

سوال اول تمرین سری یک درس میکروپروسسور
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برای پیاده‌سازی این ماشین حساب از الگوریتم و جدول تصمیم‌گیری ذکر شده در فایل ضمیمه استفاده شده است. روال برنامه به این ترتیب است که ابتدا ورودی از کاربر گرفته می‌شود. سپس ورودی را از کاراکتر اول تا آخر پرمایش می‌کنیم. در این پرمایش چنانچه به کاراکتر عددی برسیم آنرا در بافر ذخیره می‌کنیم. چنانچه به کاراکتر غیر عددی برسیم، براساس عملگری که در بالای **پشته موقت** قرار دارد، در مورد آن عملگر تصمیم‌گیری می‌کنیم. این تصمیم‌گیری که همان‌طور که در فایل ضمیمه ذکر شده است به پنج حالت منجر می‌شود. پس از پرمایش کل ورودی چنانچه خطایی در ورودی وجود نداشته باشد وارد مرحله‌ی بعد که برآورد کردن **صف** مورد نظر است، می‌شویم.

در این بخش **صف** را از ابتدا تا انتها پرمایش می‌کنیم. چنانچه به عدد برسیم آنرا در داخل **یک پشته موقت دیگر** قرار می‌دهیم تا نخستین عملگری که بعد از این می‌خوانیم روی داده‌های موجود در **پشته موقت** عمل کرده و نتیجه را در بالای پشته بگذارد. پس اینکه همه‌ی عملگرها اثر دادیم نتیجه مورد نظر در بالای **پشته موقت** قرار دارد. این نتیجه 32 بیتی را به دسیمال علامت دار تبدیل کرده و خروجی نمایش می‌دهیم. پس از این تبدیل برنامه خاتمه می‌یابد.

the calculation. This form requires an offset field in the instruction large enough to hold an address, of course, so it is less efficient than doing it the other way; however, it is nevertheless frequently the best way.

5.4.7 Based-Indexed Addressing

Some machines have an addressing mode in which the memory address is computed by adding up two registers plus an (optional) offset. Sometimes this mode is called **based-indexed addressing**. One of the registers is the base and the other is the index. Such a mode would have been useful here. Outside the loop we could have put the address of *A* in R5 and the address of *B* in R6. Then we could have replaced the instruction at *LOOP* and its successor with

```
LOOP:      MOV R4,(R2+R5)
           AND R4,(R2+R6)
```

If there were an addressing mode for indirecting through the sum of two registers with no offset, that would be ideal. Alternatively, even an instruction with an 8-bit offset would have been an improvement over the original code since we could set both offsets to 0. If, however, the offsets are always 32 bits, then we have not gained anything by using this mode. In practice, however, machines that have this mode usually have a form with an 8-bit or 16-bit offset.

5.4.8 Stack Addressing

We have already noted that making machine instructions as short as possible is highly desirable. The ultimate limit in reducing address lengths is having no addresses at all. As we saw in Chap. 4, zero-address instructions, such as *IADD*, are possible in conjunction with a stack. In this section we will look at stack addressing more closely.

Reverse Polish Notation

It is an ancient tradition in mathematics to put the operator between the operands, as in $x + y$, rather than after the operands, as in $x y +$. The form with the operator “in” between the operands is called **infix** notation. The form with the operator after the operands is called **postfix** or **reverse Polish notation**, after the Polish logician J. Lukasiewicz (1958), who studied the properties of this notation.

Reverse Polish notation has a number of advantages over infix for expressing algebraic formulas. First, any formula can be expressed without parentheses. Second, it is convenient for evaluating formulas on computers with stacks. Third, infix operators have precedence, which is arbitrary and undesirable. For example, we know that $a \times b + c$ means $(a \times b) + c$ and not $a \times (b + c)$ because multiplication has been arbitrarily defined to have precedence over addition. But does left

shift have precedence over Boolean AND? Who knows? Reverse Polish notation eliminates this nuisance.

Several algorithms for converting infix formulas into reverse Polish notation exist. The one given below is an adaptation of an idea due to E. W. Dijkstra. Assume that a formula is composed of the following symbols: variables, the dyadic (two-operand) operators $+$ $-$ $*$ $/$, and left and right parentheses. To mark the ends of a formula, we will insert the symbol \perp after the last symbol and before the first symbol.

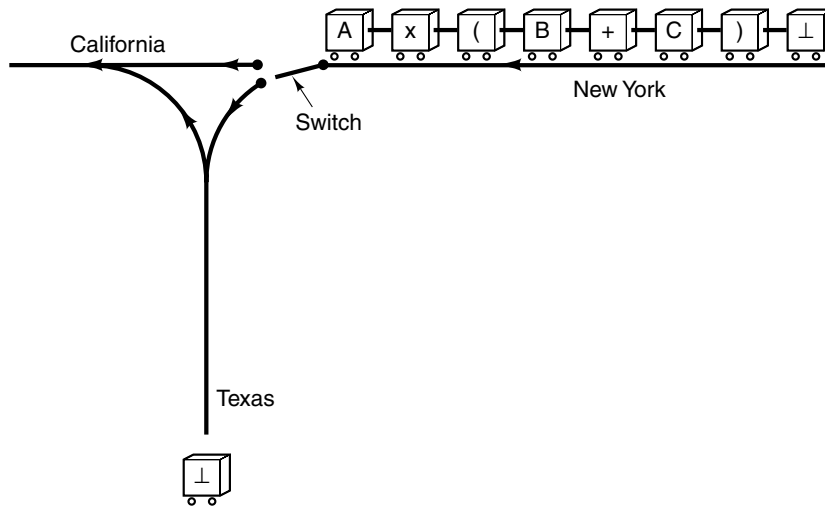


Figure 5-21. Each railroad car represents one symbol in the formula to be converted from infix to reverse Polish notation.

Figure 5-21 shows a railroad track from New York to California, with a spur in the middle that heads off in the direction of Texas. Each symbol in the formula is represented by one railroad car. The train moves westward (to the left). When each car arrives at the switch, it must stop just before it and ask if it should go to California directly or take a side trip to Texas. Cars containing variables always go directly to California and never to Texas. Cars containing all other symbols must inquire about the contents of the nearest car on the Texas line before entering the switch.

The table of Fig. 5-22 shows what happens, depending on the contents of the nearest car on the Texas line and the car poised at the switch. The first \perp always goes to Texas. The numbers refer to the following situations:

1. The car at the switch heads toward Texas.
2. The most recent car on the Texas line turns and goes to California.
3. Both the car at the switch and the most recent car on the Texas line are diverted and disappear (i.e., both are deleted).

- 4. Stop. The symbols now in California represent the reverse Polish notation formula when read from left to right.
- 5. Stop. An error has occurred. The original formula was not correctly balanced.

		Car at the switch						
		⊥	+	-	x	/	()
Most recently arrived car on the Texas line	⊥	4	1	1	1	1	1	5
	+	2	2	2	1	1	1	2
	-	2	2	2	1	1	1	2
	x	2	2	2	2	2	1	2
	/	2	2	2	2	2	1	2
	(5	1	1	1	1	1	3
)							

Figure 5-22. Decision table used by the infix-to-reverse Polish notation algorithm

After each action is taken, a new comparison is made between the car currently at the switch, which may be the same car as in the previous comparison or may be the next car, and the car that is now the last one on the Texas line. The process continues until step 4 is reached. Notice that the Texas line is being used as a stack, with routing a car to Texas being a push operation, and turning a car already on the Texas line around and sending it to California being a pop operation.

Infix	Reverse Polish notation
$A + B \times C$	$A B C \times +$
$A \times B + C$	$A B \times C +$
$A \times B + C \times D$	$A B \times C D \times +$
$(A + B) / (C - D)$	$A B + C D - /$
$A \times B / C$	$A B \times C /$
$((A + B) \times C + D) / (E + F + G)$	$A B + C \times D + E F + G + /$

Figure 5-23. Some examples of infix expressions and their reverse Polish notation equivalents.

The order of the variables is the same in infix and in reverse Polish notation. The order of the operators, however, is not always the same. Operators appear in reverse Polish notation in the order they will actually be executed during the evaluation of the expression. Figure 5-23 gives several examples of infix formulas and their reverse Polish notation equivalents.

;8086 PROGRAM HW1_2_89101089.ASM

;ABSTRACT : This program evaluates entered expression and displays result

;REGISTERS : Uses CS , DS , ES , SS , SP , BP , AX , BX , CX , DX ,

;PORTS : ---

;PROCEDURES: ins_char : inserts read character into char_buffer

; push_var : calculates decimal chars in char_buff and puts result in queue

; push_op_t : pushes operator to temporary operator stack (i.e. texas!)

; move_stck : moves top element of stack to queue

; GetRes : evaluates queue

; GetOp : calculates result of desired operator on two operators which are in stack

; UnSign : used to unsign operands before division & multiplication and set isSigned flag

; DispRes : convert 32-bit binary result to signed decimal & display it

;-----

DATA SEGMENT

Sen_Max DB 127

Sen_Real DB ?

Sentence DB 127 DUP(?)

; stack pointer backup

bcup_op_t DW ?

;character read part

isSigned DB 0

isNegative DB 0

char_fnd DB 0

char_buff DB 5 DUP(0)

char_coun DB 0

char_sum DW 0

;calculation part

operandL DW ?

operandH DW ?

Divisor DW 10

ResWord0 DW 0

ResWord1 DW 0

TempWord DW 0

queue DW 200 DUP(0) ; queue of double words

ptr_queue DW offset queue ; pointer to current double word in queue

calc_queue DW offset queue ; pointer used in calculating result

MESS1 DB 0DH , 0AH , 'Error in expression \$'

MESS2 DB 0DH , 0AH , 11 DUP(0) , '\$'

MESS3 DB 0DH , 0AH , 'Enter expression to evaluate : (Specify negative numbers with #)' , 0DH , 0AH , '\$'

DATA ENDS

;-----

STACK SEGMENT Stack

DW 100 DUP(0)
sp_func Label Word

DW 100 DUP(0) ; stack of 100 words for ascii code of operators (+ , * , / , -)
sp_op_t Label Word

DW 100 DUP(0) ; " " double words for operands during evaluation
sp_oprnd Label Word

STACK ENDS

;-----

CODE SEGMENT

ASSUME CS:CODE , DS:DATA , ES:DATA , SS:STACK

START: MOV AX , DATA
 MOV DS , AX ; initialize segments
 MOV ES , AX

 MOV AX , STACK
 MOV SS , AX
 MOV SP , offset sp_func

 MOV DX , offset MESS3 ; display message
 MOV AH , 09H
 INT 21H

read: MOV DX , offset Sen_Max
 MOV AH , 0AH
 INT 21H ; read input

 SUB BH , BH
 MOV BL , Sen_Real
 MOV Sentence[BX] , '\$' ; specify end of input

 MOV AL , '\$' ; specify first element in temporary operator stack
 CALL push_op_t

 SUB CH , CH
 MOV CL , Sen_Real
 INC CL
 ADC CH , 00H
 MOV BX , offset Sentence ; point BX to start of input

evaluate: MOV AL , [BX] ; is read character number?

char: AND AL , 0F0H

```

    CMP AL , 30H
    JNE not_number
    CALL ins_char
    JMP next                ; read next input

not_number: MOV AL , [BX]
            CMP AL , '#'
            JNE not_char

            MOV char_fnd , 1
            MOV isNegative , 1
            JMP next        ; read next input

not_char:  CMP char_fnd , 0
            JE peek_op
            CALL push_var

peek_op:  MOV BP , bcup_op_t        ; peek from temporary operand stack to AX
            MOV AX , [BP]
            MOV AH , AL             ; AH holds operand at top of temporary operand stack

            MOV AL , [BX]          ; read character is not number , find operator

;-----check for $

try0:  CMP AL , '$'
       JNE try1

nxt0:  CMP AH , '$'
       JNE nxt1
       JMP next

nxt1:  CMP AH , '('
       JNE nxt2
       MOV AH , 09H
       MOV DX , offset MESS1
       INT 21H                ; report error
       HLT

nxt2:  CALL move_stck
       JMP nxt0                ; recursive

;-----check for + and -

try1:  CMP AL , '+'
       JE nxt3
       CMP AL , '-'
       JNE try2

nxt3:  CMP AH , '$'
       JNE nxt4
       CALL push_op_t
       JMP next

```

```

nxt4: CMP AH , '('
      JNE nxt5
      CALL push_op_t
      JMP next

nxt5: CALL move_stck
      JMP nxt3

;-----check for * and -

try2: CMP AL , '*'
      JE  nxt6
      CMP AL , '/'
      JNE try3

nxt6: CMP AH , '*'
      JNE nxt7
      CALL move_stck
      JMP nxt6

nxt7: CMP AH , '/'
      JNE nxt8
      CALL move_stck
      JMP nxt6

nxt8: CALL push_op_t
      JMP next

;-----check for (

try3: CMP AL , '('
      JNE try4
      CALL push_op_t
      JMP next

;-----check for )

try4: CMP AH , '$'
      JNE nxt9

      MOV DX , offset MESS1
      MOV AH , 09H
      INT 21H      ; report error
      HLT

nxt9: CMP AH , '('
      JNE nxt10
      ADD bcup_op_t , 02H
      JMP next

nxt10: CALL move_stck
      JMP try4

;-----

```



```

next: INC BX      ; read next
      LOOP evaluate

      CALL GetRes   ; calculate result
      CALL DispRes  ; display result

      HLT

```

```

;functions*****

```

```

ins_char PROC NEAR

```

```

      PUSH BX

      MOV char_fnd , 31H           ; some random none-zero number : set char_fnd

      MOV AL , [BX]
      AND AL , 0FH

      SUB BH , BH
      MOV BL , char_coun
      MOV char_buff[BX] , AL
      INC char_coun

      POP BX
      RET

```

```

ins_char ENDP

```

```

;*****

```

```

push_var PROC NEAR

```

```

      PUSH BX
      PUSH DX           ; reset char_fnd

      MOV char_fnd , 0

      CMP char_coun , 1
      JNE two_char

```

```

one_char: MOV AL , char_buff[0]
          SUB AH , AH
          MOV char_sum , AX

```

```

          JMP ENDE
          ;-----

```

```

two_char: CMP char_coun , 2
          JNE three_char

```

```
MOV AL , char_buff[1]
SUB AH , AH
MOV char_sum , AX

MOV AL , char_buff[0]
MOV AH , 10
MUL AH
ADD char_sum , AX

JMP ENDE
;-----
```

```
three_char: CMP char_coun , 3
            JNE four_char
```

```
MOV AL , char_buff[2]
SUB AH , AH
MOV char_sum , AX

MOV AL , char_buff[1]
MOV AH , 10
MUL AH
ADD char_sum , AX

MOV AL , char_buff[0]
MOV AH , 100
MUL AH
ADD char_sum , AX

JMP ENDE
;-----
```

```
four_char: CMP char_coun , 4
            JNE five_char
```

```
MOV AL , char_buff[3]
SUB AH , AH
MOV char_sum , AX

MOV AL , char_buff[2]
MOV AH , 10
MUL AH
ADD char_sum , AX

MOV AL , char_buff[1]
MOV AH , 100
MUL AH
ADD char_sum , AX

MOV AL , char_buff[0]
SUB AH , AH
MOV DX , 1000
MUL DX
```

```
            ; caution:destroys DX
            ; multiply AX by DX :first_digit*1000
```

```

    ADD char_sum , AX          ; since result can't exceed 9000 , we don't need high word

    JMP ENDE
;-----

five_char: MOV AL , char_buff[4]
           SUB AH , AH
           MOV char_sum , AX

           MOV AL , char_buff[3]
           MOV AH , 10
           MUL AH
           ADD char_sum , AX

           MOV AL , char_buff[2]
           MOV AH , 100
           MUL AH
           ADD char_sum , AX

           MOV AL , char_buff[1]
           SUB AH , AH
           MOV DX , 1000          ; caution:destroys DX
           MUL DX                ; multiply AX by DX :second_digit*1000
           ADD char_sum , AX      ; since input is assumed to be 16 bit , we don't need high word

           MOV AL , char_buff[0]
           SUB AH , AH
           MOV DX , 10000
           MUL DX                ; multiply AX by DX :first_digit*10000
           ADD char_sum , AX      ; since input is assumed to be 16 bit , we don't need high word

;-----

ENDE: MOV char_coun , 0
      MOV char_buff[0] , 0
      MOV char_buff[1] , 0
      MOV char_buff[2] , 0
      MOV char_buff[3] , 0
      MOV char_buff[4] , 0

      CMP isNegative , 1
      JNE goon
      XOR char_sum , 0FFFFH      ; form 2's complement of AX
      ADD char_sum , 1
      MOV isNegative , 0

goon: MOV AX , char_sum          ; Byte0 : charsumL
      MOV BX , ptr_queue        ; Byte1 : charsumH
      MOV [BX] , AX             ; enqueue number to queue      Byte2 : 00H
      ADD ptr_queue , 0004H      ; point ptr_queue to next double word  Byte3 : 00H

      POP DX
      POP BX

```

```
RET
push_var ENDP
```

```
;*****
```

```
push_op_t PROC NEAR
```

```
SUB bcup_op_t, 02H
MOV BP, bcup_op_t
```

```
SUB AH, AH
MOV [BP], AX ; push AX to temporary operand stack
```

```
RET
push_op_t ENDP
```

```
;*****
```

```
move_stck PROC NEAR ; pops temporary operand stack & enqueues result to operand queue
```

```
PUSH BX
PUSH CX
```

```
MOV CL, AL ; make a copy of AL
MOV BP, bcup_op_t ; pop operand from temporary operand stack
MOV AX, [BP]
ADD bcup_op_t, 02H
```

```
ADD ptr_queue, 0002H ; point ptr_queue to upper word Byte0 : 00H
MOV BX, ptr_queue ; Byte1 : 00H
MOV [BX], AX ; enqueue operand to operand queue Byte2 : operator (AL)
ADD ptr_queue, 0002H ; Byte3 : 00H (AH)
```

```
MOV BP, bcup_op_t ; peek operand from temporary operand stack
MOV AX, [BP]
MOV AH, AL
MOV AL, CL ; restore old value of AL
```

```
POP CX
POP BX
```

```
RET
move_stck ENDP
```

```
;*****
```

```
GetRes PROC NEAR
```

```
MOV BP, offset sp_oprnd
```

```
begin: MOV AX, calc_queue
CMP AX, ptr_queue
JNE calc ; calc_queue <= ptr_queue
```

```

RET

calc: ADD  calc_queue , 0002H      ; point calc_queue to upper word
      MOV  BX   , calc_queue
      MOV  CX   , [BX]             ; load upper word to AX
      CMP  CX   , 0
      JE   num_fnd                 ; numbers are in lower words

op_fnd: CALL GetOp
      ADD  calc_queue , 0002H      ; point calc_queue to next double word
      JMP  begin

num_fnd: SUB  calc_queue , 0002H    ; point calc_queue to lower word
      MOV  BX   , calc_queue
      MOV  AX   , [BX]

      SUB  BP   , 0004H            ; decrement stack pointer (BP)
      MOV  [BP+0] , AX             ; push number to operand stack
      CMP  AX   , 0
      JGE  zero
      MOV  [BP+2] , 0FFFFH
      JMP  move

zero:  MOV  [BP+2] , 0000H          ; since inputs are assumed to be 16 bit , upper word is zero

move:  ADD  calc_queue , 0004H      ; point calc_queue to next double word
      JMP  begin

GetRes ENDP

;*****

GetOp  PROC NEAR                  ; use BP as operand stack SP

      MOV  AX   , [BP+0]
      MOV  operandL , AX           ; load first operand to ( operandH operandL ) pair
      MOV  AX   , [BP+2]
      MOV  operandH , AX

      MOV  AX   , [BP+4]           ; load second operand to ( DX   AX   ) pair
      MOV  DX   , [BP+6]

;-----

tryadd: CMP  CL   , '+'
      JNE  trysub

      ADD  AX   , operandL
      ADC  DX   , operandH
      JMP  push_res

;-----

```

```
trysub: CMP CL    , '-'
        JNE trymul
```

```
        SUB AX    , operandL
        SBB DX    , operandH
        JMP push_res
```

```
;-----
```

```
trymul: CMP CL    , '*'
        JNE trydiv
```

```
        CALL UnSign
```

```
        MOV [BP+4] , AX          ; destroys operand stack
        MOV [BP+6] , DX
```

$$;[(2^{16} * DX + AX) * ((2^{16} * High + Low))] = (2^{32} * DX * High + (2^{16} * ((DX * Low) + (AX * High))) + [AX * Low]$$

=>result is limited to 32 bits so we omit term with 2^{32}

```
        MUL operandL           ; form AX*Low
```

```
        PUSH DX                ; save  DX (1)
        PUSH AX                ; save  AX (1)
```

```
        MOV AX , [BP+4]        ; form AX*operandH
        MUL operandH
```

```
                                ; save  DX (2)
        PUSH DX                ; save  AX (2)
        PUSH AX
```

```
        MOV AX , [BP+6]
        MUL operandL           ; form DX*operandL
                                ; restore AX (2)
        POP  BX                ; restore DX (2)
        POP  CX
```

```
        ADD AX , BX
        ADC DX , CX            ; form (AX*High)+(DX*Low)
```

```
        CMP DX , 0
        JNE dort
```

```
        MOV DX , AX           ; shift (DX AX) pair by 16 bits
        SUB AX , AX
        JMP cont
```

```
cont: POP  BX                ; restore AX (1)
        POP  CX                ; restore DX (1)
```

```
        ADD AX , BX
        ADC DX , CX            ; form final result
        JMP chk_sign
```

```

dort: MOV AX , 0FFFFH           ; can' represent calculated value , so representing it as max positive
number
      MOV DX , 07FFFH
      JMP chk_sign
;-----

trydiv: CMP CL , '/'
      JNE tryerr

      CALL UnSign

      MOV [BP+4] , AX           ; destroys operand stack
      MOV [BP+6] , DX

      CMP operandH , 0000H
      JNE ex_div

; (DX AX) / Low
      MOV AX , DX              ; divide high word (DX) first
      SUB DX , DX              ; convert word to double word
      DIV operandL
      PUSH AX                  ; store high word of quotient
      MOV AX , [BP+4]
      DIV operandL
                                ; low word of quotient is in AX
      POP DX                   ; restore high word of quotient
      JMP chk_sign

; (DX AX) / (High Low) =~ DX / High

ex_div: MOV DX , AX
      SUB DX , DX
      DIV operandH
      SUB DX , DX              ; destroy remainder

      JMP chk_sign

;-----

tryerr: MOV DX , offset MESS1
      MOV AH , 09H
      INT 21H
      HLT

chk_sign: CMP isSigned , 1
      JNE push_res

      CMP DX , 0               ; if DX 's MSB is one , then overflow has occurred
      JL over

      XOR AX , 0FFFFH
      XOR DX , 0FFFFH
      ADD AX , 1

```

```

        ADC DX , 0
        JMP push_res

over: MOV DX , 0000H
        MOV AX , 8000H
        JMP push_res

push_res: ADD BP , 0004H ; push result back to operands stack
        MOV [BP+0] , AX
        MOV [BP+2] , DX

        RET

GetOp ENDP

;*****

UnSign PROC NEAR

        MOV isSigned , 0 ; reset sign

check_1: CMP operandH , 0 ; is 'operand' variable negative?
        JL negate_1
        JMP check_2

negate_1: XOR operandL , 0FFFFH ; 1's complement of operandL
        XOR operandH , 0FFFFH ; 1's complement of operandH
        ADD operandL , 1 ; 2's complement of
        ADC operandH , 0 ; operand
        MOV isSigned , 1

check_2: CMP DX , 0 ; is (DX AX) pair negative?
        JL negate_2
        RET ; end of procedure

negate_2: XOR AX , 0FFFFH ; 1's complement of AX
        XOR DX , 0FFFFH ; 1's complement of DX
        ADD AX , 1 ; 2's complement of
        ADC DX , 0 ; (DX AX) pair

;MOV [BP+4] , AX ; destroys operand stack
;MOV [BP+6] , DX

        CMP isSigned , 0
        JE signIt
        MOV isSigned , 0 ; (-) * (-) = (+)
        RET ; end of procedure

signIt: MOV isSigned , 1 ; (+) * (-) = (-)
        RET

UnSign ENDP

```


;*****

DispRes PROC NEAR

MOV operandL , 0000H ; making operand zero , so it can't effect IsSign
MOV operandH , 0000H

MOV AX , [BP+0]
MOV DX , [BP+2]

CALL UnSign

MOV ResWord0 , AX
MOV ResWord1 , DX

MOV CX , 10
MOV BX , 12 ; 3+9

;-----

again: SUB DX , DX
MOV AX , ResWord1
DIV Divisor

MOV ResWord1 , AX ; store AX (high word of new dividend)
MOV AX , ResWord0

; DX already contains correct value
DIV Divisor ; Move High word of new dividend to ResWord1
MOV ResWord0 , AX

OR DL , 30H ; DX contains decimal number
MOV MESS2[BX] , DL

DEC BX
LOOP again

;-----

CMP isSigned , 01H
JE N_Sign

MOV MESS2[2] , '+'
JMP done

N_Sign: MOV MESS2[2] , '-'

done: MOV DX , offset MESS2
MOV AH , 09H
INT 21H
RET

DispRes ENDP

;*****

CODE ENDS

END START