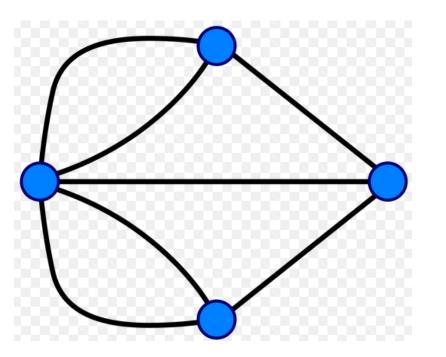
Lecture 36: Graph Theory and Eulerian Circuits

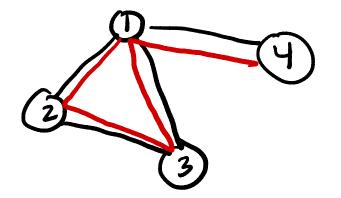
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Let G = (V, E) be a graph (directed or undirected). A walk of graph G is a sequence of alternating vertices (in V) and edges (in E) such that

- we start on any vertex and end on any vertex
- a single step in the walk proceeds upon an outgoing edge from the current vertex -

A walk has to respect edge direction in a directed graph. In an undirected graph, it doesn't matter.



Example: Consider the directed graph below:

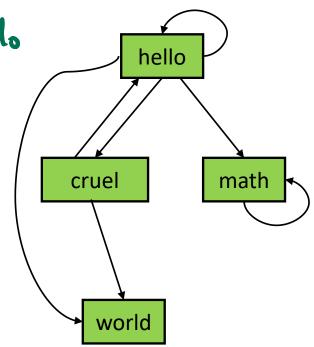
Valid walk: hello → hello → cruel → hello

Valid walk: hello-math

Valid walk: hello - cruel - world

Invalid walk: would > woeld

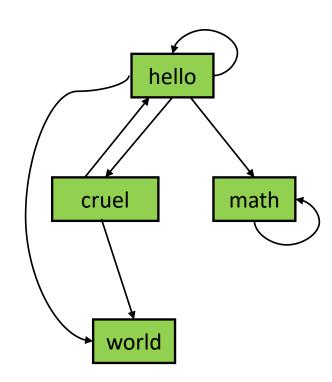
Invalid walk: world -> crue



Let G = (V, E) be a graph (directed or undirected). A **path** of a graph is a walk in which no vertex is repeated. The **length** of a path is the number of edges that it traverses.

★ Valid path: hello → cruel → world (path of length 2)

invalid path: hello -> cenel -> hello



named after the mathematician Euler

Let G = (V, E) be a graph (directed or undirected). An **Eulerian Circuit** of a graph is a special kind of walk which starts and ends at the same vertex and traverses each edge exactly once.

- A walk is pretty much anything, going from vertex to vertex.
- A **path** is a walk that does not repeat vertices.
- An Eulerian Circuit is a walk that does not repeat edges and starts and ends at the same vertex.

and traverses each edge

Example: Can there possibly be an Eulerian circuit on the graph below?

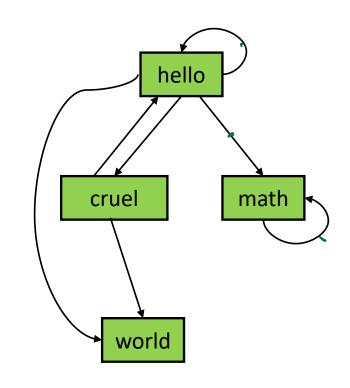
l ho!

Prove this by cases

Case 1: Statet at [hello]

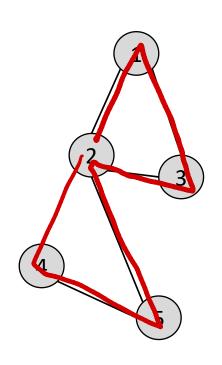
Case Z: Start at Cruel

case 3: Stant at [world]



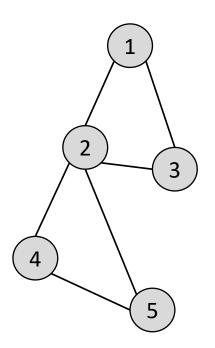
Case: Start at math)

Example: Can there possibly be an Eulerian circuit on the graph below?

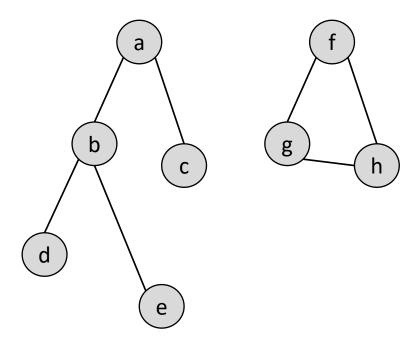


An undirected graph G is **connected** if there is a path between every pair of distinct vertices in the graph. An undirected graph is called disconnected if it is not connected.

This graph is connected.



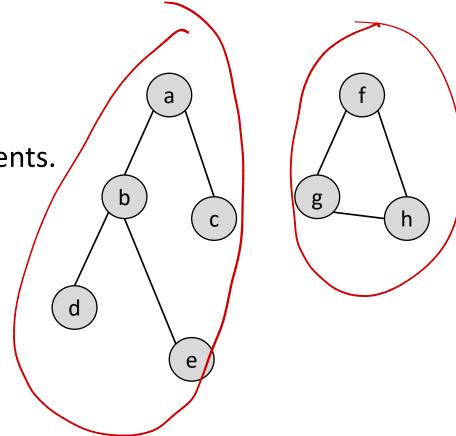
This graph is disconnected.



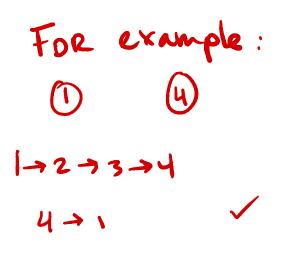
A subgraph of a disconnected graph G is called a **connected component** of G if it is a maximally connected subgraph of G.

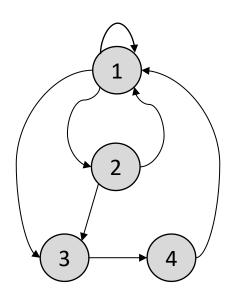
 $\{a,b,c,d,e\}$ and $\{f,g,h\}$ are connected components.

{a,b,c,d} are connected

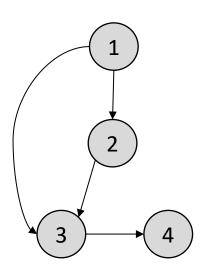


A directed graph (digraph) G is **strongly connected** if for each pair of distinct vertices a and b there is a path from a to b AND a path from b to a.





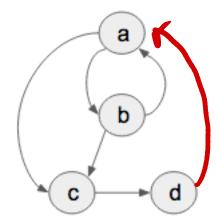
A directed graph (digraph) G is **weakly connected** if for each pair of distinct vertices a and b there is a path from a to b OR a path from b to a.



A subgraph of a digraph G is a strongly connected component if it is a maximal strongly connected subgraph of G.

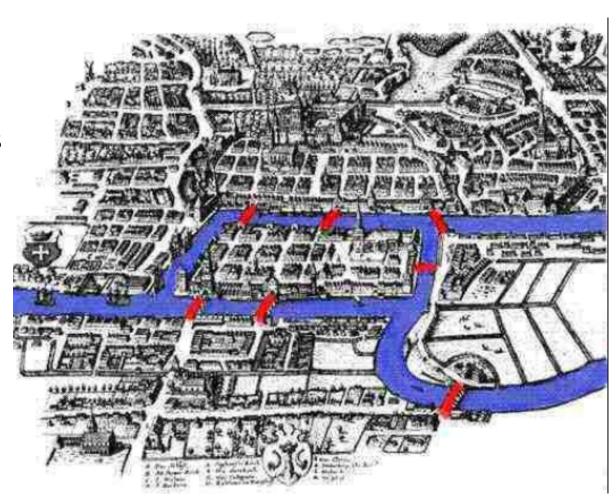
What is the strongly connected component of the graph G to the right? $G \cap G$

What is the smallest number of directed edges you could add to G to make it a strongly connected graph?



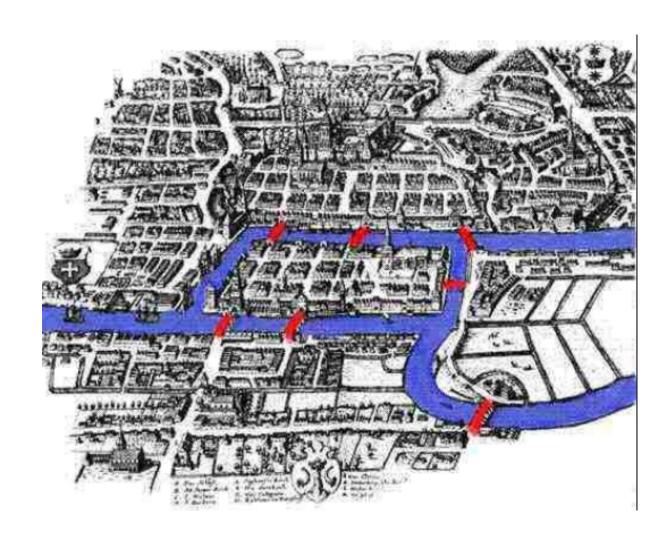
Puzzle: The city of Konigsberg has two islands formed by a river with seven bridges connecting the islands and the mainland.

Is there a circuit that traverses each bridge exactly once?



Is there a circuit that traverses each bridge exactly once?

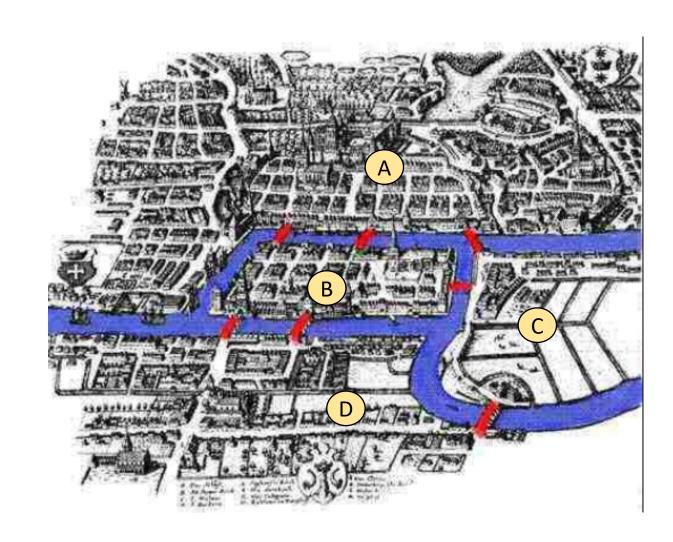
We can represent this as a graph, and solve the riddle by checking if it has an Eulerian circuit!

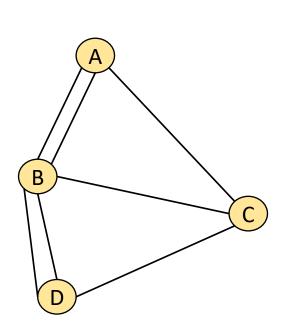


Let's formalize this as a graph.

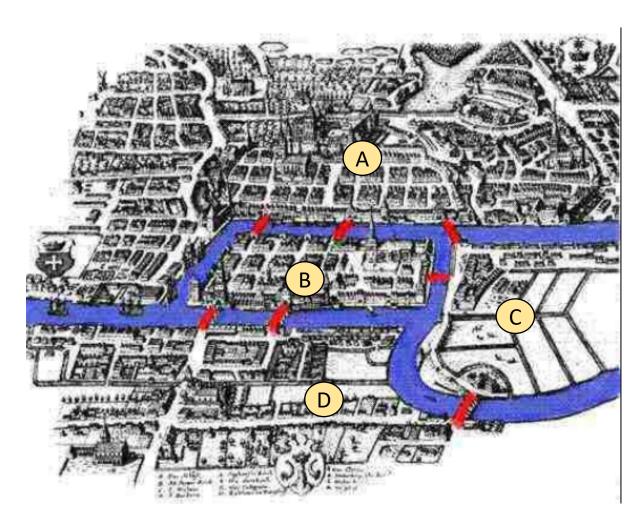
We will make the nodes (aka vertices) the separate landmasses.

The edges will be the bridges.

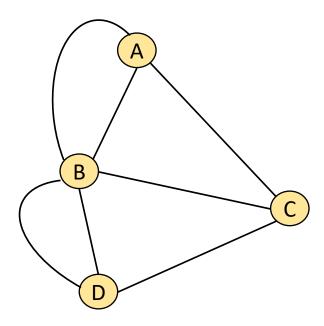




This is called a *multigraph*, because pairs of vertices have multiple edges connecting them.



Does this have an Eulerian circuit?



Theorem: A connected multigraph with at least two vertices has an Euler circuit if and only if each of its vertices has even degree.

Seven Bridges of Konigsberg: Since the graph has at least one vertex with odd degree, there cannot be an Eulerian circuit.

