 - 2. How are the selectable segments specified?
 - 3. Is execution flow through the structure restricted to include just a single selectable segment?
 - 4. What is done about unrepresented expression values?



- Early multiple selectors:
 - FORTRAN arithmetic IF (a three-way selector)

```
IF (arithmetic expression) N1, N2, N3
```

- Segments require GOTOs
- Not encapsulated (selectable segments could be anywhere)



- Modern multiple selectors
 - C's switch statement

```
switch (expression) {
    case const_expr_1: stmt_1;
    ...
    case const_expr_n: stmt_n;
    [default: stmt_n+1]
}
```



- o Design choices for C's switch statement
 - 1. Control expression can be only an integer type
 - 2. Selectable segments can be statement sequences, blocks, or compound statements
 - 3. Any number of segments can be executed in one execution of the construct (there is no implicit branch at the end of selectable segments)
 - 4. **default** clause is for unrepresented values (if there is no **default**, the whole statement does nothing)



• The Ada case statement

```
case expression is
   when choice list => stmt_sequence;
...
   when choice list => stmt_sequence;
   when others => stmt_sequence;]
end case;
```

 More reliable than C's switch (once a stmt_sequence execution is completed, control is passed to the first statement after the case statement



MULTIPLE-WAY SELECTION USING IF

• Multiple Selectors can appear as direct extensions to two-way selectors, using else-if clauses, for example in Ada:

```
if ...
  then ...
elsif ...
then ...
elsif ...
elsif ...
elsif ...
end if
```



ITERATIVE STATEMENTS

- The repeated execution of a statement or compound statement is accomplished either by iteration or recursion
- General design issues for iteration control statements:
 - 1. How is iteration controlled?
 - 2. Where is the control mechanism in the loop?



COUNTER-CONTROLLED LOOPS

- A counting iterative statement has a loop variable, and a means of specifying the *initial* and *terminal*, and *stepsize* values
- o Design Issues:
 - 1. What are the type and scope of the loop variable?
 - 2. What is the value of the loop variable at loop termination?
 - 3. Should it be legal for the loop variable or loop parameters to be changed in the loop body, and if so, does the change affect loop control?
 - 4. Should the loop parameters be evaluated only once, or once for every iteration?



• FORTRAN 90 syntax

```
DO label var = start, finish [, stepsize]
```

- Stepsize can be any value but zero
- Parameters can be expressions
- Design choices:
 - 1. Loop variable must be INTEGER
 - 2. Loop variable always has its last value
 - 3. The loop variable cannot be changed in the loop, but the parameters can; because they are evaluated only once, it does not affect loop control
 - 4. Loop parameters are evaluated only once



• FORTRAN 95: a second form:

Loop variable must be an INTEGER



ITERATIVE STATEMENTS

Pascal's for statement

```
for variable := initial (to|downto) final do
    statement
```

- Opening the property of the
 - 1. Loop variable must be an ordinal type of usual scope
 - 2. After normal termination, loop variable is undefined
 - 3. The loop variable cannot be changed in the loop; the loop parameters can be changed, but they are evaluated just once, so it does not affect loop control



Ada

```
for var in [reverse] discrete_range loop
...
end loop
```

- A discrete range is a sub-range of an integer or enumeration type
- Scope of the loop variable is the range of the loop
- Loop variable is implicitly undeclared after loop termination



o C's for statement

```
for ([expr_1] ; [expr_2] ; [expr_3])
  statement
```

- The expressions can be whole statements, or even statement sequences, with the statements separated by commas
 - The value of a multiple-statement expression is the value of the last statement in the expression
- There is no explicit loop variable
- Everything can be changed in the loop
- The first expression is evaluated once, but the other two are evaluated with each iteration



1-24

- C++ differs from C in two ways:
 - 1. The control expression can also be Boolean
 - 2. The initial expression can include variable definitions (scope is from the definition to the end of the loop body)
- Java and C#
 - Differs from C++ in that the control expression must be Boolean



ITERATIVE STATEMENTS: LOGICALLY-CONTROLLED LOOPS

- Repetition control is based on a Boolean
- o Design issues:
 - Pre-test or post-test?
 - Should the logically controlled loop be a special case of the counting loop statement? expression rather than a counter
- General forms:

```
while (ctrl_expr) do
  loop body
  while (ctrl_expr)
```



ITERATIVE STATEMENTS: LOGICALLY-CONTROLLED LOOPS: EXAMPLES

- Pascal has separate pre-test and post-test logical loop statements (while-do and repeat-until)
- o C and C++ also have both, but the control expression for the post-test version is treated just like in the pretest case (while-do and do-while)
- Java is like C, except the control expression must be Boolean (and the body can only be entered at the beginning -- Java has no goto



ITERATIVE STATEMENTS: LOGICALLY-CONTROLLED LOOPS: EXAMPLES

- Ada has a pretest version, but no post-test
- FORTRAN 77 and 90 have neither
- Perl has two pre-test logical loops, while and until, but no post-test logical loop



ITERATIVE STATEMENTS: USER-LOCATED LOOP CONTROL MECHANISMS

- Sometimes it is convenient for the programmers to decide a location for loop control (other than top or bottom of the loop)
- Simple design for single loops (e.g., break)
- Design issues for nested loops
 - 1. Should the conditional be part of the exit?
 - 2. Should control be transferable out of more than one loop?



ITERATIVE STATEMENTS: USER-LOCATED LOOP CONTROL MECHANISMS BREAK AND

CONTINUE

- o C, C++, and Java: break statement
- Unconditional; for any loop or switch; one level only
- Java and C# have a labeled break statement: control transfers to the label
- An alternative: **continue** statement; it skips the remainder of this iteration, but does not exit the loop



ITERATIVE STATEMENTS: ITERATION BASED ON DATA STRUCTURES

- Number of elements of in a data structure control loop iteration
- Control mechanism is a call to an *iterator* function that returns the next element in some chosen order, if there is one; else loop is terminate
- C's **for** can be used to build a user-defined iterator:

```
for (p=root; p==NULL; traverse(p)){
}
```



ITERATIVE STATEMENTS: ITERATION BASED ON DATA STRUCTURES (CONTINUED)

• C#'s foreach statement iterates on the elements of arrays and other collections:

```
Strings[] = strList = {"Bob", "Carol", "Ted"};
foreach (Strings name in strList)
    Console.WriteLine ("Name: {0}", name);
```

• The notation {0} indicates the position in the string to be displayed



Unconditional Branching

- Transfers execution control to a specified place in the program
- Represented one of the most heated debates in 1960's and 1970's
- Well-known mechanism: goto statement
- Major concern: Readability
- Some languages do not support goto statement (e.g., Module-2 and Java)
- C# offers goto statement (can be used in switch statements)
- Loop exit statements are restricted and somewhat camouflaged goto's



GUARDED COMMANDS

- Suggested by Dijkstra
- Purpose: to support a new programming methodology that supports verification (correctness) during development
- Basis for two linguistic mechanisms for concurrent programming (in CSP and Ada)
- Basic Idea: if the order of evaluation is not important, the program should not specify one



SELECTION GUARDED COMMAND

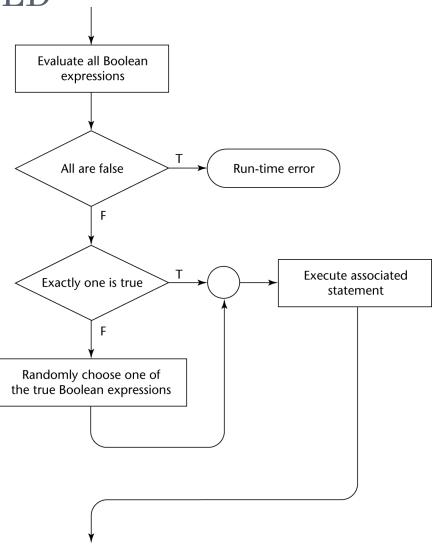
• Form

```
if <Boolean exp> -> <statement>
[] <Boolean exp> -> <statement>
    ...
[] <Boolean exp> -> <statement>
fi
```

- Semantics: when construct is reached,
 - Evaluate all Boolean expressions
 - If more than one are true, choose one nondeterministically
 - If none are true, it is a runtime error



SELECTION GUARDED COMMAND: ILLUSTRATED





SELECTION GUARDED COMMAND EXAMPLE

```
if i = 0 -> sum := sum + i
[] i > j -> sum := sum + j
[] j > i -> sum := sum + i
fi

if x >= y -> max := x
[] y >= x -> max := y
fi
```

• How would you do with ordinary coding?



LOOP GUARDED COMMAND

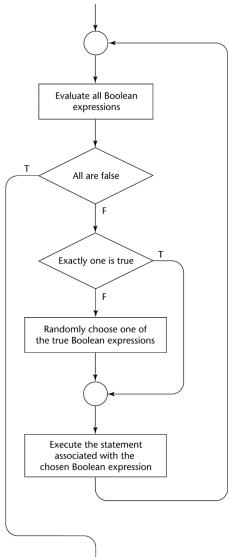
Form

```
do <Boolean> -> <statement>
[] <Boolean> -> <statement>
    ...
[] <Boolean> -> <statement>
od
```

- Semantics: for each iteration
 - Evaluate all Boolean expressions
 - If more than one are true, choose one nondeterministically; then start loop again
 - If none are true, exit loop



LOOP GUARDED COMMAND: ILLUSTRATED





LOOP GUARDED COMMAND EXAMPLE

```
do q1 > q2 -> temp := q1; q1 := q2; q2 := temp; [] q2 > q3 -> temp := q2; q2 := q3; q3 := temp; [] q3 > q4 -> temp := q3; q3 := q4; q4 := temp; od
```

• How would you do with ordinary coding?



GUARDED COMMANDS: RATIONALE

- Connection between control statements and program verification is intimate
- Verification is impossible with goto statements
- Verification is possible with only selection and logical pretest loops
- Verification is relatively simple with only guarded commands



CONCLUSION

- Variety of statement-level structures
- Choice of control statements beyond selection and logical pretest loops is a trade-off between language size and writability
- Functional and logic programming languages are quite different control structures

