

Imperative Programming

T2.1:

What does the program print if the language uses dynamic scope? Why?

y>0:

print: 1

y<0:

print: 0

because when $y > 0$, the function write () executed, the association for x is sought in the same block, which has been defined as 1, in line 2. but in the function f1(), which has been called in $y > 0$ block, there exists a sentence ($x = 0$;). according to the dynamic scope, the x in function f1() has been defined in the block , which is $\text{int } x=4$, and x has been destroyed in function f2(), $\text{int } x=5$. So x was equal to 4 and changed to 0. so the output should be 1, because the $x=1$ doesn't change.

When $y < 0$, the function write() executed, the association for x is sought in the same block, which has been defined as 1, in line 2. according to previous steps, the x in function f1() can not find a valid association in $y < 0$ block, then it will find the x has been defined as 1, in line 2. so $x=1$ has been change to 0. the output will be 0.

T2.2:

What does the program print if the language uses static scope and passes parameters by reference. Why?

Print:

5

7

the first write print the value 5, while the second one prints the value 7.

This is because: The procedure f1() is declared in the same block . Thus, $\text{int } x=5$ and f1(x) effectively associates x to 7 in first block.

In the case of the fist call of the write, we look if there exists a valid association for x in such nested block, it does exist and it is 5.

In the case of the second call of the write, we look if there exists a valid association for x in such nested block, it does exist and it is the first definition for x, and in Function f1(), the

value of x has been changed due to the static scope. So $x = x + y$ is executed, $y=2$ based on the passing parameter by reference, so $x=7$.

T2.3:

What does the program print if the language uses static scope and passes parameters by value. Why?

Print:

6

7

in the case of the first call of the function `write()`, we look if there exists a valid association for x in such nested block, it does exist and it is 6 due to $x++$, which will increment it by one.

In the case of the second call of the function `write()`, we look if there exists a valid association for x in find . -type f -exec du -a {} + | sort -n -r | less such nested block, it exists and it is the first definition of x. but in function $f1(x++)$, the parameter x is equal to 5 after calling function increments it by one. In other word, firstly, the programming call function $f1(x)$ and then execution $x=x+1$. So we can get the first definition x is $x = x + y \Rightarrow x = 2 + 5 \Rightarrow x = 7$

T2.4

What does the program print if the language uses static scope and passes parameters by name. Why?

Print:

5

7

In the case of the first call of the function `write()`, we look if there exists a valid association for x in such nested block, it does exist and it is 5 due to $f1(x++)$ passing parameter by name. In call-by-name, the actual parameter is passed unevaluated, and the address is worked out each time the formal parameter is used. In other words, in function $f1(x++)$,

$x = x + y \Rightarrow x = x + y++ \Rightarrow x = x + y; y = y + 1;$

the result will be $x = 5$ and $y = 6$, but the function `f1()` doesn't return any variables. So the output will be 5

in the case of the first call of the function `write()`, we look if there exists a valid association for x in such nested block, it does exist and has been change by the formula $x = x + y$, so the result will be 7.