WTF are CRDTs?

Adventures in distributed systems theory

Ben Tyler - YAPC::EU 2015

Follow along!

https://github.com/kanatohodets/crdt-talk

Disclaimer!

Zero practical experience

It looked cool

So I decided to try to talk about it

Sorry, no code slides

Overview

- Why this talk?
- AP systems
- The trouble with ^ (hint: missing 'C')
- Dealing with data conflicts
- Conflict free/convergent data types
- Data types: G-Counter, OR-Set
- 555

Why?

Peer pressure learning group.

Scaling reads is "easy."

High availability is "easy."

Horizontally scaling writes in a useful way?
 Really hard.

Scaling writes means...

 Multiple "masters" – many nodes that can accept writes (see: sharding).

 No 'stop the world' synchronization (that is, asynchronous replication).

Say goodbye to SQL transactions

Scaling writes means...

(Distributed systems theory edition)

Masterless

Eventually consistent

AP-ish

Aside: WAT? CAP?

Consistency, availability, partition tolerance....

- Works like a single copy (not ACID!)
- Available for updates
- Readable under netsplit

AP ...ish?

- Consistency, availability, partition tolerance...
- Pick 2
- But more like 0.7/0.8/0.5
- Riak, Cassandra: "NoSQL" distributed key-value stores.

Cassandra



What flavor of distsys?

AP







AP: The Good

A(vailability)

P(artition tolerance)

(potentially) Linear write scalability

AP: The Bad

 Eventually consistent (I just PUT it there, where'd it go?)

Network overhead (maybe)

AP: The Wookiee

Data conflicts



Resolving data conflicts

Last write wins (bet the farm on ntpd)

 Causality tracking (logical clocks, version vectors/vector clocks, interval tree clocks)

Give up and ask the client (siblings)

Last Write Wins



Time, LWW, and the farm

- "All objects with a lower timestamp will be silently deleted until GC removes the tombstone record—which means that a rogue client or node can cause the destruction of every write to a record for days to weeks afterwards"
 - https://aphyr.com/posts/299-the-trouble-withtimestamps

Causality tracking?

 Attempt to establish 'causality' – who has the more up to date data?

 Can't resolve every conflict (create a 'total ordering' – only partial).

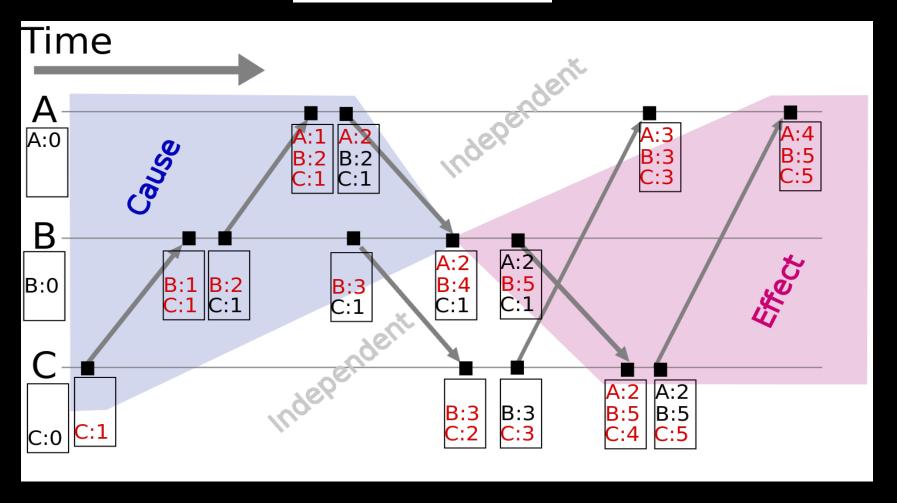
Causality tracking?

Vector clocks

- Every node has a counter. It increments whenever something happens.
- "Something" == receive message or "internal event"
- "Internal event" == client PUT something
- Any message includes all the clocks a node knows about

Causality tracking?

Vector clocks



Give up and ask the client

Merge, then tell database the new state

- "In all such systems, we find developers spend a significant fraction of their time building extremely complex and error-prone mechanisms to cope with eventual consistency and handle data that may be out of date."
 - The GOOG (F1 paper)



Can't someone do that for me?

Merging data is hard, let's not do it

Enter CRDTs

Commutative/convergent/conflict-free Replicated

Data

Types

Commutative/convergent/conflict-free Replicated

Data

Types

Commutative/convergent/conflict-free

Replicated

Data

Types

Commutative/convergent/conflict-free

Replicated

Data

Types

operation based: SBR

Commutative/convergent/conflict-free Replicated

Data

Types

state based: RBR

(with a monotonic/idempotent merge)

Aside: WAT

Monotonic:

"only moves in one direction" e.g. entirely increasing, or entirely decreasing

Idempotent

"just run it again, don't worry" e.g. add an item to a set

Commutative/convergent/conflict-free Replicated

Data

Types

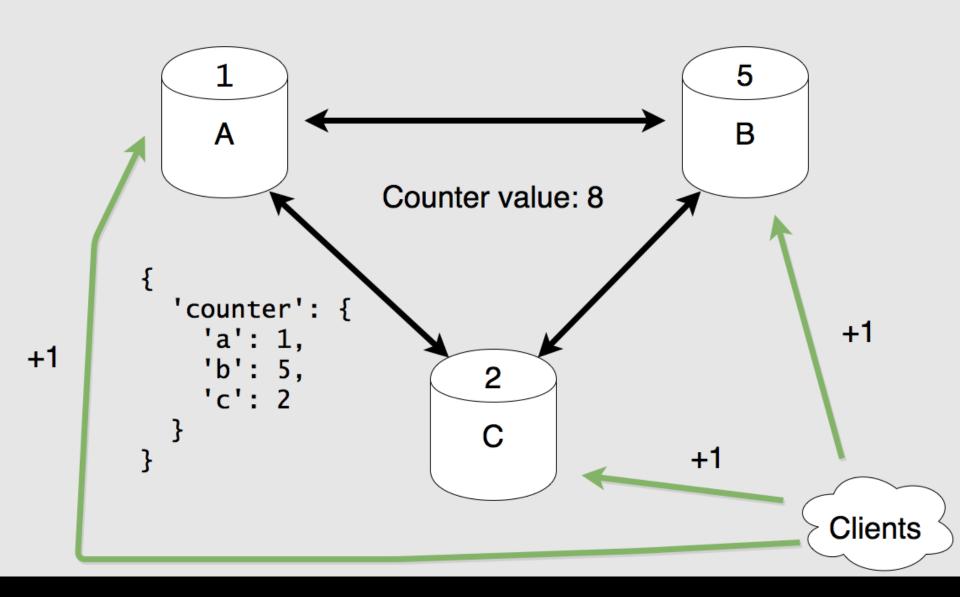
the whole gang (they're theoretically equivalent)

Like what?

Grow-only counter

- Client says:
 - 'yo, increment!'
- Track how many increments/replica: total is sum of all replica counts
- Trivially commutative operation
- ...not all that interesting
- (or I lack imagination)
- (I totally lack imagination: Youtube view count)

G-Counter cluster



What about negatives?

Two grow-only counters: increments and decrements.

Value is increments – decrements.

 If your building blocks are all CRDTs, the result is also a CRDT.

Non-negative counter?

E.g. gold in the bank in a game

...global invariant! Nooooo!

Requires synchronization

Okay, show me something I can use...

Set of insertions, each with unique tag

Set of deletions, each with unique tag

On conflict, add > delete

Merge: union of insertion, union of tombstones

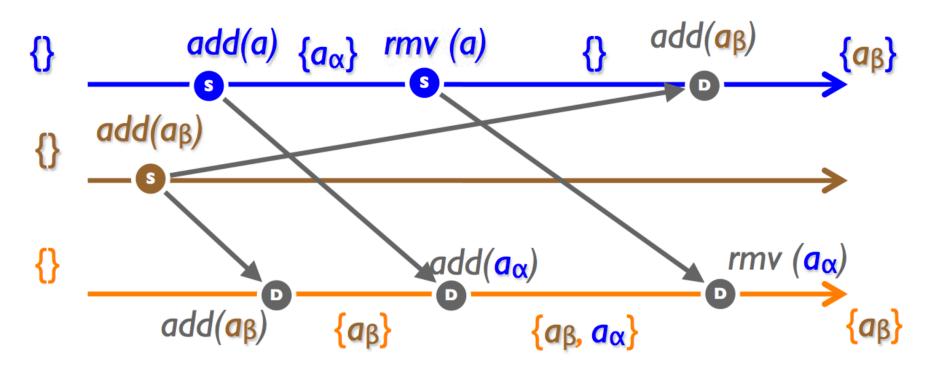
```
'type': 'or-set',
'е': Г
  ['foo', [1]],
  ['bar', [1], [1]],
  ['baz', [1, 2], [2, 3]]
```

```
'type': 'or-set',
'е': Г
 ['foo', [1]], // foo exists
  ['bar', [1], [1]],
  ['baz', [1, 2], [2, 3]]
```

```
'type': 'or-set',
'е': Г
 ['foo', [1]],
  ['bar', [1], [1]], // no bar
  ['baz', [1, 2], [2, 3]]
```

```
'type': 'or-set',
'е': Г
  ['foo', [1]],
  ['bar', [1], [1]],
  ['baz', [1, 2], [2, 3]]
] //baz exists
```





Caveats

Still eventual consistency (but stronger)

Not all client operations will be respected;
 applications need to be aware.

Garbage collection can be tricky (potential for unbounded garbage growth)

In the real world...

Riak data types (Counters, Flags, Sets, Maps, etc.)

- Soundcloud stream: LWW element set
 - https://github.com/soundcloud/roshi

- League of Legends chat: friend list
 - Friends in general
 - Online/offline friends

Bonus! DAG

Directed Acyclic Graph (e.g. tree)

Global invariant??

 Locally enforced: only add edges in existing directions

Sources & Papers

- Marc Shapiro, Nuno Preguica, Carlos Baquero, Marek Zawirski. A comprehensive study of Convergent and Commutative Replicated Data Types. [Research Report] RR-7506, Inria — Centre Paris-Rocquencourt. 2011, pp.50.
 - https://hal.inria.fr/inria-00555588
- Lots of papers
 - http://christophermeiklejohn.com/crdt/2014/07/22/readings-incrdts.html
- http://research.microsoft.com/apps/video/dl.aspx?id=153540
- https://github.com/aphyr/meangirls
- https://aphyr.com/tags/Distributed-Systems
- http://blog.acolyer.org/
- Riak docs:
 - http://docs.basho.com/riak/latest/theory/concepts/context/
 - http://docs.basho.com/riak/latest/theory/concepts/crdts/

The End