

Validity and Satisfiability T

Decide which sentences are

- valid
- satisfiable
- unsatisfiable
- neither

a. $\text{light} \Rightarrow \text{light}$

Answer

Consider the Truth Table.

light	$\text{light} \Rightarrow \text{light}$	$\neg \text{light} \vee \text{light}$
F	T	T
T	T	T

it is valid!

b. $\text{Sun} \Rightarrow \text{light}$

Answer

Consider the Truth Table:

Sun	light	$\text{Sun} \Rightarrow \text{light}$	$\neg \text{Sun} \vee \text{light}$
F	F	T	T
F	T	T	T
T	F	F	F
T	T	T	T

it is satisfiable!

$$c. ((\text{Cloudy} \vee \text{Rain}) \Rightarrow \text{Dark}) \Leftrightarrow (\text{Cloudy} \Rightarrow \text{Dark}) \vee (\text{Rain} \Rightarrow \text{Dark})$$

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A

B

$S : A \Leftrightarrow B$
 $B = B_1 \vee B_2$

Answer

Consider the truth Table:

Cloudy	Rain	Dark	A	B_1	B_2	B	S
F	F	F	T	T	T	T	T
F	F	T	T	T	T	T	T
F	T	F	T	T	F	T	T
F	T	T	T	T	T	T	T
T	F	F	T	F	T	T	T
T	F	T	F	T	F	T	T
T	T	F	F	F	F	F	T
T	T	T	T	T	T	T	T

Consider the truth Table:
 $\neg(\text{Cloudy} \vee \text{Rain}) \vee \text{Dark}$
 $\neg(\text{Cloudy} \vee \text{Dark})$
 $\neg(\text{Rain} \vee \text{Dark})$

valid

$$d. \neg(\text{fat} \vee \text{happy} \vee (\text{fat} \Rightarrow \text{happy}))$$

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Answer Consider the Truth Table

fat	happy	$\text{fat} \vee \text{happy}$	$\text{fat} \Rightarrow \text{happy}$	$\neg(\text{fat} \vee \text{happy})$	$\neg(\text{fat} \Rightarrow \text{happy})$
F	F	F	T	T	F
F	T	T	T	T	F
T	F	T	F	F	F
T	T	T	T	T	F

Diagram showing the final row of the truth table is circled in green, indicating it is unsatisfiable.

unsatisfiable

Entailment Example

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Prove that sentence S_1 entails sentence S_2

where: $S_1: (\underbrace{\text{Tall} \Rightarrow \text{handsome}}_{S_A}) \wedge (\underbrace{\neg \text{Tall} \wedge \text{skinny} \Rightarrow}_{\text{skinny}} \underbrace{\neg \text{skinny}}_{S_B})$

$S_2: \text{Tall} \wedge \text{Skinny} \Rightarrow \text{Handsome}$. $\downarrow S_B$

Answer: Consider the Truth Table.

				S_A		S_B		S_1		S_2		$\neg \text{Tall} \wedge \neg \text{Skinny} \Rightarrow \neg \text{Handsome}$	
				$\neg \text{Tall} \vee \text{Handsome}$		$\neg \text{Tall} \vee \text{Skinny}$		$\neg \text{Tall} \vee \neg \text{Skinny}$		$\neg \text{Tall} \vee \neg \text{Handsome}$		$\neg \text{Tall} \vee \neg \text{Skinny} \vee \neg \text{Handsome}$	
Tall	Skinny	Handsome		S_A		S_B		S_1		S_2		$S_1 = S_A \wedge S_B$	
F	F	F	F	T	F	F	F	F	F	T	T	F	F
F	F	F	T	T	F	F	F	F	F	T	T	F	F
F	T	F	F	T	F	F	F	F	F	T	T	F	F
F	T	T	F	T	F	F	F	F	F	T	T	F	F
T	F	F	F	F	T	T	T	F	T	F	T	F	F
T	F	T	T	T	F	F	T	F	T	F	T	F	F
T	T	F	F	T	T	F	T	T	F	F	T	F	F
T	T	T	T	T	T	T	T	T	T	T	T	T	T

S_1 entails S_2 is true
because S_2 is models where
 S_1 is true.

Same entailment, a different answer, using inference rules instead of model checking. 5

We want to prove that $S_1 \models S_2$,
with $S_1: (\text{Tall} \Rightarrow \text{Handsome}) \wedge (\neg \text{Tall} \wedge \text{Skinny} \Rightarrow \text{Skinny})$
 $S_2: \text{Tall} \wedge \text{Skinny} \Rightarrow \text{Handsome}.$

Answer:
• let us use the inference rule: $(a \Rightarrow b) = \neg a \vee b$

$S_1 \models S_2$ becomes:
 $(\neg \text{Tall} \vee \text{Handsome}) \wedge (\neg(\neg \text{Tall} \wedge \text{Skinny}) \vee \text{Skinny})$
 $\vdash (\neg(\neg \text{Tall} \wedge \text{Skinny}) \vee \text{Handsome})$

• let us use the de Morgan inference rule:
 $(\neg \text{Tall} \vee \text{Handsome}) \wedge (\text{Tall} \vee \neg \text{Skinny} \vee \text{Skinny}) \vdash$
 $(\neg \text{Tall} \vee \neg \text{Skinny} \vee \text{Handsome}).$

• But $\neg \text{Skinny} \vee \text{Skinny} = T$
 $(\neg \text{Tall} \vee \text{Handsome}) \wedge (\text{Tall} \vee T) \vdash (\neg \text{Tall} \vee \neg \text{Skinny} \vee \text{Handsome})$

• But $\text{Tall} \vee T$ shall always be TRUE
 $(\neg \text{Tall} \vee \text{Handsome}) \vdash (\neg \text{Tall} \vee \text{Handsome} \vee \neg \text{Skinny})$

because if $x = \neg \text{Tall} \vee \text{Handsome}$, this

is $x \vdash x \vee \neg \text{Skinny}$
therefore, whatever X is True, $x \vee \neg \text{Skinny}$ shall also be true.

CNF Translation

Consider the sentence:

$(\text{Unicorn} \Rightarrow \text{Mythical}) \Leftrightarrow (\text{Immortal} \vee \text{Imaginary})$

translate this sentence in CNF

Answer

1. Eliminate \Leftrightarrow using the inference rule:
 $\alpha \Leftrightarrow \beta = (\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)$

$((\text{Unicorn} \Rightarrow \text{Mythical}) \Rightarrow (\text{Immortal} \vee \text{Imaginary})) \wedge$

$(\text{Immortal} \vee \text{Imaginary}) \Rightarrow (\text{Unicorn} \Rightarrow \text{Mythical})$

2. Eliminate \Rightarrow using the inference rule
 $\alpha \Rightarrow \beta = \neg \alpha \vee \beta$

$(\neg(\neg \text{Unicorn} \vee \text{Mythical}) \vee (\text{Imortal} \vee \text{Imaginary})) \wedge$

$(\neg(\text{Immortal} \vee \text{Imaginary}) \vee (\neg \text{Unicorn} \vee \text{Mythical}))$

3. Push \neg in: using the deMorgan inference rules

3. Push \neg in: using the deMorgan inference rules

$((\text{Unicorn} \wedge \neg \text{Mythical}) \vee (\text{Immortal} \vee \text{Imaginary})) \wedge$

$\wedge ((\neg \text{Immortal} \wedge \neg \text{Imaginary}) \vee (\neg \text{Unicorn} \vee \text{Mythical}))$

4. Apply the distributivity inference rules.

$(\text{Unicorn} \vee \text{Immortal} \vee \text{Imaginary}) \wedge (\neg \text{Mythical} \vee \text{Immortal} \vee \text{Imaginary})$

$\wedge (\neg \text{Immortal} \vee \neg \text{Unicorn} \vee \text{Mythical}) \wedge (\neg \text{Imaginary} \vee \neg \text{Unicorn} \vee \text{Mythical})$

$D_1 \oplus D_2 \oplus D_3 \oplus D_4$
 Conjunctions

$D_1, D_2, D_3 \text{ and } D_4$ are
 in Normal Form
 (only disjunctions)