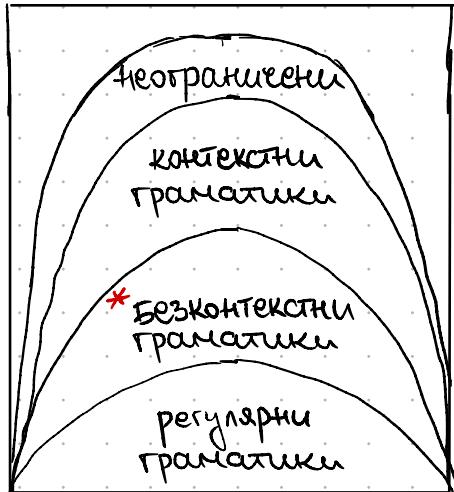


$$\textcircled{1} \quad \text{X+1} \quad \text{Sort}(L) = \{ a^l b^m c^n \mid l, m, n \in \mathbb{N} \} \quad \Sigma = \{ a, b, c \}$$

$$\text{Възможне } L = \{ a, b, c \}^* \quad \text{Тогава } \text{Sort}(L) = a^n b^n c^n$$

Класове граматики от ляво надясно се разделят на правилата им



$$\lambda \rightarrow B \quad | \quad a \rightarrow b$$

$$A \rightarrow \lambda \quad | \quad A \alpha \rightarrow \lambda \alpha \quad (\text{изиска символи до пром.)}$$

$$A \rightarrow a$$

$$A \rightarrow aB$$

$$\text{релација } X \Delta \alpha, \text{ където } X \in V \cup \Sigma \quad \alpha \in (V \cup \Sigma)^*$$

$X, Y$  - променливи/букви

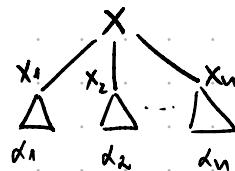
$A, B, C$  - променливи

$a, b, c$  - букви

$$0) \quad \underline{\underline{X \Delta X}}$$

$$1) \quad \underline{\underline{X \rightarrow x_1 \dots x_n}}, \quad \underline{\underline{x_1 \Delta d_1, \dots, x_n \Delta d_n}} \quad X \Delta d_1 \dots d_n$$

$$l = 1 + \sup \{ l_1, \dots, l_n \}$$



$$\text{За } n=0: \quad X \rightarrow x_1 \dots x_n = \epsilon \quad \text{1) } \underline{\underline{X \rightarrow x_1 \dots x_n}}, \quad \underline{\underline{x_1 \Delta d_1, \dots, x_n \Delta d_n}} \quad X \Delta d_1 \dots d_n$$

$$l = 1 + \sup \{ l_1, \dots, l_n \} = \underline{\underline{1+0=1}}$$

$\emptyset$

Пример:

$$L = \{ a^n b^m c^k \mid n+m \leq k \}$$

$$C = cC \mid c$$

$$S = aSC \mid bBC \mid \epsilon$$

$$B = Bbc \mid \epsilon$$

$$C = cC \mid \epsilon$$

$$S \rightarrow AC \mid \epsilon$$

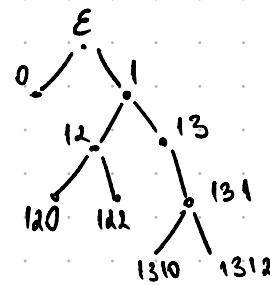
$$A \rightarrow aAc \mid bBc \mid \epsilon$$

$$B \rightarrow bBc \mid \epsilon$$

$$\Sigma = \{0, 1, \dots, b-1\}$$

$$T \subseteq \Sigma^* \text{ е } \text{гбв} \stackrel{\text{def}}{\iff} \text{Pref}(T) = \Pi$$

МН-БО от думи засв. относно префикси



$$T = \{E, 0, 1, 12, 13, 120, 122, 131, 1310, 1312\}$$

Оптимално гбв  $\Pi$ :  $d_i \in T \quad \left\{ \begin{array}{l} \Rightarrow d_j \in \Pi \\ j < i \end{array} \right.$

$$\{0\}^*. \{1\}^* \cup E \quad \epsilon \cdot \xrightarrow{\quad} \overset{0}{\underset{|}{\cdot}} \xrightarrow{\quad} \overset{00}{\underset{|}{\cdot}} \xrightarrow{\quad} \overset{000}{\underset{|}{\cdot}} \xrightarrow{\quad} \cdots$$

$P = (\Pi, \lambda)$  - гбв та извод, съвместимо с граматиката  $G$

$\Pi$ -оптимално гбв  $\{0, \dots, b-1\}^*$  където  $b = \max \{|\alpha| : A \rightarrow \alpha \text{ е правило в } G\}$

$\lambda$ -ЕТИКЕТ  $\lambda: \Pi \rightarrow V \cup \Sigma$

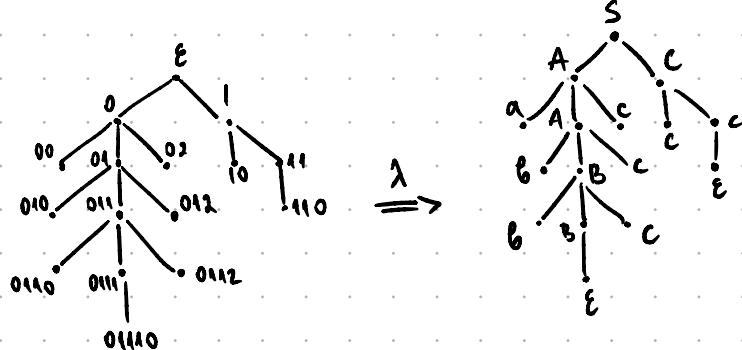
Фекса  $d \in \Pi$  и  $K+1 = |\{d_i \in \Pi \mid i < b\}|$

Тогава  $\lambda(d) \geq \lambda(d_0), \lambda(d_1), \dots, \lambda(d_K)$  е правило в  $G$

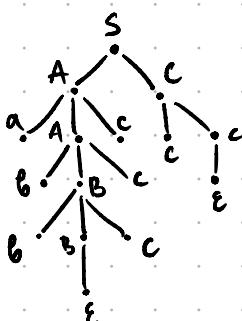
$$\lambda(\epsilon) = S \quad \lambda(011) = B$$

$$\text{height}(\Pi) = \max \{ |d| : d \in \Pi \}$$

$$\lambda(01) \xrightarrow{G} \underbrace{\lambda(010)}_a \underbrace{\lambda(011)}_A \underbrace{\lambda(012)}_C$$



$$\xrightarrow{\lambda}$$



$$\begin{aligned} S &\rightarrow Ac \\ C &\rightarrow cCc \mid \epsilon \\ A &\rightarrow aAc \mid bBc \mid \epsilon \\ B &\rightarrow bBc \mid \epsilon \end{aligned}$$

$A \trianglelefteq_G d \iff$  същ. гбв та извод  $P = (\Pi, \lambda)$  съвместимо с  $G$

$$\text{height}(\Pi) = l \quad \text{leaves}(\Pi) = d$$

$$\text{root}(P) = A$$

$A \xrightarrow{G} \alpha \Rightarrow |\alpha| \leq b^l$ , когдa  $b = \max\{|A|, A \xrightarrow{G} \alpha\}$  е правило в  $G$

$|\alpha| > b^l \Rightarrow A \xrightarrow{G^{l+1}} \alpha$   
за  $\alpha \in \tilde{\Delta}(G)$

$$\tilde{\Delta}(G) = \{\alpha \in \Sigma^* \mid S \xrightarrow{G} \alpha\}$$

$$\tilde{\Delta}(q) = \{\alpha \in \Sigma^* \mid \tilde{\Delta}_A^*(q, \alpha) \in F\} \quad \tilde{\Delta}_A^*(q) = \{\alpha \in \Sigma^* \mid |\alpha| \leq l \text{ и } \tilde{\Delta}_A^*(q, \alpha) \in F\}$$

$$\tilde{\Delta}_G(X) = \{\alpha \in \Sigma^* \mid X \xrightarrow{G} \alpha\} \quad \tilde{\Delta}_G^l(X) = \{\alpha \in \Sigma^* \mid X \xrightarrow{G} \alpha \text{ и } |\alpha| \leq l\}$$

за  $X \in V \cup \Sigma$

$$X \in \Sigma^*: \quad \tilde{\Delta}_G^0(X) = \emptyset$$

т.б.  $\tilde{\Delta}_G^{l+1}(A) = \bigcup \{\tilde{\Delta}_G^l(x_i) \mid A \xrightarrow{G} x_1 \dots x_n \text{ е правило в } G\}$

$$\tilde{\Delta}_G^l(X) \subseteq \tilde{\Delta}_G^{l+1}(X)$$

док: Используя по л

$$\tilde{\Delta}_G^l(X) = \bigcup \tilde{\Delta}_G^l(X)$$

$$\cdot l=0 \quad \tilde{\Delta}_G^0(A) = \{\alpha \in \Sigma^* \mid A \xrightarrow{G} \alpha\} \quad A \xrightarrow{G} x_1 \dots x_n \text{ и}$$

$$\alpha \in \tilde{\Delta}_G^0(x_i) \quad x_i \xrightarrow{G} \alpha_i \text{ и } \alpha_i \in \tilde{\Delta}_G^0(\alpha_i)$$

$$S \xrightarrow{a} Sb \mid \varepsilon$$

$$\alpha = \alpha_1 \dots \alpha_n$$

$$\tilde{\Delta}_G^{l+1}(S) = \tilde{\Delta}_G^l(a) \cdot \tilde{\Delta}_G^l(S) \cdot \tilde{\Delta}_G^l(b)$$

$\underbrace{\tilde{\Delta}_G^l(a)}_{\{a\}} \cdot \underbrace{\tilde{\Delta}_G^l(S)}_{\{S\}} \cdot \underbrace{\tilde{\Delta}_G^l(b)}_{\{\varepsilon\}}$

$$\alpha \in \tilde{\Delta}_G^0(x_1) \dots \tilde{\Delta}_G^0(x_n)$$

$$\tilde{\Delta}_G^l(A) \subseteq \bigcup \{\tilde{\Delta}_G^0(x_1) \dots \tilde{\Delta}_G^0(x_n)$$