Linux Interrupt Handlers

The Interrupt Handler API; Concurrency Considerations

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Overview

Roadmap:

- Terminology
- Examples
- · Registering, unregistering handlers
- proc/interrupts

Overview

Roadmap:

- Shared handlers
- Handler data
- Concurrency considerations

An interrupt:

```
"A signal that the hardware can send when it wants the processor's attention..." \label{eq:signal_processor}
```

-- Alessandro Rubini, Linux Device Drivers. O'Reilly. 2005.

Lots of devices:

- Serial ports (UARTS, SSP, I2C, ...)
- Network controllers (Ethernet, USB, CAN, ...)
- · Timers and counters, GPIO, ...

Just about everything!

An interrupt handler:

- Code that responds to a device's interrupt request
- A.k.a. "Interrupt service routine"

"Interrupt handler" != "device driver":

- Device drivers might involve interrupts, handlers
- It really depends on the device!
- Linux itself doesn't associate the two

Minimal Example

```
static int gpio_pb29_irgs;
2
    static irgreturn_t
3
    gpio_pb29_irg (int irg, void *unused)
4
5
      printk("%s: %d\n", __FUNCTION__, qpio_pb29_irqs);
6
      if (++gpio_pb29_irgs > 10) {
7
        printk(KERN ERR "%s: disabling\n", FUNCTION );
8
        disable_irq(qpio_to_irq(AT91_PIN_PB29));
10
      return IRO HANDLED;
11
12
```

```
request_irq()
```

- · Registers an interrupt handler
- Returns 0 on success

unsigned int irq

- Interrupt channel identifier
- Enumerations are architecture-specific
- Defined in <asm/irq.h>

```
irq_handler_t handler
```

- Pointer to the handler function
- · Handler is invoked when interrupt is detected

```
typedef int irqreturn_t;
typedef irqreturn_t (*irq_handler_t)(int, void *);
```

unsigned long irqflags

- Various channel-specific parameters
- Some channels might not support all flags

- IRQF_TRIGGER_RISING Trigger on rising edge IROF TRIGGER FALLING IROF_TRIGGER_HIGH — Trigger on high level IROF TRIGGER LOW
 - IRQF_DISABLED Disable interrupts during handling
 - IROF SAMPLE RANDOM Interrupt is a source of randomness
 - IROF_SHARED Allow sharing of the interrupt line

 - Trigger on falling edge
 - Trigger on low level

const char *name

• Plaintext name, for /proc/interrupts

void *dev_id

- Pointer to handler's private data
- Returned to handler during interrupts

```
free_irq()
```

- Unregisters an interrupt handler
- Disables the interrupt channel if appropriate
- Use the same dev_id passed to request_irq()

```
void free_irq(unsigned int irq, void *dev_id);
```

The /proc/interrupts Interface

The /proc filesystem:

- Pseudo-files of processor-specific information
- Pseudo-directories of process-specific information

```
$ cat /proc/interrupts
```

. . .

```
irqreturn_t handler(int irq, void *dev_id)
```

Interrupt handler signature

```
static irqreturn_t
gpio_pb29_irq (int irq, void *unused)
```

int irq

• Interrupt channel number

```
void *dev_id
```

- A.k.a. "device identifier"
- Used to distinguish devices on shared channels
- Provided in request_irq()
- Must be unique for IRQF_SHARED

```
struct irq_data {
      int data;
2
    };
    struct irg_data irg_data;
4
5
    static int
    example_open (struct inode *inode, struct file *pfile)
8
      int ret;
      ret = request_irq(IRQNUM, irq_handler, 0,
10
11
                          "example", &irg data);
      return ret;
12
13
```

```
irqreturn_t
```

· Return value from handler

```
\label{eq:local_continuity} \begin{array}{lll} & \text{IRQ\_NONE} & \textbf{— Did not handle the interrupt request} \\ & \text{IRQ\_HANDLED} & \textbf{— Handled the interrupt request} \\ & \text{IRQ\_RETVAL}(\textbf{x}) & \textbf{— Helper macro that returns NONE or HANDLED} \\ & \text{\#define } & \text{IRQ\_RETVAL}(\textbf{x}) & ((\textbf{x}) & != 0) \end{array}
```

"irq255: nobody cared"

- All registered handlers returned IRQ_NONE, or
- There were no registered handlers
- Followed by an OOPS output

Things to investigate:

- Is my return value correct?
- Is my interrupt handler detecting all possible requests?
- Am I registering with the wrong channel?
- Is it the wrong type of signal (level vs. edge)?

No Sleeping!

Interrupt handlers must be atomic:

- No pending on a semaphore
- No waiting on a completion
- No sleep-based delays

Be especially careful during:

- Memory allocation (!)
- User memory access

No Sleeping!

Memory allocation:

- Avoid it in an interrupt handler!
- Preallocate in your interface code
- Pass a pointer via dev_id

No Sleeping!

User memory access:

- Target page might be swapped out
- Use a preallocated intermediate buffer
- The mmap() APIs can lock pages
- Consider DMA

Enabling and Disabling Interrupts

void enable_irq(unsigned int irq)

Enables an interrupt channel

void disable_irq(unsigned int irq)

- Disables the requested interrupt request line
- Can be called from an interrupt handler
- Affects all registered handlers
- (You probably won't do this often)

Enabling and Disabling Interrupts

Linux uses "lazy" interrupt disables:

- The disable_irq() only sets a flag
- Pending interrupts are still serviced
- Interrupts are physically disabled later
- (This laziness is usually well-hidden)

"When do I enable interrupts?"

```
request_irq():
```

- · Registers the interrupt handler
- Enables the interrupt channel in the controller
- Does NOT enable the device to assert interrupts!

```
enable_irq():
```

- Unmasks an interrupt line
- Does NOT enable the device to assert interrupts!

"When do I enable interrupts?"

Follow this sequence:

- Probe for the device, if possible
- Disable interrupts at the device, if possible
- Initialize the associated driver, if there is one
- Initialize the device

"When do I enable interrupts?"

And then:

- Register the interrupt handler via request_irq()
- Enable interrupts at the device, as appropriate

Polling Interrupt Handlers

Polled vs. interrupt-driven:

- Your interrupt handler might support both!
- Done in some ethernet drivers, elsewhere

```
/* set up a transmit */
...;
/* send it out */
while (interrupt_handler(&data) == IRQ_HANDLED)
{
}
```

Shared Data Issues

Sharing data with interrupt handlers:

- Potentially leads to race conditions (bad!)
- Demands attention to concurrency issues

Code very cautiously!

Perfectly safe if done properly

Shared Data Issues

Commonly-used facilities:

- Completions
- Spinlocks

Also:

- Work queues
- Tasklets

Shared Data Issues

Forget about:

- Mutexes, semaphores
- Disabling interrupts

Completions

The completion API:

- One-way communication between contexts
- "The requested job is done"

Conceptually similar to a semaphore, but:

- Semaphores are optimized for the "available" case
- Semaphores are expensive during contention
- Semaphores are going away (!)

struct completion

```
#include <linux/completion.h>
2
    struct completion c;
3
4
    irqreturn_t
    example_irq_handler(int irq, void *unused)
      /* get data... */
8
      . . . ;
      complete(&c);
11
      return IRQ HANDLED;
13
```

struct completion

Recap

The interrupt API:

- request_irq()
- Interrupt handlers
- No sleeping!

Recap

Concurrency:

- Sharing data with interrupt handlers
- Demands use of completions, etc.
- Failure to do so leads to races

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