

Patterns

- ① outer loop , count no. of rows $\begin{matrix} * & * & * \\ * & * & * \\ * & * & * \end{matrix}$
- ② inner loop , focus on column
and connect them with rows $\begin{matrix} * & * & * & * \\ * & * & * & * \end{matrix}$
- ③ print then '*' inside the
inner for loop. $\begin{matrix} * & * & * & * \\ * & * & * & * \end{matrix}$

```
for( i=0 ; i<4 ; i++ )
```

{

```
    for( j=0 ; j<4 ; j++ )
```

{

```
        cout << "*";
```

cout << endl;

}

```
for( i=0 ; i<5 ; i++ )
```

{

```
    for( j=0 ; j<=i ; j++ )
```

{

```
        cout << '*';
```

{

```
    cout << endl;
```

}

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*

for(int i=0, i<5, i++)

1 2

{ for(int j=1, j<i, j++)

1 2 3

{ cout << j;

1 2 3 4

}

cout << endl;

1 2 3 4 5

}

for every row column
are going
from 1 to row itself

for (int i=1, i<5, i++)

1

{

for(j=1, j<i, j++)

2 2

{

cout << i;

3 3 3

{

cout << endl;

4 4 4 4

}

5 5 5 5 5

for every row
print row no.

for(int i=0 ; i<5 ; i++)

* * * * *

{

for(j=0; j<n-i+1 ; j++)

* * * *

{ cout << "*";

* * *

cout << endl;

* (n-row+1)

}

1 row → 5 stars

2 rows → 4 stars

3 rows → 3 stars

4 row → 2 stars

5 row → 1 stars

```
for (int i=1; i<n; i++)
```

1	2	3	4	5
*	2	3	4	

```
{ for (int j=1; j<(n-i); j++)
```

1	2	3
1	2	
1		

```
{ cout << j;
```

```
} cout << endl;
```

```
}
```

```
for (i=0; i<n; i++)
```

```
{ lspace
```

```
for (int j=0; j<(n-i-1); j++)
```

*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*

```
{ cout << " ";
```

```
*star
```

```
for (int j=0; j<(2*i+1); j++)
```

```
{ cout << "*" ;
```

```
}
```

```
lspacu
```

```
for (int j=0; j<(n-i-1); j++)
```

```
{ cout << " " ;
```

```
} cout << endl;
```

9 columns

[4space, 1star, 4space]

[3, 3, 3]

[2, 5, 2]

[1, 7, 1]

[0, 9, 0]

space. ($n - i - 1$)

* (2 $i + 1$)

$$(2^n - 2 + 1)$$

for (i=0; i<n; i++)

* * * * *

{ uspac

{ for (j=0; j<i; j++)

0 → 0, 5, 0

{ cout << " "

1 → 1, 3, 1

2 → 2, 1, 2

ustar

for (j=0; j<2n-2i+1; j++)

{ cout << "*" ;

}

// space

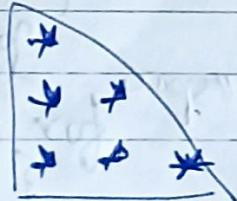
for (j=0; j<i; j++)

{ cout << " " ;

}

~~cout << endl;~~

for (int i=1; i<2n-1; i++)



{ int star = 1 ;

if (i>n) stars = 2*n - i ;

for (int j=1; j<stars; j++)

{ cout << "*" ;

}

cout << endl ;

```
int start = 1;
```

```
for( i=0; i < n; i++)
```

 for(j=0; j <= i; j++)

 if(i%2 == 0) start = 1;

 else start = 0;

```
    for( int j=0; j < i; j++) {
```

 cout << start;

 start = 1 - start;

}

cout << endl;

}

⇒ space = 2 * (n-1)

```
for( i=1; i < n; i++)
```

 // number

```
    for( j=1; j <= n; j++)
```

 cout << j;

}

// space

```
for( j=1; j < space; j++)
```

 cout << " ";

 // number

```
    for( j=i; j>=1; j--)
```

 cout << j;

}

cout << endl;

space -= 2;

1
0 1
1 0 1
0 1 0 1
1 0 1 0 1

1 2
1 2 3
1 2 3 4 4 3 2 1

1 6
1 6 5
1 6 5 4

1 . . . 1
1 2 . . . 2 1
1 2 3 . . . 3 2 1

1 2 3 4 4 3 2 1

1 6 1
2 4 2
3 2 3
4 0 4

num = 1

for (int i=1; i<=n; i++)

{ for (int j=1; j<=i; j++)

{~~for (int j=1; j<=i; j++)~~

cout << num;
num++;

cout << endl;

}

1
2 3

4 5 6

7 8 9 10

11 12 13 14 15

~~int ans[10]~~

~~int i,~~

for (i=0; i<n; i++)

A

A B

A B C

A B C D

A B C D E

{ for (char ch='A'; ch <='A'+9; ch++)

{ cout << ch;

cout << endl;

}

for (i=0; i<n; i++)

ABCDE

A B C D

A B C

A B

A

{ for (char ch='A'; ch<='A'+9-i; ch++)

{ cout << ch;

i

} cout << endl;

- ABCDE,

 A + 4 n = 4

- ABC
 A + 2 n =

- ABCD
 A + 3 (n - i - 1)

```

for (int i = 0; i < n; i++)
{
    for (char ch = 'A'; ch <='A' + (n - i); ch++)
    {
        char ch = 'A' + i;
        for (int j = 0; j < i; j++)
        {
            cout << ch << " ";
            if (cout << endl;
        }
    }
}

```

A
B B
C C C
D D D D
E E E E

```

for (int i = 0; i < n; i++)
{
    // space
    for (int j = 0; j < n - i - 1; j++)
    {
        cout << " ";
        ch = 'A';
        int breakpoint = (2 * i + 1) / 2;
        if (j <= breakpoint) ch++;
        else
            ch--;
    }
}

```

A B A
A B C B A
A B C D R B A
A B C D E D C B A

space $n - i - 1$
char $2 * i + 1$
 $i / 2 + 1 \rightarrow$ increment
breakpoint

```

//char
for (int j = 0; j < n - i - 1; j++)
{
    cout << ch;
}

```

if (cout << endl;

for (int i=0; i<n; i++) E
 { for (char ch='E'-i; ch='E'; D E
 ch++) C D E
 B DE
 A B C D E

{ cout << endl;

} }
} }
} }

0 → E E-i
1 → DE

Pairs → part of utility library

pair < int, int > = p = {1, 3};

ID access

p.first, p.second

Vectors → containers which is dynamic in nature

Array - arr[] = {1, 2, 3, 4} → we cannot modify size of array

- we can always set size of vector whenever you wish to.

vector < int > v;

→ it will contain empty container.
{}
3

v.push_back[1];

v.emplace_back[2];

~~push~~ 2 ways
to push
element in
vector

vector < int > v(5); → {0, 0, 0, 0, 0}

vector < int > v(5, 100); → {100, 100, 100, 100, 100}

vector < int > v1(v2);

push back
emplace back

for (vector < int >:: iterator it = v.begin();

it != v.end(); it++) {

cout << *it << " ";

faster

}

Q7.

→ automatically assign datatype

```
for (auto it = v.begin(); it != v.end(); it++) {  
    cout << *it << " ";  
}
```

or using for each loop.

```
for (auto it : v) {  
    cout << it << " ";  
}
```

[10, 20, 30, 40]

vec.erase (st, end)

↳ end address
after the element

~~vec.insert~~ vec.insert (v.begin() + 1, 2, 10)

occurrence

O/P: {10, 10, 10, 20, 30, 40}

• find() → searches first occurrence of a value in a range
• Iterator → searches specific element within container range
→ points the address,
can access value using *

(LIFO)	(FIFO)
Stack.	Queue
stack<int> st;	queue<int> q;
st.push(1);	q.push(1);
st.pop();	q.pop();
st.top();	q.front();
	q.back();

Priority Queue

largest value stay at the top.
push, pop, top

set → sorted, unique

Multiset → sorted, (duplication allowed)

Unordered → unique.

Map → map<int, int> mpp;
 unique keys in sorted order

MultiMap → duplicate keys, sorted.
 $\{1, 2\}, \{1, 3\}$

UnorderedMap → unique key, not sorted

✓ works in constant time
 $O(1)$

Extra $\text{sort}(a, a+n)$

↑
start

↑
end

$[1, 2, 3, 4, 5]$
↑
 $a+n$

automatically sort in ascending

Basic Maths . Concept

extraction of digit

$$7789 \% 10 = 9$$

$$110 \left(\begin{array}{r} 70080 \\ \hline 700 \end{array} \right)$$

$$778 \% 10 = 8$$

$$110 \left(\begin{array}{r} 70 \\ \hline 7 \end{array} \right)$$

$$7 \% 10 = 7$$

$$110 \left(\begin{array}{r} 0 \\ \hline 0 \end{array} \right)$$

$$10 \left(\begin{array}{r} 778 \\ \hline 7789 \\ - 778 \\ \hline 9 \end{array} \right)$$

Division

/ → return quotient

Floor Division

// → return largest integer <= to the quotient
↳ round off to nearest integer

Modulo

% → returns remainder

while ($N \neq 0$)

{

last digit = $N \% 10$

$N = N / 10$

}

Time complexity
 $O(\log_{10}(n))$

count of digits

$$\log_{10}(7789) = 3.89$$

$$\frac{+1}{4.89} \rightarrow (\underline{\underline{\text{int}}})$$

```
#include <bits/stdc++.h>
```

```
int count (int n) {  
    int cnt = (int)(log10(n)) + 1;  
    return cnt;  
}
```

Palindrome: ~~(2)~~

~~reverse no.
and check
condition~~ 1331 → reverse. 1331

Armstrong:

$$371 = 3^3 + 7^3 + 1^3 = 371$$

$$1634 = 1^3 + 6^3 + 3^3 + 4^3 = 1634.$$

dup = N
sum = 0

while ($N > 0$)

{

$$\text{lastdigit} = N \% 10;$$

$$\text{sum} = \text{sum} + (\text{lastdigit} \times \text{lastdigit} \times \text{lastdigit})$$

$$N = N / 10$$

}

If ($\text{sum} == \text{dup}$) \Rightarrow armstrong

Print all divisor

36 → 1, 2, 3, 4, 6, 9, 12, 18, 36

we can say that it lies b/t {1 to N}

for ($i=1$; $i <= N$; $i++$)

{

// completely dividing
if ($N \% i == 0$)
print (i)

}

Alternative

$N = 36$

$1 \times 36 \leftarrow N_1$
 $2 \times 18 \leftarrow N_2$
 $3 \times 12 \leftarrow N_3$
 $4 \times 9 \leftarrow N_4$
 $6 \times 6 \leftarrow N_5$

same

9×4
 12×3
 18×2
 36×1

we will move till \sqrt{N}

$$\sqrt{36} = 6$$

6 should not come 2 time

take a note.

O/P: 1, 36, 2, 18, 3, 12, 4, 9, 6 ← not in sorted manner

```

for (i=1 ; i<=sqrt(N) ; i++)
{
    if (N % i == 0)
        { print(i)
        if ((n/i) != i)
            print(n/i); }
}

```

Prime No: \rightarrow exactly two factors
 $(1 \& \text{ itself})$

counter = 0

for (i=1 ; i<=N ; i++)

{ if (N % i == 0)

counter++;

Note:
 we can
 also get
 factor using
 $\text{sqrt}(N)$

}

if (counter == 2)

printf ("Prime")

else

printf ("Not Prime")

of using \sqrt{n}

$cnt = 0$

for ($i=1$; $i*i <= n$; $i++$)

{ if ($n \% i == 0$)

{ $cnt++$;

if ($(n/i) != i$)

$cnt++$;

$O(\sqrt{n})$

} if ($cnt == 2$)

Prime

else

Not Prime

GCD or HCF

greatest common division
highest common factor

$$GCD(9, 12) = 3$$

$N=9$

$N=12$

~~(1, 3)~~ 9

(1, 2, 3, 6, 12) 4

Highest common = 3

Logic
Let $N_1 = 10$ and $N_2 = 12$

$N_1 = 9$ $N_2 = 12$

Loop 1 to $\min(N_1, N_2)$

check for all i value
where it is divisible by both
 N_1 and N_2

and replace it with GCD
value

$O(\min(n_1, n_2))$

for ($i=1$; $i \leq \min(n_1, n_2)$; $i+1$)

{ if ($n_1 \% i == 0$ & $n_2 \% i == 0$)

$$\gcd = i;$$

}

Euclidean Algorithm Time complexity
 $O(\log(\min(a, b)))$

$$\gcd(a, b) = \gcd(a-b, b)$$

where $a > b$

$$\gcd(n_1, n_2) = \gcd(n_1 - n_2, n_2) \quad n_1 > n_2$$

$$\begin{aligned} \gcd(52, 10) &\rightarrow \gcd(42, 10) \xrightarrow{\textcircled{1}} \gcd(32, 10) \xrightarrow{\textcircled{2}} \gcd(22, 10) \xrightarrow{\textcircled{3}} \\ &\quad \cancel{\gcd(12, 10)} \xrightarrow{\textcircled{4}} \cancel{\gcd(2, 10)} \end{aligned}$$

$\cancel{52/10 = 5}$
 $\cancel{52 \% 10 = 2}$

$$\gcd(a, b) = \gcd(a \% b, b)$$

greater % smaller one of them zero the other is.

$\gcd(0, n_2)$