

Edge Computing and UAV Swarm Cooperative Task Offloading in Vehicular Networks

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Abstract—Recently, unmanned aerial vehicle (UAV) swarm has been advocated to provide diverse data-centric services including data relay, content caching and computing task offloading in vehicular networks due to their flexibility and conveniences. Since only offloading computing tasks to edge computing devices (ECDs) can not meet the real-time demand of vehicles in peak traffic flow, this paper proposes to combine edge computing and UAV swarm for cooperative task offloading in vehicular networks. Specifically, we first design a cooperative task offloading framework that vehicles' computing tasks can be executed locally, offloaded to UAV swarm, or offloaded to ECDs. Then, the selection of offloading strategy is formulated as a mixed integer nonlinear programming problem, the object of which is to maximize the utility of the vehicle. To solve the problem, we further decompose the original problem into two subproblems: minimizing the completion time when offloading to UAV swarm and optimizing the computing resources when offloading to ECD. For offloading to UAV swarm, the computing task will be split into multiple subtasks that are offloaded to different UAVs simultaneously for parallel computing. A Q-learning based iterative algorithm is proposed to minimize the computing task's completion time by equalizing the completion time of its subtasks assigned to each UAV. For offloading to ECDs, a gradient descent algorithm is used to optimally allocate computing resources for offloaded tasks. Extensive simulations are lastly conducted to demonstrate that the proposed scheme can significantly improve the utility of vehicles compared with conventional schemes.

Index Terms—Vehicular networks, UAV swarm, edge computing, task offloading.

I. INTRODUCTION

With the rapid development of vehicular networks in recent years, vehicles have become smarter and connected more closely. Consequently, various vehicle applications (e.g., entertainments, navigation, augmented/virtual reality applications) have increased significantly [1], [2]. However, a large number of computing resources will be occupied when processing these applications' tasks, resulting in a high delay due to the limited computing resources in vehicles.

To reduce the computing delay, edge computing has been applied to realize computing task offloading. Although edge computing presents significant benefits, it still has some deficiencies [3]–[5]. On one hand, edge computing usually relies on fixed edge servers with limited coverage and low flexibility. On the other hand, during peak traffic periods, edge computing devices (ECDs) are prone to overload and oversaturation. Therefore, a flexible and efficient offloading scheme for vehicles should be investigated.

Fortunately, unmanned aerial vehicle (UAV) swarm has been advocated for providing efficient computing task offloading services for vehicles, by harnessing its advantages including low cost, high mobility, and line-of-sight (LoS) communication. Zhang *et al.* [6] use multiple UAVs to execute computing tasks offloaded from mobile devices. Meanwhile, Zhu *et al.* [7] devise a UAVs-assisted computation offloading paradigm where multiple UAVs exist. Besides, Zhang *et al.* [8] exploit multiple UAVs to provide offloading services for mobile devices. Although the above works can realize the computing task offloading, the load balance among UAVs is not sufficiently considered. In addition, rather than offloading the computing task to the same UAV, the computing task can be split into multiple subtasks and executed simultaneously to enjoy the low latency of parallel computing. Keeping that in view, offloading vehicles' computing tasks to UAV swarm is still an open and vital issue.

In this paper, an edge computing and UAV swarm assisted cooperative task offloading scheme in vehicular networks is studied. Specifically, a cooperative task offloading framework where vehicles' computing tasks can be executed locally, offloaded to UAV swarm, or offloaded to ECDs is first designed. In order to maximize the utility of the vehicle, the selection of offloading strategy is then formulated as a mixed integer nonlinear programming problem. Afterwards, the original problem is decomposed into two subproblems: minimizing the completion time when offloading to UAV swarm and optimizing the computing resources when offloading to ECD. Furthermore, when offloading to UAV swarm, a Q-learning based iterative algorithm is proposed to minimize the task completion time. When offloading to ECD, a gradient descent algorithm is utilized to find the optimal computing resources allocated for offloaded tasks. At last, extensive experimental results demonstrate that the proposed scheme outperforms the conventional schemes in terms of the vehicle's utility.

The remainder of this article is organized below. Section II reviews the related work and the system model is described in Section III. Section IV formulates the problem followed by the solution to the subproblems in Section V. The performance of the proposed scheme is evaluated in Section VI. Finally, Section VII closes this paper with conclusions.

II. RELATED WORK

Various works have been done to improve the efficiency of task offloading in vehicular networks. He *et al.* [9] de-

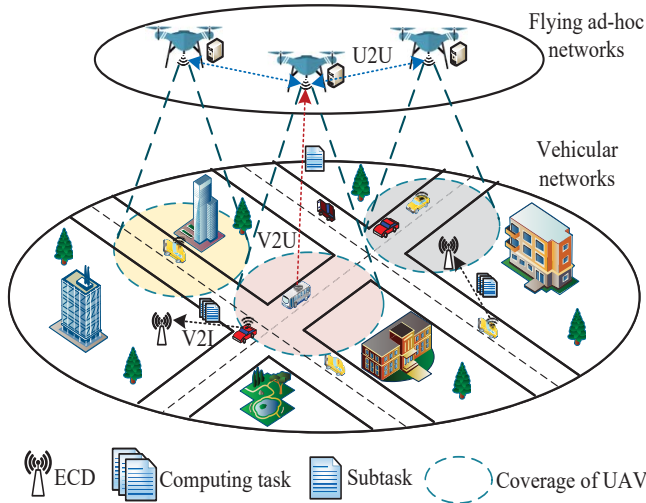


Fig. 1. Network architecture of edge computing and UAV swarm cooperative task offloading in vehicular networks.

signed an edge network that combines caching spaces, edge computational capacities and communication bandwidths with task offloading to enhance vehicles' satisfaction of Quality of Experience (QoE). Misra *et al.* [10] proposed a mobility-aware task offloading scheme, which aims to minimize the task computation time in a software-defined vehicular network. Yadav *et al.* [11] offloaded user tasks to a nearby vehicular node and minimized service latency and energy consumption.

However, the above works only offload vehicles' computing tasks to ECDs, which cannot meet the real-time demand of vehicles in the peak traffic period. In our work, we use edge computing and UAV swarm cooperative task offloading to guarantee vehicles' real-time demand.

Besides, many works use UAVs to help the offloading. Xu *et al.* [12] focused on the computing delay of offloaded tasks executed by multiple UAVs with the aim of minimizing the task completion time. Zhao *et al.* [13] proposed a Software-Defined Networking (SDN) enabled UAV-assisted vehicular computation offloading framework to minimize the system cost of vehicles' computing tasks. Haber *et al.* maximized the rate of served requests by optimizing UAVs' positions, the allocated resources, and the offloading decisions in [14].

However, the above works only offload vehicles' computing tasks to a single UAV, which cannot further reduce the latency. In our work, the computing task is split into several subtasks that are offloaded to different UAVs simultaneously for parallel execution to reduce the completion time.

III. SYSTEM MODEL

A. Network Model

As shown in Fig. 1, we consider an edge computing and UAV swarm cooperative vehicular network. The flying ad-hoc networks consists of several UAVs that can communicate with each other through UAV-to-UAV (U2U) links. Besides, the vehicular network is established by multiple vehicles and ECDs. There are N UAVs, and the set is

denoted by $\mathbf{U} = \{u_1, \dots, u_n, \dots, u_N\}$. Each vehicle has a computing task to execute and the set of vehicles is denoted as $\mathbf{V} = \{v_1, \dots, v_m, \dots, v_M\}$, where its total number is M . Besides, there exists several ECDs, denoted as $\mathbf{E} = \{e_1, \dots, e_s, \dots, e_S\}$. With Vehicle-to-UAV (V2U) and Vehicle-to-Infrastructure (V2I) links, vehicles can offload their computing tasks to UAV swarm or ECDs.

The task of vehicle v_m is denoted as T_m , which can be described by a tuple: $\{D_m, c_m, \tau_m, F_m\}$, where D_m represents the data size (bit), c_m represents the number of CPU cycles needed per bit, τ_m represents the maximum delay tolerance (s) and F_m represents the total number of CPU cycles required by T_m . Apparently, $F_m = c_m D_m$.

Each computing task can be executed with three strategies: local execution, offload to UAV swarm and offload to ECD. Especially, when v_m chooses to offload T_m to UAV swarm for execution, it can be divided into multiple subtasks for parallel computing (described in Sect. V. A). The reason is that subtasks executed in parallel can reduce the completion time of computing tasks. For data-partition oriented applications, the input data is known in advance and can be partitioned arbitrarily for parallel processing because of bit independence. The typical examples are file/graphics compression, virus scanning, visual applications, and object recognition [15], [16].

B. Communication Model

We consider a three-dimensional coordinate system, the coordinates of vehicle v_m , UAV u_n and ECD e_s are $g_m = (x_m(t), y_m(t), 0)$, $q_n = (x_n(t), y_n(t), h_n(t))$ and $o_s = (x_s, y_s, 0)$, respectively. The Orthogonal Frequency Division Multiple Access (OFMDA) technology is adopted in this paper [17]. The bandwidth allocated to each vehicle and each UAV is B_0 and B_1 , respectively. Based on [8], the channel gain of the communication link between v_m and u_n is

$$h_{m,n} = \frac{\beta_0}{\|q_n - g_m\|^2}, \quad (1)$$

where β_0 is the channel power gain at the reference distance $d_0 = 1m$. According to Shannon theorem, the communication rate between v_m and u_n is

$$r_{v_m, u_n} = B_0 \log_2 \left(1 + \frac{p_m h_{m,n}}{B_0 N_0} \right), \quad (2)$$

where p_m is the transmission power of v_m and N_0 is the noise power density.

We view the U2U link between UAV u_n and UAV u_j as LoS link, so the communication rate can be calculated by

$$r_{u_n, u_j} = B_1 \log_2 \left(1 + \frac{p_n \beta_1}{\|q_n - q_j\|^2 B_1 N_0} \right), \quad (3)$$

where B_1 is the bandwidth allocated to each UAV, β_1 is the channel power gain at the reference distance $d_0 = 1m$.

The communication rate between v_m and e_s is derived as

$$r_{v_m, e_s} = B_0 \log_2 \left(1 + \frac{p_m h^2}{\|g_m - o_s\|^2 \sigma} \right), \quad (4)$$

where h is uplink channel fading coefficient and σ is Gaussian noise power.

IV. PROBLEM FORMULATION

A. Vehicle's Utility When Computing Locally

The limited local computing resources will result in a long task completion time. The time needed for local computing is

$$t_{T_m,L} = \frac{F_m}{f_m^L}, \quad (5)$$

where f_m^L is the computing resources assigned to T_m from v_m . For v_m , if v_m obtains short task completion time, v_m will have more utility. Consequently, v_m 's utility is

$$s_m^L = C - e^{-(\tau_m - t_{T_m,L})}, \quad (6)$$

where C is a constant which makes the utility non-negative.

B. Vehicle's Utility When Offloading to UAV Swarm

In pursuit of low latency, T_m can be decomposed into several subtasks, each of them will be executed by different UAVs simultaneously, i.e., $T_m = \{t_{m,1}, \dots, t_{m,i}, \dots, t_{m,I}\} (I \leq N)$ will be offloaded to $U_m = \{u_{m,1}, \dots, u_{m,i}, \dots, u_{m,I}\}$. The data size set of the subtask is $D_{T_m} = (d_{m,1}, \dots, d_{m,i}, \dots, d_{m,I})$, where $d_{m,i}$ is $t_{m,i}$'s data size. The total number of CPU cycles required by $t_{m,i}$ is $f_{m,i}$ and $f_{m,i} = c_m d_{m,i}$. According to the above definition, we can obtain the following formula:

$$\sum_{i=1}^I d_{m,i} = D_m, \sum_{i=1}^I f_{m,i} = F_m, \forall v_m \in \mathbf{V}. \quad (7)$$

When subtask $t_{m,i}$ is offloaded to UAV $u_{m,i}$, $t_{m,i}$ can be transmitted directly if v_m is within $u_{m,i}$'s communication range. Otherwise, $t_{m,i}$ needs a relay UAV. We introduce a binary variable x_m^i . $x_m^i = 1$ denotes that v_m is connected with $u_{m,i}$, otherwise $x_m^i = 0$. Since UAVs can communicate with each other, when $x_m^i = 0$, the relay UAV is

$$u_{m,i}^{relay} = \arg \max_{u_j \in \mathbf{U}, u_j \neq u_{m,i}} x_{m,j} \frac{1}{\frac{1}{r_{v_m, u_j}} + \frac{1}{r_{u_j, u_{m,i}}}}, \quad (8)$$

where $\frac{1}{r_{v_m, u_j}} + \frac{1}{r_{u_j, u_{m,i}}}$ is the time required to transmit one bit of data from v_m to $u_{m,i}$ through u_j , i.e., $\frac{1}{r_{v_m, u_{m,i}}}$. Therefore, the transmission time is

$$t_{t_{m,i}, u_{m,i}}^{tra} = \frac{d_{m,i}}{r_{v_m, u_{m,i}}}. \quad (9)$$

If subtasks are assigned randomly without considering the load of each UAV, the computing task's completion time will be prolonged. We use $F_{m,i}^{wait}$ to represent the CPU cycles required for the remaining tasks of $u_{m,i}$, then $u_{m,i}$ needs to wait before executing its next subtask. The waiting time is

$$t_{u_{m,i}}^{wait} = \frac{F_{m,i}^{wait}}{f_u}, \quad (10)$$

where f_u is the amount of computing resources allocated from each UAV to computing tasks. With the help of $t_{u_{m,i}}^{wait}$, the current load of $u_{m,i}$ can be obtained. The computing time required for $u_{m,i}$ to execute subtask $t_{m,i}$ is

$$t_{u_{m,i}, t_{m,i}}^{cmt} = \frac{f_{m,i}}{f_u}. \quad (11)$$

Because the data size of computing results is relatively small, we omit the transmit time of results. Since $u_{m,i}$ is still executing the remaining tasks during $t_{m,i}$'s transmission, $t_{m,i}$ can be executed immediately if $t_{t_{m,i}, u_{m,i}}^{tra}$ is greater than $t_{u_{m,i}}^{wait}$. Therefore, the total time required to execute subtask $t_{m,i}$ is

$$t_{t_{m,i}, u_{m,i}} = \max(t_{t_{m,i}, u_{m,i}}^{tra}, t_{u_{m,i}}^{wait}) + t_{u_{m,i}, t_{m,i}}^{cmt}. \quad (12)$$

As T_m is divided into several subtasks, the execution time of T_m is determined by the largest execution time, i.e.,

$$t_{T_m, U} = \{t_{t_{m,i}, u_{m,i}} | \forall t_{m,k} \in T_m : t_{t_{m,i}, u_{m,i}} \geq t_{t_{m,k}, u_{m,k}}\}. \quad (13)$$

As executing offloaded tasks consumes considerable battery energy of UAVs, the vehicle has to pay the UAVs. Similar to [18], the communication energy is ignored since it's too small compared to the computing energy. Therefore, the payment for offloading computing task T_m to UAV swarm is

$$p_{T_m, U} = \sum_{i=1}^I \gamma_U \delta d_{m,i} (f_u)^2 = \gamma_U \delta F_m (f_u)^2, \quad (14)$$

where γ_U is UAV's unit price of energy and δ is a constant related to the hardware architecture. We use the satisfaction of the vehicle and the payment to the UAV swarm to measure the utility of the vehicle. Thereby, v_m 's utility is

$$s_m^U = \alpha (C - e^{-(\tau_m - t_{T_m, U})}) - p_{T_m, U}, \quad (15)$$

where α is the offloading willingness.

C. Vehicle's Utility When Offloading to ECD

According to the principle of proximity, the nearest ECD will be selected. The time required for transmitting T_m from v_m to e_s is $t_{T_m, e_s}^{tra} = \frac{D_m}{r_{v_m, e_s}}$. The ECD e_s will compute the task as soon as it arrives, the time of e_s calculating T_m is

$$t_{e_s, T_m}^{cmt} = \frac{F_m}{f_m^{e_s}}, \quad (16)$$

where $f_m^{e_s}$ is the allocated computing resources from e_s to T_m . Therefore, T_m 's completion time is

$$t_{T_m, e_s} = t_{T_m, e_s}^{tra} + t_{e_s, T_m}^{cmt}. \quad (17)$$

Similar to offloading to UAV swarm, the vehicle has to pay the ECD. Denote γ_E as the unit price of ECD's energy, then the utility of v_m when offloading to e_s is

$$s_m^E = \alpha (C - e^{-(\tau_m - t_{T_m, e_s})}) - \gamma_E \delta F_m (f_m^{e_s})^2. \quad (18)$$

D. Problem Formulation

We introduce $d_m = \{d_m^L, d_m^U, d_m^E\}$ to describe the computing strategy. Binary variable d_m^L, d_m^U, d_m^E indicate that whether T_m is executed locally, offloaded to UAV swarm, or offloaded to ECD, respectively. $d_m^L = 1$ means that T_m is executed locally, $d_m^L = 0$ means other situations. d_m^U and d_m^E are similar to d_m^L . This paper aims to maximize the utility of vehicle, and the optimization problem is formulated as:

$$\begin{aligned} P_1 : & \max_{d_m} (d_m^L s_m^L + d_m^U s_m^U + d_m^E s_m^E) \\ \text{s.t.} \quad & C_1 : d_m^L + d_m^U + d_m^E \leq 1 \\ & C_2 : d_m^L t_{T_m, L} + d_m^U t_{T_m, U} + d_m^E t_{T_m, e_s} \leq \tau_m \end{aligned} \quad (19)$$

where C_1 indicates that the computing task can only be executed locally, offloaded to UAV swarm, or offloaded to ECD. C_2 means that T_m 's completion time should be less than the tolerable delay τ_m .

V. PROPOSED OPTIMIZATION APPROACH

P_1 is NP-hard since it is a 0-1 integer programming problem that cannot be solved in linear time. By observing the problem, we decompose it into three subproblems: 1) maximizing the utility when offloading to UAV swarm, 2) maximizing the utility when offloading to ECD, and 3) maximizing the utility when the task is computed locally.

A. Maximize Vehicle's Utility When Offloading to UAV swarm

According to Eq. (15), s_m^U is determined by $t_{T_m,U}$. As long as $t_{T_m,U}$ is minimized, s_m^U is maximized. Therefore, the subproblem of maximizing s_m^U can be transformed into minimizing $t_{T_m,U}$, i.e.,

$$P_2 : \min_{D_{T_m}} t_{T_m,U} \quad \text{s.t. } t_{T_m,U} \leq \tau_m \quad (20)$$

Obviously, the minimum latency $t_{T_m,U}^{\min}$ is reached when the parallel subtasks take the same completion time. Therefore, the first step is to determine UAVs allocation state, i.e., which UAVs are selected to execute T_m 's subtasks. As a kind of reinforcement learning, Q-learning is an effective solution when the number of actions and states of a decision problem is limited [19]. Consequently, the Q-learning approach is used to find the optimal UAVs allocation state s^* and the Q-learning can be expressed as a tuple $\{S, A, R(s, a)\}$. Specially,

1) $S = \{s_1, \dots, s_i, \dots, s_I\}$ is the allocation states set of UAVs, where $s_i = (u_{m,1}, u_{m,2}, \dots, u_{m,i})$ represents the set of UAVs allocated to execute T_m 's subtasks. To realize load balance, the UAVs in status s_i are selected from U in ascending order of waiting time, i.e., $t_{u_{m,1}}^{\text{wait}} < t_{u_{m,2}}^{\text{wait}} < \dots < t_{u_{m,i}}^{\text{wait}}$.

2) $A = \{a_1, \dots, a_i, \dots, a_I\}$ represents the set of replacement actions for each state, in which $a_i = (u_i^-, u_{i+1}^+)$ represents the replacement action can be taken in state s_i . $a_i = (1, 0)$ indicates that $u_{m,i}$ is removed from the allocation set to reach state s_{i-1} . $a_i = (0, 1)$ indicates that $u_{m,i+1}$ is added to the allocation set to reach status s_{i+1} .

3) $R(s, a) = \{R_1(s_1, a_1), \dots, R_i(s_i, a_i), \dots, R_I(s_I, a_I)\}$ is the reward set of performing action a_i in state s_i .

The Q-value is described as the state-action pair $Q(s_i, a_i)$, and the update of the Q-value can be expressed as:

$$Q(s_i, a_i) = (1 - \nu)Q(s_i, a_i) + \nu[R_i(s_i, a_i) + \gamma \max_{a_i'} Q(s_{i'}, a_{i'})], \quad (21)$$

where s_i is the current allocation state of UAVs. $s_{i'}$ is the new allocation state of UAVs after the execution of action a_i , and $a_{i'}$ is the action that can be taken in state $s_{i'}$. ν is the learning rate and γ is the discount factor used to determine the impact of future rewards on present rewards.

Algorithm 1 : Q-Learning based Iterative Algorithm

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1: Initialize Sort  $U$  decremently by waiting time;
2: Repeat
3:   Selects a UAV replacement action  $a_i$ ;
4:   Calculate the reward according to Eq. (22);
5:   Update Q-value according to Eq. (21);
6:    $s_{i'} \leftarrow s_i$ ;
7:    $a_{i'} \leftarrow a_i$ ;
8:   Until Q-value converges
9:   Find  $s^*$  according to Q-value;
10:  Calculate  $t_{sub}^{\min}, t_{sub}^{\max}$  in  $s^*$  and set  $t_{T_m,U}^{\min} = t_{sub}^{\min}$ ;
11:  while  $t_{T_m,U}^{\min} < t_{sub}^{\max}$  do
12:    Calculate  $d_{m,k}$  in  $s^*$  according to Eq. (24) and Eq. (26);
13:    if  $\sum_1^{I^*} d_{m,i} - D_m < \varepsilon_1$  then
14:      Break
15:    end if
16:     $t_{T_m,U}^{\min} = t_{T_m,U}^{\min} + \varepsilon_2$ ;
17:  end while
18: Calculate  $s_m^U$  according to Eq. (15);

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In order to find the optimal UAVs allocation state and achieve load balance, we formulate $R_i(s_i, a_i)$ as follows:

$$R_i(s_i, a_i) = \frac{F_m}{\sum_{k=1}^{i'} d_{m,k} \Big|_{t_{T_m,U}=t_{u_{m,i'}+1}^{\text{wait}}} - F_m}, \quad (22)$$

where the first term in the denominator represents the sum of data size of T_m 's subtasks when the completion time of each subtask is $t_{u_{m,i'}+1}^{\text{wait}}$ in state $s_{i'}$. The implication of Eq. (22) is that if state s_i can complete task T_m within the time of $t_{u_{m,i+1}}^{\text{wait}}$, there is no need to transform from s_i to s_{i+1} . The reason is that the time to complete T_m in s_{i+1} must be longer than $t_{u_{m,i+1}}^{\text{wait}}$. In this way, the task completion time is reduced and load balance can be ensured. In the next, we inversely calculate $d_{m,k}$ of each subtask assigned to UAVs in state s_i according to the task completion time $t_{T_m,U}$.

In status s_i , there are two possible situations for each UAV $u_{m,k}$ ($1 \leq k \leq i$) after assigning subtasks: $t_{t_{m,k},u_{m,k}}^{\text{tra}} \leq t_{u_{m,k}}^{\text{wait}}$ or $t_{t_{m,k},u_{m,k}}^{\text{tra}} > t_{u_{m,k}}^{\text{wait}}$. The boundary condition of these two situations is that $t_{t_{m,k},u_{m,k}}^{\text{tra}} = t_{u_{m,k}}^{\text{wait}}$. Then, the subtask's data size under this boundary condition is $t_{u_{m,k}}^{\text{wait}} r_{v_m u_{m,k}}$, and the total completion time required for this part of data is

$$t_{t_{m,k},u_{m,k}}^{\text{bor}} = t_{u_{m,k}}^{\text{wait}} + \frac{t_{u_{m,k}}^{\text{wait}} r_{v_m u_{m,k}} F_m}{D_m f_u}, \quad (23)$$

where the last item is the calculation time.

Therefore, if $t_{T_m,U} \leq t_{t_{m,k},u_{m,k}}^{\text{bor}}$, it will be the first case and the data size of the subtask offloaded to $u_{m,k}$ is

$$d_{m,k} = D_m \frac{(t_{T_m,U} - t_{u_{m,k}}^{\text{wait}}) f_u}{F_m}. \quad (24)$$

If $t_{T_m,U} > t_{t_{m,k},u_{m,k}}^{\text{bor}}$, it will be the second case, where $d_{m,k}$ can be divided into two parts. The transmitting time of the first part overlaps with the waiting time, so the data size of the first part is $d_{m,k}^f = t_{u_{m,k}}^{\text{wait}} r_{v_m u_{m,k}}$, and the total time required for the first part is shown in Eq. (23). Therefore, the

total time consumed by the second part is $t_{T_m,U} - t_{t_m,k,u_m,k}^{bor}$. The second part consumes not only additional transmission time, but also computation time. As a result, the data size of the second part is

$$d_{m,k}^s = \frac{t_{T_m,U} - t_{t_m,k,u_m,k}^{bor}}{\frac{1}{rv_{m,u_m,k}} + \frac{1}{D_m} \frac{F_m}{f_u}}, \quad (25)$$

where the denominator is the transmission time and computation time required for one bit of data. As a result, in the second case:

$$d_{m,k} = d_{m,k}^f + d_{m,k}^s. \quad (26)$$

After achieving $d_{m,k}$, the optimal UAVs allocation state s^* can be obtained through Q-learning. The next step is to find each subtask's data size and the match between subtasks and UAVs, i.e., the data size of each subtask assigned to each UAV in s^* . First, we set up subtasks of equal data size for each UAV, i.e., $d_{m,i} = \frac{D_m}{I^*}$, ($1 \leq i \leq I^*$). I^* is the number of UAVs in s^* . Then, we figure out the real completion time of subtasks according to Eq. (12). The shortest completion time and the longest completion time among them are denoted as t_{sub}^{\min} and t_{sub}^{\max} , respectively, i.e., $t_{sub}^{\min} = \min(t_{t_m,i,u_m,i} | 1 \leq i \leq I^*)$ and $t_{sub}^{\max} = \max(t_{t_m,i,u_m,i} | 1 \leq i \leq I^*)$.

Because $t_{T_m,U}^{\min}$ must exist between t_{sub}^{\min} and t_{sub}^{\max} and there is not much numerical difference between t_{sub}^{\min} and t_{sub}^{\max} , so we use the iterative method to find out $t_{T_m,U}^{\min}$. Obviously, after $t_{T_m,U}^{\min}$ is determined, we can figure out the data size of each subtask like $d_{m,k}$.

The concrete implementation process is shown in Algorithm 1. We first sort UAVs in ascending order of load so that low load UAVs would be preferentially selected to execute subtasks, which is conducive to load balance. In Line 2 to Line 9, Q-learning is used to find the optimal UAVs allocation state s^* . Afterwards, the minimum task completion time is obtained by iteration in Line 10 to Line 18. Since the time required for each UAV to complete its respective subtask is equal, the load balance among UAVs is further satisfied.

B. Maximize Vehicle's Utility When Offloading to ECD

The subproblem of maximizing the vehicle's utility when the computing task offloaded to ECD can be formulated as

$$P_3 : \max_{f_m^s} s_m^E \quad (27)$$

$$s.t. \quad t_{T_m,e_s} \leq \tau_m$$

According to Eq. (18), the utility of v_m is determined by f_m^s . We first calculate the first-order partial derivative of s_m^E with respect to f_m^s :

$$\frac{\partial s_m^E}{\partial f_m^s} = \frac{\alpha F_m}{(f_m^s)^2} e^{\frac{F_m}{f_m^s} + t_{T_m,e_s}^{tra} - \tau_m} - 2\gamma_E \delta F_m f_m^s, \quad (28)$$

Then, we calculate the second-order partial derivative of s_m^E with respect to f_m^s :

$$\frac{\partial^2 s_m^E}{\partial (f_m^s)^2} = -(2 + \frac{F_m}{f_m^s}) \frac{\alpha F_m}{(f_m^s)^3} e^{\frac{F_m}{f_m^s} + t_{T_m,e_s}^{tra} - \tau_m} - 2\gamma_E \delta F_m. \quad (29)$$

Algorithm 2 : Gradient Descent Algorithm

- 1: **Initialize** $f_m^{e_s(0)}$, $k = 0$, *step*;
- 2: $s(f_m^{e_s}) = -s_m^E$;
- 3: **while** $s(f_m^{e_s(k)}) - s(f_m^{e_s(k+1)}) > \varepsilon$ **do**
- 4: *step* = $\frac{(\nabla s(f_m^{e_s(k)}))^T \nabla s(f_m^{e_s(k)})}{\nabla s(f_m^{e_s(k)})^T \nabla^2 s(f_m^{e_s(k)}) \nabla s(f_m^{e_s(k)})}$;
- 5: $f_m^{e_s(k+1)} = f_m^{e_s(k)} - \textit{step} \times \nabla s(f_m^{e_s(k)})$;
- 6: **end while**
- 7: Calculate s_m^E according to Eq. (18);

Due to the negativity of the second derivative, s_m^E is concave, i.e., there exists the maximum utility for the vehicle. We then conduct the following operations:

$$\lim_{f_m^s \rightarrow 0^+} \frac{\partial s_m^E}{\partial f_m^s} = +\infty, \quad \lim_{f_m^s \rightarrow +\infty} \frac{\partial s_m^E}{\partial f_m^s} = -\infty. \quad (30)$$

Therefore, let Eq. (28) be equal to zero, we can get the optimal computing resource $f_m^{e_s^*}$. Since it can not be solved directly, we use the gradient descent algorithm to approximate $f_m^{e_s^*}$ as shown is Algorithm 2. In Line 2, we first find the negative of s_m^E because the maximum value of s_m^E is the minimum value of $-s_m^E$. The gradient descent algorithm is then used to find the optimal computing resource from Line 3 to Line 6, where $\nabla s(f_m^{e_s(k)}) = -\frac{\partial s_m^E}{\partial f_m^s}$ and $\nabla^2 s(f_m^{e_s(k)}) = -\frac{\partial^2 s_m^E}{\partial (f_m^s)^2}$.

C. Maximize Vehicle's Utility When Computing Locally

Since s_m^L is a monotonic function of f_m^L , the greater f_m^L , the greater s_m^L . s_m^L is a fixed value when f_m^L is determined, so we do not need to consider this subproblem.

VI. PERFORMANCE EVALUATIONS

A. Simulation Setup

We consider a 500m×500m simulation area, in which there are 9 uniformly distributed UAVs and 4 uniformly distributed ECDs. Vehicles' coordinates are randomly selected from the area and each vehicle has a computing task to execute. The computing capacity of the vehicle is $f_m^L = 1\text{GHZ}$, the computing resource that each UAV can provide is $f_u = 4\text{GHZ}$. The CPU cycles required for one bit is $c_m = 1000$ [8] and channel power gain $\beta_0 = \beta_1 = -50\text{dB}$ at reference distance $d_0 = 1\text{m}$. For the computation tasks, the data size ranges from 100KB to 500KB [15] and the flying height of UAV is 20m [17]. Besides, $\delta = 1 \times 10^{-28}$ [18].

Since [6]–[8] offload the computing task to a single UAV, we compare our proposal with the scheme that offloads tasks to a greedily selected UAV (OTU). In OTU, vehicles' computing tasks are offloaded to a UAV selected from a greedy algorithm to ensure the vehicle obtains more utility. Another scheme is computing the task locally (CTL) that vehicles' computing tasks are executed locally.

B. Performance Evaluation

Fig. 2 shows the vehicle's utility obtains from three schemes when the data size of the computing task changes. For our

proposed scheme, subtasks are executed in parallel and the optimal computing resources are used. Consequently, the utility obtained from our proposed scheme is the largest. When the data size of the computing task becomes larger, the limited computing resources of a single UAV cannot complete the computing task in a short time. Therefore, the utility obtained from OTU is becoming less and less.

Fig. 3 shows the vehicle's utility obtained from offloading computing tasks to ECD when the ECD allocates different computing resources. In Fig. 3, there are three schemes to determine $f_m^{e_s}$ that e_s allocates to the offloaded computing task: optimal computing resources $f_m^{e_s*}$ determined by our scheme, $f_m^{e_s} = 4\text{GHz}$ and $f_m^{e_s} = 5\text{GHz}$. As seen from Fig. 3, our scheme is always the most effective. The reason is that the vehicle's utility when offloading to the ECD is mainly determined by the computing resources allocated to the computing task.

VII. CONCLUSION

In this paper, we have proposed an edge computing and UAV swarm assisted cooperative task offloading scheme in vehicular networks to guarantee vehicles' real-time demand in peak traffic. Specifically, a cooperative task offloading framework is first designed in which vehicles' computing tasks can be executed locally, offloaded to UAV swarm, or offloaded to ECDs. A mixed integer nonlinear programming problem is then formulated to maximize the utility of the vehicle by finding the optimal offloading strategy. Furthermore, the original problem is decomposed into two subproblems: minimizing the completion time when the computing task is offloaded to UAV swarm and optimizing the computing resources when offloaded to ECD. When offloading to UAV swarm, a Q-learning based iterative algorithm that can realize the load balance among UAVs is designed to minimize the completion time. When offloading to ECDs, a gradient descent algorithm is utilized to find the optimal computing resources allocated for computing tasks. At last, numerical simulations have demonstrated the effectiveness and superiority of our proposed scheme in comparison with conventional schemes. In future work, the communication security between vehicles and UAVs will be studied.

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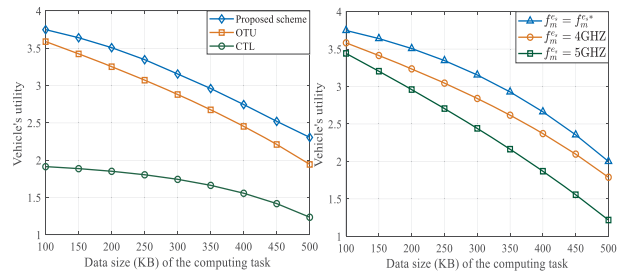


Fig. 2. Utility of the vehicle under different schemes.

Fig. 3. Utility of the vehicle under different $f_m^{e_s}$.

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