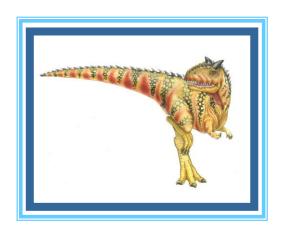
# Operating System LAB-03 Prepared By

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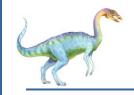


#### **EXPERIMENT NO-3**

# CPU SCHEDULINGALGORITHMS ROUND ROBIN

AIM: To simulate the CPU scheduling algorithm round-robin





#### **EXPERIMENT NO-03**

### DESCRIPTION

To aim is to calculate the average waiting time. There will be a time slice, each process should be executed within that time-slice and if not it will go to the waiting state so first check whether the burst time is less than the time-slice. If it is less than it assign the waiting time to the sum of the total times. If it is greater than the burst-time then subtract the time slot from the actual burst time and increment it by time-slot and the loop continues until all the processes are completed.





**Step 1: Start the process** 

Step 2: Accept the number of processes in the ready Queue and time quantum (or) time slice

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Calculate the no. of time slices for each process where No. of time slice for process (n) = burst time process (n)/time slice

Step 5: If the burst time is less than the time slice then the no. of time slices =1.

Step 6: Consider the ready queue is a circular Q, calculate



- a) Waiting time for process (n) = waiting time of process(n-1)+ burst time of process(n-1) + the time difference in getting the CPU from process(n-1)
- b) Turnaround time for process(n) = waiting time of process(n) + burst time of process(n)+ the time difference in getting CPU from process(n).

Step 7: Calculate

- c) Average waiting time = Total waiting Time / Number of process
- d) Average Turnaround time = Total Turnaround Time / Number of process Step
- 8: Stop the process





### Round Robin (RR)

- □ Each process gets a small unit of CPU time (time quantum q), usually 10-100 milliseconds. After this time has elapsed, the process is preempted and added to the end of the ready queue.
- □ If there are n processes in the ready queue and the time quantum is q, then each process gets 1/n of the CPU time in chunks of at most q time units at once. No process waits more than (n-1)q time units.
- ☐ Timer interrupts every quantum to schedule next process
- Performance
  - □ q large  $\Rightarrow$  FIFO
  - $q \text{ small} \Rightarrow q \text{ must be large with respect to context switch,}$ otherwise overhead is too high

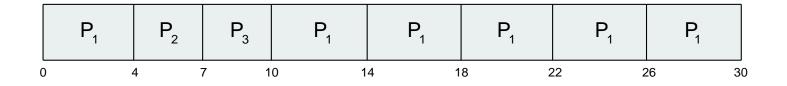




## **Example of RR with Time Quantum = 4**

<u>Process</u>	<b>Burst Time</b>
$P_1$	24
$P_2$	3
$P_3$	3

The Gantt chart is:

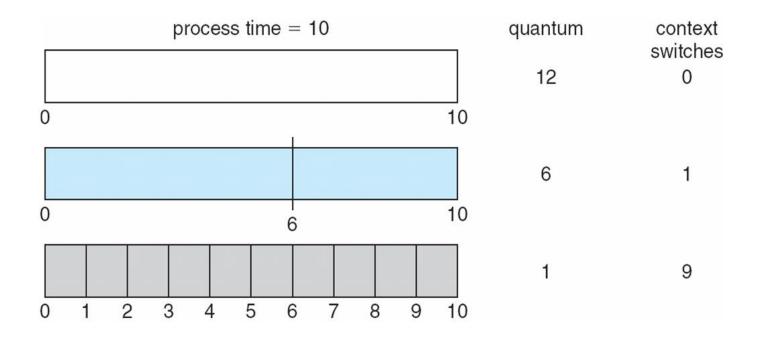


- Typically, higher average turnaround than SJF, but better response
- □ q should be large compared to context switch time
- □ q usually 10ms to 100ms, context switch < 10 usec





#### **Time Quantum and Context Switch Time**







# **EXPERIMENT NO-03: C PROGRAM**

```
#include<stdio.h>
int main()
  int n,i,qt,count=0,temp,sq=0,bt[10],wt[10],tat[10],rem_bt[10];
  float awt=0,atat=0;
  printf("Enter the Number of process: ");
  scanf("%d",&n);
  for(i=0;i<n;i++)
    printf("Enter burst time for Process P%d~~",i+1);
    scanf("%d",&bt[i]);
    rem_bt[i]=bt[i];
  printf("Enter quantum time: ");
  scanf("%d",&qt);
```





## **EXPERIMENT NO-03: C PROGRAM**

```
while(1)
    for(i=0,count=0;i<n;i++)</pre>
       temp = qt;
       if(rem_bt[i]==0)
          count++;
          continue;
       if(rem_bt[i]>qt)
          rem_bt[i]=rem_bt[i]-qt;
       else
          if(rem_bt[i]>=0)
            temp=rem_bt[i];
            rem_bt[i]=0;
```





# **EXPERIMENT NO-03: C PROGRAM**

```
sq=sq+temp;
         tat[i]=sq;
    if(n==count)
       break:
  printf("\n\tPROCESS\t BURST TIME \t WAITING TIME\tTURNAROUND TIME\n");
  for(i=0;i<n;i++)
    wt[i]=tat[i]-bt[i];
    awt=awt+wt[i];
    atat=atat+tat[i];
    printf("\t%d \t %d \t\t %d \t\t %d \n",i+1,bt[i],wt[i],tat[i]);
  awt=awt/n;
  atat=atat/n;
  printf("\nThe Average Turnaround time is -- %f",atat);
  printf("\nThe Average Waiting time is -- %f ",awt);
  return 0;
```



## End of Lab 02

