

$$F(t) = 1 - e^{-\lambda t} \Rightarrow f(t) = \lambda e^{-\lambda t}$$

Let $\lambda = 2$ $f(1) = 0.27$, $p(1) = 0$
 $f(0.1) = 2e^{-2(0.1)} = 1.63 > 0$

PDF's are not probabilities. They are probability "densities" i.e. how fast the CDF is changing. Its value doesn't matter, only
 (1) its integrals matter
 (2) its ratios matter

$$\frac{f(0.1)}{f(1)} = \frac{1.63}{0.27} \approx 6$$

This ratio means that realizations "near" 0.1 are 6x more likely than realizations near 1.

$$\lim_{\varepsilon \rightarrow 0} \frac{P(T \in [0.1, 0.1 + \varepsilon])}{P(T \in [1, 1 + \varepsilon])} = \frac{F(0.1 + \varepsilon) - F(0.1)}{F(1 + \varepsilon) - F(1)}$$

$$= \lim_{\varepsilon \rightarrow 0} \frac{\frac{F(0.1 + \varepsilon) - F(0.1)}{\varepsilon}}{\frac{F(1 + \varepsilon) - F(1)}{\varepsilon}} = \frac{f(0.1)}{f(1)}$$

$$P(T \in (-\infty, \infty)) = 1 \Rightarrow \int_{-\infty}^{\infty} f(t) dt = 1$$

equiv of $\sum_{x \in \text{supp}(X)} p(x) = 1$ for discrete rv

Properties of continuous rv X

$$\textcircled{1} \text{supp}[X] = |\mathbb{R}|$$

$$\textcircled{2} f(x) \geq 0$$

$$\textcircled{3} \int_{\mathbb{R}} f(x) dx = 1$$

$$\textcircled{4} \text{supp}[X] = \{x : f(x) > 0\}$$

$$\textcircled{5} Q[X, q] = F^{-1}(q) = \arg\min_x \{F(x) \geq q\}$$

$$\textcircled{6} \text{Definition of equals in distribution } X_1 \stackrel{d}{=} X_2 \Rightarrow f_{X_1}(x) = f_{X_2}(x)$$

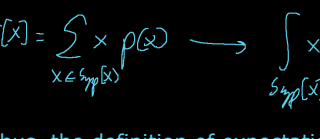
or $F_{X_1}(x) = F_{X_2}(x)$

How do we define expectation of a continuous rv, $E[X]$?

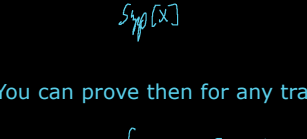
Recall for a discrete rv,

$$E[X] = \sum_{x \in \text{supp}[X]} x \cdot p(x)$$

For a continuous rv, there are no "probability masses", only density:



We can approximate any continuous rv as a discrete rv by defining a "bin width" and taking integrals over the grid via a left/right/mid rect-angle approx.



This is called a Riemann sum approximation. If the width of the rectangle goes to 0 in its limit, then the sum becomes the definition of the integral.

$$E[X] = \sum_{x \in \text{supp}[X]} x \cdot p(x) \longrightarrow \int_{\text{supp}[X]} x \cdot f(x) dx$$

Thus, the definition of expectation for a continuous rv is

$$\mu := E[X] := \int_{\text{supp}[X]} x \cdot f(x) dx$$

You can prove then for any transformation $g(X)$ that

$$E[g(X)] = \int_{\text{supp}[X]} g(x) \cdot f(x) dx$$

Thus, we can define variance which is nothing but an expectation of a transformation of X , the squared error loss,

$$\sigma^2 := \text{Var}[X] := E[(X - \mu)^2] = \int_{\text{supp}[X]} (x - \mu)^2 \cdot f(x) dx$$

We proved lots of rules for expectation and variance of discrete rv's. You can show all those rules are the same for continuous rv's. Here's a list:

$$E[aX + c] = a E[X] + c = a\mu + c$$

$$\text{Var}[aX + c] = a^2 \text{Var}[X] = a^2 \sigma^2$$

$$\Downarrow$$

$$\text{SD}[aX + c] = |a| \sqrt{\text{Var}[X]} = |a| \sigma$$

$$E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i] = n\mu$$

if X_1, \dots, X_n equally distributed

$$\text{Var}\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n \text{Var}[X_i] = n\sigma^2$$

if X_1, \dots, X_n ind. if X_1, \dots, X_n iid

When we "scrunched up" the geometric last class, we invented a new rv, the "exponential rv"

$$P(X > x) = e^{-\lambda x}$$

$$X \sim \text{Exp}(\lambda) := \underbrace{\lambda e^{-\lambda x}}_{f(x)}, \quad F(x) = 1 - \underbrace{e^{-\lambda x}}_{P(X \leq x)}$$

$\lambda > 0$ from its def. with n, p last class

$$\text{supp}[X] = [0, \infty)$$

$$E[X] = \int_0^{\infty} x \cdot \lambda e^{-\lambda x} dx = \lambda \int_0^{\infty} x e^{-\lambda x} dx$$

Recall integration by parts: $\int u dv = uv - \int v du$

$$\text{Let } u = x \Rightarrow du = dx, \quad dv = e^{-\lambda x} dx \Rightarrow v = \int e^{-\lambda x} dx = -\frac{1}{\lambda} e^{-\lambda x}$$

$$= \lambda \left(\left[x \left(-\frac{1}{\lambda} e^{-\lambda x} \right) \right]_0^{\infty} - \int_0^{\infty} -\frac{1}{\lambda} e^{-\lambda x} dx \right)$$

$$= \lambda \left(\underbrace{\left(-\frac{1}{\lambda} \left[\frac{x}{e^{\lambda x}} \right]_0^{\infty} \right)}_{0-0} - \frac{1}{\lambda^2} \underbrace{\left[e^{-\lambda x} \right]_0^{\infty}}_{0-1} \right) = \frac{1}{\lambda}$$

We said the geometric rv was "memoryless". Is the exponential rv memoryless?

$$P(X > a+b \mid X > b) = \frac{P(X > a+b \text{ and } X > b)}{P(X > b)} = \frac{P(X > a+b)}{P(X > b)}$$

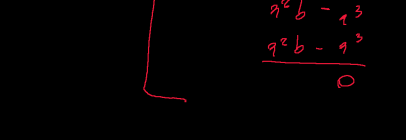
$$= \frac{e^{-\lambda(a+b)}}{e^{-\lambda b}} = \frac{e^{-\lambda a} e^{-\lambda b}}{e^{-\lambda b}}$$

$$= e^{-\lambda a} = P(X > a)$$

Recall the uniform discrete rv. e.g. $X \sim U(\{1, 7, 19\}) = \begin{cases} 1 & \text{w.p. } \frac{1}{3} \\ 7 & \text{w.p. } \frac{1}{3} \\ 19 & \text{w.p. } \frac{1}{3} \end{cases}$

There is a uniform continuous rv called the "uniform rv"

$X \sim U(a, b)$ where $a < b$ and $a, b \in \mathbb{R}$, $\text{Supp}[X] = [a, b]$



$$1 = \int_a^b c dx = c [x]_a^b = c(b-a) \Rightarrow c = \frac{1}{b-a}$$

$$F(x) = \int f(x) dx + c = \int \frac{1}{b-a} dx + c = \frac{x}{b-a} + c, \quad c = ?$$

$$F(a) = 0, F(b) = 1 \Rightarrow \frac{a}{b-a} + c = 0 \Rightarrow c = -\frac{a}{b-a}$$

$$\Rightarrow F(x) = \frac{x-a}{b-a} = \frac{x}{b-a} - \frac{a}{b-a}$$

$$E[X] = \int_a^b x \cdot \frac{1}{b-a} dx = \frac{1}{b-a} \int_a^b x dx = \frac{1}{b-a} \left[\frac{x^2}{2} \right]_a^b = \frac{b^2 - a^2}{2(b-a)}$$

$$= \frac{(b-a)(b+a)}{2(b-a)} = \boxed{\frac{a+b}{2}}$$

$$\text{Var}[X] = \int_a^b \left(x - \frac{a+b}{2} \right)^2 \frac{1}{b-a} dx$$

let's use a different formula...

$$= E[X^2] - \mu^2$$

$$= \int_a^b x^2 \frac{1}{b-a} dx - \left(\frac{a+b}{2} \right)^2$$

$$= \frac{1}{b-a} \int_a^b x^2 dx - \frac{a^2 + 2ab + b^2}{4} = \frac{1}{b-a} \left[\frac{x^3}{3} \right]_a^b - \frac{a^2 + 2ab + b^2}{4}$$

$$= \frac{b^3 - a^3}{3(b-a)} - \frac{a^2 + 2ab + b^2}{4} = \frac{b^2 + ab + a^2}{3} - \frac{a^2 + 2ab + b^2}{4}$$

$$= \frac{4b^2 + 4ab + 4a^2 - 3a^2 - 6ab - 3b^2}{12} = \frac{a^2 - 2ab + b^2}{12}$$

$$= \frac{(b-a)^2}{12}$$

$$\Rightarrow \text{SD}[X] = \frac{b-a}{\sqrt{12}}$$

Let's compute the median of X , $\text{Med}[X] := Q[X, 0.5] = F^{-1}(0.5)$

$$F^{-1}(q) = ? \quad F(x) = q \Rightarrow \frac{x-a}{b-a} = q \Rightarrow x-a = q(b-a)$$

$$\Rightarrow x = q(b-a) + a$$

$$= qb - qa + a$$

$$= qb + a - qa$$

also known as the "quantile function"

$$\boxed{F^{-1}(q) = qb + (1-q)a}$$

$$\text{Med}[X] = F^{-1}(0.5) = 0.5b + (1-0.5)a = \frac{a+b}{2} = E[X]$$

in this case

There is a "standard uniform" rv where $a = 0$ and $b = 1$

$$X \sim U(0, 1) = \frac{1}{1-0} = 1 = f(x)$$

$\text{supp}[X] = [0, 1]$

$$F(x) = x$$

$$E[X] = \frac{1}{2} = \text{Med}[X]$$

This is one of the most widely-used rv's. It's used in every programming language e.g. `Math.random()` in Java realized a standard uniform rv.