Quertion - 31

Given Context free grammas es = s -> xax × -> ax/bx/n

a) Given Content Free Grammas (CFG) is unambiguous because of how only one pause tree for given grammas

tox exomple:

 $S \rightarrow XaX \quad (S \rightarrow XaX)$

 $\rightarrow a \times a \times (\times \rightarrow a \times)$

 $\rightarrow aax (x \rightarrow \wedge)$

 $\rightarrow aa \quad (x \rightarrow 1)$

"aabba"

 $s \rightarrow x \alpha \chi \quad (s \rightarrow x \alpha x)$

 $\rightarrow a \times a \times (x \rightarrow a \times)$

 $\rightarrow aaxax (x \rightarrow ax)$

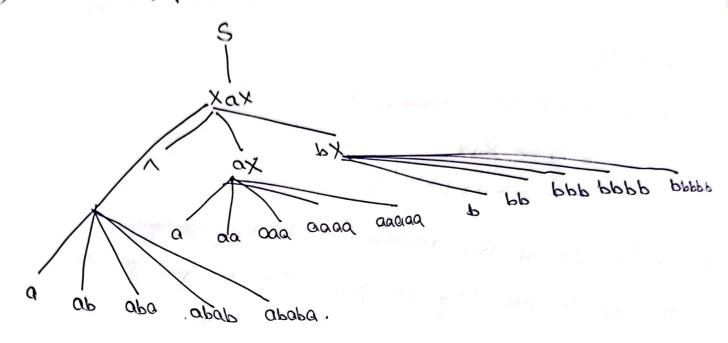
→aabxax (x→bx)

→oabbxax (x→bx)

 $\rightarrow aappax (x \rightarrow v)$

 $\rightarrow aabba (x \rightarrow n)$

There is only one way to devise them wing given gramma so it is not ambiguous



The quen geammas is unambiguous.