

# Posit Arithmetic Representation

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# Deliverables

- Analysing Posit-Arithmetic Representation.
- Documentation and Evaluation.

# Posit Arithmetics

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## What's wrong with IEEE 754?

- It's a guideline, not a standard.
- No guarantee of identical results across systems.
- Invisible rounding errors; the inexact flag is useless.
- Breaks algebra laws, like  $a+(b+c) = (a+b)+c$
- Overflows to infinity, underflows to zero.
- No way to express most of the real number line.

# Example of invisible Rounding errors

## Why worry about floating point?

Find the scalar product  $a \cdot b$ :

$$\begin{aligned}a &= (3.2e8, 1, -1, 8.0e7) \\ b &= (4.0e7, 1, -1, -1.6e8)\end{aligned}$$

**Note:** All values are integers that can be expressed *exactly* in the IEEE 754 Standard floating-point format (single or double precision)

Single Precision, 32 bits:  $a \cdot b = 0$

Double Precision, 64 bits:  $a \cdot b = 0$

Correct answer:  $a \cdot b = 2$

**Most** linear algebra is unstable with floats!

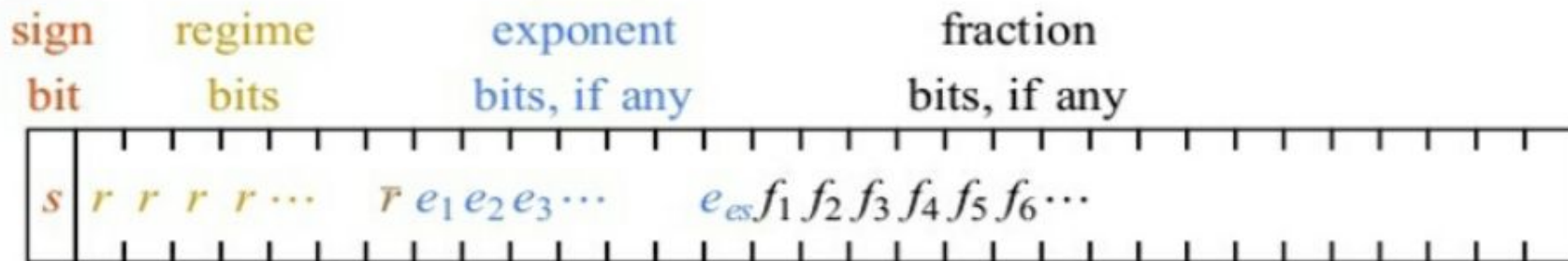
# What is a Posit?

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Posit is a unum that behaves much like a floating-point number of fixed size, rounding to the nearest expressible value if the result of a calculation is not expressible exactly; however, the posit representation offers more accuracy and a larger dynamic range than floats with the same number of bits, as well as many other advantages.

# The Posits format

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# Regime bits

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$K = -m$  if  $m$  bits are 0

$= m-1$  if  $m$  bits are 1

Binary	0001	001x	01xx	10xx	110x	1110
Run-length, $k$	-3	-2	-1	0	1	2

The regime contributes to a scaling factor of  $\text{used}^k$ , where  $\text{used} = 2^{2^{\text{es}}}$ .

es	0	1	2	3	4
used	2	$2^2 = 4$	$4^2 = 16$	$16^2 = 256$	$256^2 = 65536$

- Posit Dynamic Range :

Width		Posit <i>es</i>	Posit Dynamic Range
16		1	$4 * 10^{-9} - 3 * 10^8$
32		3	$6 * 10^{-73} - 2 * 10^{72}$
64		4	$2 * 10^{-299} - 4 * 10^{298}$

- Decoding a posit:

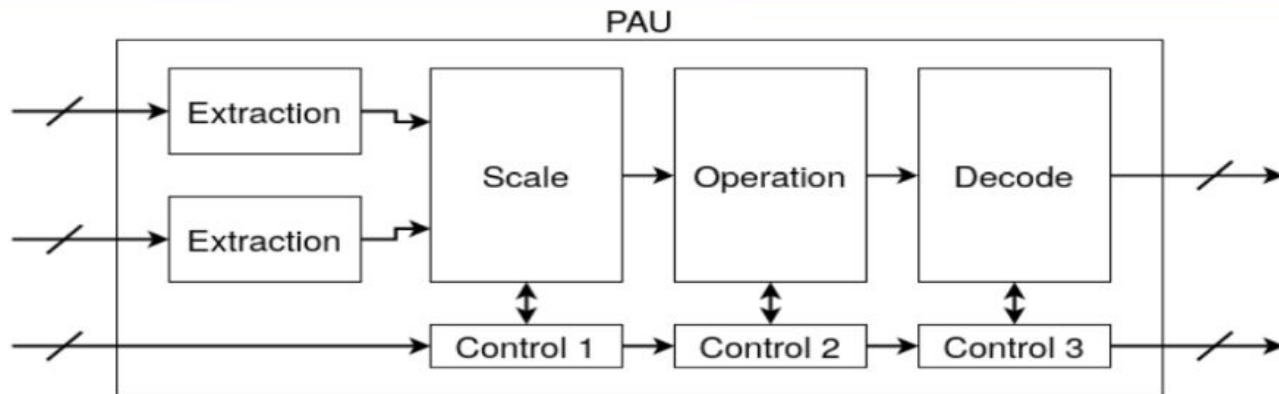
$$X_{10} = \begin{cases} 0, & p=0 \\ \pm\infty, & p=-2^{n-1} \\ (-1)^S \times useed^k \times 2^{exp} \times (1.Fraction), & \text{all other } p. \end{cases}$$



# Posit Core Design

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This effectively splits the design into several stages:



- Extraction.
- Scaling
- Operation
- Decoding along with the necessary control logic.

# Extraction

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## Positive Extraction Algorithm

If negative

Take 2's complement

Check for exception values

Remove sign bit

If regime > 0

Count leading 0

regime = zero\_count - 1

else

negate input

Count leading zeros

regime = zero\_count

temp = input << (zero\_count - 1)

exp = temp[top:top - es + 1]

frac = {1'b1, temp[top - es:end]}

# Scale

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Scale factor = (regime << es) + exp

## Addition/Subtraction Scaling

shit\_value=|sf\_a - sf\_b|

If op\_a > op\_b

    greater\_frac = a\_frac

    smaller\_frac = b\_frac

    greatest\_scaling\_factor = sf\_a

else

    greater\_frac = b\_frac

    smaller\_frac = a\_frac

    greatest\_scaling\_factor=sf\_b

# Operations

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We can do many operation in posit here we have done fout of them:

- 1) Addition
- 2) Subtraction
- 3) Multiplication
- 4) Division

**Decoding block**

**Control block**

# Floating Point to Posit Conversion

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check for exception values

if ( $x < 0$ )

negate  $x$

convert decimal value to binary

remove sign bit

if (MSB=1)

count leading ones

$k = \text{ones} - 1$

else

Count leading zeros

$k = -\text{ones}$

regime =  $\text{used}^k$

if ( regime is entire bit pattern )

return max value

remove regime from binary string

$\text{exp} = 2^{\text{MSB's of binary string}}$

fraction = remaining bits

compute exact fraction

$\text{pos} = \text{sign} * \text{regime} * \text{exp} * \text{frac}$

# Addition

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```
static struct unpacked_t add(struct unpacked_t a, struct unpacked_t b, bool neg)
{
    struct unpacked_t r;

    POSIT_LUTYPE afrac = HIDDEN_BIT(a.frac);
    POSIT_LUTYPE bfrac = HIDDEN_BIT(b.frac);
    POSIT_LUTYPE frac;

    if (a.exp > b.exp) {
        r.exp = a.exp;
        bfrac = RSHIFT(bfrac, a.exp - b.exp);
    } else {
        r.exp = b.exp;
        afrac = RSHIFT(afrac, b.exp - a.exp);
    }

    frac = afrac + bfrac;
    if (RSHIFT(frac, POSIT_WIDTH) != 0) {
        r.exp++;
        frac = RSHIFT(frac, 1);
    }

    r.neg = neg;
    r.frac = LSHIFT(frac, 1);

    return r;
}
```

# Subtraction

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```
static struct unpacked_t sub(struct unpacked_t a, struct unpacked_t b, bool neg)
{
    struct unpacked_t r;

    POSIT_UTYPE afrac = HIDDEN_BIT(a.frac);
    POSIT_UTYPE bfrac = HIDDEN_BIT(b.frac);
    POSIT_UTYPE frac;

    if (a.exp > b.exp || (a.exp == b.exp && a.frac > b.frac)) {
        r.exp = a.exp;
        bfrac = RSHIFT(bfrac, a.exp - b.exp);
        frac = afrac - bfrac;
    } else {
        neg = !neg;
        r.exp = b.exp;
        afrac = RSHIFT(afrac, b.exp - a.exp);
        frac = bfrac - afrac;
    }

    r.neg = neg;
    r.exp -= CLZ(frac);
    r.frac = LSHIFT(frac, CLZ(frac) + 1);

    return r;
}
```

# Result of Addition&Subtraction

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	Float	Posit
Exact	18.5%	25.0%
Inexact	70.1%	75.0%
NaN	10.6%	0.00153%
Overflow	0.757%	0.0%



# Multiplication

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```
struct unpacked_t op2_mul(struct unpacked_t a, struct unpacked_t b)
{
    struct unpacked_t r;

    POSIT_LUTYPE afrac = HIDDEN_BIT(a.frac);
    POSIT_LUTYPE bfrac = HIDDEN_BIT(b.frac);
    POSIT_UTYPE frac = RSHIFT(afrac * bfrac, POSIT_WIDTH);
    POSIT_STYPE exp = a.exp + b.exp + 1;

    if ((frac & POSIT_MSB) == 0) {
        exp--;
        frac = LSHIFT(frac, 1);
    }

    r.neg = a.neg ^ b.neg;
    r.exp = exp;
    r.frac = LSHIFT(frac, 1);

    return r;
}
```

# Division

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```
struct unpacked_t op2_div(struct unpacked_t a, struct unpacked_t b)
{
    struct unpacked_t r;

    POSIT_LUTYPE afrac = HIDDEN_BIT(a.frac);
    POSIT_LUTYPE bfrac = HIDDEN_BIT(b.frac);
    POSIT_STYPE exp = a.exp - b.exp;

    if (afrac < bfrac) {
        exp--;
        bfrac = RSHIFT(bfrac, 1);
    }

    r.neg = a.neg ^ b.neg;
    r.exp = exp;
    r.frac = LSHIFT(afrac, POSIT_WIDTH) / bfrac;

    return r;
}
```

# Result of Multiplication&Division

— — —

	Float	Posit
Exact	22.3%	18.0%
Inexact	51.2%	82.0%
NaN	10.7%	0.00305%
Overflow	12.5%	0.0%
Underfow	3.34%	0.0%

# Posit to Floating point Conversion

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//Here, we can find the Posit to Floating-Point converter module. It includes the following files.

//created a Top-module which takes N (posit word size), E (FP exponent size) and es (posit exponent size) as parameters.

//after we created the Dynamic right shifter sub-module and Dynamic left shifter sub-module.

//and then we created the Leading-One-Detector sub-module and Leading-Zero-Detector sub-module.

//and we already define all theories in the above section how we find regime, exponent, mantissa in posit number so here we computed all these things like that.

//and then decode the output and calculate the final answer.

# Disadvantages of Posit

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- Real-world hardware doesn't support posits (so far).
- Posits don't distinguish between positive infinity, negative infinity and NaN.
- Since there is no NaN, “the calculation is interrupted, and the interrupt handler can be set to report the error”.

# Contribution

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## Karan Jain (B20AI016):

Conversion to floating point, Arithmetic operation on posits

## Kethireddy Harshith Reddy (B20AI018):

Conversion to posit numbers, advantage of posit over floating point numbers

# Reference:-

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1. <https://posithub.org/docs/Posits4.pdf>
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***THANK YOU***