

CSPE32: Combinatorics and Graph Theory
Assignment-1

Finding the γ_r (Mixed Roman Domination Number) of $P_3 \circ P_n$

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CSE-B

The weak Roman domination number, denoted by $\gamma_r(G)$, is the minimum weight of a WRDF in G

Weak Roman Dominating Function requirements

We need to assign weights of 0, 1, or 2 to each vertex in such a way that:

- Every vertex with weight 0 has a neighboring vertex with weight > 0 .
- If a vertex with weight 0 is changed to 1 then a weight of 1 is reduced from one of the adjacent vertices with weight > 0 and still point one satisfies while also not affecting the edges not involved

Lexicographic Product

The **lexicographic product** of two graphs $G = (V(G), E(G))$ and $H = (V(H), E(H))$, denoted $G[H]$, is a graph defined as follows:

- **Vertices:** The vertex set of $G[H]$ is $V(G) \times V(H)$, where each vertex is represented as a pair (g, h) with $g \in V(G)$ and $h \in V(H)$.
- **Edges:** There is an edge between two vertices (g_1, h_1) and (g_2, h_2) in $G[H]$ if:
 1. $g_1 = g_2$ and $\{h_1, h_2\} \in E(H)$, or
 2. $\{g_1, g_2\} \in E(G)$.

Now consider P_3 and P_n , where $P_3 = \{u_1, u_2, u_3\}$ is the path graph on 3 vertices with edges $\{u_1u_2, u_2u_3\}$, and $P_n = \{v_1, v_2, \dots, v_n\}$ is the path graph on n vertices with edges $\{v_i v_{i+1} : 1 \leq i < n\}$.

- **Vertex Set:** The vertices of $P_3[P_n]$ are:

$$\{(u_1, v_1), (u_1, v_2), \dots, (u_1, v_n), (u_2, v_1), (u_2, v_2), \dots, (u_2, v_n), (u_3, v_1), \dots, (u_3, v_n)\}.$$

- **Edges:** Two vertices (u_i, v_j) and (u_k, v_l) are adjacent if:
 1. $u_i = u_k$ and $\{v_j, v_l\} \in E(P_n)$, or
 2. $\{u_i, u_k\} \in E(P_3)$ (i.e., $|i - k| = 1$).

Thus, the graph P_3OP_n has the following structure:

- Three copies of P_n , corresponding to u_1, u_2 , and u_3 .
- **Cross connections:**
 - Every vertex in the first copy (corresponding to u_1) is connected to every vertex in the second copy (corresponding to u_2).
 - Every vertex in the second copy (corresponding to u_2) is connected to every vertex in the third copy (corresponding to u_3).

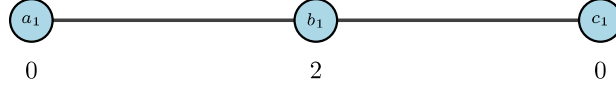
Strategy for Assigning Weights

- We need to assign weight of 0 to each of the vertex.
- Then we start by assigning a weight of 1 to each vertex with max degree, if not satisfied we bump it to 2.
- If still not satisfied we go through the b_n vertices as they are only the ones affected in terms of adjacent weight and try to assign a weight of 1 to any of them and see if it satisfies.
- If not satisfied then we just remove the extra weight given and just assign one of the vertices of a_n or c_n (either is fine) a weight of 2 as it will cover for all the b_n .

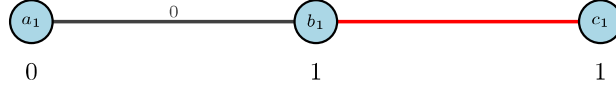
Proof

1 $P_3 \circ P_1$

Starting with $P_3 = \{a, b, c\}$ and $P_1 = \{1\}$, the lexicographic product of the two gives two vertices $V = \{a_1, b_1\}$ and edges $E = \{(a_1, b_1)\}$.



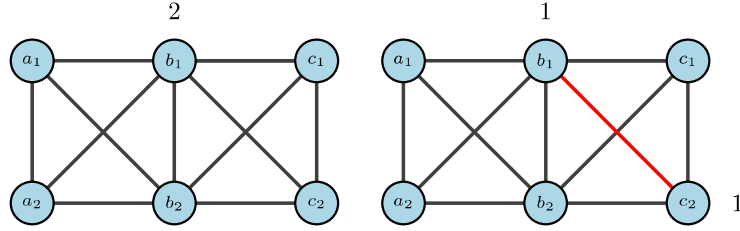
it satisfies the first function which is that b_1 with weight 2 is adjacent to c_1 , a_1 with weight 0 (i.e 1)



The γ_r for $P_3 \circ P_1=2$

2 $P_3 \circ P_2$

$P_3 = \{a, b, c\}$ and $P_2 = \{1, 2\}$, the lexicographic product of the two gives a complete graph with 4 vertices $V = \{a_1, b_1, a_2, b_2\}$ and edges $E = \{(a_1, b_1), (a_1, b_2), (a_2, b_1), (a_2, b_2), (a_1, a_2), (b_1, b_2)\}$

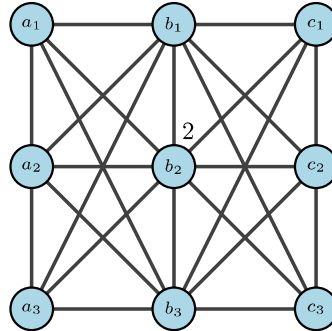


it satisfies as b_1 is connected to all vertices and when moved to another position it still has a weight of 1 to cover for the other vertices.

The γ_r for $P_3 \circ P_2=2$

3 $P_3 \circ P_3$

$P_2 = \{a, b\}$ and $P_3 = \{1, 2, 3\}$, the lexicographic product of the two gives as follows



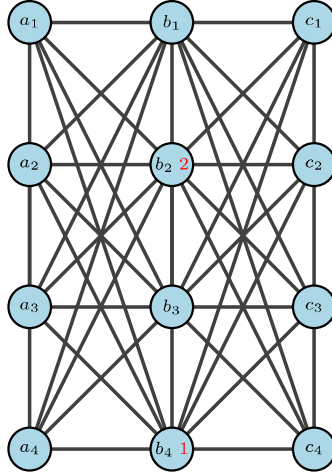
All the vertices are covered by b_2 .

We can observe that b_2 covers for all vertices with its weight of 2. If we put 1 instead of 2 then when moved from b_2 to a_1 , c_3 remains undefended.

The γ_r for $P_3 \circ P_3=2$

4 $P_3 \circ P_4$

$P_3 = \{a, b, c\}$ and $P_4 = \{1, 2, 3, 4\}$, the lexicographic product of the two gives as follows

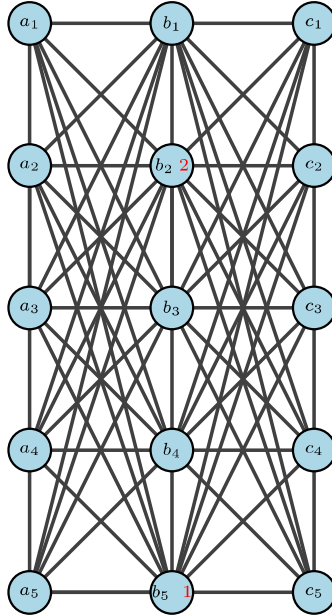


Here b_4 remains undefended so we just assign it a weight of 1 so it satisfies the first and second function.

The γ_r for $P_3 \circ P_4=3$

5 $P_3 \circ P_5$

$P_3 = \{a, b, c\}$ and $P_5 = \{1, 2, 3, 4, 5\}$, the lexicographic product of the two gives as follows



We assign b_5 the weight of 1 to cover for both the undefended b_4 and b_5 . This also ensures coverage when weight moves from b_5 to any of the possible ways.

After this for $P_3 \circ P_n$ where $n > 5$ we have b_n vertices upto 6 where if we try to assign 1 to b_5 . If moved to b_4 then b_6 remains undefended and when moved to b_6 , b_4 remains undefended so we instead of giving assigning 2 to b_5 we instead give one of the a_n with weight of 2 which will be repeated.

The γ_r ; for $P_3 \circ P_5=3$

Summary

The γ_r for $P_3 \circ P_n$ is given by:

The γ_r for $P_3 \circ P_1 = 2$,

The γ_r for $P_3 \circ P_2 = 2$,

The γ_r for $P_3 \circ P_3 = 2$,

The γ_r for $P_3 \circ P_4 = 3$,

The γ_r for $P_3 \circ P_5 = 3$,

The γ_r for the rest = 4,

Bonus next page

Bonus

Finding the γ_R (Mixed Roman Domination Number) of $P_3 \odot P_n$:

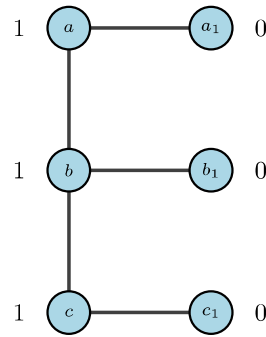
Corona Product

One copy of P_3 and 3 copies of P_n is taken. Every vertex in P_3 is connected to all the vertices of the corresponding copy of P_n .

Strategy for Assigning Weights

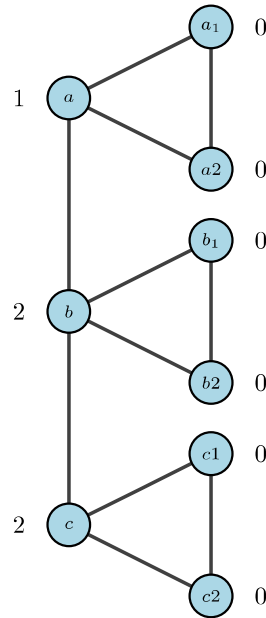
- Assign 2 to the vertices a, b, c .
- Assignment to the edges of each copy of P_n follows this pattern.
- 0, 1, 0, 2, 0, 2, 0, 0, 2, 0, 1, 0, 2, 0, 2, 0, 2, 0, 0, 2, 0, 2, 0, 1 etc.

1 $P_3 \odot P_1$



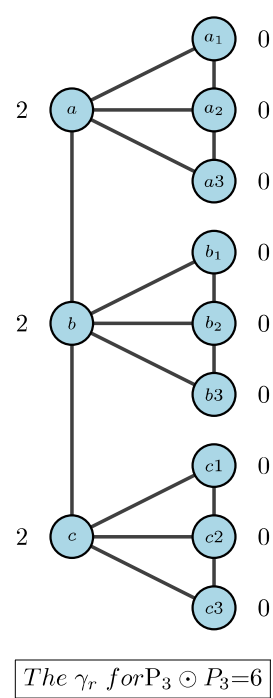
The γ_r for $P_3 \odot P_1 = 3$

2 $P_3 \odot P_2$



The γ_r for $P_3 \odot P_2 = 1$

3 $P_3 \odot P_3$



Summary

The γ_r for $P_3 \odot P_n$ is given by:

- The γ_r for $P_3 \odot P_1 = 3$,
- The γ_r for $P_3 \odot P_2 = 3$,
- The γ_r for the rest = 6.

Thank You