HW1 examples

Source: Andy

Destination: Alice

	Andy	Bob	Claire	David	Alice
Andy	0	1	1	0	0
Bob	1	0	5	4	0
Claire	1	5	0	0	1
David	0	4	0	0	0
Alice	0	0	1	0	0

HW1 examples

Source: Andy

Destination: Alice

	Andy	Bob	Claire	David	Alice
Andy	0	1	1	0	0
Bob	1	0	5	4	0
Claire	1	5	0	0	1
David	0	4	0	0	0
Alice	0	0	1	0	0

Issues

- Reference implementation created by the TAs did not follow the homework specifications exactly.
- "standard" algorithms are usually under-determined, e.g., they do not specify ordering hence they may not be directly applicable.
- "pure" algorithms, e.g., BFS strictly uses a FIFO queue, are under-determined most likely those will not work. E.g., your BFS answer could use a FIFO queue, plus some loop detection logic, plus some extra logic to enforce the alphabetical popping.
- "equivalent tricks", e.g., "to achieve alphabetical popping, use reverse alphabetical pushing in DFS" may not work use only if you are sure your trick truly is equivalent to the stated specifications.

Queue: add successors to queue back;

empty queue from front (top)

#	state	depth	path cost	parent #
$\overline{1}$	Andy	0	0	



Andy

Claire

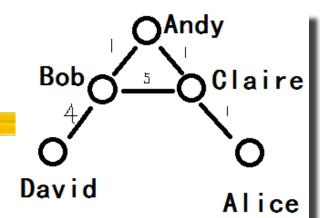
Alice

Bob

David

Queue: add successors to queue back;

empty queue from front (top)



#	state	depth	path cost	parent #	
1	Andy	0	0		_
2	Bob	1	1	1	
3	Claire	1	1	1	

Queue: add successors to queue back; empty queue from front (top)

#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1

Bob

David

Note: When the costs of two or more nodes are equal, you need to make sure these nodes are **popped** off the search queue in alphabetical order. This will resolve ambiguity and ensure that there is only one correct solution for each problem.

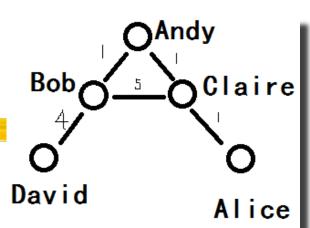
Log: Andy

Claire

Alice

Queue: add successors to queue back;

empty queue from front (top)



#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
4	David	2	2	2

Children of #2 (Bob): Andy, Claire, David

Andy: not enqueued (case 3) Claire: not enqueued (case 2) David: enqueued (case 1)

Note 2: Your algorithm should perform *loop detection*. As studied in lectures 2-4, do not enqueue a child that has a state already visited, unless the child has a better cost than when we previously visited that state (see slides about "A clean robust algorithm" in session02-04.pptx lecture slides).

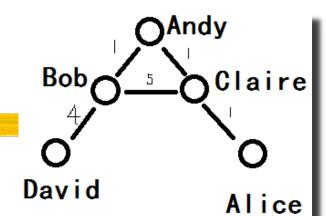
Log: Andy - Bob

A Clean Robust Algorithm

```
[... see previous slide ...]
         children ← Expand(currnode, Operators[problem])
         while children not empty
                  child ← Remove-Front(children)
                  if no node in open or closed has child's state
      Case 1
                           open ← Queuing-Fn(open, child)
                  else if there exists node in open that has child's state
      Case 2
                           if PathCost(child) < PathCost(node)</pre>
                                    open ← Delete-Node(open, node)
                                    open ← Queuing-Fn(open, child)
                  else if there exists node in closed that has child's state
      Case 3
                           if PathCost(child) < PathCost(node)</pre>
                                    closed ← Delete-Node(closed, node)
                                    open ← Queuing-Fn(open, child)
         end
[... see previous slide ...]
                                CS 561, Sessions 2-4
```

Queue: add successors to queue back;

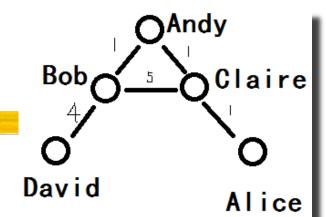
empty queue from front (top)



#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
4	David	2	2	2
5	Alice	2	2	3

Queue: add successors to queue back;

empty queue from front (top)



#	state	depth	path cost	parent #	
1	Andy	0	0		
2	Bob	1	1	1	
3	Claire	1	1	1	
4	David	2	2	2	
5	Alice	2	2	3	

Queue: add successors to queue back;

empty queue from front (top)

, ,	Andy
Bob	Claire
0	\0
David	Alice

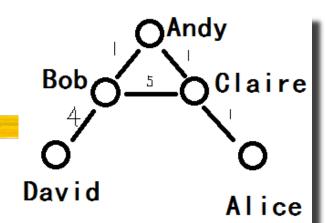
#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
4	David	2	2	2
5	Alice	2	2	3

To get the path: backtrack from solution node up the chain of parent nodes

Path: ... (#5) Alice

Queue: add successors to queue back;

empty queue from front (top)

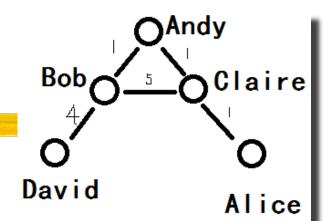


#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1 🛑
4	David	2	2	2
5	Alice	2	2	3

Path: ... (#3) Claire - (#5) Alice

Queue: add successors to queue back;

empty queue from front (top)



#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
4	David	2	2	2
5	Alice	2	2	3

Path: (#1) Andy - (#3) Claire - (#5) Alice

Queue: add successors to queue front;

empty queue from front (top)

state depth

Bob 5 Claire

O O

David

Alice

parent #

path cost

Andy

1 Andy 0 --

Log:

14

Queue: add successors to queue front;

empty queue from front (top)

state

depth

path cost

David Alice

parent #

Bob

Andy

Claire

Bob Claire Andy

Log: Andy

Queue: add successors to queue front;

empty queue from front (top)

state depth

path cost

parent #

Bob

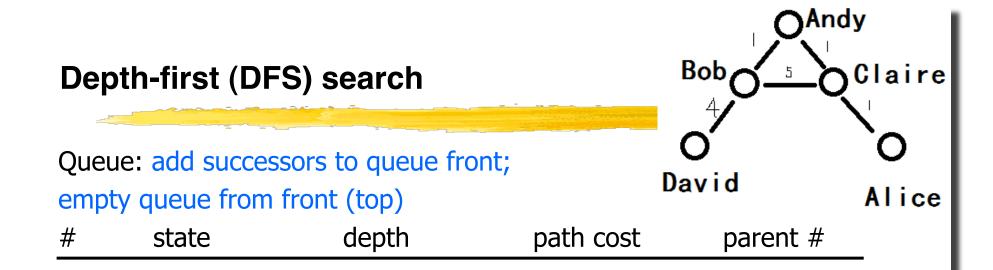
David

Andy

Claire

Alice

Log: Andy



Enqueue Claire again? No because of loop detection rule (node #3 in open queue already has state Claire)

Note 2: Your algorithm should perform *loop detection*. As studied in lectures 2-4, do not enqueue a child that has a state already visited, unless the child has a better cost than when we previously visited that state (see slides about "A clean robust algorithm" in session02-04.pptx lecture slides).

4	David	2	2	2
2	Bob	1	1	1
3	Claire	1	1	1
1	Andv	0	0	

Log: Andy - Bob

A Clean Robust Algorithm

```
[... see previous slide ...]
         children ← Expand(currnode, Operators[problem])
         while children not empty
                  child ← Remove-Front(children)
                  if no node in open or closed has child's state
                           open ← Queuing-Fn(open, child)
                  else if there exists node in open that has child's state
                           if PathCost(child) < PathCost(node)</pre>
                                    open ← Delete-Node(open, node)
                                    open ← Queuing-Fn(open, child)
                  else if there exists node in closed that has child's state
                           if PathCost(child) < PathCost(node)</pre>
                                    closed ← Delete-Node(closed, node)
                                    open ← Queuing-Fn(open, child)
         end
[... see previous slide ...]
                                 CS 561, Sessions 2-4
                                                                                18
```

Queue: add successors to queue front;

empty queue from front (top)

state depth

path cost

parent #

Bob

David

Andy

Claire

Alice

4	David	2	2	2
2	Bob	1	1	1
3	Claire	1	1	1
1	Andv	0	0	

Log: Andy - Bob

Queue: add successors to queue front;

empty queue from front (top)

state depth path cost parent

 4
 David
 2
 2

 2
 Bob
 1
 1
 1

 3
 Claire
 1
 1
 1

 1
 Andy
 0
 0
 -

Log: Andy - Bob - David

Andy

Claire

Alice

Bob

David

Queue: add successors to queue front;

empty queue from front (top)

state depth

path cost

parent #

Bob

David

Andy

Claire

Alice

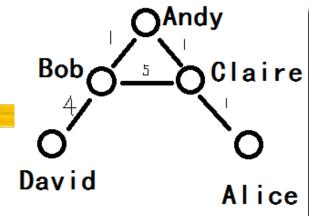
5	Alice	2	2	3
4	David	2	2	2
2	Bob	1	1	1
3	Claire	1	1	1
1	Andy	0	0	

Log: Andy - Bob - David - Claire

Queue: add successors to queue front;

empty queue from front (top)

state depth path cost



parent #

Path: (#1) Andy - (#3) Claire - (#5) Alice

5	Alice	2	2	3
4	David	2	2	2
2	Bob	1	1	1
3	Claire	1	1	1
1	Andy	0	0	

Log: Andy - Bob - David - Claire - Alice

Queue: keep queue sorted by path cost; empty queue from front (top)

Andy			
Bob	Claire		
4/			
0	0		
David	Alice		

#	state	depth	path cost	parent #	
1	Andy	0	0		_

Log:

Queue: keep queue sorted by path cost;

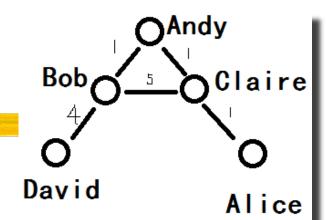
empty queue from front (top)

OAndy			
Bob O-	Claire		
0	\		
David	Alice		

#	state	depth	path cost	parent #	
1	Andy	0	0		_
2	Bob	1	1	1	
3	Claire	1	1	1	

Queue: keep queue sorted by path cost;

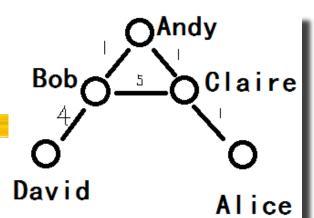
empty queue from front (top)



#	state	depth	path cost	parent #	_
1	Andy	0	0		•
2	Bob	1	1	1	
3	Claire	1	1	1	
4	David	2	5	2	

Queue: keep queue sorted by path cost;

empty queue from front (top)

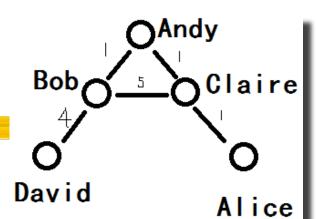


#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
5	Alice	2	2	3
4	David	2	5	2

Note the sorting by path cost (#5 goes above #4)

Queue: keep queue sorted by path cost;

empty queue from front (top)



#	state	depth	path cost	parent #
1	Andy	0	0	
2	Bob	1	1	1
3	Claire	1	1	1
5	Alice	2	2	3
4	David	2	5	2

Path: (#1) Andy - (#3) Claire - (#5) Alice

Summary

Prof Itti

Original HW1

HW1 Erratum

Log: Andy-Bob-Claire-Alice

Path: Andy-Bob-Claire-Alice

BFS

Log: Andy - Bob - Claire - Alice

Path: Andy - Claire - Alice Log: Andy-Bob-Claire-David-Alice

Path: Andy-Claire-Alice

DFS

Log: Andy - Bob - David - Claire - Alice

Path: Andy - Claire - Alice
Log: Andy-Bot

Log: Andy-Bob-David-Claire-Alice

Path: Andy-Claire-Alice

UCS

Log: Andy - Bob - Claire - Alice

Path: Andy - Claire - Alice Log: Andy-Bob-Claire-Alice

Path: Andy-Claire-Alice

Summary

Prof Itti

Original HW1

HW1 Erratum

BFS

Log: Andy - Bob - Claire - Alice

Path: Andy - Claire - Alice Log: Andy-Bob-Claire-David-Alice

Path: Andy-Claire-Alice

DFS

Log: Andy - Bob - David - Claire - Alice

Path: Andy - Claire - Alice

Log: Andy-Bob-David-Claire-Alice

Path: Andy-Claire-Alice

Log: Andy Bob-Claire Alice Path: Andy-Bob-Claire Alice

UCS

Log: Andy - Bob - Claire - Alice

Path: Andy - Claire - Alice Log: Andy-Bob-Claire-Alice

Path: Andy-Claire-Alice

Recommendation

- As already recommended earlier, please make sure your algorithm complies with the specifications given in the HW handout. Following the specs overrides following the examples.
- "standard" algorithms are usually under-determined, e.g., they do not specify ordering hence they may not be directly applicable.
- "pure" algorithms, e.g., BFS strictly uses a FIFO queue, are under-determined most likely those will not work. E.g., your BFS answer could use a FIFO queue, plus some loop detection logic, plus some extra logic to enforce the alphabetical popping.
- "equivalent tricks", e.g., "to achieve alphabetical popping, use reverse alphabetical pushing in DFS" may not work use only if you are sure your trick truly is equivalent to the stated specifications.
- please do not try to convince the TAs that some standard implementation should give the correct result for this problem standard implementations usually are not 100% compliant with this specific problem.
- if/where possible ambiguity remains, try your best to comply with the specifications and let us know your thoughts in a README file.