

# SIG Proceedings Paper in LaTeX Format\*

Extended Abstract<sup>†</sup>

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This paper provides a sample of a  $\LaTeX$  document which conforms, somewhat loosely, to the formatting guidelines for ACM SIG Proceedings.<sup>1</sup>

CCS Concepts: • **Computer systems organization** → **Embedded systems**; *Redundancy*; Robotics; • **Networks** → Network reliability;

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## 1 PATHS AND ALIASING

Model dynamics using judgement of the form:

$$V, H, R, F \vdash_{P:\Sigma} e \Downarrow v, H', F'$$

Where  $V : \text{Var} \rightarrow \text{Val} \times \text{State}$ ,  $H : \text{Loc} \rightarrow \text{Val}$ ,  $R \subseteq \text{Loc}$ ,  $F \subseteq \text{Loc}$ , and  $\Sigma : \text{Var} \rightarrow \text{FTypes}$ . This can be read as: under stack  $V$ , heap  $H$ , roots  $R$ , freelist  $F$ , and program  $P$  with signature  $\Sigma$ , the expression  $e$  evaluates to  $v$ , and engenders a new heap  $H'$  and freelist  $F'$ .

A *program* is then a  $\Sigma$  indexed map  $P$  from  $\text{Var}$  to pairs  $(y_f, e_f)_{f \in \Sigma}$ , where  $\Sigma(y_f) = A \rightarrow B$ , and  $\Sigma; y_f : A \vdash e_f : B$  (typing rules are discussed in 6). We write  $P : \Sigma$  to mean  $P$  is a program with signature  $\Sigma$ . Because the signature  $\Sigma$  for the mapping of function names to first order functions does not change during evaluation, we drop the subscript  $\Sigma$  from  $\vdash_\Sigma$  when the context of evaluation is clear. It is convenient to think of the evaluation judgement  $\vdash$  as being indexed by a family of signatures  $\Sigma$ 's, each of which is a set of “top-level” first-order declarations to be used during evaluation.

For a partial map  $f : A \rightarrow B$ , we write  $\text{dom}$  for the defined values of  $f$ . Sometimes we shorten  $x \in \text{dom}(f)$  to  $x \in f$ . We write  $f[x \mapsto y]$  for the extension of  $f$  where  $x$  is

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<sup>†</sup>The full version of the author's guide is available as `acmart.pdf` document

<sup>‡</sup>This author is the one who did all the really hard work.

<sup>1</sup>This is an abstract footnote

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BTypes	$\tau ::=$		
	nat	nat	naturals
	unit	unit	unit
	bool	bool	boolean
	prod( $\tau_1; \tau_2$ )	$\tau_1 \times \tau_2$	product
	list( $\tau$ )	$L(\tau)$	list
FTypes	$\rho ::=$		
	arr( $\tau_1; \tau_2$ )	$\tau_1 \rightarrow \tau_2$	first order function
Exp	$e ::=$		
	$x$	$x$	variable
	nat[ $n$ ]	$\bar{n}$	number
	unit	()	unit
	T	T	true
	F	F	false
	if( $x; e_1; e_2$ )	if $x$ then $e_1$ else $e_2$	if
	lam( $x : \tau.e$ )	$\lambda x : \tau.e$	abstraction
	ap( $f; x$ )	$f(x)$	application
	tpl( $x_1; x_2$ )	$\langle x_1, x_2 \rangle$	pair
	case( $x_1, x_2.e_1$ )	case $p \{ (x_1; x_2) \hookrightarrow e_1 \}$	match pair
	nil	[]	nil
	cons( $x_1; x_2$ )	$x_1 :: x_2$	cons
	case{l}( $e_1; x, xs.e_2$ )	case $l \{ \text{nil} \hookrightarrow e_1 \mid \text{cons}(x; xs) \hookrightarrow e_2 \}$	match list
	let( $e_1; x : \tau.e_2$ )	let $x = e_1$ in $e_2$	let
	share( $x; x_1, x_2.e$ )	share $x$ as $x_1, x_2$ in $e$	share
Val	$v ::=$		
	val( $n$ )	$n$	numeric value
	val(T)	T	true value
	val(F)	F	false value
	val(Null)	Null	null value
	val(cl( $V; x.e$ ))	$(V, x.e)$	function value
	val( $l$ )	$l$	loc value
	val(pair( $v_1; v_2$ ))	$\langle v_1, v_2 \rangle$	pair value
State	$s ::=$		
	alive	alive	live value
	dead	dead	dead value
Loc	$l ::=$		
	loc( $l$ )	$l$	location
Var	$l ::=$		
	var( $x$ )	$x$	variable

mapped to  $y$ , with the constraint that  $x \notin \text{dom}(f)$ .

Roots represents the set of locations required to compute the continuation *excluding* the current expression. We can think of roots as the heap allocations necessary to compute the context with a hole that will be filled by the current expression.

In order to prove soundness of the type system, we need some auxiliary judgements to defining properties of a heap. Below we define  $\text{reach} : \text{Val} \rightarrow \{\{\text{Loc}\}\}$  that maps stack values to the root *multiset*, the multiset of locations that's already on the stack. Next we define reachability of values:

$$\begin{aligned}\text{reach}_H(\langle v_1, v_2 \rangle) &= \text{reach}_H(v_1) \uplus \text{reach}_H(v_2) \\ \text{reach}_H(l) &= \{l\} \uplus \text{reach}_H(H(l)) \\ \text{reach}_H(\_) &= \emptyset\end{aligned}$$

For a multiset  $S$ , we write  $\mu_S : S \rightarrow \mathbb{N}$  for the multiplicity function of  $S$ , which maps each element to the count of its occurrence. If  $\mu_S(x) \geq 1$  for a multiset  $S$ , then we write  $x \in S$  as in the usual set membership relation. If for all  $s \in S$ ,  $\mu(s) = 1$ , then  $S$  is a property set, and we denote it by  $\text{set}(S)$ . Additionally,  $A \uplus B$  denotes counting union of sets where  $\mu_{A \uplus B}(s) = \mu_A(s) + \mu_B(s)$ , and  $A \cup B$  denotes the usual union where  $\mu_{A \cup B}(s) = \max(\mu_A(s), \mu_B(s))$ . For the disjoint union of sets  $A$  and  $B$ , we write  $A \sqcup B$ .

Next, we define the predicates `no_alias`, `stable`, and `disjoint`:

`no_alias(V, H)`:  $\forall x, y \in V, x \neq y$ . Let  $r_x = \text{reach}_H(V(x))$ ,  $r_y = \text{reach}_H(V(y))$ .

Then:

(1)  $\text{set}(r_x), \text{set}(r_y)$

(2)  $r_x \cap r_y = \emptyset$

`stable(R, H, H')`:  $\forall l \in R. H(l) = H'(l)$ .

`safe(V, H, F)`:  $\forall x \in V. \text{reach}_H(V(x)) \cap F = \emptyset$

`disjoint(C)`:  $\forall X, Y \in C. X \cap Y = \emptyset$

For a stack  $V$  and a heap  $H$ , whenever `no_alias(V, H)` holds, visually, one can think of the situation as the following: the induced graph of heap  $H$  with variables on the stack as additional leaf nodes is a forest: a disjoint union of arborescences (directed trees); consequently, there is at most one path from a live variable on the stack  $V$  to a location in  $H$  by following the pointers.

First, we define  $FV^\star(e)$ , the multiset of free variables of  $e$ . As the usual  $FV$ , it is defined inductively over the structure of  $e$ ; the only unusual thing is that multiple occurrences of a free variable  $x$  in  $e$  will be reflected in the multiplicity of  $FV^\star(e)$ .

Next, we define  $\text{locs}_{V, H}$  using the previous notion of reachability.

$$\text{locs}_{V, H}(e) = \bigcup_{x \in FV(e)} \text{reach}_H(V(x))$$

*size* calculates the *literal size* of a value, e.g. the size to store its address.

$$\begin{aligned} \text{size}(\langle v_1, v_2 \rangle) &= \text{size}(v_1) + \text{size}(v_2) \\ \text{size}(\_) &= 1 \end{aligned}$$

Let  $\text{card}(S)$  denote the number of unique elements, e.g. the cardinality of a multiset  $S$ . We write  $\|v\|_H$  for  $\text{card}(\text{reach}_H(v))$

As usual, we extend it to stacks  $V$ :  $\|V\|_H = \sum_{V(x)=v} \|v\|_H$

$\text{copy}(H, L, v)$  takes a heap  $H$ , a set of locations  $L$ , and a value  $v$ , and returns a new heap  $H'$  and a location  $l$  such that  $l$  maps to  $v$  in  $H'$ .

$$\begin{aligned} \text{copy}(H, L, \langle v_1, v_2 \rangle) &= \\ \text{let } L_1 \sqcup L_2 &\subseteq L \\ \text{where } |L_1| &= \|v_1\|_H, |L_2| = \|v_2\|_H \\ \text{let } H_1, v'_1 &= \text{copy}(H, L_1, v_1) \\ \text{let } H_2, v'_2 &= \text{copy}(H_1, L_2, v_2) \text{ in} \\ H_2, \langle v'_1, v'_2 \rangle & \\ \text{copy}(H, L, l) &= \\ \text{let } l' \in L &\text{ in} \\ \text{let } H', v &= \text{copy}(H, L \setminus \{l'\}, H(l)) \text{ in} \\ H' \{l' \mapsto v\}, l' & \\ \text{copy}(H, L, v) &= \\ H, v & \end{aligned}$$



## 2 GARBAGE COLLECTION SEMANTICS

$$\begin{array}{c}
\frac{V(x) = v}{V, H, R, F \vdash x \Downarrow v, H, F} (S_1) \qquad \frac{}{V, H, R, F \vdash \bar{n} \Downarrow \text{val}(n), H, F} (S_2) \\
\\
\frac{}{V, H, R, F \vdash \top \Downarrow \text{val}(\top), H, F} (S_3) \qquad \frac{}{V, H, R, F \vdash F \Downarrow \text{val}(F), H, F} (S_4) \\
\\
\frac{}{V, H, R, F \vdash () \Downarrow \text{val}(\text{Null}), H, F} (S_5) \\
\\
\frac{V = V'[x \mapsto \top] \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V, H}(e_1)\} \quad V', H, R, F \cup g \vdash e_1 \Downarrow v, H', F'}{V, H, R, F \vdash \text{if}(x; e_1; e_2) \Downarrow v, H', F'} (S_6) \\
\\
\frac{V = V'[x \mapsto F] \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V, H}(e_2)\} \quad V', H, R, F \cup g \vdash e_2 \Downarrow v, H', F'}{V, H, R, F \vdash \text{if}(x; e_1; e_2) \Downarrow v, H', F'} (S_7) \\
\\
\frac{V(x) = v' \quad P(f) = (y_f, e_f) \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V, H}(e_f)\} \quad [y_f \mapsto v'], H, R, F \cup g \vdash e_f \Downarrow v, H', F'}{V, H, R, F \vdash f(x) \Downarrow v, H', F'} (S_8) \\
\\
\frac{V(x_1) = v_1 \quad V(x_2) = v_2}{V, H, R, F \vdash \langle x_1, x_2 \rangle \Downarrow \langle v_1, v_2 \rangle, H, F} (S_9) \\
\\
\frac{V = V'[x \mapsto \langle v_1, v_2 \rangle] \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V, H}(e)\} \quad V'' = V'[x_1 \mapsto v_1, x_2 \mapsto v_2] \quad V'', H, R, F \cup g \vdash e \Downarrow v, H', F'}{V, H, R, F \vdash \text{case } x \{ (x_1; x_2) \hookrightarrow e \} \Downarrow v, H', F'} (S_{10}) \\
\\
\frac{}{V, H, R, F \vdash \text{nil} \Downarrow \text{val}(\text{Null}), H, F} (S_{11}) \\
\\
\frac{v = \langle V(x_1), V(x_2) \rangle \quad H' = H\{l \mapsto v\}}{V, H, R, F \vdash \text{cons}(x_1; x_2) \Downarrow l, H', F'} (S_{12}) \\
\\
\frac{V = V'[x \mapsto v'] \quad g = \text{reach}_H(v') \quad V', H, R, F \cup g \vdash e \Downarrow v, H', F'}{V, H, R, F \vdash \text{drop}(x; e) \Downarrow v, H', F'} (S_{13}) \\
\\
\frac{V = V'[x \mapsto v'] \quad L \subseteq F \quad |L| = \|v'\|_H \quad H', v'' = \text{copy}(H, L, v') \quad V'[x_1 \mapsto v', x_2 \mapsto v''], H', R, F \setminus L \vdash e \Downarrow v, H'', F'}{V, H, R, F \vdash \text{shareCopy } x \text{ as } x_1, x_2 \text{ in } e \Downarrow v, H'', F'} (S_{14}) \\
\\
\frac{V(x) = \text{Null} \quad V' \subseteq V \quad \text{dom}(V') = FV(e_1) \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V, H}(e_1)\} \quad V', H, R, F \cup g \vdash e_1 \Downarrow v, H', F'}{V, H, R, F \vdash \text{case } x \{ \text{nil} \hookrightarrow e_1 \mid \text{cons}(x_h; x_t) \hookrightarrow e_2 \} \Downarrow v, H', F'} (S_{15}) \\
\\
\frac{V(x) = l \quad H(\hat{l}) = \langle v_h, v_t \rangle \quad V' \subseteq V \quad \text{dom}(V') = FV(e_2) \setminus \{x_h, x_t\} \quad V'' = V'[x_h \mapsto v_h, x_t \mapsto v_t] \quad g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V'', H}(e_2)\} \quad V'', H, R, F \cup g \vdash e_2 \Downarrow v, H', F'}{V, H, R, F \vdash \text{case } x \{ \text{nil} \hookrightarrow e_1 \mid \text{cons}(x_h; x_t) \hookrightarrow e_2 \} \Downarrow v, H', F'} (S_{16}) \\
\\
\frac{}{V = V_1 \sqcup V_2 \quad \text{dom}(V_1) = FV(e_1)}
\end{array}$$

### 3 OPERATIONAL SEMANTICS

In order to prove the soundness of the type system, we also define a simplified operational semantics that does not account for garbage collection.

$$\boxed{V, H \vdash e \Downarrow v, H'}$$

This can be read as: under stack  $V$ , heap  $H$  the expression  $e$  evaluates to  $v$ , and engenders a new heap  $H'$ . We write the representative rules.

$$\frac{v = \langle V(x_1), V(x_2) \rangle \quad (L \sqcup \{l\}) \cap \text{dom}(H) = \emptyset \quad H', l = \text{copy}(H, L, v)}{V, H \vdash \text{cons}(x_1; x_2) \Downarrow l, H'} \text{(S}_{18}\text{)}$$

$$\frac{\begin{array}{l} V(x) = l \quad H(l) = \langle v_h, v_t \rangle \quad V' \subseteq V \quad \text{dom}(V') = FV(e_2) \setminus \{x_h, x_t\} \\ V'' = V'[x_h \mapsto v_h, x_t \mapsto v_t] \quad V'', H \vdash e_2 \Downarrow v, H' \end{array}}{V, H \vdash \text{case } x \{ \text{nil} \hookrightarrow e_1 \mid \text{cons}(x_h; x_t) \hookrightarrow e_2 \} \Downarrow v, H'} \text{(S}_{19}\text{)}$$

$$\frac{\begin{array}{l} V = V_1 \sqcup V_2 \quad \text{dom}(V_1) = FV(e_1) \quad \text{dom}(V_2) = FV(\text{lam}(x : \tau.e_2)) \\ V_1, H \vdash e_1 \Downarrow v_1, H_1 \quad V'_2 = V_2[x \mapsto v_1] \quad V'_2, H_1 \vdash e_2 \Downarrow v_2, H_2 \end{array}}{V, H \vdash \text{let}(e_1; x : \tau.e_2) \Downarrow v_2, H_2} \text{(S}_{20}\text{)}$$

$$\frac{\begin{array}{l} V = V'[x \mapsto v'] \quad L \text{ fresh} \\ |L| = \|v'\|_H \quad H', v'' = \text{copy}(H, L, v') \quad V'[x_1 \mapsto v', x_2 \mapsto v''], H' \vdash e \Downarrow v, H'' \end{array}}{V, H \vdash \text{shareCopy } x \text{ as } x_1, x_2 \text{ in } e \Downarrow v, H''} \text{(S}_{21}\text{)}$$

### 4 WELL DEFINED ENVIRONMENTS

In order to define the potential for first-order types, we need a notion of well-define environments, one that relates heap values to semantic values of a type. We first give a denotational semantics for the first-order types:

$$\begin{aligned} () &\in \llbracket \text{unit} \rrbracket \\ \perp &\in \llbracket \text{bool} \rrbracket \\ \top &\in \llbracket \text{bool} \rrbracket \\ 0 &\in \llbracket \text{nat} \rrbracket \\ n + 1 &\in \llbracket \text{nat} \rrbracket \text{ if } n \in \llbracket \text{nat} \rrbracket \\ [] &\in \llbracket L(A) \rrbracket \\ \pi(a, l) &\in \llbracket L(A) \rrbracket \text{ if } a \in \llbracket A \rrbracket \text{ and } l \in \llbracket L(A) \rrbracket \end{aligned}$$

Where semantic set for each type is the least set such that the above holds. Note  $\pi(x, y)$  is the usual set-theoretic pairing function, and write  $[a_1, \dots, a_n]$  for  $\pi(a_1, \dots, \pi(a_n, []))$ .

Now we give the judgements relating heap values to semantic values, in the form  $\boxed{H \models v \mapsto a : A}$ , which can be read as: under heap  $H$ , heap value  $v$  defines the semantic value  $a \in \llbracket A \rrbracket$ .

$$\begin{array}{c}
\frac{n \in \mathbb{Z}}{H \models n \mapsto n : \text{nat}} (\text{V:ConstI}) \qquad \frac{}{H \models \text{Null} \mapsto n : \text{unit}} (\text{V:ConstI}) \\
\\
\frac{A \in \text{BType}}{H \models \text{Null} \mapsto n : L(A)} (\text{V:Nil}) \qquad \frac{}{H \models \top \mapsto \top : \text{bool}} (\text{V:True}) \\
\\
\frac{}{H \models \text{F} \mapsto \perp : \text{bool}} (\text{V:False}) \\
\\
\frac{l \in \text{Loc} \quad H(l) = \langle v_h, v_t \rangle \quad H \models v_h \mapsto a_1 : A \quad H \models v_t \mapsto [a_2, \dots, a_n] : L(A)}{H \models l \mapsto [a_1, \dots, a_n] : L(A)} (\text{V:Cons})
\end{array}$$

## 5 STACK VS HEAP ALLOCATED TYPES

In order to share variables, we need to distinguish between types that are allocated on the stack and the heap. We write  $\boxed{\text{stack}(A)}$  to denote that values of type  $A$  will be allocated *entirely* on the stack at run time (no references into the heap).

$$\begin{array}{c}
\frac{A \in \{\text{unit}, \text{bool}, \text{nat}\}}{\text{stack}(A)} (\text{S:Const}) \qquad \frac{\text{stack}(A_1) \quad \text{stack}(A_2)}{\text{stack}(A_1 \times A_2)} (\text{S:Product})
\end{array}$$

## 6 LINEAR GARBAGE COLLECTION TYPE RULES

The linear version of the type system takes into account of garbage collected cells by returning potential locally in a match construct. Since we are interested in the number of heap cells, all constants are assumed to be nonnegative. The second let rule expresses the fact that since stack types don't reference heap cells, any heap cells used in the evaluation of  $e_1$  can be deallocated, as there are no longer references to them in  $v_1$ .



$$\begin{array}{c}
\frac{n \in \mathbb{Z}}{\Sigma; \emptyset \mid \frac{q}{q} n : \text{nat}} \text{(L:ConstI)} \qquad \frac{}{\Sigma; \emptyset \mid \frac{q}{q} () : \text{unit}} \text{(L:ConstU)} \\
\\
\frac{}{\Sigma; \emptyset \mid \frac{q}{q} \top : \text{bool}} \text{(L:ConstT)} \qquad \frac{}{\Sigma; \emptyset \mid \frac{q}{q} \text{F} : \text{bool}} \text{(L:ConstF)} \\
\\
\frac{}{\Sigma; x : B \mid \frac{q}{q} x : B} \text{(L:Var)} \qquad \frac{\Sigma(f) = A \xrightarrow{q/q'} B}{\Sigma; x : A \mid \frac{q}{q'} f(x) : B} \\
\\
\frac{\Sigma; \Gamma \mid \frac{q}{q'} e_t : B \quad \Sigma; \Gamma \mid \frac{q}{q'} e_f : B}{\Sigma; \Gamma, x : \text{bool} \mid \frac{q}{q'} \text{if } x \text{ then } e_t \text{ else } e_f : B} \text{(L:Cond)} \\
\\
\frac{}{\Sigma; x_1 : A_1, x_2 : A_2 \mid \frac{q}{q} \langle x_1, x_2 \rangle : A_1 \times A_2} \text{(L:Pair)} \\
\\
\frac{\Sigma; \Gamma, x_1 : A_1, x_2 : A_2 \mid \frac{q}{q} e : B}{\Sigma; \Gamma, x : (A_1, A_2) \mid \frac{q}{q'} \text{case } x \{ (x_1; x_2) \hookrightarrow e \} : B} \text{(L:MatP)} \\
\\
\frac{}{\Sigma; \emptyset \mid \frac{q}{q} \text{nil} : L^p(A)} \text{(L:Nil)} \\
\\
\frac{}{\Sigma; x_h : A, x_t : L^p(A) \mid \frac{q+p+1}{q} \text{cons}(x_h; x_t) : L^p(A)} \text{(L:Cons)} \\
\\
\frac{\Sigma; \Gamma \mid \frac{q}{q'} e_1 : B \quad \Sigma; \Gamma, x_h : A, x_t : L^p(A) \mid \frac{q+p+1}{q'} e_2 : B}{\Sigma; \Gamma, x : L^p(A) \mid \frac{q}{q'} \text{case } x \{ \text{nil} \hookrightarrow e_1 \mid \text{cons}(x_h; x_t) \hookrightarrow e_2 \} : B} \text{(L:MatL)} \\
\\
\frac{\Sigma; \Gamma_1 \mid \frac{q}{p} e_1 : A \quad \Sigma; \Gamma_2, x : A \mid \frac{p}{q'} e_2 : B}{\Sigma; \Gamma_1, \Gamma_2 \mid \frac{q}{q'} \text{let}(e_1; x : \tau.e_2) : B} \text{(L:Let)} \\
\\
\frac{\Sigma; \Gamma_1 \mid \frac{q}{p} e_1 : A \quad \text{stack}(A) \quad \Sigma; \Gamma_2, x : A \mid \frac{\max(p, q)}{q'} e_2 : B}{\Sigma; \Gamma_1, \Gamma_2 \mid \frac{q}{q'} \text{let}(e_1; x : \tau.e_2) : B} \text{(L:LetS)} \\
\\
\frac{\Sigma; \Gamma \mid \frac{q}{q'} e : B}{\Sigma; \Gamma, x : A \mid \frac{q}{q'} \text{drop}(x; e) : B} \text{(L:Drop)} \\
\\
\frac{\text{stack}(A) \quad \Sigma; \Gamma, x_1 : A, x_2 : A \mid \frac{q}{q'} e : B}{\Sigma; \Gamma, x : A \mid \frac{q}{q'} \text{share } x \text{ as } x_1, x_2 \text{ in } e : B} \text{(L:Share)}
\end{array}$$

Now if we take  $\dagger : L^p(A) \mapsto L(A)$  as the map that erases resource annotations, we obtain a simpler typing judgement  $\boxed{\Sigma^\dagger; \Gamma^\dagger \vdash e : B^\dagger}$ .

## 7 TYPE RULES FOR SHARING

$$\begin{aligned} L^p(A) - n &= L^{\max(p-n, 0)}(A - n) \\ A_1 \times A_2 - n &= A_1 - n \times A_2 - n \\ A - n &= A \end{aligned}$$

$$\frac{A \curlyvee A_1, A_2, 1 \quad \Sigma; \Gamma, x_1 : A_1, x_2 : A_2 \mid \frac{q}{q'} e : B}{\Sigma; \Gamma, x : A \mid \frac{q}{q'} \text{ share } x \text{ as } x_1, x_2 \text{ in } e : B} \text{(M:Share)}$$

$$\frac{\Sigma; \Gamma_1 \mid \frac{q}{p} e_1 : A \quad \Sigma; \Gamma_1 \mid \text{cf} e_1 : A' \quad \Sigma; \Gamma_2, x : (A' - 1) \mid \frac{p}{q'} e_2 : B}{\Sigma; \Gamma_1, \Gamma_2 \mid \frac{q}{q'} \text{ let}(e_1; x : \tau.e_2) : B} \text{(M:Let)}$$

Where  $A \curlyvee A_1, A_2, n$  is the sharing relation defined as:

$$\begin{aligned} L^p(A) \curlyvee L^q(A_1), L^r(A_2), n & \quad \text{if } p = q + r + n \text{ and } A \curlyvee A_1, A_2, n \\ A \curlyvee A_1, A_2 & \quad \text{if } \text{stack}(A), \text{stack}(A_1), \text{stack}(A_2) \text{ and } A \equiv A_1 \equiv A_2 \end{aligned}$$

## 8 SOUNDNESS FOR LINEAR GC

We simplify the soundness proof of the type system for the general metric to one with monotonic resource. (No function types for now)

**Definition 8.1** (Well-formed computation). When considering the input mode arguments of a evaluation judgment  $V, H, R, F \vdash e \Downarrow v, H', F'$ , we say the 5-tuple  $(V, H, R, F, e)$  is a *well-formed computation* given the following:

- (1)  $\text{dom}(V) = \text{FV}(e)$
- (2)  $\text{no\_alias}(V, H)$ , and
- (3)  $\text{disjoint}(\{R, F, \text{loc}_{SV, H}(e)\})$

And we write  $\text{wfc}(V, H, R, F, e)$  to denote this fact.

LEMMA 1.1. *If  $\Sigma; \Gamma \mid \frac{q}{q} e : B$ , then  $\Sigma^\dagger; \Gamma^\dagger \vdash e : B^\dagger$ .*

LEMMA 1.2. *If  $\Sigma; \Gamma \mid \frac{q}{q} e : B$ , then  $\text{set}(\text{FV}^\star(e))$  and  $\text{dom}(\Gamma) = \text{FV}(e)$ .*

PROOF. Induction on the typing judgement. □

LEMMA 1.3. *Let  $H \models v \mapsto a : A$ . For all sets of locations  $R$ , if  $\text{reach}_H(v) \subseteq R$  and  $\text{stable}(R, H, H')$ , then  $H' \models v \mapsto a : A$  and  $\text{reach}_H(v) = \text{reach}_{H'}(v)$ .*

PROOF. Induction on the structure of  $v$ . □

COROLLARY 1.3.1. *Let  $H \models V : \Gamma$ . For all sets of locations  $R$ , if  $\bigcup_{x \in V} \text{reach}_H(V(x)) \subseteq R$  and  $\text{stable}(R, H, H')$ , then  $H' \models V : \Gamma$ .*

PROOF. Follows from Lemma 1.3. □

LEMMA 1.4. *Let  $H \models v \mapsto a : A$ . If  $\text{stack}(A)$ , then  $\Phi_H(v : A) = 0$ .*

PROOF. Induction on  $H \models v \mapsto a : A$ . □

LEMMA 1.5 (HEAP CONSERVATION). *Let  $\text{wfc}(V, H, R, F, e)$ ,  $V, H, R, F \vdash e \Downarrow v, H', F'$ , and  $g = \text{gc}(H', R, F')$ . Then  $\|V\|_H + |F| \leq \|v\|_{H'} + |F' \cup g|$ .*

PROOF. Induction on evaluation.

### Case 1: E:Var

$$\begin{aligned}
 V &= [x \mapsto v] && (\text{since } \text{dom}(V) = \text{FV}(e) = \{x\}) \\
 \|V\|_H &= \|v\|_H && (\text{def of } \|\cdot\|_H) \\
 \|V\|_H + |F| &\leq \|v\|_{H'} + |F \cup g|
 \end{aligned}$$

### Case 2: E:Const\* Due to similarity, we show only for E:ConstI

$$\begin{aligned}
 V &= \emptyset && (\text{since } \text{dom}(V) = \text{FV}(e) = \emptyset) \\
 \|V\|_H &= \|v\|_H && (\text{def of } \|\cdot\|_H) \\
 \|V\|_H + |F| &\leq \|v\|_{H'} + |F \cup g|
 \end{aligned}$$

### Case 4: E:App

**Case 5: E:CondT** Similar to E:MatNil

**Case 6: E:CondF** Similar to E:CondT

**Case 7: E:Let**

$$\|V_1\|_H + |F| \leq \|v_1\|_{H_1} + |F_1 \cup g| \quad (\text{IH on first premise})$$

$$\text{Let } g' = \text{gc}(H_2, R, F_2)$$

$$\|V'_2\|_{H_1} + |F_1 \cup g| \leq \|v_2\|_{H_2} + |F \cup g'| \quad (\text{IH on second premise})$$

$$\|V'_2\|_{H_1} = \|V_2\|_{H_1} + \|v_1\|_{H_1} \quad (\text{definition of semantic size})$$

$$= \|V_2\|_H + \|v_1\|_{H_1} \quad (\text{main lemma})$$

$$\|V_2\|_H + \|v_1\|_{H_1} + |F_1 \cup g| \leq \|v_2\|_{H_2} + |F \cup g'|$$

$$\|V_1\|_H + \|V_2\|_H + \|v_1\|_{H_1} + |F| + |F_1 \cup g| \leq \|v_1\|_{H_1} + \|v_2\|_{H_2} + |F_1 \cup g| + |F \cup g'|$$

$$\|V\|_H + |F| \leq \|v_2\|_{H_2} + |F \cup g'|$$

**Case 8: E:Pair** Similar to E:Var

**Case 9: E:MatP** Similar to E:MatCons

**Case 10: E:Nil** Similar to E:Const\*

**Case 11: E:Cons**

$$V = [x_1 \mapsto v_1, x_2 \mapsto v_2] \quad (\text{since } \text{dom}(V) = FV(e) = \{x_1, x_2\})$$

$$\|V\|_H = \|v_1\|_H + \|v_2\|_H \quad (\text{def of } \|\cdot\|_H)$$

$$\|l\|_{H'} = 1 + \|H'(l)\|_{H'} = 1 + \|v\|_{H''} = 1 + \|v_1\|_{H''} + \|v_2\|_{H''} \quad (\text{def of semantic size})$$

$$= 1 + \|v_1\|_H + \|v_1\|_H$$

$$= 1 + \|V\|_H$$

$$L \sqcup \{l\} \subseteq g \quad (R \cap F = \emptyset \text{ and } L \sqcup \{l\} \subseteq H'')$$

$$|g| \geq |L \sqcup \{l\}| = \text{size}(v) + 1$$

$$|F' \cup g| \geq |F|$$

$$\|V\|_H + |F| \leq \|v\|_{H'} + |F \cup g|$$

**Case 12: E:MatNil**

**Case 13: E:MatCons**

$$\text{Let } g' = \text{gc}(H', R, F')$$

$$\|V''\|_H + |F \cup g| \leq |F' \cup g'| \quad (\text{IH (wfc from main lemma)})$$

$$\|V''\|_H = \|V'[x_h \mapsto v_h, x_t \mapsto v_t]\|_H$$

$$= \|V'\|_H + \|v_h\|_H + \|v_t\|_H$$

$$= \|V'\|_H + \|l\|_H - 1$$

$$= \|V\|_H - 1$$

$$\|V\|_H - 1 + |F \cup g| \leq |F' \cup g'|$$

$$\begin{aligned} \|v\|_H - 1 + |F| + |g| &\leq |F' \cup g'| & (F \cap g = \emptyset) \\ \|v\|_H + |F| &\leq |F' \cup g'| & (|g| \geq 1 \text{ from main lemma}) \end{aligned}$$

**Case 13: E:Drop**

$$\begin{aligned} \text{Let } g' &= \text{gc}(H', R, F') \\ \|V'\|_H + |F \cup g| &\leq \|v\|_{H'} + |F' \cup g'| & (\text{IH}) \\ HV &= \|+\|_{V'} \| \\ v' \|V\|_H - \|+\|_{v'} |F \cup \text{reach}_H(v')| &\leq \|v\|_{H'} + |F' \cup g'| \\ \|V\|_H - \|+\|_{v'} |F| + |\text{reach}_H(v')| &\leq \|v\|_{H'} + |F' \cup g'| \\ \|V\|_H + |F| &\leq \|v\|_{H'} + |F' \cup g'| \end{aligned}$$

**Case 13: E:ShareCopy**

$$\begin{aligned} e &= \text{shareCopy } x \text{ as } x_1, x_2 \text{ in } e & (\text{case}) \\ \text{Let } g' &= \text{gc}(H', R, F') \\ \|V'[x_1 \mapsto v', x_2 \mapsto v'']\|_{H'} + |F \setminus L| &\leq \|v\|_{H'} + |F' \cup g'| & (\text{IH, well-formedness from main lemma}) \\ \|V'[x_1 \mapsto v', x_2 \mapsto v'']\|_{H'} + |F \setminus L| &= \|V\|_H + \|v''\|_{H'} + |F| - |L| & (\text{stability lemma for copy}) \\ &= \|V\|_H + |L| + |F| - |L| & (\text{lemma about copy}) \\ &= \|V\|_H + |F| \end{aligned}$$

□

LEMMA 1.6. *Let  $\Sigma; \Gamma \vdash \frac{q}{q'} e : B$  and  $V, H, R, F \vdash e \Downarrow v, H', F'$ . Then  $\|V\|_H - \|v\|_{H'} + q \geq q'$ .*

LEMMA 1.7 (MAIN LEMMA). *For all stacks  $V$  and heaps  $H$ , let  $V, H, R, F \vdash e \Downarrow v, H', F'$  and  $\Sigma; \Gamma \vdash e : B$ . Then given the following:*

- (1)  $\text{dom}(V) = \text{FV}(e)$
- (2)  $\text{no\_alias}(V, H)$ , and
- (3)  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$

*We have the following:*

- (1)  $\text{set}(\text{reach}_{H'}(v))$
- (2)  $\text{disjoint}(\{R, F', \text{reach}_{H'}(v)\})$ , and
- (3)  $\text{stable}(R, H, H')$

PROOF. Nested induction on the evaluation judgement and the typing judgement.

**Case 1: E:Var**

Suppose  $H \models V : \Gamma, \text{dom}(V) = \text{FV}(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$

$$\begin{array}{ll}
\text{set}(\text{reach}_H(v)) & (\text{no\_alias}(V, H)) \\
\text{disjoint}(\{R, F, \text{reach}_H(v)\}) & (\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})) \\
\text{no\_alias}(V, H) & (\text{Sp.}) \\
\text{stable}(R, H, H') & (H = H')
\end{array}$$

**Case 2: E:Const\*** Due to similarity, we show only for E:ConstI

$$\begin{array}{ll}
\text{Suppose } H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\}) & \\
\text{set}(\text{reach}_H(v)) & (\text{reach}_H(v) = \emptyset) \\
\text{disjoint}(\{R, F, \emptyset\}) & (\text{disjoint}(R, F)) \\
\text{no\_alias}(V, H) & (\text{Sp.}) \\
\text{stable}(R, H, H') & (H = H')
\end{array}$$

**Case 4: E:App**

**Case 5: E:CondT** Similar to E:MatNil

**Case 6: E:CondF** Similar to E:CondT

**Case 7: E:Let**

$$\begin{array}{ll}
V, H, R, F \vdash \text{let}(e_1; x : \tau.e_2) \Downarrow v_2, H_2, F_2 & (\text{case}) \\
V, H, R', F \vdash e_1 \Downarrow v_1, H_1, F_1 & (\text{ad.}) \\
\Sigma; \Gamma_1, \Gamma_2 \vdash \text{let}(e_1; x : \tau.e_2) : B & (\text{case}) \\
\Sigma; \Gamma_1 \vdash e_1 : A & (\text{ad.}) \\
\text{Suppose } H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\}) & \\
H \models V_1 : \Gamma_1 & (\text{def of W.D.E and Lemma 1.2})
\end{array}$$

By IH, we have invariant on  $J_1$

NTS (1) - (3) to instantiate invariant on  $J_1$

$$\begin{array}{ll}
(1) \quad \text{dom}(V_1) = FV(e_1) & (\text{def of } V_1) \\
(2) \quad \text{no\_alias}(V_1, H) & (\text{no\_alias}(V, H) \text{ and } V_1 \subseteq V) \\
(3) \quad \text{disjoint}(R', F, \text{locs}_{V, H}(e_1)) & \\
F \cap R' = \emptyset & (F \cap \text{locs}_{V, H}(e) = \emptyset \text{ and } \text{locs}_{V_2, H}(\text{lam}(x : \tau.e_2)) \subseteq \text{locs}_{V, H}(e)) \\
FV(e_1) \cap FV(\text{lam}(x : \tau.e_2)) = \emptyset & (\text{Lemma 1.2}) \\
\text{locs}_{V, H}(e_1) \cap \text{locs}_{V_2, H}(\text{lam}(x : \tau.e_2)) = \emptyset & (\text{no\_alias}(V, H)) \\
R' \cap \text{locs}_{V, H}(e_1) = \emptyset & (\text{disjoint}(\{R, \text{locs}_{V, H}(e)\})) \\
F \cap \text{locs}_{V, H}(e_1) = \emptyset & (\text{Sp.})
\end{array}$$

Thus we have  $\text{disjoint}(R', F, \text{locs}_{V, H}(e_1))$

By IH,

$$\begin{array}{ll}
(1) \quad \text{set}(\text{reach}_{H_1}(v_1)) & \\
(2) \quad \text{disjoint}(\{R', F_1, \text{reach}_{H_1}(v_1)\}) & \\
(3) \quad \text{stable}(R', H, H_1) & \\
V'_2, H_1, R, F_1 \cup g \vdash e_2 \Downarrow v_2, H_2, F_2 & (\text{ad.})
\end{array}$$

$$\Sigma; \Gamma_2, x : A \vdash e_2 : B \quad (\text{ad.})$$

$$H_1 \models V'_2 : (\Gamma_2, x : A) \quad (???)$$

By IH, we have invariant on  $J_2$

NTS (1) - (3) to instantiate invariant on  $J_2$

$$(1) \quad \text{dom}(V'_2) = \text{FV}(e_2) \quad (\text{def of } V'_2)$$

$$(2) \quad \text{no\_alias}(V'_2, H_1)$$

Let  $x_1, x_2 \in V'_2, x_1 \neq x_2$  be arb.

**case:**  $x_1 \neq x, x_2 \neq x$

$$\text{reach}_H(V'_2(x_1)) \subseteq R' \quad (\text{reach}_H(V'_2(x_1)) \subseteq \text{locs}_{V'_2, H}(\text{lam}(x : \tau.e_2)))$$

$$\text{reach}_H(V'_2(x_2)) \subseteq R' \quad (\text{reach}_H(V'_2(x_2)) \subseteq \text{locs}_{V'_2, H}(\text{lam}(x : \tau.e_2)))$$

$$\text{reach}_H(V'_2(x_1)) = \text{reach}_{H_1}(V'_2(x_1)), \text{reach}_H(V'_2(x_2)) = \text{reach}_{H_1}(V'_2(x_2))$$

(stable( $R', H, H_1$ ) and Lemma 1.3)

$$\text{reach}_{H_1}(V'_2(x_1)) = \text{reach}_H(V(x_1)), \text{reach}_{H_1}(V'_2(x_2)) = \text{reach}_H(V(x_2))$$

(stable( $R', H, H_1$ ) and Lemma 1.3)

$$\text{no\_alias}(V'_2, H_1) \quad (\text{no\_alias}(V, H))$$

**case:**  $x_1 = x, x_2 \neq x$

$$\text{reach}_{H_1}(V'_2(x_1)) = \text{reach}_{H_1}(v_1) \quad (\text{def of } V'_2)$$

$$\text{reach}_{H_1}(V'_2(x_2)) \subseteq R' \quad (\text{same as above})$$

$$\text{set}(\text{reach}_{H_1}(v_1)) \quad (\text{IH 1.1})$$

$$\text{reach}_{H_1}(V'_2(x_2)) = \text{reach}_H(V(x_2)) \quad (\text{same as above})$$

$$\text{set}(\text{reach}_{H_1}(V'_2(x_2))) \quad (\text{no\_alias}(V, H))$$

$$\text{reach}_{H_1}(V'_2(x_1)) \cap \text{reach}_{H_1}(V'_2(x_2)) = \emptyset \quad (\text{disjoint}(\{R', \text{reach}_{H_1}(v_1)\}))$$

Thus we have  $\text{no\_alias}(V'_2, H_1)$

$$(3) \quad \text{disjoint}(\{R, F_1 \cup g, \text{locs}_{V'_2, H_1}(e_2)\})$$

$$R \cap F_1 = \emptyset \quad (\text{disjoint}(\{R', F_1\}) \text{ from 1.2 and } R \subseteq R')$$

$$R \cap (F_1 \cup g) = \emptyset \quad (\text{def of } g)$$

$$\text{NTS } (F_1 \cup g) \cap \text{locs}_{V'_2, H_1}(e_2) = \emptyset$$

Let  $l \in \text{locs}_{V'_2, H_1}(e_2)$  be arb.

$$l \in \text{reach}_{H_1}(V'_2(x')) \text{ for some } x' \in V'_2$$

**case:**  $x' \neq x$

$$\text{reach}_H(V'_2(x')) = \text{reach}_{H_1}(V'_2(x')) \quad (\text{same as above})$$

$$\text{reach}_{H_1}(V'_2(x')) \subseteq R' \quad (\text{def of } R')$$

$$\text{reach}_{H_1}(V'_2(x')) \cap F_1 = \emptyset \quad (\text{disjoint}(\{R', F_1\}) \text{ from 1.2})$$

**case:**  $x' = x$

$$\text{reach}_{H_1}(V'_2(x')) = \text{reach}_{H_1}(v_1) \quad (\text{def of } V'_2)$$

$$\text{reach}_{H_1}(V'_2(x')) \cap F_1 = \emptyset \quad (\text{disjoint}(\{F_1, \text{reach}_{H_1}(v_1)\}) \text{ from 1.2})$$

$$reach_{H_1}(V'_2(x')) \subseteq locs_{V'_2, H_1}(e_2) \quad (\text{def of } locs_{V, H})$$

$$reach_{H_1}(V'_2(x')) \cap g = \emptyset \quad (\text{def of } g)$$

$$\text{Thus } reach_{H_1}(V'_2(x')) \cap (F_1 \cup g) = \emptyset$$

$$\text{NTS } R \cap locs_{V'_2, H_1}(e_2) = \emptyset$$

Let  $l \in locs_{V'_2, H_1}(e_2)$  be arb.

$$l \in reach_{H_1}(V'_2(x')) \text{ for some } x' \in V'_2$$

**case:**  $x' \neq x$

$$reach_H(V_2(x')) = reach_{H_1}(V'_2(x')) \quad (\text{same as above})$$

$$l \in locs_{V, H}(e) \quad (\text{def of } locs_{V, H})$$

$$l \notin R \quad (\text{disjoint}(\{R, locs_{V, H}(e)\}) \text{ from 0.3})$$

**case:**  $x' = x$

$$reach_{H_1}(V'_2(x')) = reach_{H_1}(v_1) \quad (\text{def of } V'_2)$$

$$reach_{H_1}(V'_2(x')) \cap R = \emptyset \quad (\text{disjoint}(\{R', reach_{H_1}(v_1)\}) \text{ from 1.2 and } R \subseteq R')$$

$$\text{Thus } reach_{H_1}(V'_2(x')) \cap R = \emptyset$$

$$\text{Hence we have (3) } \text{disjoint}(R, F_1 \cup g, locs_{V'_2, H_1}(e_2))$$

By instantiating the invariant on  $J_2$ , we have

$$(1) \text{ set}(reach_{H_2}(v_2))$$

$$(2) \text{ disjoint}(\{R, F_2, reach_{H_2}(v_2)\})$$

$$(3) \text{ stable}(R, H_1, H_2)$$

Lastly, showing (1) - (3) holds for the original case  $J_0$  :

$$(1) \text{ set}(reach_{H_2}(v_2)) \quad (\text{By 2.1})$$

$$(2) \text{ disjoint}(\{R, F_2, reach_{H_2}(v_2)\}) \quad (\text{By 2.2})$$

$$(3) \text{ stable}(R, H_1, H_2)$$

Let  $l \in R$  be arb.

$$H(l) = H_1(l) \quad (\text{stable}(R', H, H_1) \text{ from 1.3})$$

$$H_1(l) = H_2(l) \quad (\text{stable}(R, H_1, H_2) \text{ from 2.3})$$

$$H(l) = H_2(l)$$

$$\text{Hence stable}(R, H, H_2)$$

**Case 8: E:Pair** Similar to E:Var

**Case 9: E:MatP** Similar to E:MatCons

**Case 10: E:nil** Similar to E:Const\*

**Case 11: E:Cons**

$$V, H, R, F \vdash e \Downarrow l, H'', F' \quad (\text{case})$$

$$\text{Suppose } H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, locs_{V, H}(e)\})$$

NTS (1) - (3) holds after evaluation

$$(1) \text{ set}(reach_{H''}(l))$$



$\text{stable}(\{\text{locs}_{V,H}(e)\}, H, H'')$   
 $(\text{disjoint}(\{F, \text{locs}_{V,H}(e)\}) \text{ and } \text{copy} \text{ only updates } l \in L \subseteq F)$   
 $\text{reach}_H(V(x_i)) = \text{reach}_{H''}(V(x_i)) \quad (\text{reach}_H(V(x_i)) \subseteq \text{locs}_{V,H}(e) \text{ and } \text{1.3 for } i = 1, 2)$   
 $\text{reach}_{H''}(l) = \{l\} \cup \text{reach}_{H''}(V(x_1)) \cup \text{reach}_{H''}(V(x_2)) \quad (\text{def of } \text{reach}_H)$   
 $\text{set}(\text{reach}_{H''}(l)) \quad (l \notin \text{locs}_{V,H}(e) \text{ and } \text{no\_alias}(V, H))$   
 $(2) \quad \text{disjoint}(\{R, F', \text{reach}_{H''}(l)\})$   
 $R \cap F' = \emptyset \quad (F' \subseteq F \text{ and } \text{disjoint}(\{R, F'\}))$   
 $R \cap \text{reach}_{H''}(l) = \emptyset \quad (l \in F \text{ and } \text{disjoint}(\{R, \text{locs}_{V,H}(e)\}))$   
 $F' \cap \text{reach}_{H''}(l) = \emptyset \quad (F' \subseteq F \text{ and } \text{disjoint}(\{F, \text{locs}_{V,H}(e)\}))$   
 Thus we have  $(2) \quad \text{disjoint}(\{R, F', \text{reach}_{H''}(l)\})$   
 $(3) \quad \text{stable}(R, H, H'') \quad (\text{since copy only updates } l \in L \subseteq F \text{ and } F \cap R = \emptyset)$

### Case 12: E:MatNil

Suppose  $H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, \text{locs}_{V,H}(e)\})$   
 $\Sigma; \Gamma' \vdash e_1 : B \quad (\text{ad.})$   
 $V, H, R, F \cup g \vdash e_1 \Downarrow v, H', F' \quad (\text{ad.})$   
 $H \models V' : \Gamma' \quad (\text{def of W.D.E})$

By IH, we have invariant on  $J_1$

NTS (1) - (3) to instantiate invariant on  $J_1$

$(1) \quad \text{dom}(V') = FV(e_1) \quad (\text{def of } V')$   
 $(2) \quad \text{no\_alias}(V', H) \quad (\text{no\_alias}(V, H) \text{ and } V' \subseteq V)$   
 $(3) \quad \text{disjoint}(\{R, F, \text{locs}_{V',H}(e_1)\})$   
 $(\text{disjoint}(\{R, F, \text{locs}_{V,H}(e)\}) \text{ and } \text{locs}_{V',H}(e_1) \subseteq \text{locs}_{V,H}(e))$

Instantiating invariant on  $J_1$ ,

$(1) \quad \text{set}(\text{reach}_{H'}(v))$   
 $(2) \quad \text{disjoint}(\{R, F_1, \text{reach}_{H'}(v)\})$   
 $(3) \quad \text{stable}(R, H, H')$

### Case 13: E:MatCons

$V(x) = l \quad (\text{ad.})$   
 $H(l) = \langle v_h, v_t \rangle \quad (\text{ad.})$   
 $\Gamma = \Gamma', x : L(A) \quad (\text{ad.})$   
 $\Sigma; \Gamma', x_h : A, x_t : L(A) \vdash e_2 : B \quad (\text{ad.})$   
 $V'', H, R, F \cup g \vdash e_2 \Downarrow v_2, H_2, F' \quad (\text{ad.})$   
 Suppose  $H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{F, R, \text{locs}_{V,H}(e)\})$   
 $H \models V(x) : L(A) \quad (\text{def of W.D.E})$   
 $H'' \models v_h : A, H'' \models v_t : L(A) \quad (\text{ad.})$

$$H \models v_h : A, H \models v_t : L(A) \quad (???)$$

$$H \models V'' : \Gamma', x_h : A, x_t : L(A) \quad (\text{def of W.D.E})$$

By IH, we have invariant on  $J_1$

NTS (1) - (3) to instantiate invariant on  $J_1$

$$(1) \quad \text{dom}(V'') = FV(e_2) \quad (\text{def of } V'')$$

$$(2) \quad \text{no\_alias}(V'', H)$$

Let  $x_1, x_2 \in V'', x_1 \neq x_2, r_{x_1} = \text{reach}_H(V''(x_1)), r_{x_2} = \text{reach}_H(V''(x_2))$

**case:**  $x_1 \notin \{x_h, x_t\}, x_2 \notin \{x_h, x_t\}$

(1), (2) from  $\text{no\_alias}(V, H)$

**case:**  $x_1 = x_h, x_2 \notin \{x_h, x_t\}$

$\text{set}(r_{x_1})$  ( since  $\text{set}(\text{reach}_H(V(x)))$  from  $\text{no\_alias}(V, H)$ )

$\text{set}(r_{x_2})$  (since  $\text{no\_alias}(V, H)$ )

$x_2 \in FV(e)$  (def of  $FV$ )

$\text{reach}_H(V(x)) \cap r_{x_2} = \emptyset$  (def of  $\text{reach}$  and  $\text{no\_alias}(V, H)$ )

hence  $r_{x_1} \cap r_{x_2} = \emptyset$

**case:**  $x_1 = x_h, x_2 = x_t$

$\text{set}(r_{x_1})$  since  $\text{set}(\text{reach}_H(V(x)))$  from  $\text{no\_alias}(V, H)$

$\text{set}(r_{x_2})$  since  $\text{set}(\text{reach}_H(V(x)))$  from  $\text{no\_alias}(V, H)$

$r_{x_1} \cap r_{x_2} = \emptyset$  ( $\text{set}(\text{reach}_H(V(x)))$ )

**case: otherwise**

similar to the above

Thus we have  $\text{no\_alias}(V'', H)$

$$(3) \quad \text{disjoint}(\{R, F \cup g, \text{locs}_{V'', H}(e_2)\})$$

$$(F \cup g) \cap R = \emptyset \quad (\text{since } F \cap R = \emptyset \text{ and by def of } g)$$

$$\text{NTS } R \cap \text{locs}_{V'', H}(e_2) = \emptyset$$

Let  $l' \in \text{locs}_{V'', H}(e_2)$  be arb.

**case:**  $l' \in \text{reach}_H(V''(x'))$  for some  $x' \in FV(e_2)$  where  $x' \notin \{x_h, x_t\}$

$x' \in V$  (def of  $V''$ )

$l' \in \text{reach}_H(V(x'))$

$x' \in FV(e)$  (def of  $FV$ )

$l' \in \text{locs}_{V, H}(e)$  (def of  $\text{locs}_{V, H}$ )

$l' \notin R$  ( $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$ )

**case:**  $l' \in \text{reach}_H(V''(x_h))$

*tom*  $l' \in \text{reach}_H(v_h)$

$l' \in \text{reach}_H(V(x))$  (def of  $\text{reach}$ )

$l' \in \text{locs}_{V, H}(e)$  (def of  $\text{locs}_{V, H}$ )

$l' \notin R$  (since  $\text{disjoint}(\{F, R, \text{locs}_{V,H}(e)\})$ )  
**case:**  $l' \in \text{reach}_H(V''(x_l))$   
 similar to above  
 Hence  $R \cap \text{locs}_{V'',H}(e_2) = \emptyset$   
 $F \cap \text{locs}_{V'',H}(e_2) = \emptyset$  (Similar to above)  
 $g \cap \text{locs}_{V'',H}(e_2) = \emptyset$  (def. of  $g$ )  
 $(F \cup g) \cap \text{locs}_{V'',H}(e_2) = \emptyset$   
 Thus  $\text{disjoint}(\{R, F \cup g, \text{locs}_{V'',H}(e_2)\})$   
 Instantiating invariant on  $J_1$ ,  
 (1)  $\text{set}(\text{reach}_{H'}(v))$   
 (2)  $\text{disjoint}(\{R, F', \text{reach}_{H'}(v)\})$   
 (3)  $\text{stable}(R, H, H')$

### Case 13: E:Drop

$e = \text{drop}(x; e')$  (case)  
 $V', H, R, F \cup g \vdash e' \Downarrow v, H', F'(\mathcal{J}_1)$  (ad.)  
 $\Gamma = \Gamma', x : A$  (case)  
 $\Sigma; \Gamma' \mid \frac{q}{q'} e' : B$   
 Suppose  $\text{dom}(V) = FV(e)$ ,  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V,H}(e)\})$   
 By IH, we have invariant on  $\mathcal{J}_1$   
 NTS (1) - (3) for  $\mathcal{J}_1$   
 (1)  $\text{dom}(V') = FV(e')$  ( $\text{dom}(V) = FV(e)$  and def of  $FV$ )  
 (2)  $\text{no\_alias}(V', H)$  ( $\text{no\_alias}(V, H)$  and  $V' \subseteq V$ )  
 (3)  $\text{disjoint}(\{R, F \cup g, \text{locs}_{V',H}(e')\})$   
 $g = \text{reach}_H(v')$  (case)  
 $g \subseteq \text{locs}_{V,H}(e)$  (def of  $\text{locs}_{V,H}$ )  
 $R \cap (F \cup g) = \emptyset$  ( $\text{disjoint}(\{R, F\})$  and  $\text{disjoint}(\{R, \text{locs}_{V,H}(e)\})$ )  
 $R \cap \text{locs}_{V',H}(e') = \emptyset$  ( $\text{disjoint}(\{R, \text{locs}_{V,H}(e)\})$  and  $\text{locs}_{V',H}(e) \subseteq \text{locs}_{V,H}(e)$ )  
 $F \cap \text{locs}_{V',H}(e') = \emptyset$  ( $\text{disjoint}(\{F, \text{locs}_{V,H}(e)\})$  and  $\text{locs}_{V',H}(e) \subseteq \text{locs}_{V,H}(e)$ )  
 $g \cap \text{locs}_{V',H}(e') = \emptyset$  ( $\text{no\_alias}(V, H)$ )  
 Instantiating invariant on  $\mathcal{J}_1$ ,  
 (1)  $\text{set}(\text{reach}_{H'}(v))$   
 (2)  $\{R, F', \text{reach}_{H'}(v)\}$   
 (3)  $\text{stable}(R, H, H')$

### Case 13: E:Share

$e = \text{shareCopy } x \text{ as } x_1, x_2 \text{ in } e'$  (case)

Suppose  $H \models V : \Gamma, \text{dom}(V) = FV(e), \text{no\_alias}(V, H), \text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$   
(def. of wtf)

Let  $V_2 = V'[x_1 \text{mapstov}', x_2 \text{mapstov}']$

We show the subsequent computation is also well-formed to invoke the IH:

(1)  $\text{dom}(V_2) = FV(e')$  ( $\text{dom}(V) = FV(e)$  and def of  $FV$ )

(2)  $\text{no\_alias}(V'[x_1 \mapsto v', x_2 \mapsto v''], H)$   $\text{no\_alias}(V'[x_1 \mapsto v', x_2 \mapsto v''], H)$   
( $\text{no\_alias}(V, H)$ )

Let  $x' \mapsto v''' \in V'[x_1 \mapsto v']$ .  $\text{STS } \text{reach}_{H'}(v''') \cap \text{reach}_{H'}(v'') = \emptyset$

$\text{reach}_{H'}(v'') \subseteq L \subseteq F$  (lemma about copy)

$\text{reach}_{H'}(v''') \subseteq \text{locs}_{V'[x_1 \mapsto v'], H'}(e') \subseteq \text{locs}_{V, H}(e)$  (stability lemma for copy)

(wfc( $V, H, R, F$ ))

(3)  $\text{disjoint}(\{R, F \setminus L, \text{locs}_{V_2, H'}(e')\})$

By IH:

(1)  $\text{set}(\text{reach}_{H''})$

(2)  $\{R, F', \text{reach}_{H''}(v)\}$

(3)  $\text{stable}(R, H', H'')$

$\text{STS stable}(R, H, H')$ , which follows from  $L \cap R = \emptyset$  and stability for copy

□

**TASK 1.8 (SOUNDNESS).** *let  $H \models V : \Gamma, \Sigma; \Gamma \mid \frac{q}{q'} e : B$ , and  $V, H \vdash e \Downarrow v, H'$ . Then  $\forall C \in \mathbb{Q}^+$  and  $\forall F, R \subseteq \text{Loc}$ , if we have the following (existence lemma):*

- (1)  $\text{dom}(V) = FV(e)$
- (2)  $\text{no\_alias}(V, H)$
- (3)  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$ , and
- (4)  $|F| \geq \Phi_{V, H}(\Gamma) + q + C$

*then there exists  $F' \subseteq \text{Loc}$  s.t.*

- (1)  $V, H, R, F \vdash e \Downarrow v, H', F'$
- (2)  $|F'| \geq \Phi_{H'}(v : B) + q' + C$

**PROOF.** Nested induction on the evaluation judgement and the typing judgement.

### Case 1: E:Var

$V, H, R, F \vdash x \Downarrow V(x), H, F$  (admissibility)

$\Sigma; x : B \mid \frac{q}{q'} x : B$  (admissibility)

$|F| - |F'|$  (1)

$= |F| - |F|$  (ad.)

$$= 0 \quad (2)$$

$$\Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q') \quad (3)$$

$$= \Phi_{V,H}(x : B) + q - (\Phi_H(V(x) : B) + q) \quad (\text{ad.})$$

$$= \Phi_H(V(x) : B) + q - (\Phi_H(V(x) : B) + q) \quad (\text{def. of } \Phi_{V,H})$$

$$= 0 \quad (4)$$

$$|F| - |F'| \leq \Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q') \quad ((3),(5))$$

**Case 2: E:Const\*** Due to similarity, we show only for E:ConstI

$$|F| - |F'| = |F| - |F| \quad (\text{ad.})$$

$$= 0$$

$$\Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q') = \Phi_{V,H}(\emptyset) + q - (\Phi_H(v : \text{int}) + q) \quad (\text{ad.})$$

$$= 0 \quad (\text{def of } \Phi_{V,H})$$

$$|F| - |F'| \leq \Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q')$$

**Case 4: E:App**

**Case 5: E:CondT**

$$\Gamma = \Gamma', x : \text{bool} \quad (\text{ad.})$$

$$H \models V : \Gamma' \quad (\text{def of W.F.E})$$

$$\Sigma; \Gamma' \mid \frac{q}{q'} e_t : B \quad (\text{ad.})$$

$$V, H, R, F \cup g \vdash e_t \Downarrow v, H', F' \quad (\text{ad.})$$

$$|F \cup g| - |F'| \leq \Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q') \quad (\text{IH})$$

$$|F| - |F'| \leq \Phi_{V,H}(\Gamma) + q - (\Phi_{H'}(v : B) + q')$$

**Case 6: E:CondF** Similar to E:CondT

**Case 7: E:Let**

$$V, H \vdash e \Downarrow v_2, H_2 \quad (\text{case})$$

$$V, H \vdash e_1 \Downarrow v_1, H_1 \quad (\text{ad.})$$

$$\Sigma; \Gamma_1 \mid \frac{q}{p} e_1 : A \quad (\text{ad.})$$

$$H \models V_1 : \Gamma_1 \quad (\text{def of W.D.E})$$

Let  $C \in \mathbb{Q}^+$ ,  $F, R \subseteq \text{Loc}$  be arb.

Suppose  $\text{dom}(V) = FV(e)$ ,  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V,H}(e)\})$ , and  $|F| \geq \Phi_{V,H}(\Gamma) + q + C$

NTF  $F'$  s.t.

$$1. V, H, R, F \vdash e \Downarrow v_2, H_2, F' \text{ and}$$

$$2. |F'| \geq \Phi_{H_2}(v_2 : B) + q' + C$$

Let  $R' = R \cup \text{locs}_{V,H}(\text{lam}(x : \tau.e_2))$

$\text{disjoint}(\{R', F, \text{locs}_{V,H}(e_1)\})$  (Similar to case in Lemma 1.7)

Instantiate IH with  $C = C + \Phi_{V_2,H}(\Gamma_2)$ ,  $F = F, R = R'$ , we get existence lemma on  $J_1$  :

NTS (1) - (4) to instantiate existence lemma on  $J_1$

- (1)  $\text{dom}(V_1) = FV(e_1)$
- (2)  $\text{no\_alias}(V_1, H)$
- (3)  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$  ((1) - (3) all verbatim as in Lemma 1.7)
- (4)  $|F| \geq \Phi_{V_1, H}(\Gamma_1) + q + C + \Phi_{V, H}(\Gamma_2)$   
 $(|F| \geq \Phi_{V, H}(\Gamma) + q + C \text{ and } \Phi_{V, H}(\Gamma) \geq \Phi_{V_1, H}(\Gamma_1) + \Phi_{V, H}(\Gamma_2))$

Instantiating existence lemma on  $J_1$ , we get  $F''$  s.t.

- 1.  $V, H, R', F \vdash e_1 \Downarrow v_1, H_1, F''$  and
- 2.  $|F''| \geq \Phi_{H_1}(v_1 : A) + p + C + \Phi_{V_2, H_1}(\Gamma_2)$

For the second premise:

- $\Sigma; \Gamma_2, x : A \mid \frac{p}{q} e_2 : B$  (ad.)
- $H_1 \models v_1 : A$  and (Theorem 3.3.4)
- $H_1 \models V : \Gamma_2$  (???)
- $H_1 \models V' : \Gamma_2, x : A$  (def of  $\models$ )
- $V', H_1 \vdash e_2 \Downarrow v_2, H_2$  (ad.)

Let  $g = \{l \in H_1 \mid l \notin F_1 \cup R \cup \text{locs}_{V', H_1}(e_2)\}$

Instantiate IH with  $C = C, F = F'' \cup g, R = R$ , we get existence lemma on  $J_2$  :

NTS (1) - (4) to instantiate existence lemma on  $J_1$

- (1)  $\text{dom}(V'_2) = FV(e_2)$
- (2)  $\text{no\_alias}(V'_2, H_1)$
- (3)  $\text{disjoint}(\{R, F'' \cup g, \text{locs}_{V'_2, H_1}(e_2)\})$  ((1) - (3) all verbatim as in Lemma 1.7)
- (4)  $|F'' \cup g| \geq \Phi_{V'_2, H_1}(\Gamma_2, x : (A - 1)) + p + C$   
 $\text{STS } |F'' \cup g| \geq \Phi_{V_2, H_1}(\Gamma_2) + \Phi_{H_1}(v_1 : (A - 1)) + p + C$   
 $|F'' \cup g| \geq \|V_1\|_H + |F| - \|v_1\|_{H_1}$  (conservation lemma)  
 $\geq \Phi_{V, H}(\Gamma) + q + C + \|V_1\|_H - \|v_1\|_{H_1}$  ( $|F| \geq \Phi_H(V) + q + C$ )  
 $\text{STS } \Phi_{V_1, H}(\Gamma_1) + q + C + \|V_1\|_H - \|v_1\|_{H_1} \geq \Phi_{H_1}(v_1 : (A - 1)) + p + C$   
 $\Phi_{V_1, H}(\Gamma_1) \geq \Phi_{H_1}(v_1 : (A - 1))$  (lemma about cf typing)  
 $\text{STS } \|V_1\|_H - \|v_1\|_{H_1} + q \geq p$  (done by aux lemma)

Instantiating existence lemma on  $J_2$ , we get  $F^{(3)}$  s.t.

- 1.  $V'_2, H_1, R, F'' \cup g \vdash e_2 \Downarrow v_2, H_2, F^{(3)}$
- 2.  $|F^{(3)}| \geq \Phi_{H_2}(v_2 : B) + q' + C$

Take  $F' = F^{(3)}$

- $V, H, R, F \vdash e \Downarrow v_2, H_2, F'$  and (E:Let)
- $|F'| \geq \Phi_{H_2}(v_2 : B) + q' + C$  (from IH)

**Case 14: E:Let1**

$$V, H \vdash e \Downarrow v_2, H_2 \quad (\text{case})$$

$$V, H \vdash e_1 \Downarrow v_1, H_1 \quad (\text{ad.})$$

$$\Sigma; \Gamma_1 \mid \frac{q}{p} e_1 : A \quad (\text{ad.})$$

$$H \models V_1 : \Gamma_1 \quad (\text{def of W.D.E})$$

Let  $C \in \mathbb{Q}^+$ ,  $F, R \subseteq \text{Loc}$  be arb.

Suppose  $\text{dom}(V) = FV(e)$ ,  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$ , and  $|F| \geq \Phi_{V, H}(\Gamma) + q + C$   
NTF  $F'$  s.t.

$$1. V, H, R, F \vdash e \Downarrow v_2, H_2, F' \text{ and}$$

$$2. |F'| \geq \Phi_{H_2}(v_2 : B) + q' + C$$

Let  $R' = R \cup \text{locs}_{V, H}(\text{lam}(x : \tau.e_2))$

$\text{disjoint}(\{R', F, \text{locs}_{V, H}(e_1)\})$  (Similar to case in Lemma 1.7)

Instantiate IH with  $C = C + \Phi_{V_2, H}(\Gamma_2)$ ,  $F = F$ ,  $R = R'$ , we get existence lemma on  $J_1$  :

NTS (1) - (4) to instantiate existence lemma on  $J_1$

$$(1) \text{ dom}(V_1) = FV(e_1)$$

$$(2) \text{ no\_alias}(V_1, H)$$

$$(3) \text{ disjoint}(\{R, F, \text{locs}_{V, H}(e)\}) \quad ((1) - (3) \text{ all verbatim as in Lemma 1.7})$$

$$(4) |F| \geq \Phi_{V_1, H}(\Gamma_1) + q + C + \Phi_{V, H}(\Gamma_2) \\ (|F| \geq \Phi_{V, H}(\Gamma) + q + C \text{ and } \Phi_{V, H}(\Gamma) \geq \Phi_{V_1, H}(\Gamma_1) + \Phi_{V, H}(\Gamma_2))$$

Instantiating existence lemma on  $J_1$ , we get  $F''$  s.t.

$$1. V, H, R', F \vdash e_1 \Downarrow v_1, H_1, F'' \text{ and}$$

$$2. |F''| \geq \Phi_{H_1}(v_1 : A) + p + C + \Phi_{V_2, H_1}(\Gamma_2)$$

For the second premise:

$$\Sigma; \Gamma_2, x : A \mid \frac{\max(p, q)}{q'} e_2 : B \quad (\text{ad.})$$

$$H_1 \models v_1 : A \text{ and} \quad (\text{Theorem 3.3.4})$$

$$H_1 \models V : \Gamma_2 \quad (???)$$

$$H_1 \models V' : \Gamma_2, x : A \quad (\text{def of } \models)$$

$$V', H_1 \vdash e_2 \Downarrow v_2, H_2 \quad (\text{ad.})$$

Let  $g = \{l \in H_1 \mid l \notin F'' \cup R \cup \text{locs}_{V', H_1}(e_2)\}$

Instantiate IH with  $C = C$ ,  $F = F'' \cup g$ ,  $R = R$ , we get existence lemma on  $J_2$  :

NTS (1) - (4) to instantiate existence lemma on  $J_1$

$$(1) \text{ dom}(V'_2) = FV(e_2)$$

$$(2) \text{ no\_alias}(V'_2, H_1)$$

$$(3) \text{ disjoint}(\{R, F'' \cup g, \text{locs}_{V'_2, H_1}(e_2)\}) \quad ((1) - (3) \text{ all verbatim as in Lemma 1.7})$$

$$(4) |F'' \cup g| \geq \Phi_{V'_2, H_1}(\Gamma_2, x : A) + q + C$$

$$\begin{aligned}
|F'' \cup g| &\geq |F''| \\
&\geq \Phi_{H_1}(v_1 : A) + p + C + \Phi_{V_2, H}(\Gamma_2) & (\text{IH}) \\
&= \Phi_{H_1}(v_1 : A) + p + C + \Phi_{V'_2, H_1}(\Gamma_2) & (\text{Lemma 4.3.3}) \\
&= \Phi_{V'_2, H_1}(\Gamma_2, x : A) + p + C & (\text{def of } \Phi)
\end{aligned}$$

Instantiating existence lemma on  $J_2$ , we get  $F^{(3)}$  s.t.

$$\begin{aligned}
1. & V'_2, H_1, R, F'' \cup g \vdash e_2 \Downarrow v_2, H_2, F^{(3)} \\
2. & |F^{(3)}| \geq \Phi_{H_2}(v_2 : B) + q' + C
\end{aligned}$$

Take  $F' = F^{(3)}$

$$V, H, R, F \vdash e \Downarrow v_2, H_2, F' \text{ and} \quad (\text{E:Let})$$

$$|F'| \geq \Phi_{H_2}(v_2 : B) + q' + C \quad (\text{from IH})$$

**Case 8: E:Pair** Similar to E:Const\*

**Case 9: E:MatP** Similar to E:MatCons

**Case 10: E:Nil** Similar to E:Const\*

**Case 11: E:Cons**

$$V, H \vdash \text{cons}(x_1; x_2) \Downarrow l, H' \quad (\text{case})$$

Let  $C \in \mathbb{Q}^+$ ,  $F, R \subseteq \text{Loc}$  be arb.

Suppose  $\text{dom}(V) = FV(e)$ ,  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$ ,  $|F| \geq \Phi_{V, H}(\Gamma) + q + C$

NTF  $F'$  s.t.

$$\begin{aligned}
1. & V, H, R, F \vdash e \Downarrow v, H', F' \text{ and} \\
2. & |F'| \geq \Phi_{H'}(v : B) + q' + C
\end{aligned}$$

Let  $F' = F$

**Case 12: E:MatNil** Similar to E:Cond\*

**Case 13: E:MatCons**

$$V(x) = (l, \text{alive}) \quad (\text{ad.})$$

$$H(l) = \langle v_h, v_t \rangle \quad (\text{ad.})$$

$$\Gamma = \Gamma', x : L^P(A) \quad (\text{ad.})$$

$$\Sigma; \Gamma', x_h : A, x_t : L^P(A) \Big| \frac{q+p+1}{q'} e_2 : B \quad (\text{ad.})$$

$$V'', H \vdash e_2 \Downarrow v, H' \quad (\text{ad.})$$

Let  $C \in \mathbb{Q}^+$ ,  $F, R \subseteq \text{Loc}$  be arb.

$$H \models V(x) : L^P(A) \quad (\text{def of W.D.E})$$

$$H'' \models v_h : A, H'' \models v_t : L^P(A) \quad (\text{ad.})$$

$$H \models v_h : A, H \models v_t : L^P(A) \quad (???)$$

$$H \models V'' : \Gamma', x_h : A, x_t : L^P(A) \quad (\text{def of W.D.E})$$

Suppose  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\})$ , and  $|F| \geq \Phi_{V, H}(\Gamma) + q + C$



NTF  $F'$  s.t.

$$1. V, H, R, F \vdash e \Downarrow v, H', F' \text{ and}$$

$$2. |F'| \geq \Phi_{H'}(v : B) + q' + C$$

Let  $g = \{l \in H \mid l \notin F \cup R \cup \text{locs}_{V'', H}(e_2)\}$

We want to  $g$  nonempty, in particular, that  $l \in g$

$$l \notin F \cup R \quad (\text{disjoint}(\{R, F, \text{locs}_{V, H}(e)\}))$$

$$\text{AFSOC } l \in \text{locs}_{V'', H}(e_2)$$

Then  $l \in \text{reach}_H(\bar{V}''(x'))$  for some  $x' \neq x$

$$x' \in \{x_h, x_t\} \quad (\text{since } \text{reach}_H(\bar{V}(x')) \cap \text{reach}_H(\bar{V}x) = \emptyset \text{ from no\_alias}(V, H))$$

WLOG let  $x' = x_h$

But then  $\mu_{\text{reach}_H(\bar{V}(x))}(l) \geq 2$  and  $\text{set}(\text{reach}_H(\bar{V}(x)))$  doesn't hold

$$l \notin \text{locs}_{V'', H}(e_2)$$

Hence  $l \in g$

Next, we have  $\text{no\_alias}(V'', H)$  and  $\text{disjoint}(\{R, F \cup g, \text{locs}_{V'', H}(e_2)\})$   
(similar to case in Lemma 1.2)

By IH with  $C' = C, F'' = F \cup g$  and the above conditions, we have:  $F^{(3)}$  s.t.

$$1. V'', H, R, F \cup g \vdash e_2 \Downarrow v, H', F^{(3)}$$

$$2. |F^{(3)}| \geq \Phi_{H'}(v : B) + q' + C$$

Where we also verify the precondition that

$$|F''| \geq \Phi_{V'', H}(\Gamma', x_h : A, x_t : L^p(A)) + q + p + 1 + C' :$$

$$\begin{aligned} |F''| &= |F \cup g| \\ &= |F| + |g| && (F \text{ and } g \text{ disjoint}) \\ &\geq \Phi_{V, H}(\Gamma) + q + C + |g| && (\text{Sp.}) \\ &= \Phi_{V, H}(\Gamma', x_h : A, x_t : L^p(A)) + p + q + C + |g| && (\text{Lemma 4.1.1}) \\ &= \Phi_{V, H}(\Gamma', x_h : A, x_t : L^p(A)) + p + q + C + 1 && (g \text{ nonempty}) \end{aligned}$$

Now take  $F' = F^{(3)}$

$$V, H, R, F \vdash e \Downarrow v, H', F' \quad (\text{E:MatCons})$$

$$|F'| \geq \Phi_{H'}(v : B) + q' + C \quad (\text{From the IH})$$

### Case 13: E:Share

$$V, H \vdash e \Downarrow v, H'' \quad (\text{case})$$

$$V'[x_1 \mapsto v', x_2 \mapsto v''], H' \vdash e' \Downarrow v, H'' \quad (\text{ad})$$

$$\Sigma; \Gamma, x : A \mid \frac{q}{q'} e : B \quad (\text{case})$$

$$A \curlyvee A_1, A_2, 1 \quad (\text{ad.})$$

$$\Sigma; \Gamma, x_1 : A_1, x_2 : A_2 \mid \frac{q}{q} e : B \quad (\text{ad.})$$

Let  $C \in \mathbb{Q}^+$ ,  $F, R \subseteq \text{Loc}$  be arb.

Suppose  $\text{no\_alias}(V, H)$ ,  $\text{disjoint}(\{R, F, \text{locs}_{V,H}(e)\})$ , and  $|F| \geq \Phi_{V,H}(\Gamma, x : A) + q + C$   
NTF  $F''$  s.t.

1.  $V, H, R, F \vdash e \Downarrow v, H'', F''$  and
2.  $|F''| \geq \Phi_{H''}(v : B) + q' + C$

We need to show the freelist is sufficient for the subsequent computation to invoke the IH:

Instantiate with  $C, F \setminus L$ , and  $R$

STS  $|F \setminus L| \geq \Phi_{V_2, H'}(\Gamma, x_1 : A_1, x_2 : A_2) + q + C$

$$\iff |F| - |L| \geq \Phi_{V_2, H'}(\Gamma) + \Phi_{V_2, H'}(x_1 : A_1) + \Phi_{V_2, H'}(x_2 : A_2) + q + C$$

$$\iff |F| \geq \Phi_{V_2, H'}(\Gamma) + \Phi_{V_2, H'}(x_1 : A_1) + \Phi_{V_2, H'}(x_2 : A_2) + \|v'\|_H + q + C$$

$$\iff |F| \geq \Phi_{V_2, H'}(\Gamma) + \Phi_{V, H}(x : A) + q + C \quad (\text{definition of sharing relation})$$

$$\iff |F| \geq \Phi_{V, H}(\Gamma, x : A) + q + C \quad (\text{stability of copying})$$

done from assumption

By IH, we get  $F''$  fulfilling the previous two points for the case.

□

## 9 COPY-FREE GARBAGE COLLECTION SEMANTICS

Consider the GC semantics (from now on copy semantics) above, with the share rule replaced with the following:

$$\frac{V = V'[x \mapsto v'] \quad V'[x_1 \mapsto v', x_2 \mapsto v'], H', R, F \vdash e \Downarrow v, H'', F'}{V, H, R, F \vdash \text{share } x \text{ as } x_1, x_2 \text{ in } e \Downarrow v, H'', F'} (S_{22})$$

Call this new semantics free semantics for copy-free. It is easy to see that any terminating computation in copy has a corresponding one in free that can be instantiated with an equal or smaller freelist. This is expressed as the following lemma:

**LEMMA 1.9.** *Let  $V, H, R, F \vdash^{\text{copy}} v, H', F'$ . Then there exists an  $F''$  s.t.  $V, H, R, F \vdash^{\text{free}} v, H', F''$  and  $|F''| \geq |F'|$ .*