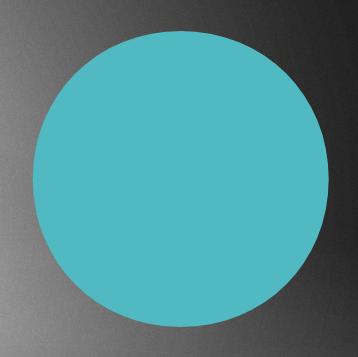
ARRAYS

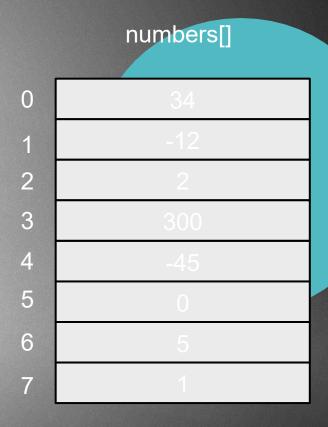
ARRAY



Arrays: data structures

A collection of elements / values each identified by an array index or key !!!

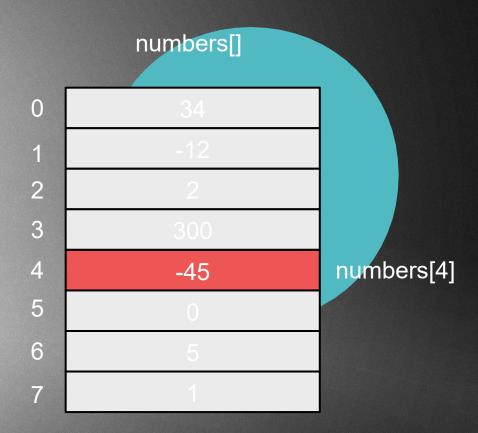
- index starts at zero
- because of the indexes: random access is possible



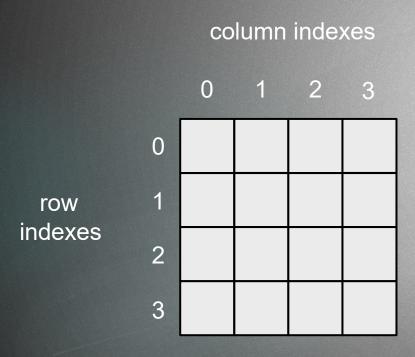
Arrays: data structure

A collection of elements / values each identified by an array index or key !!!

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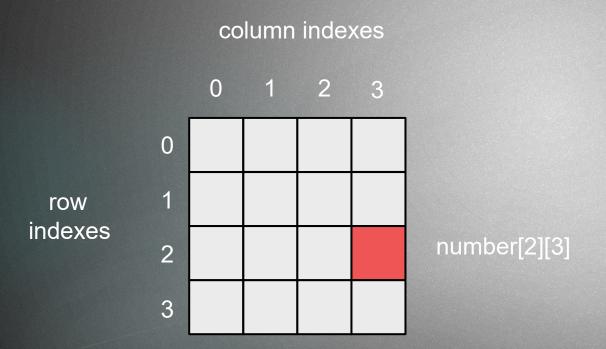
Multidimensional arrays: it can prove to be very important in mathematical related computations (matrixes)



numbers[][] two dimensional array

First paramter: row index
Second parameter: column index

Multidimensional arrays: it can prove to be very important in mathematical related computations (matrixes)



numbers[][] two dimensional array

First paramter: row index
Second parameter: column index

<u>Arrays</u>

- Arrays are data structures in order to store items of the same type
- We use indices as keys !!!
- Arrays can have as many dimensions as we want: one or two dimensional arrays are quite popular
- For example: storing a matrix → two dimensional array
- Dynamic array: when the size of the array is changing dynamically
- Applications: lookup tables / hashtables, heaps

Advantages

- We can use random access because of the keys: getItem(int index) will return the value with the given key very fast // O(1)
- Very easy to implement and to use
- Very fast data structure
- We should use arrays in applications when we want to add items over and over again and we want to take items with given indexes!!! ~ it will be fast

<u>Disadvantages</u>

- We have to know the size of the array at compile-time: so it is not so dynamic data structure
- If it is full: we have to create a bigger array and have to copy the values one by one // reconstructing an array is O(N) operation
- It is not able to store items with different types

Arrays operation: add

We can keep adding values to the array as far as the array is not full

We can keep adding values to the array as far as the array is not full

add(34)



We can keep adding values to the array as far as the array is not full

add(34)



We can keep adding values to the array as far as the array is not full

add(12)



We can keep adding values to the array as far as the array is not full

add(120)

34	0
12	1
120	2
	3
	4
	5
	6
	7

We can keep adding values to the array as far as the array is not full

add(-5)

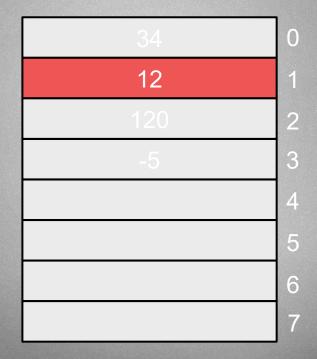
34	0
12	1
120	
-5	3
	4
	5
	6
	7

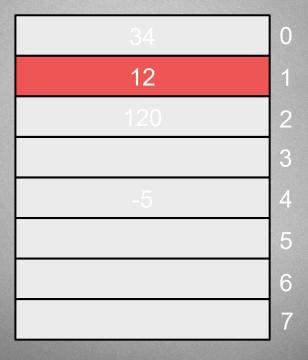
We can keep adding values to the array as far as the array is not full

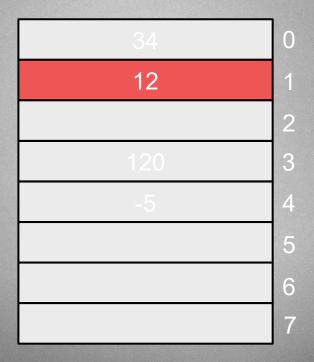
34	0
12	1
120	2
-5	2 3
	4
	5
	6
	7

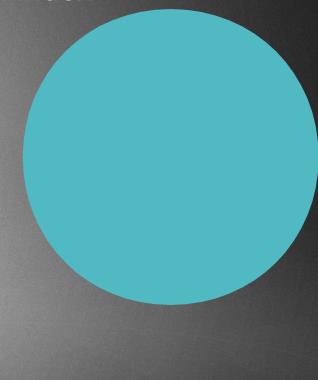
So: when adding new values to the list, we just have to insert it with the next index → very fast **O(1)** operation

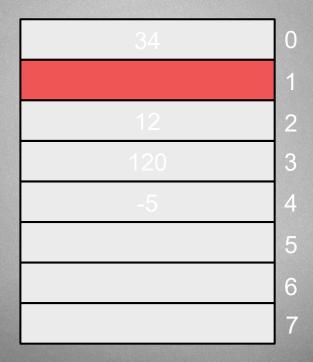
34	0
12	1
120	2
-5	3
	4
	5
	6
	7

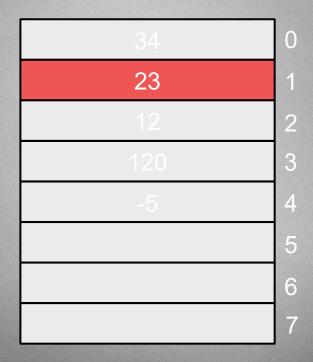




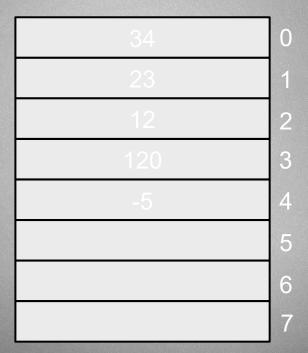








insert(23,1);



So: it is a bit more problematic, sometime we have to shift lots of values in order to be able to insert the new one $!!! \sim O(N)$ time complexity

Arrays operation: insert item

We would like to insert a given value with a given index

Add new item: O(1)
Insert item to a given index: O(N)

Arrays operation: remove items

We would like to remove the last item, it is very simple,

just remove it // O(1) time complexity

34	0
23	1
12	2
120	3
-5	4
	5
	6
	7

Arrays operation: remove items

We would like to remove the last item, it is very simple,

just remove it // O(1) time complexity

removeLast();

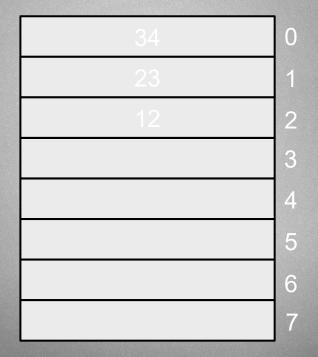
34	0
23	1
12	2
120	3
	4
	5
	6
	7

Arrays operation: remove items

We would like to remove the last item, it is very simple,

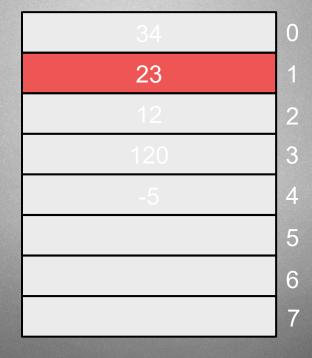
just remove it // O(1) time complexity

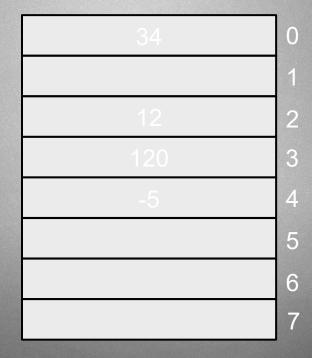
removeLast();

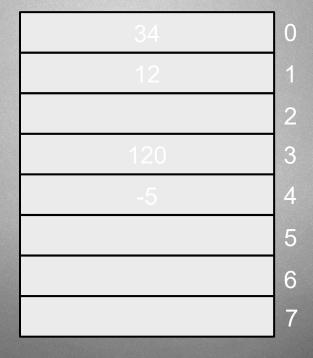


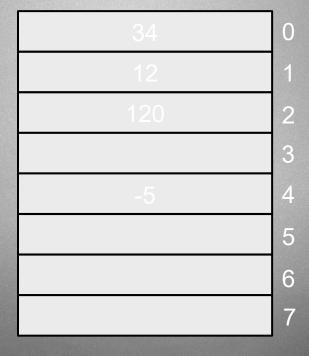
-		
	34	0
	23	1
	12	2
	120	2 3
	-5	4
		5
		6
		7

34	0
23	1
12	2
120	3
-5	4
	5
	6
	7

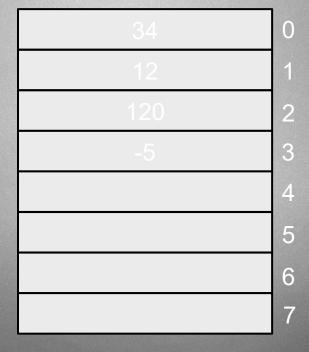








remove(1);



So: overall complecity will be linear O(N)

Removing last item: O(1)
Removing f.e. middle item: O(N)