

Inference in FOL

1

Examples of finding the Most General Unifier (MGU)

Example 1: Find the MGU of:

$$X = f(g(*, u(y, B), n(B, y)))$$

and

$$Y = f(g(x, z, w))$$

Solution: Observation: Both X and Y are FOL atomic sentences that use the same predicates f and g

Let us look more carefully at X & Y

$$X = f(g(\cancel{\text{arg}}_1^g, \cancel{\text{arg}}_2^g, \cancel{\text{arg}}_3^g))$$

$$Y = f(g(\cancel{\text{arg}}_1^g, \cancel{\text{arg}}_2^g, \cancel{\text{arg}}_3^g))$$

wee: $\left\{ \begin{array}{l} \cancel{\text{arg}}_1^g = x \\ \cancel{\text{arg}}_3^g = n(B, y) \\ \cancel{\text{arg}}_1^g = x \quad \cancel{\text{arg}}_2^g = z \quad \cancel{\text{arg}}_3^g = w \end{array} \right.$

Therefore :

$$\text{arg}_1^f = \alpha$$

$$\text{arg}_1^{f'} = \alpha$$

$$\text{arg}_2^f = m(y, B)$$

$$\text{arg}_2^{f'} = z$$

$$\text{arg}_3^f = n(B, y)$$

$$\text{arg}_3^{f'} = w$$

How do we decide on the substitution?

$$\theta = \{ z / m(y, B); w / n(B, y) \}$$

How did the UNIFY function work?

X and Y
are compounds

having the
same operators (f and g)

$$X. \text{Args} = \{ \alpha, m(y, B), n(B, y) \}$$

$$Y. \text{Args} = \{ \alpha, z, w \}$$

$$\begin{aligned} &\Downarrow \\ &\text{UNIFY}(\alpha, \alpha) \Rightarrow \theta = \{ \} \\ &\text{UNIFY}(m(y, B), z) \Rightarrow \end{aligned}$$

$$\theta = \{ z / m(y, B) \}$$

$$\begin{aligned} &\text{UNIFY-VR}(z, m(y, B)) \\ &\text{occur-check}(z, w) \quad \cancel{\theta} \\ &\text{UNIFY-VR}(w, n(B, y)) \end{aligned}$$

$$\begin{aligned} &\text{UNIFY}(n(B, y), w) \\ &\theta = \{ z / m(y, B), w / n(B, y) \} \end{aligned}$$

$$\text{MGCS} = f(f(\alpha, m(y, B), n(B, y)))$$

$$\Downarrow \text{UNIFY}(X, Y) = \theta$$

Example 2: Find the MGV of:

$$X = C(f(x), y, g(B))$$

$$Y = C(f(x), f(x), *)$$

$$MGU = \{y/f(x); x/g(B)\}$$

$$MGS = C(f(g(B)), f(g(B)), g(B))$$

How did the UNIFY function work?

X and Y are compounds!

\Rightarrow operator f is relying on its unifat
 $UNIFY-VAR \Rightarrow y/f(x) \} MGU$
 $UNIFY-VAR \Rightarrow x/g(B) \} MGU$

[4]

Example 3 Find the MOU of:

Mother (child(x), Jane, Nephew (Cousin
(Uncle (Jack))))

and

Mother (child (John), y, Nephew (Cousin(y)))

x / John ; y / Jane ; y / Uncle (Jack)
→ occurs - check = FAIL

Standardize Apart:

Mother (child(x), Jane, Nephew (Cousin (Uncle (Jack))))

Mother (child (John), y, Nephew (Cousin (z)))

MOU = {x / John ; y / Jane ; z / Uncle (Jack)}

Forward Chaining Examples

5

Example 1

To eat a ~~dog~~ mouse, a cat should be hungry and chase the mouse instead of fighting dogs, which do not eat mice.

P1: $\text{cat}(x) \wedge \text{mouse}(y) \wedge \text{dog}(z) \wedge \text{chase}(x, y) \wedge \text{eat}(z, y) \wedge \text{fight}(x, z) \Rightarrow \text{eat}(x, y)$

To chase the mouse, both the eat and the mouse should be in the house and the cat should be healthy.

P2: $\text{cat}(x) \wedge \text{mouse}(y) \wedge \text{in-house}(x) \wedge \text{in-house}(y) \wedge \text{healthy}(x) \wedge \text{not}(\text{dead}(x)) \Rightarrow \text{chase}(x, y)$

Cats that are friends with dogs do not fight.

P3: $\text{cat}(x) \wedge \text{dog}(z) \wedge \text{friends}(x, z) \Rightarrow \text{not}(\text{fight}(x, z))$

kitty is a cat in the house.

P4: $\text{Cat}(\text{kitty})$

P5: $\text{in-house}(\text{kitty})$

Fido is our in house dog.

P6: $\text{dog}(\text{Fido})$

P7: $\text{in-house}(\text{Fido})$

Mickey, the mouse, entered the house.

P8: $\text{mouse}(\text{Mickey})$

P9: $\text{in-house}(\text{Mickey})$.

kitty is healthy.

P10: healthy(kitty)

Fido and kitty are friends.

P11: friends(kitty, Fido)

Dogs do not eat mice.

P12: dog(\exists) \Rightarrow eat(\exists, \forall)
 \wedge mouse(\forall)
 mouse

Prove that kitty eat Micky.

Q: eat(kitty, Micky).

Forward chaining

a) Which sentences have implications?

P1, P2, P3, P12

b) Find sentences from the KB
 which unify with premises
 of P1, P2, P3 or P12

Iteration 1: $\frac{P_1 \wedge P_4}{P_1}$: $\theta_1 \models x / \text{kitty}$

$\text{cat}(\text{kitty}) \not\models P_1$ has unsatisfied
 premises.

$\frac{P_2 \wedge P_8}{P_2}$: $\theta_2 \models y / \text{Micky}$
 $\text{mouse}(\text{Micky}) \not\models P_2$ has unsatisfied
 premises.

P₃ & P₆ : θ₃ : $\exists z/Fido)$
dog(Fido)

P₃ has unsatisfied premises

P₁₂ \Rightarrow all premises satisfied.

dog(Fido) \wedge house(Mickey)

$\boxed{Q_1: \exists eat(Fido, Mickey)}$

Iteration 2: • P₁ \wedge θ₁ \wedge θ₂ \wedge θ₃

$\exists Q_1$

cat(kitty) \wedge mouse(Mickey) \wedge dog(Fido)

\wedge chase(kitty, Fido)

\wedge eat(Fido, Mickey)

\wedge fight(kitty, Fido)

P₁ has unsatisfied premises.

• P₂ \wedge P₅ \wedge θ₃

cat(kitty) \wedge in-house(kitty)

\wedge healthy(kitty) \wedge in-house(Mickey)

P₂ has unsatisfied premises

• P₃ \wedge P₁₁ cat(kitty) \wedge dog(Fido) \wedge Preyed(kitty, Fido)

P₃ has all premises satisfied!

$\boxed{Q_2: \exists fight(kitty, Fido)}$

Iteration 3

- P_1 premises:

$\text{cat(kitty)} \wedge \text{mouse(Mickey)} \wedge \text{dog(Fido)}$

$\wedge \underline{\text{chase(kitty, Fido)}}$

$\wedge \underline{\text{eat(Fido, Mickey)}} \wedge \text{fight(kitty, Fido)}$

still has 1 unsatisfied premise!

- $P_2 \wedge P_9$

$\text{cat(kitty)} \wedge \text{in-house(kitty)} \wedge \text{in-house(Mickey)}$

$\wedge \underline{\text{healthy(kitty)}}$

$\hookrightarrow P_2$ still has 1 unsatisfied premise!

Iteration 4: P_1 has unsatisfied premise:

$\wedge \text{chase(kitty, Fido)}$

$P_2 \wedge P_{10} \Rightarrow P_2$ has all premises satisfied!

$\boxed{Q_3: \text{chase(kitty, Mickey)}}$

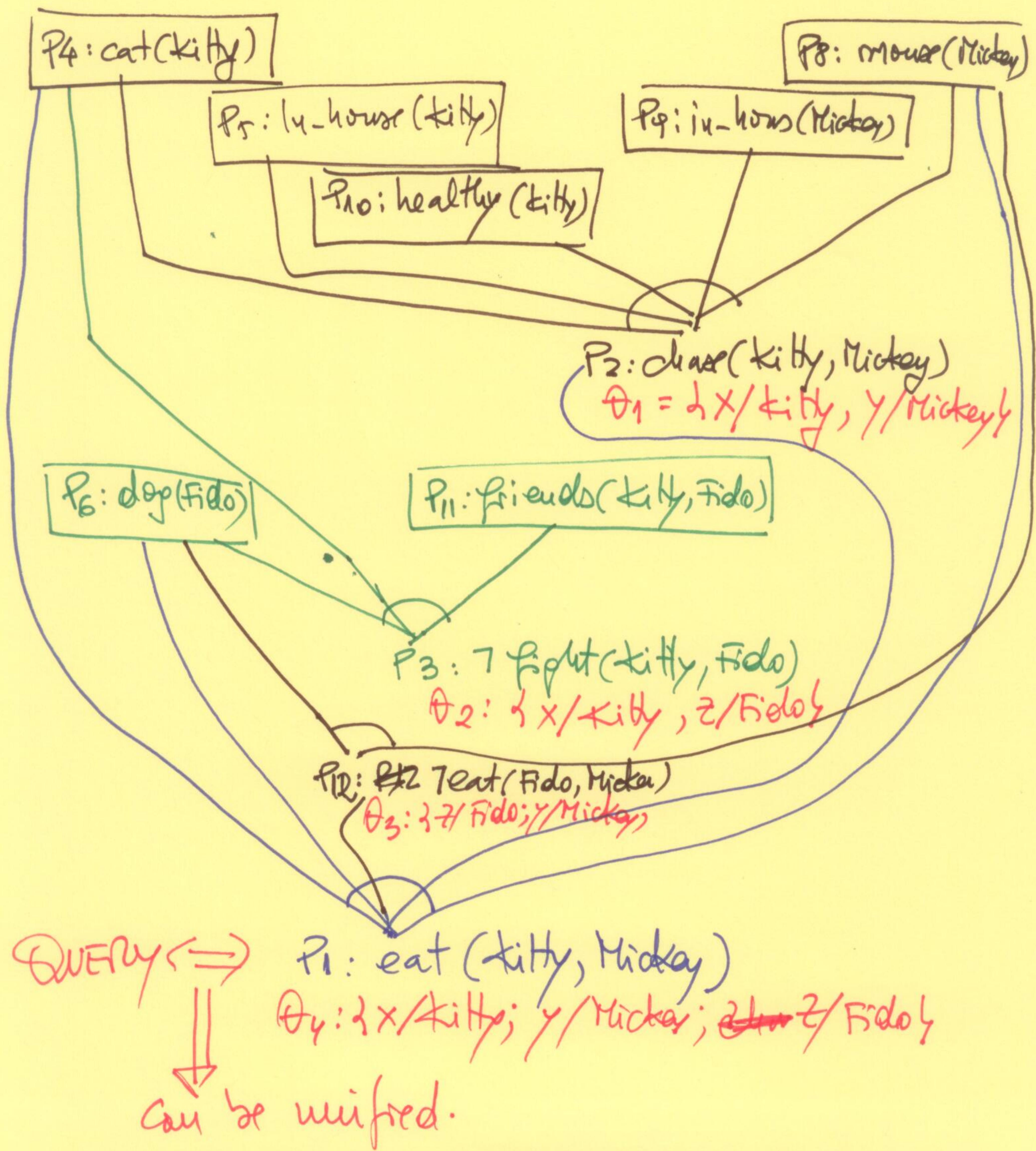
Iteration 5: $P_1 \wedge Q_3 \rightarrow P_1$ has all premises satisfied!

$\boxed{Q_4: \text{eat(kitty, Mickey)}}$

Can unify with the query Δ

FOL Feed forward Chaining up using AND-OR Graphs + Substitutions

9



2nd Example of FOL Forward Chaining

10

1. Only athletes that win gold medals at the Olympics are considered Olympic champions.
2. The swimmer Michael Phelps won six gold medals at the 2004 Summer Olympics.
3. Is Michael Phelps an Olympic champion?

Encoding in FOL

1. $\text{athlete}(a) \wedge \text{win}(a, \text{Gold-Medal}, \theta) \rightarrow \text{Olympic-Champion}(a)$
 $\wedge \text{Olympics}(\theta)$
 2. swimmer(MF)
 3. $\text{win(MF, Gold-Medal, 2004)} \wedge \text{Olympics(2004)}$
 - Q: $\text{Olympic-champion(MF)}$
- We also need 4. $\text{swimmer}(x) \Rightarrow \text{athlete}(x)$
encoding world knowledge
(commonsense knowledge)

Solution:

Transform 3 into

2. swimmer(MF)

4. athlete(MF)
 $\theta_1 \& x / MF$

3. a. $\text{win(MF, Gold-Medal, 2004)}$
3. b. Olympics(2004)

3. a. $\text{win(MF, Gold-Medal, 2004)}$
3. b. Olympics(2004)

$\text{Olympic-Champion(MF)} \Leftarrow Q$
 $\theta_2 \& a / MF, \theta / 2004$

FOL Forward Chaining Example #3

Qual-level Problem!

Suppose your KB contains the following Horn clauses:

1. Descendant(child(x), x)
2. Descendant(x , y) \wedge Descendant(y , z) \Rightarrow Descendant(x , z)

Consider the Forward-Chaining algorithm, that on the j -th iteration terminates if the KB contains a sentence that unifies with the query, else it adds to the KB every atomic sentence that can be inferred from the sentences already in the KB after iteration $j-1$.

For each of the following queries, say whether the Forward-Chaining algorithm will give an answer, and if so, write down the answer:

Q1: Descendant(child(y), John)

Answer: YES : A: { y/x ; x/John }
 \rightarrow Descendant(child(John), John)

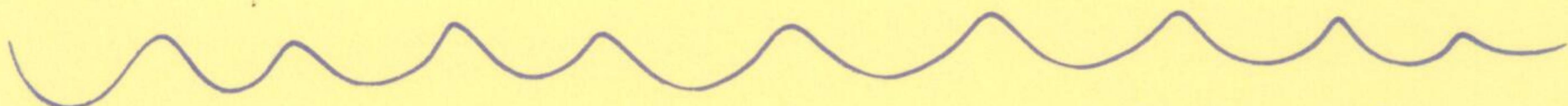
Q2: Descendant(child(child(y)), John)

YES ; After 1st Iteration: x/John
 After 2nd Iteration, it stops as $\text{John} = \text{child}(y)$

Q3: Descendant(child(Mary), child(~~child~~(Mary)))

Q3 : Descendant (Child (Mary),
Child (Child (Mary)))

Answer: $\theta: \exists x/Mary^y$ \Rightarrow FC does not
wind sentence 1 terminate



FOL Backward Chaining Examples

let us consider again the Colonel West example.

KB

P1: American(x) \wedge Weapon(y) \wedge Nation(z) \wedge
 \vdots \wedge Hostile(z) \wedge Sells(x, y, z) \Rightarrow Criminal(y)

P2a Owns (Nono, N1)
b Missile (M1)
c Nation (Nono)

P3. Missile (x) \wedge Owns (Nono, x) \Rightarrow Sells (West, x, Nono)

P4: $\forall x$ Missile(x) \Rightarrow Weapons(x)

P5: $\forall x$ Enemy(x, America) \Rightarrow Hostile(x)

P6: American (West)

P7: Enemy (Nono, America)

Query : Criminal (West)

The Backward Chaining Algorithm

13

FOL - BC - ASK (KB , goals, θ)

a list of conjuncts
forming a query
current substitution

goals: Criminal(West)? ~~or not~~ ← first(Goals)

Goals 1

$g' \leftarrow$ a) where do I have in KB an
axiom with the conclusion unifiable
to first(Goals)?

* P1
b) what is the substitution that
would enable the unification?

$\theta(\text{?}/\text{West})$

c) Are the preconditions of P_1
standardized-apart?

YES NO YES

\Rightarrow new-goals $\leftarrow \{ \text{Weapon}(y), \text{Nation}(z), \text{Hostile}(z), \text{Sells}(\text{West}, y, z) \}$

answers \leftarrow FOL - BC - ASK (KB , new-goals,
 $\text{Compound}(\theta, \theta')$)
values.

↓
goals \leftarrow new-goals.

Goals 2

$p' \leftarrow \text{Weapon}(y)$

a) where do I have in KB an
axiom with conclusion
 $\text{Weapon}(y)$?
* P4

b) What substitution would enable the unification of P' and P_4 ?

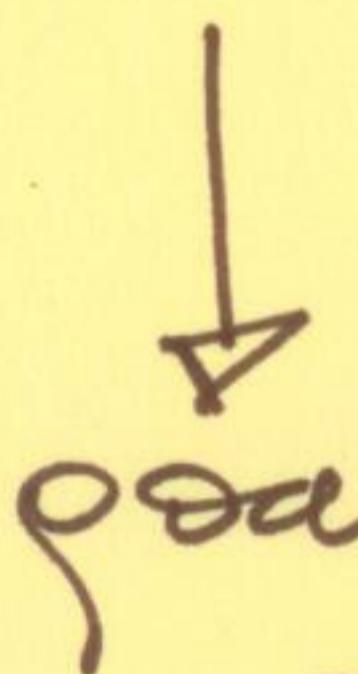
$$\theta' = \{y/x\}$$

c) Are the preconditions of P_1 standardised apart?

YES

\Rightarrow new-goals $\leftarrow \{ \text{Missile}(x), \text{Nation}(z), \theta \in \theta \cup \theta' \text{ Hostile}(z), \text{Sells}(\text{West}, y, z) \}$

answers $\leftarrow \text{FOL-BC-ASK}(\text{KB}, \text{new-goals}, \text{cnpurp}(\theta', \theta)) \setminus \text{vacuous}$



goals \leftarrow new-goals

GOALS

$P' \leftarrow \text{Missile}(x)$

a) Where do I have in KB an axiom with conclusion $\text{Missile}(x)$?

b) What substitution will enable the unification of P' and $P_2.b$?

$$\theta' = \{x/M1\}$$

c) Are these preconditions to $P_2.b \Rightarrow \text{No}$
 But we need to standardise x into $M1$ Noticing to standardise apart

\Rightarrow new-goals $\leftarrow \{ \text{Nation}(z), \text{Hostile}(z), \text{Sells}(\text{West}, M1, z) \}$
 $\leftarrow \text{REST}(\text{new-goals})$

$$\theta \leftarrow \{x/\text{West}; y/M1\}$$

answers $\leftarrow \text{FOL-BC-ASK}(\text{KB}, \text{new-goals}, \text{cnpurp}(\theta', \theta)) \setminus \text{vacuous}$



\downarrow
goals \leftarrow new-goals

GOAL-4

$f' \leftarrow \text{Native}(z)$

- a). where do I have in KB an axiom
with the conclusion unifiable
to f' ? $\Rightarrow P2.C$

- b). what substitution will enable the
unification of f' with $P2.C$?

$$\theta' = \{ z / \text{Nono} \}$$

- c). Are there pre-conditions to $P2.C$? \Rightarrow No!

\Rightarrow new-goals $\leftarrow \{ \text{Hostile}(\text{Nono}), \text{Sells}(\text{West}, M1, \text{Nono}) \}$

$$\theta = \{ x / \text{West}, y / M1, z / \text{Nono} \}$$

answers $\leftarrow \text{FOL-BC-ASK}(\text{KB}, \text{new-goals},$
 $\text{unif}\theta(\theta, \theta'))$
 \downarrow
values

goals \leftarrow new-goals

GOAL-5

$f' \leftarrow \text{Hostile}(\text{Nono})$

- a). where do I have in KB an axiom
with the conclusion $\text{Hostile}(\text{Nono})$?

$\Rightarrow P5$

- b). what substitution will enable the
unification of f' with $P5$?

$$\theta' = \{ x / \text{Nono} \}$$

- c). Are there pre-conditions to $P5$
that need to be standardized
apart?

\downarrow
 $\text{Enemy}(\text{Nono}, \text{America})$

\Rightarrow new-goals $\leftarrow \{ \text{Enemy}(\text{Nono}, \text{America}), \text{Sells}(\text{West}, M1, \text{Nono}) \}$

\downarrow
 answers $\leftarrow \text{FOL-BC-Ask}(\text{KB}, \text{new-goals}, \text{compared } (\theta, \theta')) \cup \text{answers}$

\downarrow
 $\text{goals} \leftarrow \text{new-goals}$

GOKZ-6 $f^1 \leftarrow \text{Enemy}(\text{Noro}, \text{America})$

- a) • where do I find in KB an axiom with conclusion $\text{Enemy}(x, y)$?

$\Rightarrow P7$

- b) • what substitution will enable the unification of f^1 with $P7$?

c) all standardized apart! $\theta' = \{y/x\}$
 $\Rightarrow \text{new-goals} \leftarrow \text{sells}(\text{West}, \text{M1}, \text{Nano})$

\downarrow

answers $\leftarrow \text{FOL-BC-Ask}(\text{KB}, \text{new-goals}, \text{compared } (\theta', \theta)) \cup \text{answers}$

\downarrow
 $\text{goals} \leftarrow \text{new-goals}$

GOKZ-7 $f^1 \leftarrow \text{sells}(\text{West}, \text{M1}, \text{Nano})$

- a) • where do I find in KB an axiom with conclusion $\text{sells}(a, b, c)$?

b) $\Rightarrow P3$
 b) • what substitution will enable the unification of f^1 with $P3$?

$\theta' = \{a/West, b/M1, c/Nano\}$

c) all standardized apart.

$\Rightarrow \text{new-goals} = \emptyset$

\downarrow
 answers $\leftarrow \text{FOL-BC-Ask}(\text{KB}, \emptyset, \text{compared } (\theta, \theta')) \cup \text{answers}$

$\theta = \{a/West, b/M1, c/Nano\}$

\Rightarrow return $\underline{\theta} \quad \underline{\text{PROVEN}}$

2nd Example of FOL - BC Proof Tree

17

let us revisit example I used for FOL forward-chaining
(pages 5-9)

KB

P1: cat(x) \wedge mouse(y) \wedge dog(z) \wedge chase(x, y) \wedge eat(z, y)
 \wedge right(x, z) \Rightarrow eat(x, y)

P2: cat(x) \wedge mouse(y) \wedge in-house(x) \wedge in-house(y) \wedge
 \wedge healthy(x) \Rightarrow chase(x, y)

P3: cat(x) \wedge dog(z) \wedge friends(x, z) \Rightarrow right(x, z)

P4: cat(kitty)

P5: in-house(kitty)

P6: dog(Fido)

P7: in-house(Fido)

P8: mouse(Mickey)

P9: in-house(Mickey)

P10: healthy(kitty)

P11: friends(kitty, Fido)

P12: dog(z) \wedge mouse(y) \Rightarrow eat(z, y)

query: eat(kitty, Mickey)

Q: Eat(kity, Mickey) Goal +

18

