normalize / regularize (by dividing by the entire area)

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$$R(x, \lambda) + R(x, \lambda) = |$$
 $R(x, \lambda) + R(x, \lambda) = |$ 
 $R(x,$ 

W, ..., Wx 2 Exp(S), TK = W, + ... + Wx ~ Erlang(K, x)

$$F_{K} = W_{1} + ... + W_{K} \sim F_{K} - F_{K}$$

$$P(T_{\kappa} > t) = S_{T_{\kappa}} = Q(k, \lambda t) \leftarrow$$

$$(\lambda) := \frac{e^{-\lambda} \lambda^{n}}{n!}$$

$$P(N \leq n) = F_{N}(n) = Q(n+1, \lambda) \in$$

$$P(N \geq n) = S_{N}(n) = P(n+1, \lambda) \in$$

$$P(N > h) = S_N(h) = P(h, h)$$
The "Poisson Process" is the equivalence of these two events: waiting and counting. Let's start with k=1.

$$T \sim \text{Erlang}(2, \lambda)$$

$$P(N > N) = S_N(N) = P(N, N)$$
son Process" is the equivalence of these and counting. Let's start with k=1.

$$P(N > h) = S_N(h) = P(h+1, > h)$$
son Process" is the equivalence of these and counting. Let's start with k=1.

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$$k=1$$
.

Poisson Process" is the equivalence of this gand counting. Let's start with 
$$k=1$$
.

Friang  $(2, \lambda)$ 

$$\sum_{n} \alpha \text{ Erlang}(2, \lambda)$$

$$P(t_1 > 1) = S_{t_1}(t) = Q(1, \lambda) = Q(0 + 1, \lambda) = F_N(\omega) = P(N \le 0) = R(\omega)$$

$$\frac{1}{\sqrt{(1-x)}} = S_{T_1}(x) = Q(x, x) = Q(x,$$

$$P(T_1 > 1) = S_{T_1}(1) = Q(1, \lambda) = Q(0 + 1, \lambda) = F_N(0) = P(N \leq 0)$$

$$P(T_2 > 1) = S_{T_2}(1) = Q(2, \lambda) = Q(1 + 1, \lambda) = F_N(1) = P(N \leq 1)$$

$$P(N \geq 0)$$

$$P(T_{2}>1) = 5_{T_{2}}(1) = (2,3) = Q(1+1,3) = F_{N}(1) = P(N=1)$$

$$P(T_{3}>1) = 5_{T_{3}}(1) = Q(3,3) = Q(2+1,3) = F_{N}(2) = P(N=2)$$

$$P(N=4,1,23)$$

$$O(T_3 > 1) = S_{T_3}(1) = Q(S_3)$$

$$\vdots$$
 $O(T_3 > 1) = S_{T_3}(1) = Q(K_3)$ 

$$P(T_{K}>1) = 5_{T_{K}}(1) = Q(K, \lambda) = Q((k-1)+1, \lambda) = F_{K}(k-1) = P(k \leq k-1)$$

$$P(N \leq \{0, 1, ..., k-1\}, \lambda) = F_{K}(k-1) = P(N \leq \{0, ..., k-1\}, \lambda)$$

$$P(T_{k}>1) = S_{T_{k}}(1) = Q(k)$$

$$P(T_{K}>1) = 5_{T_{K}}(1) = Q(k, \lambda) = Q((k-1)+1, \lambda) = F_{W}(k-1) = P(W \in \{0,1,...\})$$

$$N \text{ models the number of events that occur in the first second.}$$

Is there analogous relationship between the Binomial and the Negative Binomial? Yes (HW)

P(WE {0,1,.., k-13)

factorial function to 
$$\mathbb{R}$$
.

$$T \sim \text{Erlang}(k, \lambda) := \frac{\lambda^{\kappa} e^{-\lambda \cdot \ell} e^{k-1}}{(k-1)!} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{k-1}}{\tau} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{k-1}}{\tau} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{k-1}}{\tau} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{k-1}}{\tau} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{-\lambda \kappa} e^{k-1}}{\tau} \, \mathbb{1}_{\ell \geq \rho} = \frac{\lambda^{\kappa} e^{-\lambda \kappa} e^{-\lambda \kappa$$

The gamma function is a commonly used extension of the

Let's "invent" two new rv's. Recall  $\Gamma(x) = (x-1)!$  or  $\Gamma(x+1) = x!$  If  $x \in \mathbb{N}$ 

T ~ NegBin(k, p) := 
$$\binom{k+\ell-1}{k-1} (1-p)^{\frac{\ell}{2}} p^{\frac{\ell}{2}} \int_{\ell \in \mathcal{N}_p} \frac{d^{\ell}(k+\ell)}{(\ell-k)^{\frac{\ell}{2}}} (1-p)^{\frac{\ell}{2}} p^{\frac{\ell}{2}} \int_{\ell \in \mathcal{N}_p} d^{\ell} d^{$$

do they make any conceptual sense? Yes. You can imagine averaging multiple waiting times for different legal k values I can also prove that f(t) and p(t) are legal for all  $k \in (0, \infty)$ . Thus we now have two new rv's just by extending the parameter space of old rv's and employing the gamma function extension.

$$\times \sim \text{Ex+NegBin}(k, \rho) := 51me$$
extended negative binomial

Transformations of Discrete rv's
$$\times \sim \text{Bern}(\rho) = \rho^*(I - \rho)^{1-\kappa} \mathbf{1}_{x \in \{\rho, 1\}} = \begin{cases} 1 & \text{up } \rho \\ 0 & \text{up } 1 - \rho \end{cases}$$

 $Y = g(X) = X + 3 \sim \begin{cases} 4 & \text{pp} \\ 3 & \text{pl-p} \end{cases} = p^{3}(1-p)^{-\frac{1}{2}(3)}$  y = g(x) = x + 3

conjecture: 
$$\rho_{Y}(y) = \gamma_{X}(y^{-1}(y))$$

$$\rho_{Y}(y) = \rho_{X}(y^{-1}(y))$$

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 $X \sim U(\{1,2,3,4,5,6,7,8,9,10\}) = 1/10$ , the old-style PMF Y = g(X) = min{X, 3} which is not a 1:1 function

What if the inverse doesn't exist? Consider:

Assume the inverse function exists

Y = {0,1,8,..

on Supp[X]

 $\times \sim \beta ih(\beta, \beta), Y = X^{3} \sim \begin{pmatrix} h \\ y^{\frac{1}{3}} \end{pmatrix} p^{y^{\frac{1}{3}}} (1-p)^{h-y^{\frac{1}{3}}} \frac{1}{y^{\frac{1}{3}} \in \{2, 1, 2, ..., h\}}$ 

If  $Y = g(X) = X^2$ , would this strategy work? Yes... because on Supp[X] it is invertible (and that's all that matters).

Transformations for continuous rv's when g-inverse exists on Supp[X] 
$$f_{Y}(y) \neq f_{X}(g^{-1}(y)) \qquad \qquad \bigvee_{i \in [0,1]} (y = g(X) = Z) \qquad \Rightarrow \qquad X = \frac{Y}{Z} = g^{-1}(Y)$$

$$f(x) = g(X) = Z \qquad \Rightarrow \qquad X = \frac{Y}{Z} = g^{-1}(Y)$$