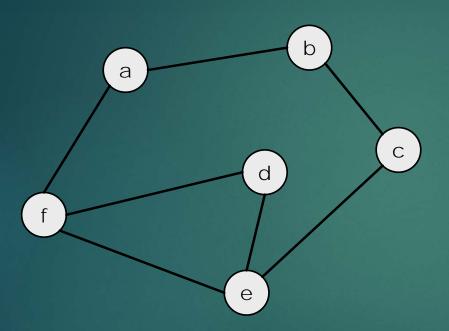
# HAMILTONIAN PATH

**BACKTRACKING** 

# Hamiltonian cycle

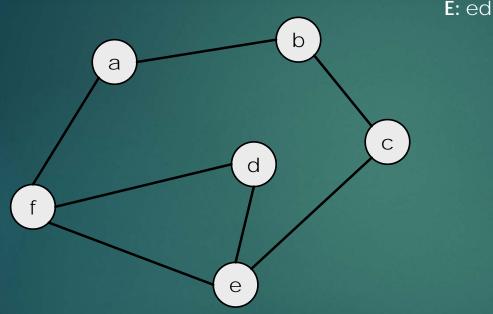
G(V,E)

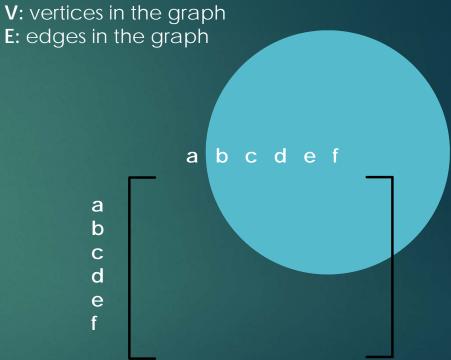


V: vertices in the graph E: edges in the graph

# Hamiltonian cycle

G(V,E)

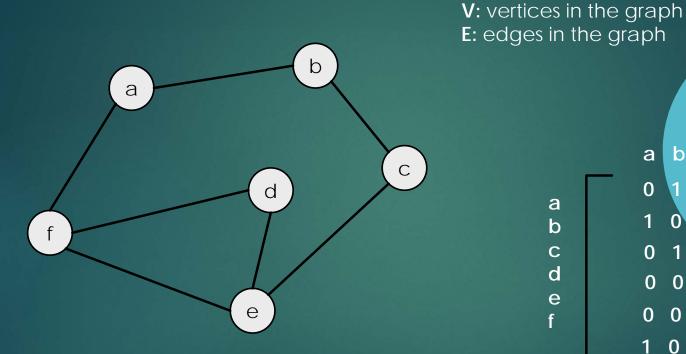




 $A(i,j) = \{ 1 - if \text{ there is a connection between } i \text{ and } j ; 0 - if \text{ no connection } \}$ 

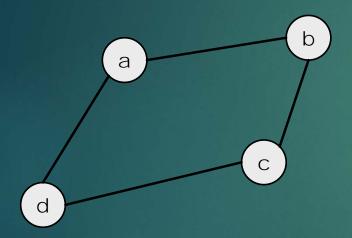
# Hamiltonian cycle

G(V,E)

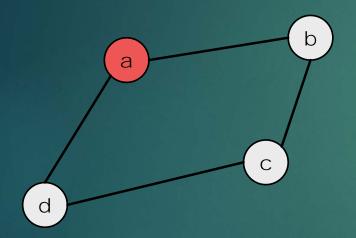




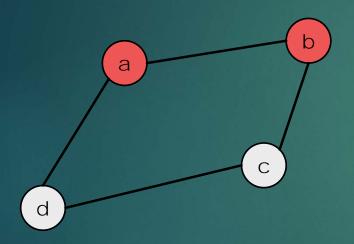
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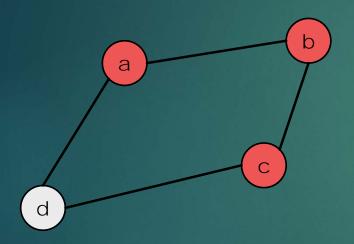
A <u>hamiltonian path</u> in an undirected graph is a path that visits every node exactly once !!!



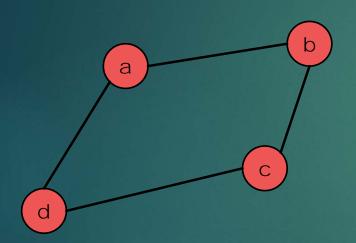
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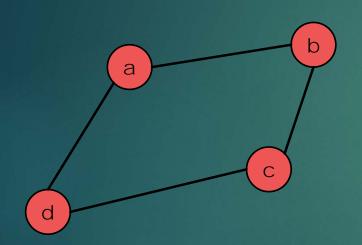
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<u>Hamiltonian cycle</u>: the first node and the last node of the path are the same vertexes

StartingPoint == EndPoint

A valid hamiltonian path is: { a b c d a }

There may be several hamiltonian path in a given graph !!!

# Hamiltonian problem

- Determining whether such paths and cycles exist in graphs is the Hamiltonian path problem
- ▶ This is an NP-complete complete problem !!!
- <u>Dirac-principle</u>: a simple graph with N vertices is hamiltonian if every vertex has degree N/2 or greater (degree is the number of edges of a vertex)
- Important fact: finding Hamiltonian path is NP-complete, but we can decide whether such path exists in linear time complexity with topological ordering

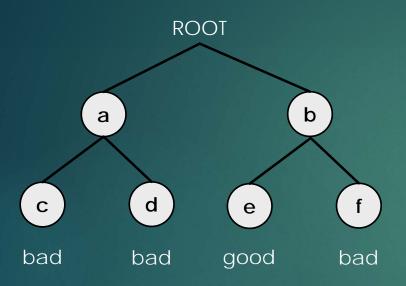
# **Solutions**

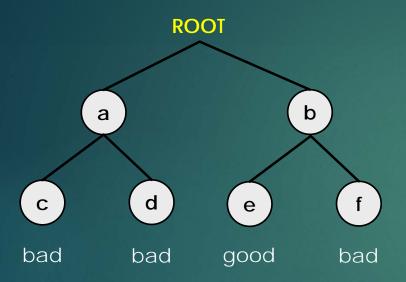
- ▶ Naive approach:
- Generate all possible configurations of the vertices and print a configuration that satisfies the given constraints
- ▶ Problem → if the graph has N vertices, there are N! configurations, so the "solution space" is enormous
- ▶ Very very inefficient !!!

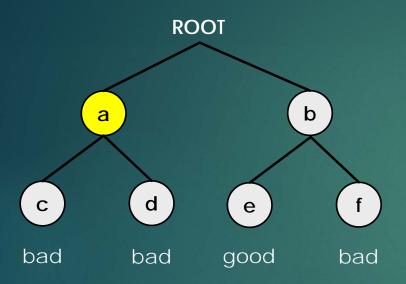
# **Solutions**

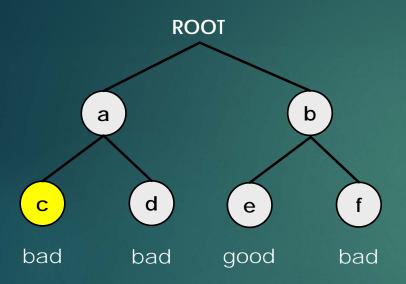
► Constructing a tree:

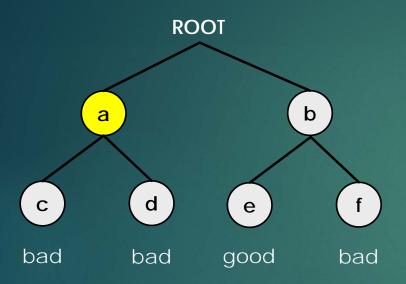


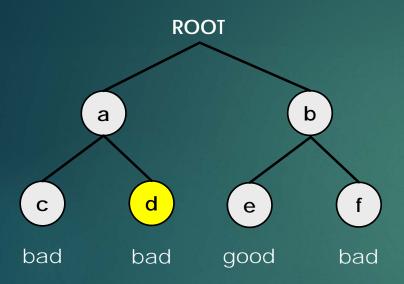


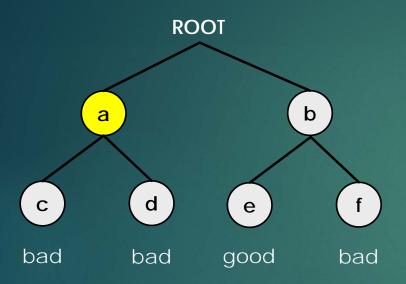


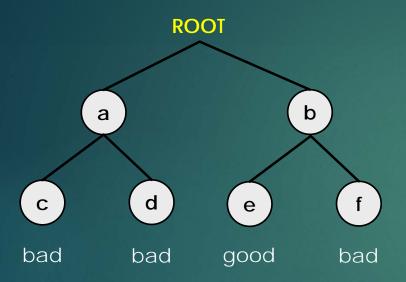


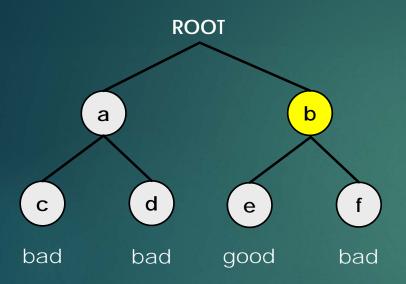


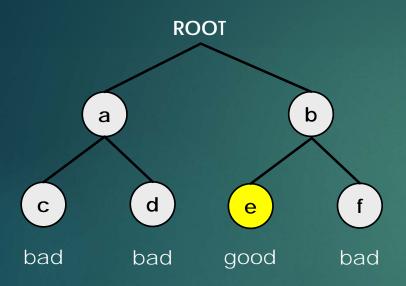


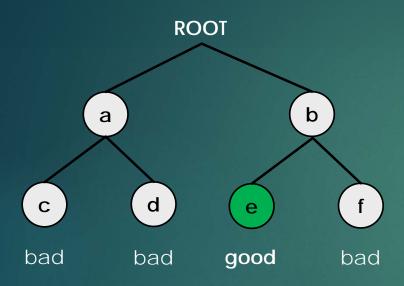












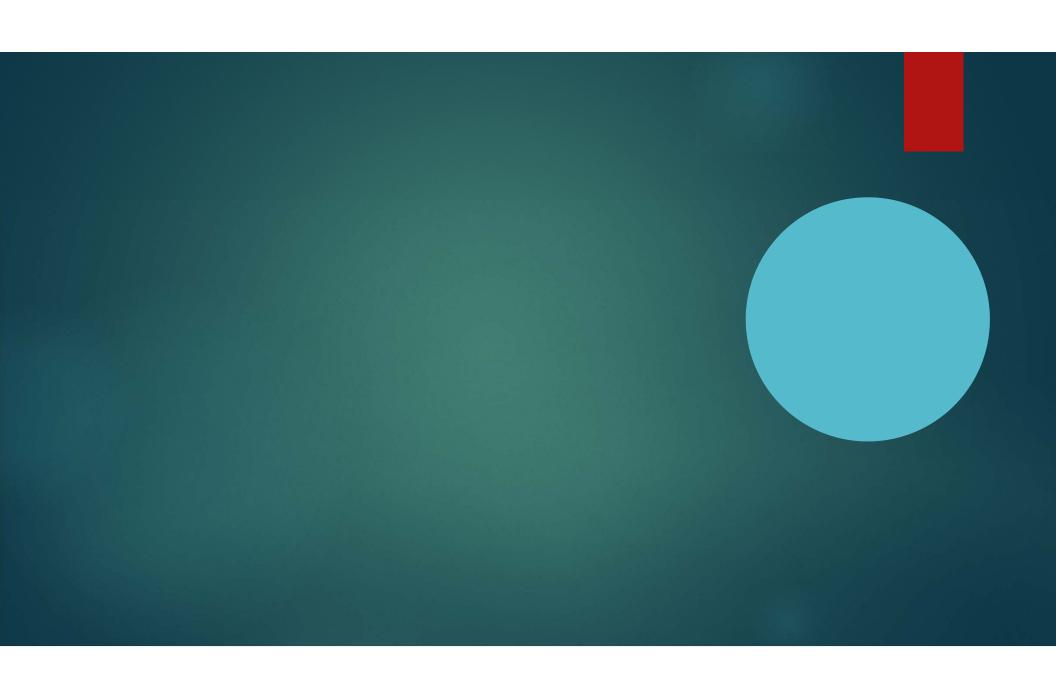
If we get to a bad leaf: we just "backtrack" and keep moving on by revoking our most recent choice

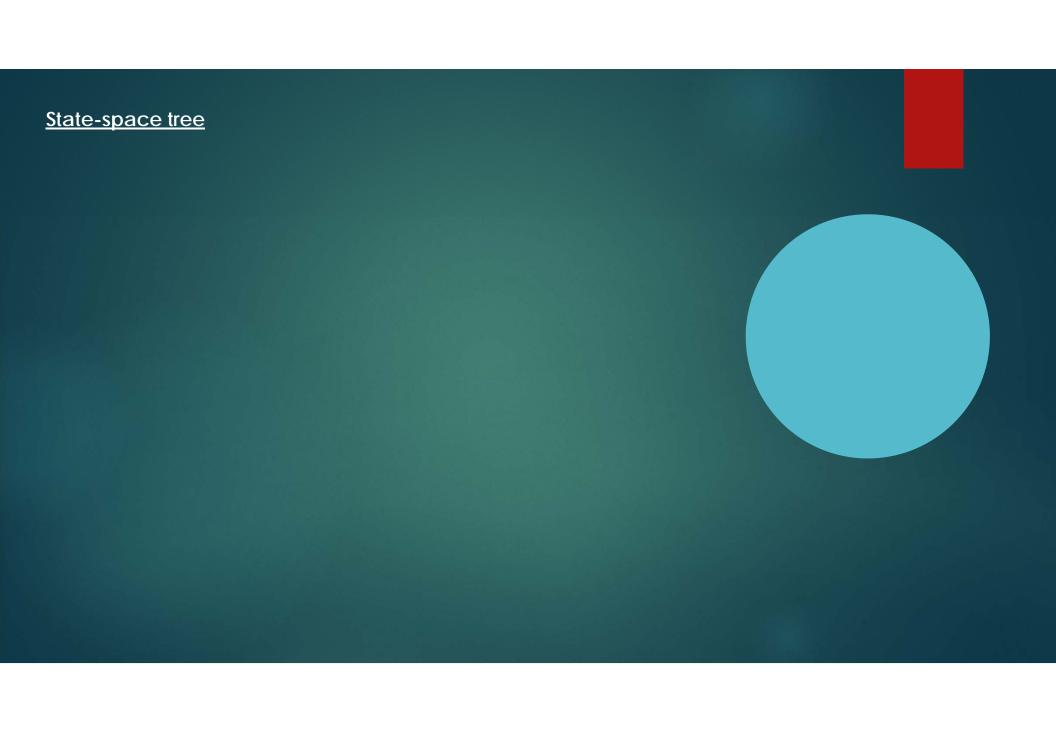
The tree is an abstract model of the possible sequences of choices we could make Here we do a depth-first search on the tree

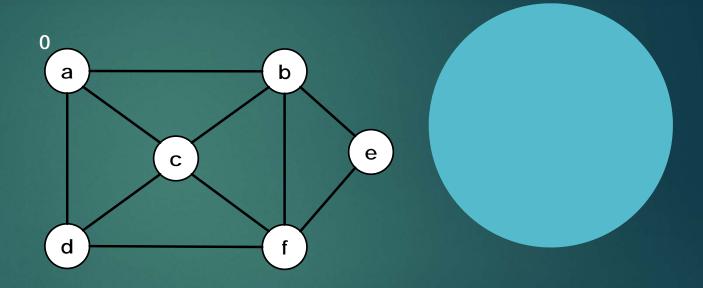
Problem: hard to construct a tree if **N** is big!!!

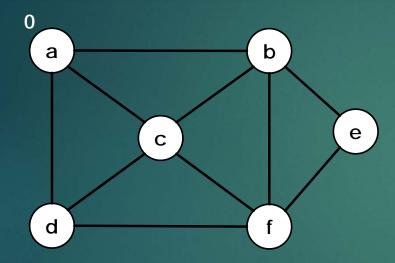
## **Solutions**

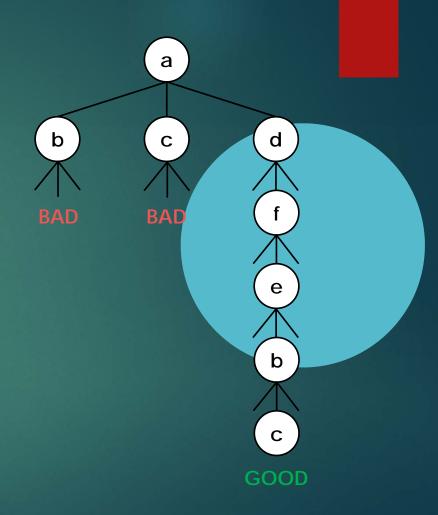
- ▶ Backtracking:
- We use recursion to solve the problem
- Create an empty path array and add vertex 0 to it as the starting vertex
- Add other vertices, starting from the vertex 1
- ▶ Before adding a vertex, check whether it is adjacent to the previously added vertex + make sure it is not already added
- ▶ If we find such a vertex → we add the vertex as part of the solution
- ▶ If we do not find a vertex → we return false

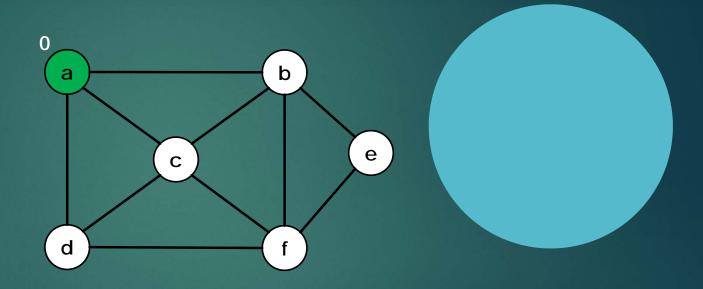


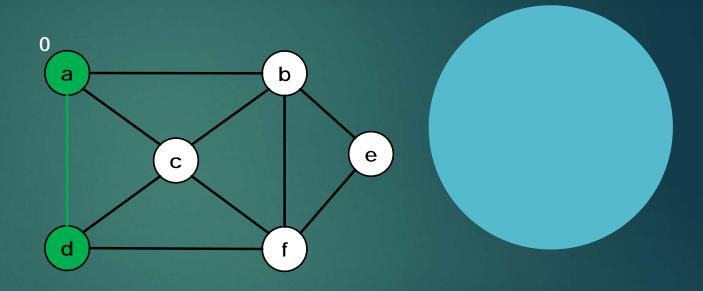


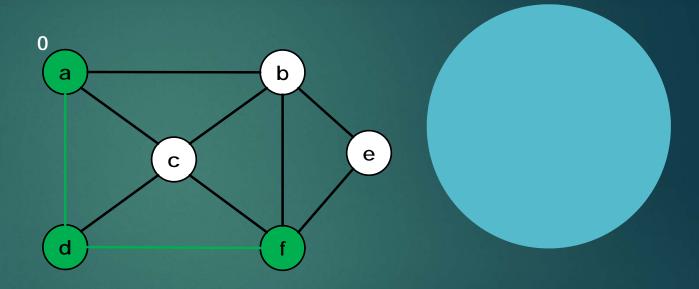


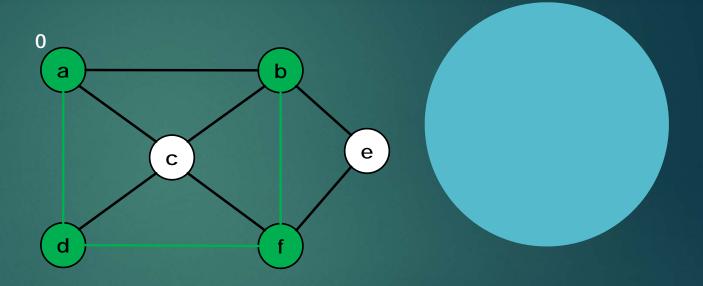


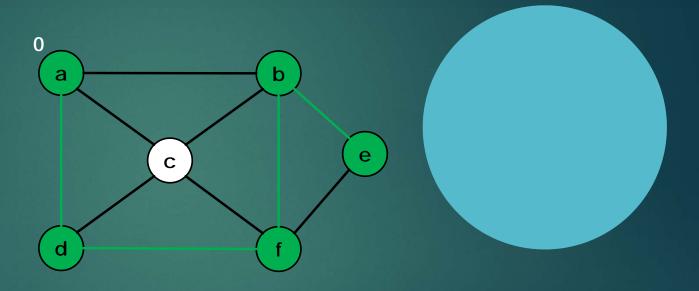


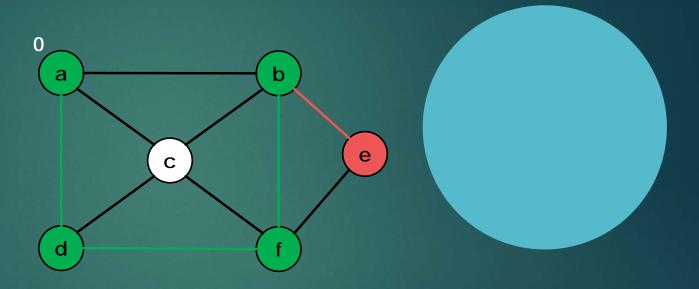


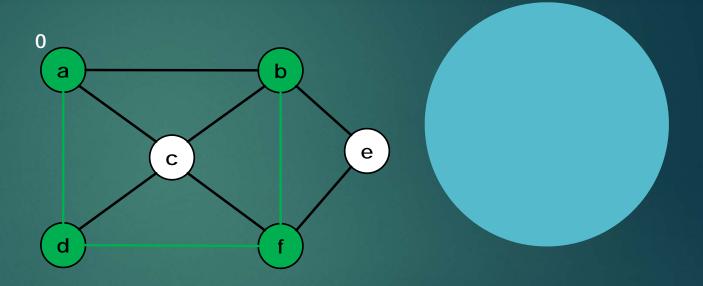


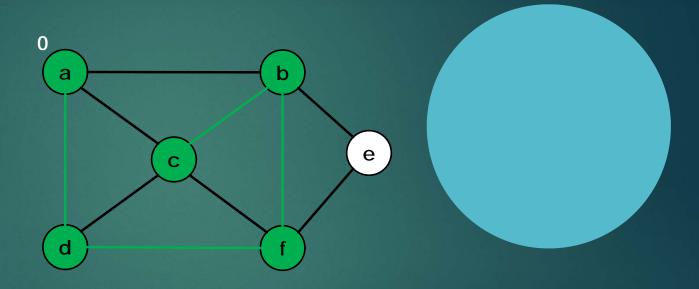


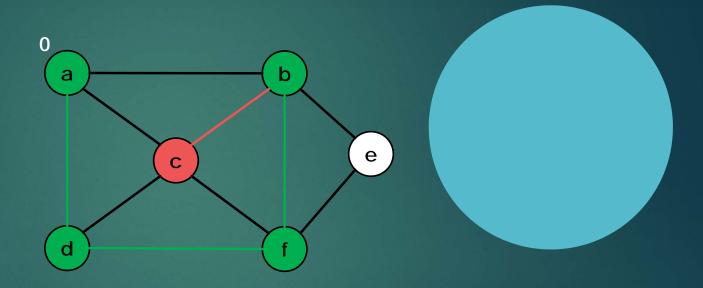


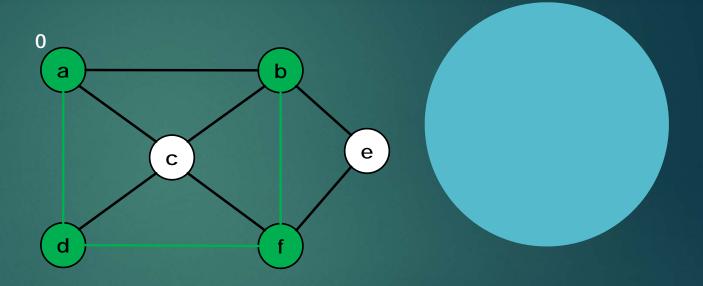


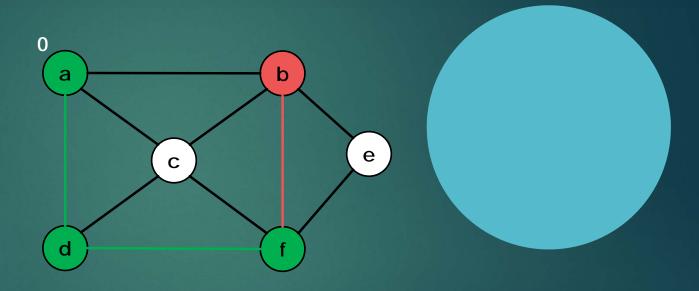


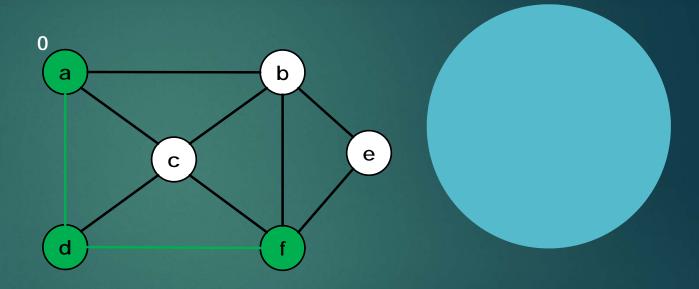


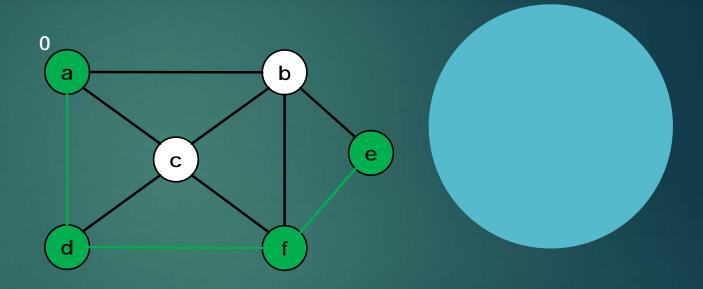


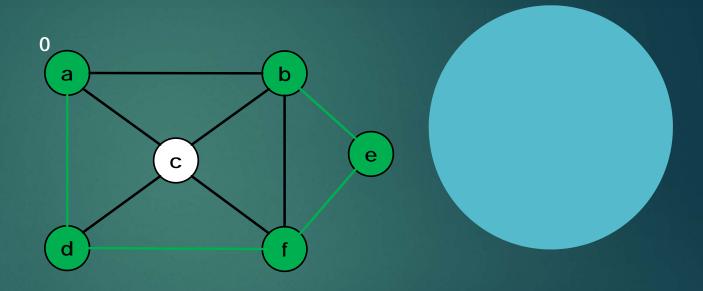


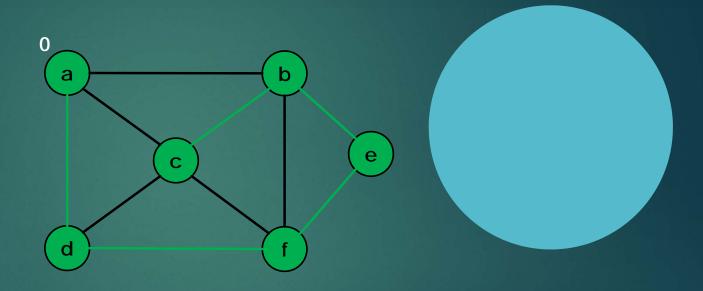


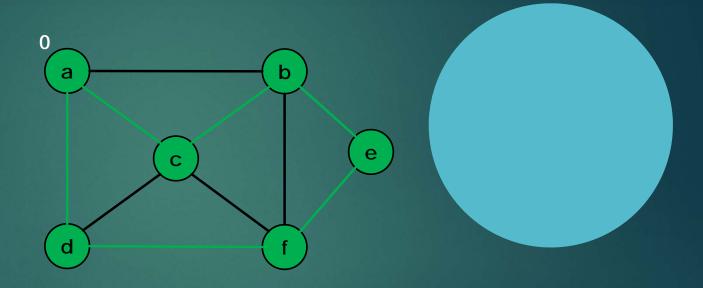


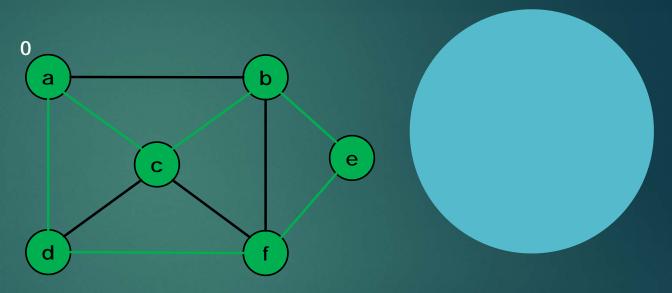












We have found the Hamiltonian-cycle in this graph

{a,d,f,e,b,c,a}