

Agenda:

Space Complexity

$$2^k = N$$

$$\underbrace{2 \times 2 \times 2 \times \dots \times 2}_{k \text{ times}} = N$$

$$\log_2 N : k$$

Introduction To Arrays.

$$4N^2 + 3N + 1$$

$$4N^2 \quad \underline{\underline{O(N^2)}}$$

Reverse the array.

$$O(N^2)$$

Rotate the array k times.

$$1 \leq N \leq 10^5$$

Dynamic Arrays.

$$10^{10} \quad 10^8$$

TLE

Space Complexity:

```
func (int N)
{
    int x; → 4 bytes (32 bits)
    long y; → 8 bytes (64 bits)
    int z;    int[] A: new int[N]
}
16 Bytes. O(1)
```

1 byte = 8 bit

0 or 1

long → 8 bytes

Input Space
(x)

Output Space
(x)

O(1)

```
fun() (A[], N)
{
    ans = A[0];
    for (i → 1 to N-1)
    {
        ans = max(ans, A[i]);
    }
    return ans;
}
```

Introduction To Arrays:

$A[0]$ $A[N-1]$

`int arr[N];`

Indexing starts from 0.
ends at $N-1$

`int` \rightarrow 4 bytes.

`int[]` \rightarrow 4

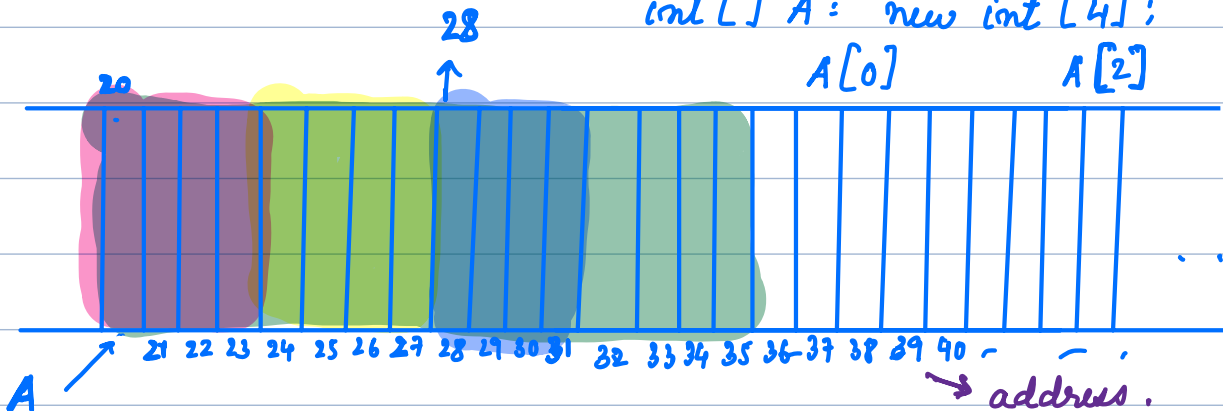
Time Complexity for accessing random element in the array. $\rightarrow O(1)$

`int[] A = new int[4];`

$A[0]$

$A[2]$

$$20 + (2 \times 4) = \underline{\underline{28}}$$



`int arr[] = new int[5];`

$O(1)$

Print all elements in the array.

```
fun (int[] A)
{
    int N = A.length;
    for (i → 0 to N-1)
    {
        println(A[i]);
    }
}
```

A.length();

Break : 8:07 AM

Swapping.

swap(1, 3)

A: ⁰5 ¹4 ²9 ³3

int temp = A[i];
A[i] = A[3];
A[3] = temp;

Question

Reverse an array.

i	res
N-1: 4	0
N-2	1

A: ⁰3 ¹4 ²2 ³8 ⁴6

res: - - - - -

A': ⁰6 ¹8 ²2 ³4 ⁴3

```
int[] reverseArray (int[] N)
{
    int[] res = new int [N.length];
    for (i = N.length-1; i ≥ 0; i--)
    {
        res[N.length-1-i] = A[i];
    }
    return res;
}
```

TC: $O(N)$

SC: $O(N)$

}

	0	1	2	i	
A:	3	4	2	3	4
				8	6
		j			

$i < j$

```
int[] reverseArray (int[] A)
{
    int i = 0; int j = A.length - 1;
    while (i < j)
    {
        swap (A[i], A[j]);
        i++; j--;
    }
    return A;
}
```

TC: $O(N)$

SC: $O(1)$

```
void swap (int[] A, int i, int j)
{
    // swap logic
}
```

Question: Reverse an array within a given range.

A: ⁰3 ¹4 ²2 ³8 ⁴6

$S = 2$ $e = 4$

Ans \rightarrow A: ⁰3 ¹4 ²6 ³8 ⁴2

initial value of
 i & j will be
 S & e

same as last program

Question: Given an array of size N , rotate the array from right to left k times. You can't use a new array.

$A: \overset{0}{1} \ \overset{1}{2} \ \overset{2}{3} \ \overset{3}{4} \ \overset{4}{5}$
 $R(1): 5 \ 1 \ 2 \ 3 \ 4$
 $R(2): 4 \ 5 \ 1 \ 2 \ 3$
 $R(3): 3 \ 4 \ 5 \ 1 \ 2$
 $R(4): 2 \ 3 \ 4 \ 5 \ 1$
 $R(5): 1 \ 2 \ 3 \ 4 \ 5$

$k = 3 \rightarrow 5$

$3 \% 5 \rightarrow \underline{\underline{3}}$

$k \% N$

$k \% N$

$4 \rightarrow \underline{\underline{3}}$

k
 $6 \rightarrow 1$
 $7 \rightarrow 2$
 $8 \rightarrow 3$
 $9 \rightarrow 4$
 $10 \rightarrow 0$

$11 \rightarrow 1$
 $12 \rightarrow 2$
 $13 \rightarrow 3$

$temp = 5$

`void rotateK(int[] A, int K)`

`{ int N = A.length;`

`K = K % N;`

`for (i = 0; i < K; i++)`

`{ int temp = A[N-1];`

`for (j = N-2; j >= 0; j--)`

`{ A[j+1] = A[j];`

`}`

`A[0] = temp;`

`}`

`}`

j
 $A: \overset{0}{5} \ \overset{1}{1} \ \overset{2}{2} \ \overset{3}{3} \ \overset{4}{4}$

$K * N \Rightarrow O(K * N)$

$TC \rightarrow \underline{\underline{O(N^2)}}$

$SC \rightarrow O(1)$

$1 \leq N \leq 10^5$

Optimized Approach:

A: ^{0 1 2 3 4 5 6}
5 6 7 1 2 3 4

^{0 1 2 3 4 5 6}
1 2 3 4 5 6 7 $k=3$

Reverse(A) → 7 6 5 4 3 2 1

Reverse(A, 0, k-1) → 5 6 7 4 3 2 1

Reverse(A, k, N-1) → 5 6 7 1 2 3 4

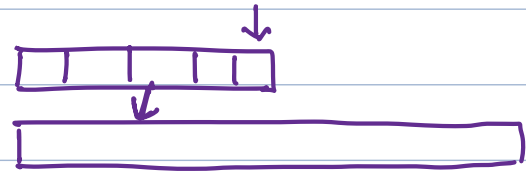
$$\begin{array}{lcl} \text{Reverse(A)} \rightarrow & O(N) & \\ \text{Reverse(A, 0, k-1)} \rightarrow & O(N) & + \\ \text{Reverse(A, k, N-1)} \rightarrow & O(N) & + \end{array} \quad \begin{array}{l} \\ \\ 3 \times O(N) \end{array}$$

$\Rightarrow \underline{O(N)}$

Dynamic Arrays:

Array \rightarrow length static.

Dynamic array \rightarrow You can add



Next Lecture:

Prefix Sum

Problems on Prefix Sum.

Doubts:

```
for(i=0; i<N; i++)  
{  
    int i  
}
```