ROD CUTTING PROBLEM

DYNAMIC PROGRAMMING



Rod cutting problem

- Given a rod with certain length I
- Given the prices of different lengths
- ► How to cut the rod in order to **maximize** the profit?
- ► This is the rod cutting problem



Rod length \rightarrow I = 5m

Prices for different lengths:

1m → \$2

2m → \$5

 $3m \rightarrow 7

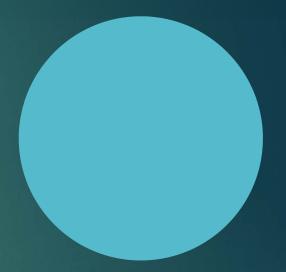
4m → \$3

Solution to the rod cutting problem: {2,3} so a cut the rod to get a 2m piece and a 3m piece

OR

{2,2,1} 2 2m pieces and a single 1m piece, it is going to be the same \$12 profit

Total value for both solutions: \$12



Recursion

- ► The naive approach is to use a simple recursive method / function
- ▶ N-1 cuts can be made in the rod of length N
- There are 2^{N-1} ways to cut the rod
- Problems: time complexity + overlapping subproblems
- Exponential time complexity: O(2^N) where N is the length of the rod in units
- ▶ (for every length we have 2 options whether to cut or not)

Dynamic programming

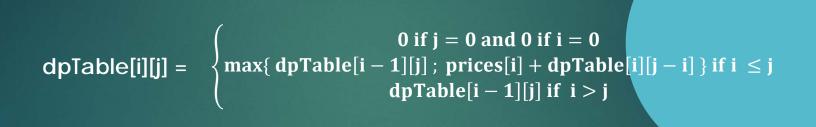
We have to create a solution matrix:

dpTable[numOfLengths+1][originalLength+1]
rows columns

We have to define the <u>base cases</u>:

- if originalLength is $0 \rightarrow 0$ is the profit
- if we do not consider any lengths \rightarrow 0 is the profit

Complexity: O(numOfLengths*originalLength)

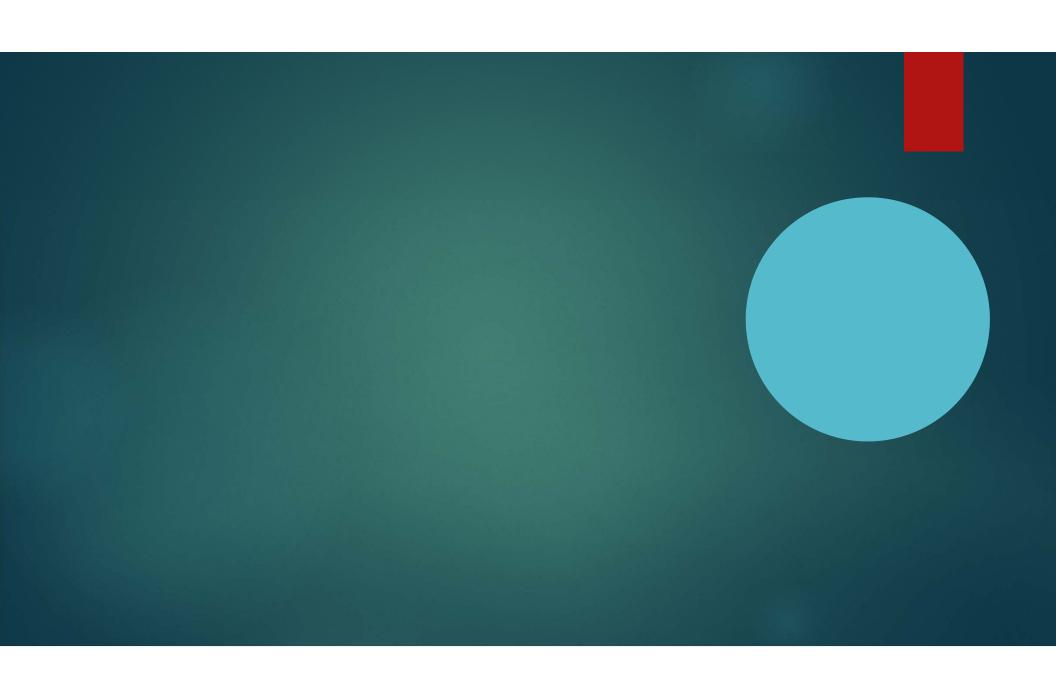


These are the base cases

$$\label{eq:dpTable} \begin{split} \text{dpTable[i][j] = } & \begin{cases} 0 \text{ if } j = 0 \text{ and } 0 \text{ if } i = 0 \\ \max\{\text{dpTable}[i-1][j] \text{ ; prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j \\ \text{dpTable}[i-1][j] \text{ if } i > j \end{cases} \end{split}$$

The total value when total length is **j** and we have the first **i** pieces

If the piece is greater than the length of the rod → we skip it



EXAMPLE



I = 5m 0m → \$0 1m → \$2 2m → \$5 3m → 7\$ 4m → 3\$ numOfRows = prices.length+1 prices[] = $\{0, 2, 5, 7, 3\}$

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]	length
\$0 – 0m							
\$2 – 1m							
\$5 – 2m							
\$7 – 3m							
3\$ - 4m							
pieces							

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]	length
No cuts	\$0 – 0m							
Just one	\$2 – 1m							
First two	\$5 – 2m							
First three	\$7 – 3m							
All	3\$ - 4m							
	pieces							

<u>Subproblems</u>: we consider the lengths {0,1,2,3,4,5} step by step when we can have {1,2,3,4} unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0					
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces						

Subproblems: we consider the lengths {0,1,2,3,4,5} step by step when we can have {1,2,3,4} unit lengths at the same time!!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0	2				
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces						

Subproblems: we consider the lengths {0,1,2,3,4,5} step by step when we can have {1,2,3,4} unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0	2	4			
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces						

Subproblems: we consider the lengths {0,1,2,3,4,5} step by step when we can have {1,2,3,4} unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0	2	4	6		
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces						

Subproblems: we consider the lengths {0,1,2,3,4,5} step by step when we can have {1,2,3,4} unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0	2	4	6	8	
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces	1000					

<u>Subproblems</u>: we consider the lengths $\{0,1,2,3,4,5\}$ step by step when we can have $\{1,2,3,4\}$ unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m

$$0m \rightarrow \$0 \ 1m \rightarrow \$2 \ 2m \rightarrow \$5 \ 3m \rightarrow 7\$ \ 4m \rightarrow 3\$$$

prices[] = $\{0, 2, 5, 7, 3\}$
numOfColumns = I +1
numOfRows = prices.length+1

		0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
No cuts	\$0 – 0m	0	0	0	0	0	0
Just one	\$2 – 1m	0	2	4	6	8	10
First two	\$5 – 2m	0					
First three	\$7 – 3m	0					
All	3\$ - 4m	0					
	pieces						

<u>Subproblems</u>: we consider the lengths $\{0,1,2,3,4,5\}$ step by step when we can have $\{1,2,3,4\}$ unit lengths at the same time !!! We solve the subproblems and combine them for the final solution

I = 5m $0m \rightarrow $0.1m \rightarrow $2.2m \rightarrow $5.3m \rightarrow 7$.4m \rightarrow 3$. num()$ $prices[] = <math>\{0, 2, 5, 7, 3\}$

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2				
\$7 – 3m	0					
3\$ - 4m	0					

 $\label{eq:continuous} \mbox{dpTable[i][j] = } \begin{cases} &0 \mbox{ if } j = 0\\ &\max\{\mbox{dpTable}[i-1][j] \mbox{; prices}[i] + \mbox{dpTable}[i][j-i]\} \mbox{ if } i \leq j\\ &\mbox{dpTable}[i-1][j] \mbox{ if } i > j \end{cases}$

dpTable[2][1] = dpTable[1][1]

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5			
\$7 – 3m	0					
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[2][2] = MAX { dpTable[1][2] ; 5 + dpTable[2][0] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7		
\$7 – 3m	0					
3\$ - 4m	0					

 $\label{eq:continuous} \begin{aligned} \text{pieces} \\ \text{dpTable[i][j] = } & \begin{cases} 0 \text{ if } j = 0 \\ \max\{\text{dpTable}[i-1][j] \text{ ; } \text{prices}[i] + \text{dpTable}[i][j-i]\} \text{ if } i \leq j \\ \text{dpTable}[i-1][j] \text{ if } i > j \end{cases} \end{aligned}$

dpTable[2][3] = MAX { dpTable[1][3] ; 5 + dpTable[2][1] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	
\$7 – 3m	0					
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[2][4] = MAX { dpTable[1][4] ; 5 + dpTable[2][2] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0					
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[2][5] = MAX { dpTable[1][5] ; 5 + dpTable[2][3] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2				
3\$ - 4m	0					

 $\label{eq:continuous} \begin{aligned} \text{pieces} \\ \text{dpTable[i][j] =} & \begin{cases} 0 \text{ if } j = 0 \\ \max\{\text{dpTable}[i-1][j] \text{; prices}[i] + \text{dpTable}[i][j-i]\} \text{ if } i \leq j \\ \text{dpTable}[i-1][j] \text{ if } i > j \end{cases} \end{aligned}$

dpTable[3][1] = dpTable[2][1]

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5			
3\$ - 4m	0					

 $\label{eq:continuous} \mbox{dpTable[i][j] = } \begin{cases} &0 \mbox{ if } j = 0 \\ &\max\{\mbox{dpTable}[i-1][j] \mbox{; prices}[i] + \mbox{dpTable}[i][j-i]\} \mbox{ if } i \leq j \\ &\mbox{dpTable}[i-1][j] \mbox{ if } i > j \end{cases}$

dpTable[3][2] = dpTable[2][2]

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7		
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[3][3] = MAX { dpTable[2][3] ; 7 + dpTable[3][0] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[3][4] = MAX { dpTable[2][4] ; 7 + dpTable[3][1] }

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0					

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[3][5] = MAX { dpTable[2][5] ; 7 + dpTable[3][2] }

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2				

 $\label{eq:continuous} \mbox{dpTable[i][j] = } \begin{cases} &0 \mbox{ if } j = 0 \\ &\max\{\mbox{dpTable}[i-1][j] \mbox{; prices}[i] + \mbox{dpTable}[i][j-i]\} \mbox{ if } i \leq j \\ &\mbox{dpTable}[i-1][j] \mbox{ if } i > j \end{cases}$

dpTable[4][1] = dpTable[3][1]

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5			

 $\label{eq:continuous} \mbox{dpTable[i][j] = } \begin{cases} &0 \mbox{ if } j = 0 \\ &\max\{\mbox{dpTable}[i-1][j] \mbox{; prices}[i] + \mbox{dpTable}[i][j-i]\} \mbox{ if } i \leq j \\ &\mbox{dpTable}[i-1][j] \mbox{ if } i > j \end{cases}$

dpTable[4][2] = dpTable[3][2]

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7		

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[4][3] = dpTable[3][3]

numOfColumns = I +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	

 $\label{eq:pieces} \text{dpTable[i][j] = } \begin{cases} &0 \text{ if } j = 0\\ &\text{max} \{ \text{dpTable}[i-1][j] \text{; } \text{prices}[i] + \text{dpTable}[i][j-i] \} \text{ if } i \leq j\\ &\text{dpTable}[i-1][j] \text{ if } i > j \end{cases}$

dpTable[4][4] = MAX { dpTable[3][4] ; 3 + dpTable[4][0] }

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

 $\label{eq:continuous} \begin{aligned} \text{pieces} \\ \text{dpTable[i][j] =} & \begin{cases} 0 \text{ if } j = 0 \\ \max\{\text{dpTable}[i-1][j] \text{; prices}[i] + \text{dpTable}[i][j-i]\} \text{ if } i \leq j \\ \text{dpTable}[i-1][j] \text{ if } i > j \end{cases} \end{aligned}$

dpTable[4][5] = MAX { dpTable[3][5] ; 3 + dpTable[4][1] }

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION:

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12 🔨

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION:

It is coming from the cell above: it means there is no 4m cut in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12 🔭
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts? It is coming from

SOLUTION:

It is coming from the cell above: it means there is no 3m cut in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12 🛌
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts? It is NOT coming from

SOLUTION: 2m,

It is NOT coming from the cell above: it means there is a 2m cut in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts? It is NOT coming from

SOLUTION: 2m, 2m,

It is NOT coming from the cell above: it means there is a 2m cut again in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2 🔨	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts? It is coming from

SOLUTION: 2m, 2m,

It is coming from the cell above: it means there is no more 2m cut in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2 🔨	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION: 2m, 2m, 1m

It is NOT coming from the cell above: it means there is a 1m cut in the solution

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2 🔨	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION: 2m, 2m, 1m

It is NOT coming from the cell above: it means there is a 1m cut in the solution

length

We have to subtract the given price (\$2) corresponding to the length (1m) from \$2 ... go to that position \$2 - \$2 = \$0 → so we go to dpTable[1][0]!!!

numOfColumns = | +1
numOfRows = prices.length+1

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2 🔨	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION: 2m, 2m, 1m

It is NOT coming from the cell above: it means there is a 1m cut in the solution

length

We have to subtract the given price (\$2) corresponding to the length (1m) from \$2 ... go to that position \$2 - \$2 = \$0 → so we go to dpTable[1][0]!!!

numOfColumns = | +1
numOfRows = prices.length+1

length

No cuts
Just one
First two
First three
All

	0 [m]	1 [m]	2 [m]	3 [m]	4 [m]	5 [m]
\$0 – 0m	0	0	0	0	0	0
\$2 – 1m	0	2	4	6	8	10
\$5 – 2m	0	2	5	7	10	12
\$7 – 3m	0	2	5	7	10	12
3\$ - 4m	0	2	5	7	10	12

pieces

OK, we can make a \$12 profit ... but what are the optimal cuts?

SOLUTION: 2m, 2m, 1m ... We will have 2 2m length cut and a single 1m cut, this is the optimal solution and we can make \$12 profit