Lecture 7 – Spatial Domain Decomposition and Data Structures for Multi-Block (Domain) Computations

- Domain Decomposition has several meanings
- There are different types of domain decomposition:
 - Spatial/grid domain decomposition is where a single domain is broken up into multiple "blocks" or sub-domains in order to
 - Make grid generation easier
 - Use parallel computing to solve larger problems and/or solve a problem faster
 - Discipline domain decomposition is where a domain is broken up into different physics (with different equation sets and/or models)
 - For different disciplines
 - For different physics
 - Algorithm domain decomposition is where different algorithms are used depending on purpose
 - Convergence acceleration (Eg multi-grid) vs fine-grid
 - Processor/Parallel block domain decomposition is where multiple blocks are distributed across multiple processors in such a way
 - Evenly distribute the work-load (load balancing) across the processors

Data Structures for Spatial/Grid and Processor/Parallel Domain Decomposition

- When developing a parallel code to solve an engineering simulation, we have to consider
 - Where data must be exchanged at block and processor boundaries during the course of the algorithm
 - What data must be exchanged
 - If barriers are necessary to synchronize the processors
 - What blocks and processors must exchange data
- Different engineering simulation codes may use different types of computational grid and differencing (or integration) stencil
 - Multi-block structured grids
 - Multi-block unstructured grids
 - Multi-blocks of hybrid, structured/unstructured grids

Multi-Block Data Structures

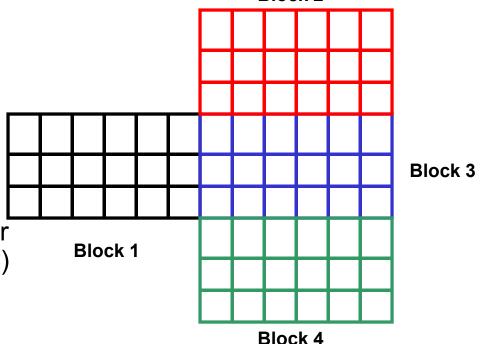
- Inter-block and inter-processor communication is dependent on block neighbor information
- Each block (of computational points or cells) must keep track of
 - Its own block and processor number
 - Block and processor numbers of its neighbors (E, W, N, S, H, Y)
 - The boundary conditions to be applied along its face (edge)
 - Additional information for special applications (such as parent/child relationships for embedding-adaptation)
- Upon execution, each block can then determine
 - The processor which has been assigned to it
 - The processors associated with its neighbors
 - If communication is necessary for each block face (edge)

Example: Spatial/Grid Decomposition Block 2

 Flow through a sudden expansion, or heat transfer in a "T" strut, etc.

Each block must store

- Its N, S, W, E, H, Y neighbor block numbers (connectivity)
 - Example: Block-1 has neighbors 0,0,0,3,0,0
- The boundary condition types (if any) at each point along is N, S, W, E, H, Y faces (edges)
 - Example: Block-1 has a BC type of 1 (solid boundary) along its southern boundary or -3 along its eastern boundary (which says that this is a block/block boundary



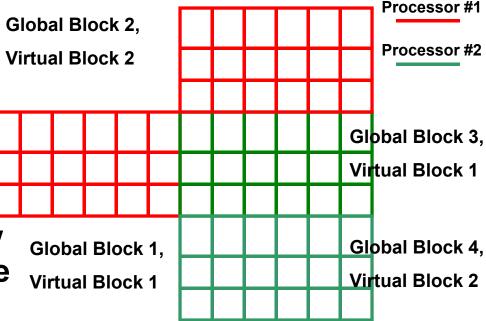
- Another way to store BC information would be to store regions of a face (sub-faces) with descriptions of the region dimensions and BC type
- Yet another way to store BC information would be point-by-point. This may make the user input description tedious, however.

Creation of Multi-Block Grids

- The creation of a multi-block structured, unstructured, or hybrid grid is generally performed with a separate grid-generation code.
 - A multi-block restart file can be written that contains the block dimensions, connectivity, and coordinates
 - PLOT3D is an "old" file format for restart files
 - CGNS is a new (modern) data-base format for restart files (you can find information regarding CGNS at: www.cgns.org
 - The initial values of the simulation code may be written in these same formats either as a separate file (PLOT3D) or as part of the data-base (CGNS)
- In the case of your project simulation code, the next step will be to write a separate code to break up the original single-block grid into multiple blocks with connectivity and BC information (spatial/grid domain decomposition) for serial execution (projects 2 and 3).
 - Note, however, that you will ultimately need to create the data structure for parallel execution (projects 4 and 5).

Example: Processor/Parallel Domain Decomposition

- Processor/parallel domain decomposition may be used to determine the processor assignments to each block
- For instance, say we knew ahead of execution that we could use 2 processors
 - Processor #1 is assigned to block 1 and 2
 - Processor #2 is assigned to block 3 and 4
- So at this point, each processor can find out its neighboring processors via the block neighbor information



For instance, block-1:

- Knows that its east boundary is a block/block interface from its BC information for that face (edge)
- Knows its east neighbor is block-3
- Can determine that block-3's processor is #2 so the east face (edge) is also an inter-processor boundary

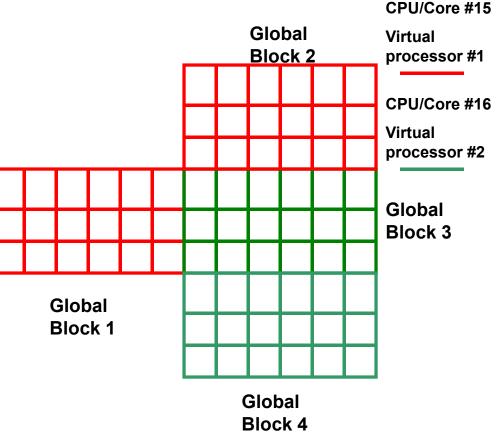
Example: Block and Processor Decomposition

- So before execution even begins, we know the following for each processor:
 - Block numbers of all blocks
 - Grid dimensions of all blocks
 - Coordinates and physical variables at nodes or cell-centers of all blocks
 - Boundary condition types for all points on all faces (edges) of all blocks
 - Neighbors of all blocks (N, S, W, E, H, Y)
 - Inter-processor boundaries (if any) for all blocks
 - The processor numbers assigned to the neighbors across interprocessor boundaries (This tells us who to communicate with)
- Most of this information could be generated by a preprocessor spatial decomposition or grid-generation code, then written out to a file for each processor that could be read at the beginning of simulation execution.
 - Determination of inter-processor boundaries and block assignments to virtual processors (last 2 steps above) would come later during processor/parallel domain decomposition

Example: Virtual to Physical Processor Mapping

Upon execution, physical processor numbers are mapped to computational processor numbers by MPI

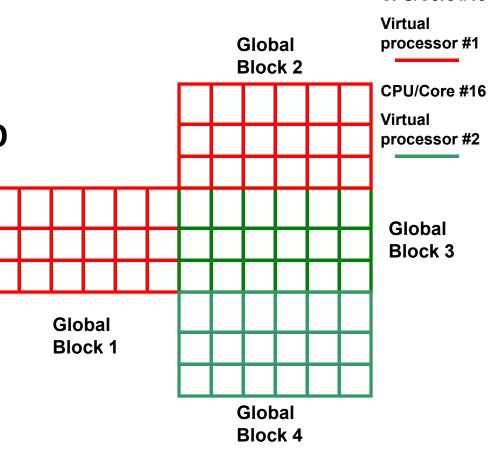
- Example: CPU/Core #15 is assigned all the blocks of computational processor #1 (red blocks) and reads the pertinent block information
- CPU/Core #16 is assigned all the blocks of computational processor #2 and reads all the pertinent block information (green blocks)
- Solution iteration begins



Example: SPMD vs MPMD Parallel Computing

 If all blocks are executing the same analysis (code), then this type of parallel computing is called SPMD (Single Process, Multiple Data)

• If the red blocks were executing one type of analysis and the green blocks were executing a different analysis (Multi-Disciplinary Simulations), then this type of parallel computing is called MPMD (Multiple Process, Multiple Data)



CPU/Core #15

Data Structures to Aid in Multi-Block Discretization and Parallel Communication

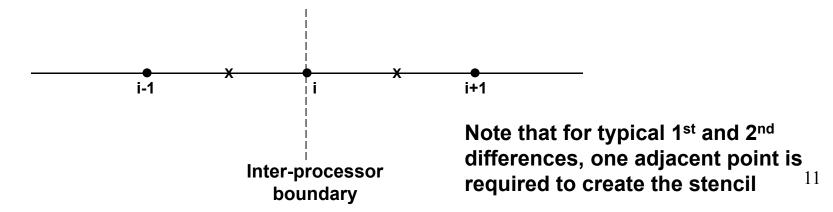
- So far, we have only addressed issues related to breaking up the grid, assigning sub-groups of blocks to processors, and making the boundary condition specification more general
- We now must address the communication that takes place between blocks and processors during iteration of the solution algorithm
- This involves the discretization stencil which is dependent on the type of algorithm being used
 - Finite difference, finite volume, finite element, other

Data Structures to Aid in Multi-Block Discretization and Parallel Communication

- The differencing (or integration) stencil in engineering simulations depends on the governing equations to be solved and the numerical algorithm used to solve them
 - For instance, most 2nd order-accurate difference operators used for the physical science simulations use a 3-point stencil

$$\left. \frac{\partial \phi}{\partial x} \right|_{i} = \frac{\phi_{i+1} - \phi_{i-1}}{2\Delta x} \quad or \quad \frac{\partial^{2} \phi}{\partial x^{2}} \right|_{i} = \frac{\phi_{i+1} - 2\phi_{i} + \phi_{i-1}}{\Delta x^{2}}$$

 One must keep the largest stencil used during the course of execution in mind when developing parallel code



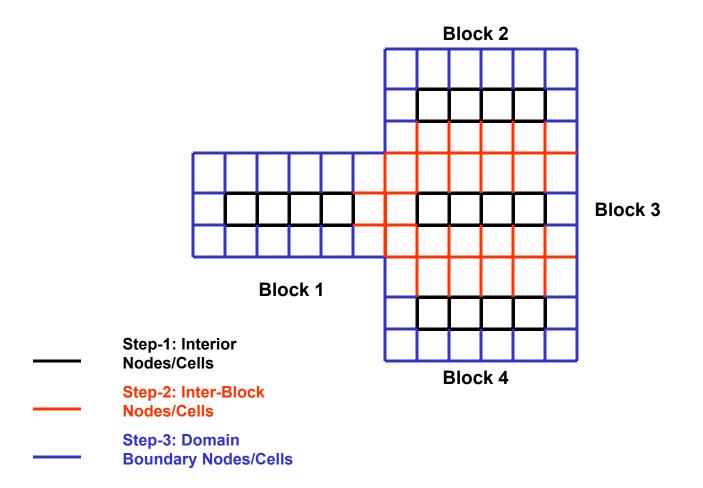
Data Structures to Aid in Multi-Block Discretization and Parallel Communication

- Different strategies may be used to create the stencil at the interblock boundary
 - Since the block neighbor and boundary condition information is known for each side of every block, we can quickly determine the adjacent block processor numbers
- Method-1: Information may be passed on the fly during the construction of the stencil at the edge
 - Tends to be inefficient due to a large number of small-size packets of information exchanged
- Method-2: Stencil may be constructed via an "accumulation" of forward and backward stencils at the edges (lecture 5)
 - Sum the contributions from different processors at the edges at the end of stencil construction (cell-vertex schemes)
 - Requires additional logic to handle various stencil operators toward edges
- Method-3: Information may be passed at the end of each iteration into buffers (halo or ghost cells/nodes) that "extend" each block
 - Requires additional storage but is efficient
 - Ghost cells/nodes may also be used to perform boundary conditions

Method-1: "On the Fly" Stencil Construction

- Pseudo-Code for multi-block, structured-grid simulations using method-1:
 - Loop over blocks
 - Loop to construct discretization of stencil at internal nodes/cells
 - Loop to construct discretization of stencil at internal north boundary nodes/cells (reach across to north neighbor)
 - Loop to construct discretization of stencil at internal south boundary nodes/cells (reach across to south neighbor)
 - Loop to construct discretization of stencil at internal west boundary nodes/cells (reach across to west neighbor)
 - Loop to construct discretization of stencil at internal east boundary nodes/cells (reach across to east neighbor)
 - Loop to construct discretization of stencil at internal hither boundary nodes/cells (reach across to hither neighbor)
 - Loop to construct discretization of stencil at internal yonder boundary nodes/cells (reach across to yonder neighbor)
 - Construct discretization of stencil at block corner nodes/cells (reach across to two appropriate neighbors for each)
 - · Perform boundary conditions on all non-inter-block faces
 - Update variables and make sure inter-block boundary nodes/cells are consistent

Example: "On the Fly" Stencil Construction

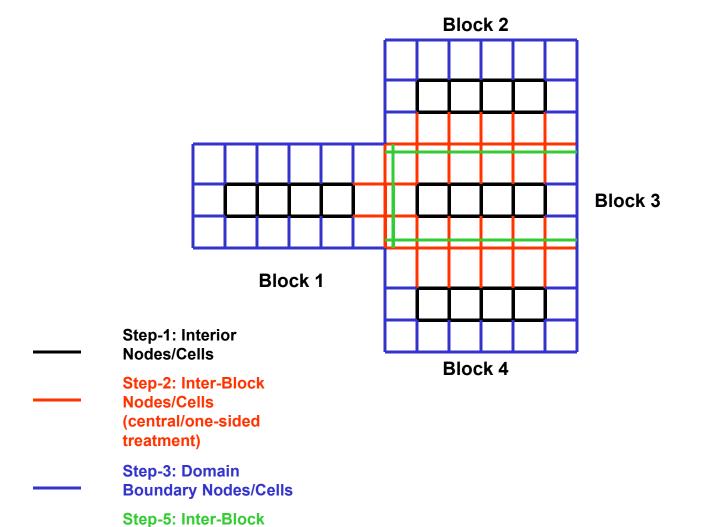


Method-2: "Accumulation" Stencil Construction

Pseudo-Code for multi-block, structured-grid simulations using method-2:

- Loop over blocks
 - Loop to construct discretization of stencil at internal nodes/cells
 - Loop to construct discretization of stencil at north boundary using combination of central and one-sided treatments
 - Loop to construct discretization of stencil at south boundary using combination of central and one-sided treatments
 - Loop to construct discretization of stencil at west boundary using combination of central and one-sided treatments
 - Loop to construct discretization of stencil at east boundary using combination of central and one-sided treatments
 - Loop to construct discretization of stencil at hither boundary using combination of central and one-sided treatments
 - Loop to construct discretization of stencil at yonder boundary using combination of central and one-sided treatments
 - Perform boundary conditions on all non-inter-block faces
- Accumulate contributions at inter-block faces
 - Message pass (send/receive) boundary contributions
- Update variables

Example: "Accumulation" Stencil Construction

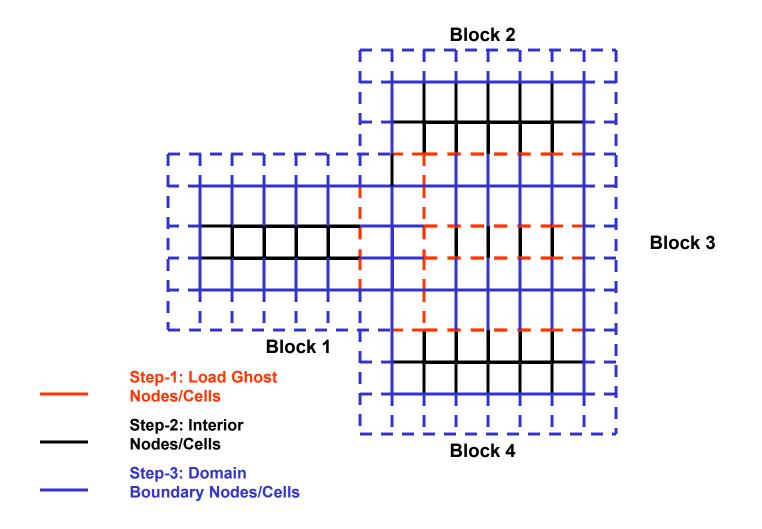


Accumulation

Method-3: "Ghost Cell/Node" Stencil Construction

- Pseudo-Code for multi-block, structured-grid simulations using method-3:
 - Create enlarged-block data-structure that has additional planes beyond each face (halo or ghost nodes/cells)
 - Number of additional planes beyond each face is (stencil_size-1)/2
 - Read original blocks into enlarged-block data structure
 - Loop over enlarged-blocks
 - At inter-block boundaries, pack halo/ghost nodes/cell information with appropriate neighbor information
 - At domain boundaries, create appropriate halo/ghost node/cell information for boundary condition treatment
 - Loop over enlarged-blocks
 - Loop to construct discretization of stencil at all original block nodes/cells
 - Perform boundary conditions on all non-inter-block faces
 - Update variables

Example: "Ghost Cell/Node" Stencil Construction

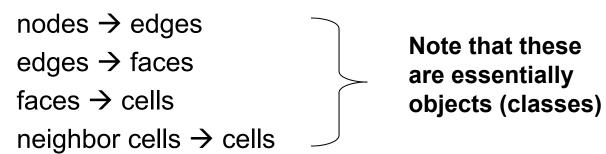


Multi-Block Unstructured-Grid Treatments

- The multi-block data structure and inter-block stencil construction for unstructured-grids is similar
 - Methods 1-3 can be constructed for unstructured-grids
 - Except instead of looping over faces/corners, the boundary is treated as an unstructured column vector of cells/nodes
 - The indirect-addressing associated with the randomness of the grid cells/nodes makes the treatments more tedious, time-consuming, and memory-consuming
 - For unstructured-grids, spatial/grid decomposition is often combined with processor/parallel decomposition
 - Unstructured grid-generation tools often create the grid over the entire domain as a single block

Hybrid Data Structures for Multi-Block Solution Procedures

- Unstructured-grid procedures have random node numbering for each cell
 - This allows for cells to be created arbitrarily to fill in a domain and straightforward cell division for grid-adaptation schemes
 - However, it requires a data structure that has:



Thus the data structure between the nodes and the cells can be multi-level which leads to expensive indirect addressing of memory

- We can greatly reduce this overhead by treating blocks in an unstructured grid format while keeping the cells inside of the block in a structured-grid format
 - This is the concept of "multi-block" structured grids

Multi-Block Structured Grids

So the advantage of multi-block structured grids are:

- Practically any geometry can be gridded using unstructured-grid like block topologies that are connected using a block-neighbor data structure
 - The resulting data structure does not need all of the overhead of a totally unstructured-grid procedure
 - Only neighbor information is necessary since the nodes on adjacent faces of different blocks can be easily mapped
 - Lends very well to both distributed- and shared-memory parallel computer architectures
- The operators inside of the block are all structured using direct addressing of memory and rapid indexing
 - Lends very well to vector and/or graphical processors
- Grid adaptation can be performed uniformly and directionally on a per-block or per-sub-block basis

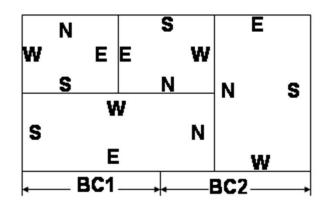
Multi-Block Structured Grids – Additional Considerations

Additional information must now be added to singleblock solvers to allow for multi-block topologies:

- For each given block, store:
 - Global block number
 - Block type (in multi-disciplinary problems, "solid" or "fluid")
 - This could also include "overlaid", "structured", "unstructured"
 - Processor number (parallel) the block belongs to
 - Possibly the orientation
- For each given block face, store:
 - The start and end indices (note that start index can be larger than end index to reverse order of processing)
 - The adjacent global block number (if one exists)
 - The face of the adjacent block (if one exists)
 - The orientation of the adjacent block OR the start and end indices in the directions of the adjacent block

Possibility of Sub-Faces

- In many cases, you will find that it is inconvenient topology-wise to create blocks in such a fashion that the block faces have a single boundary condition over the entire face
- Multiple sub-faces can be used to allow for:
 - Blocks to be staggered
 - Multiple physical boundary conditions to be applied along a face
- This de-couples the block topology from the physical boundary conditions



Sub-faces are not necessary for your projects.
However, they are handy for general purpose procedures.

Sub-Face Considerations

Incorporating sub-faces into a solver requires

- Number of sub-faces for each face (S,N,W,E,H,Y) as part of the stored block information
- Face information:
 - Start and stop indices (in each direction) of the sub-face
 - Must have the capability to traverse in positive or negative i-, j-, kdirections (ie start index can be larger than stop index)
 - Note that traversing of face information MUST be consistent at inter-block boundaries since message passing has an order to it
 - Physical boundary condition type or
 - Along with any physical information regarding that boundary condition (eg total pressure, total temperature, etc)
 - Neighboring block number
 - Along with the start and stop indices (in each direction) that is point-matched to the given sub-face. Note that traverse direction of data in neighbor must be consistent with your own face traverse direction.

Blocks and Faces as Objects (Abstraction)

- When coding up your projects, you could think of block faces abstractly (generically).
 - You could create a (linked) list of faces that have
 - The list could consist of pointers to the temperature of the face nodes
 - Physical boundary conditions
 - Inter-block boundary treatments
 - For each face in the face list, you could store
 - Start and end point indices to loop over
 - For physical boundaries, the boundary condition type
 - · For inter-block boundaries.
 - Neighbor block, face, and processor numbers
 - Indices of start and end points to loop over in each direction

This is the unstructured-grid manner of processing blocks and faces

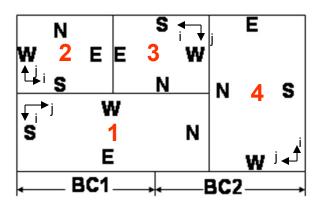
- You could also think of blocks abstractly (generically).
 - You could create a (linked) list of blocks that have
 - Block type (for your projects, you only have solid blocks)
 - Processor number
 - · Dimensions in each direction
 - Number of sub-faces in each direction

Blocks and Faces

- OR you could simply loop over the south, north, west, and east faces
 - Performing kernels for
 - Physical boundary conditions
 - Inter-block boundary conditions
 - This tends to produce repeatable sections of code, however, since the faces are doing similar tasks
- AND loop over the blocks in a Cartesian grid system (since in your problem, the grid is broken up into N x M blocks)
 - This is NOT typical, however, so the creation of a list of blocks is more general

Options for Orientation

- Keeping track of block orientation is only important in terms of neighbor information to find out what neighbor face is adjacent to your face
- Block orientation can be performed either by
 - Keeping track of how a given face maps to a standard-oriented block face
 - Let standard i- be 1 and standard j- be 2.
 - block 2 would have orientation of 1,2
 - block 1 would then have an orientation of -2,1.
 - block 3 would then have an orientation of -1,-2
 - block 4 would then have an orientation of 2,-1
 - BE CAREFUL! All blocks should be right-handed to make integrations come out correct.
 - OR Keeping track of start and end indices of each face (or sub-face) and matching them to neighbor correctly
 - Note that block 2's east face is ordered in the opposite direction as block 3's east face



TAGS for Send/Receiving Messages at Inter-Block Boundaries

- Tags of Send and Receive Messages must be unique AND must be the same on the Send and Receive sides
- Combinatorial mathematics can be used to determine a unique TAG number:
 - A combination of 2 positive, non-zero numbers A and B will be unique if A+BN is unique. For this to happen, N=max(A).
 - A combination of 3 positive, non-zero numbers A, B, and C will be unique if A+BN+CM is unique or when N=max(A) and M=max(A)+max(B)max(A).
 - A combination of 4 positive, non-zero numbers A, B, C, and D will be unique if A+BN+CM+DQ is unque or when N=max(A), M=max(A)+max(B)max(A), and Q=max(A)+max(B)max(A)+(max(A)+max(B)max(A))max(C) or Q=(max(A)+max(B)max(A))(1+max(C))
 - and so on...

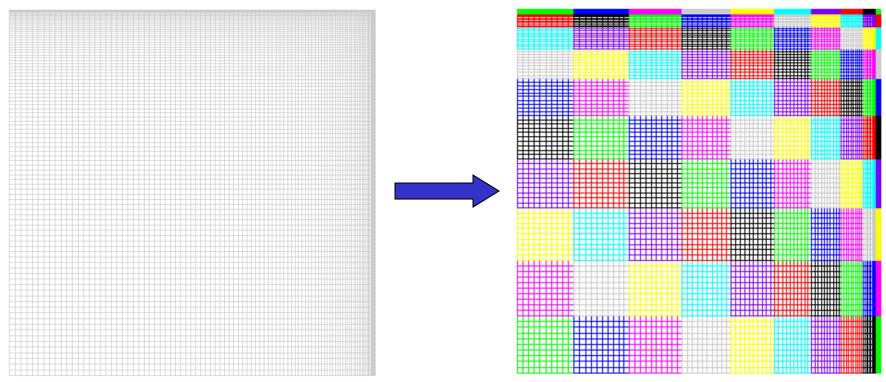
TAGS

- The number of unique tags available will depend on the number of variables that you use, A, B, C, D, etc.
- For your projects,
 - A could be your block number
 - B could be your face number (S=1, N=2, W=3, E=4)
 - C could be your neighbor's block number
 - D could be your neighbor's face number that is adjacent to you
 - Etc.
- There are other combination of variables that can be used to make up TAGS
- Try to keep the TAG formulas as simple as possible that give you the required number of unique values

- Now that we have shown that we can develop a singleblock engineering simulation code, let's break up the problem into multiple blocks.
- For your simulation code developed under project-1, you need to write a code that will break up the domain into multiple blocks that will still run on a single processor
- This is an intermediate step prior to developing a code for multiple processors

- For the Heat Conduction (Default) problem, develop a code that will divide the domain into N x M blocks, each with dimensions of 1+(IMAX-1)/N x 1+(JMAX-1)/M
 - Where IMAX = JMAX = 101 and 501 for default heat-conduction problem
 - The code should be completely general so that N and M can be arbitrary
- Develop a multi-block data structure that writes out a connectivity file with the following information for each block:
 - Global block number of each block
 - Neighbor block numbers of each block
 - Boundary conditions for each side of block
 - · Could be unique to each side
 - Could have regions for each side
 - Number of sub-regions for each side (if you decide to have sub-regions)
 - Dimensions of each sub-region (start and stop indices)
 - Boundary condition for each sub-region
 - Could be defined point-by-point
 - Block orientation (all blocks will probably have the same orientation for our problems).
- And multi-block grid and temperature files (Plot3D or other format) that has the coordinates and initial temperature along with
 - Number of global blocks
 - Block dimensions of each global block

You should end up with something similar to:



in the x', y' frame with associated connectivity and boundary condition file(s)

 Project-3 will be to write a multi-block serial code that will solve your simulation problem

Due Tuesday 10/27:

- An overview of your simulation problem
 - Describe the problem, algorithm, and boundary conditions
- Listing of your spatial/grid decomposition code
- A plot of your decomposed computational grid (or data decomposition) for two-different decompositions:
 - 10 x 10 (N=10, M=10) grid for heat-conduction problem
 - 5 x 4 (N=5, M=4) grid for heat-conduction problem
 - Using the 101 x 101 and 501 x 501 grids
- A sample listing of your connectivity/boundary condition file(s)