Worksheet 11: Graphs

Question 1.

- (A) K_n (the complete graph on n vertices) has edges and vertices.
- **(B)** $K_{m,n}$ (the complete bipartite graph on sets of sizes m, n) has edges and vertices.
- (C) W_n (wheel graph on n vertices n-gonal pyramid viewed from above) has edges and vertices.
- **(D)** Q_n (*n*-dimensional cube) has edges and vertices.

Question 2 (Rosen7e, Ch.10, Q15-Q17).

- (A) The length of the longest simple circuit in K_5 is
- **(B)** The length of the longest simple circuit in W_{10} is
- (C) The length of the longest simple circuit in $K_{4,10}$ is

Note. A simple circuit in (Rosen2019) is defined as a circular sequence of vertices $v_0, v_1, \ldots, v_n = v_0$, where each two neighboring vertices are connected by an edge (and it does not contain any edge more than once). It can return to the same vertex multiple times.

Question 3 (Rosen7e, Ch.10, Q19-Q24). In each example find the dimensions of a matrix; and number of 0s and 1s in it: Find *X*, *Y*, *Z*, *T*.

- (A) The adjacency matrix for $K_{m,n}$ has size (rows times columns) $X \times Y$; it has Z 0's and T 1's.
- **(B)** The adjacency matrix for K_n has size $X \times Y$; it has Z 0's and T 1's.
- (C) The adjacency matrix for C_n has size $X \times Y$; it has Z 0's and T 1's.
- **(D)** The adjacency matrix for Q_4 has size $X \times Y$; it has Z 0's and T 1's.
- (E) The incidence matrix for W_n has size $X \times Y$; it has Z 0's and T 1's.
- (**F**) The incidence matrix for Q_5 has size $X \times Y$; it has Z 0's and T 1's.

Note. Adjacency matrix is a square matrix of size $|V| \times |V|$, but incidence matrix is a rectangular matrix of size $|V| \times |E|$.

Question 4 (Rosen7e, Ch.10, Q28-Q31).

- (A) List all positive integers n such that K_n has an Euler circuit; what is its length in terms of n?
- **(B)** List all positive integers n such that Q_n has an Euler circuit.

Question 5 (Rosen7e, Ch.10, Q43).

If G is a planar connected graph with 12 regions and 20 edges, then G has vertices.

Question 6 (Rosen7e, Ch.10, Q44).

If G is a planar connected graph with 20 vertices, each of degree 3, then G has regions.

Question 7 (Rosen7e, Ch.10, Q45).

If a regular graph G has 10 vertices and 45 edges, then each vertex of G has degree

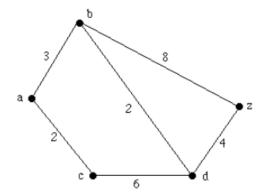
Note. A *regular graph* is a graph where all vertices have the same degree.

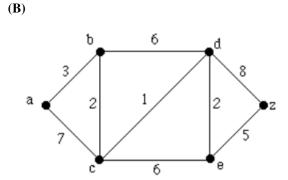
Question 8 (Rosen7e, Ch.10, Q59-Q82).

- (A) A simple graph with 6 vertices, whose degrees are 2, 2, 2, 3, 4, 4.
- **(B)** A simple graph with 8 vertices, whose degrees are 0, 1, 2, 3, 4, 5, 6, 7.
- **(C)** A simple graph with degrees 1, 2, 2, 3.
- **(D)** A simple graph with degrees 2, 3, 4, 4, 4.
- (E) A simple graph with degrees 1, 1, 2, 4.
- **(F)** A simple digraph with indegrees 0, 1, 2 and outdegrees 0, 1, 2.
- **(G)** A simple digraph with indegrees 1, 1, 1 and outdegrees 1, 1, 1.
- (H) A simple digraph with indegrees 0, 1, 2, 2 and outdegrees 0, 1, 1, 3.
- (I) A simple digraph with indegrees 0, 1, 2, 4, 5 and outdegrees 0, 3, 3, 3, 3.
- (**J**) A simple digraph with indegrees 0, 1, 1, 2 and outdegrees 0, 1, 1, 1.
- (K) A simple digraph with indegrees: 0, 1, 2, 2, 3, 4 and outdegrees: 1, 1, 2, 2, 3, 4.
- (L) A simple graph with 6 vertices and 16 edges.
- (M) A connected simple planar graph with 5 regions and 8 vertices, each of degree 3.
- (N) A graph with 4 vertices that is not planar.
- (O) A planar graph with 10 vertices.
- **(P)** A planar graph with 8 vertices, 12 edges, and 6 regions.
- (Q) A planar graph with 7 vertices, 9 edges, and 5 regions.

Question 9 (Rosen7e, Ch.10, Q108).

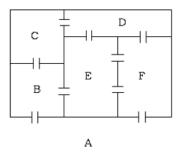
Use Dijkstras Algorithm to find the shortest path length between the vertices a and z in these weighted graphs.





Question 10 (Rosen7e, Ch.10, Q113).

The picture at the right shows the floor plan of an office. Show that it is impossible to plan a walk that passes through each doorway exactly once, starting and ending at A.



Hall's Marriage Theorem (Rosen2019, p.772)

The bipartite graph G = (V, E) with partition of vertices into 2 disjoint sets $V = X \cup Y$ has a *maximum matching* that saturates X iff for all $A \subseteq X$ we have $|X| \le |N(X)|$.

Note. A *matching* in a graph is a set of of edges such that no two edges share a common endpoint. A *maximum matching* is matching containing the greatest number of edges. And a matching *saturates* a set X, if each vertex $v \in X$ belongs to some matching edge.

Question 11 (Königs Marriage Theorem)

Prove that if all the vertices of a bipartite graph have the same degree, then it has a perfect matching. (Quines2017, p.11); https://cjquines.com/files/halls.pdf

Note. A matching is *perfect*, if it saturates all vertices (every vertex has a pair).

Question 12

We have a regular deck of 52 playing cards, with exactly 4 cards of each of the 13 ranks. The cards have been randomly dealt into 13 piles, each with 4 cards in it. Prove that there is a way to take a card from each pile so that after we take a card from every pile, we

have exactly a card of every rank. (Quines2017, p.11). **Answers**

Question 6. Answer: TBD

Question 1. Answer:

Question 7. Answer: TBD

Question 8. Answer: TBD

(A) K_n has $\frac{n(n-1)}{2}$ edges and n vertices. (B) $K_{m,n}$ has $m \cdot n$ edges and m + n vertices. (C) W_n has 2n edges and n + 1 vertices. (D) Q_n has $n \cdot 2^{n-1}$ edges and 2^n vertices.

The number of edges in Q_n could be first expressed by a recurrent formula (and then proven by mathematical induction).

Question 9. Answer: TBD

Question 2. Answer: TBD

Question 10. Answer: TBD

Question 3. Answer: TBD

Question 11. Answer: TBD

Question 4. Answer: TBD **Question 5.** Answer: TBD

Question 12. Answer: TBD