

Integer Division

Division of Unsigned Integers

- Two operands: **Divisor** and **Dividend**

Dividend (D): 274

Divisor (M): 13

$$\begin{array}{r} 13 \overline{) 274} 21 \leftarrow \text{Quotient (Q)} \\ \underline{26} \\ 14 \\ \underline{13} \\ 1 \leftarrow \text{Remainder (R)} \end{array}$$

- $D = Q \times M + R$
- Long-hand-division** (paper-and-pencil method)

Binary Division

- In binary division the only possibilities for quotient bits are 0 and 1
- Divisor is positioned appropriately with respect to the dividend and perform subtraction
- If the remainder is 0 or positive, a quotient bit of 1 is determined
- The remainder is extended by another bit of the dividend
- The divisor is repositioned and another subtraction is performed
- If the remainder is negative, a quotient bit of 0 is determined
- The dividend is restored by adding back the divisor

$$\begin{array}{r}
 1101 \overline{) 100010010(10101} \\
 \underline{1101} \\
 10000 \\
 \underline{11001} \\
 0001110 \\
 \underline{11001} \\
 0001
 \end{array}$$

n -bit Unsigned Division – Restoring Division

- For any n -bit division, we will have $n/2$ -bit divisor and n -bit dividend
- It produces n -bit quotient and n -bit remainder
- Restoring division method uses 3 registers
 - **Register M:** $m_n m_{n-1} \dots m_1 m_0$. $(n+1)$ -bit length
 - It holds $n/2$ -bit positive divisor. The MSB holds 0
 - **Register Q:** $q_{n-1} \dots q_1 q_0$
 - It holds n -bit positive dividend at the start of the operation
 - After the division is complete, it will contain n -bit quotient
 - **Register A:** $a_n a_{n-1} \dots a_1 a_0$. $(n+1)$ -bit length
 - Set to 0 at the beginning of the operation
 - After the division is complete, it will contain n -bit remainder
 - The extra bit in MSB of A and M accommodates the sign bit during subtraction
 - Subtraction is done using 2's complement arithmetic

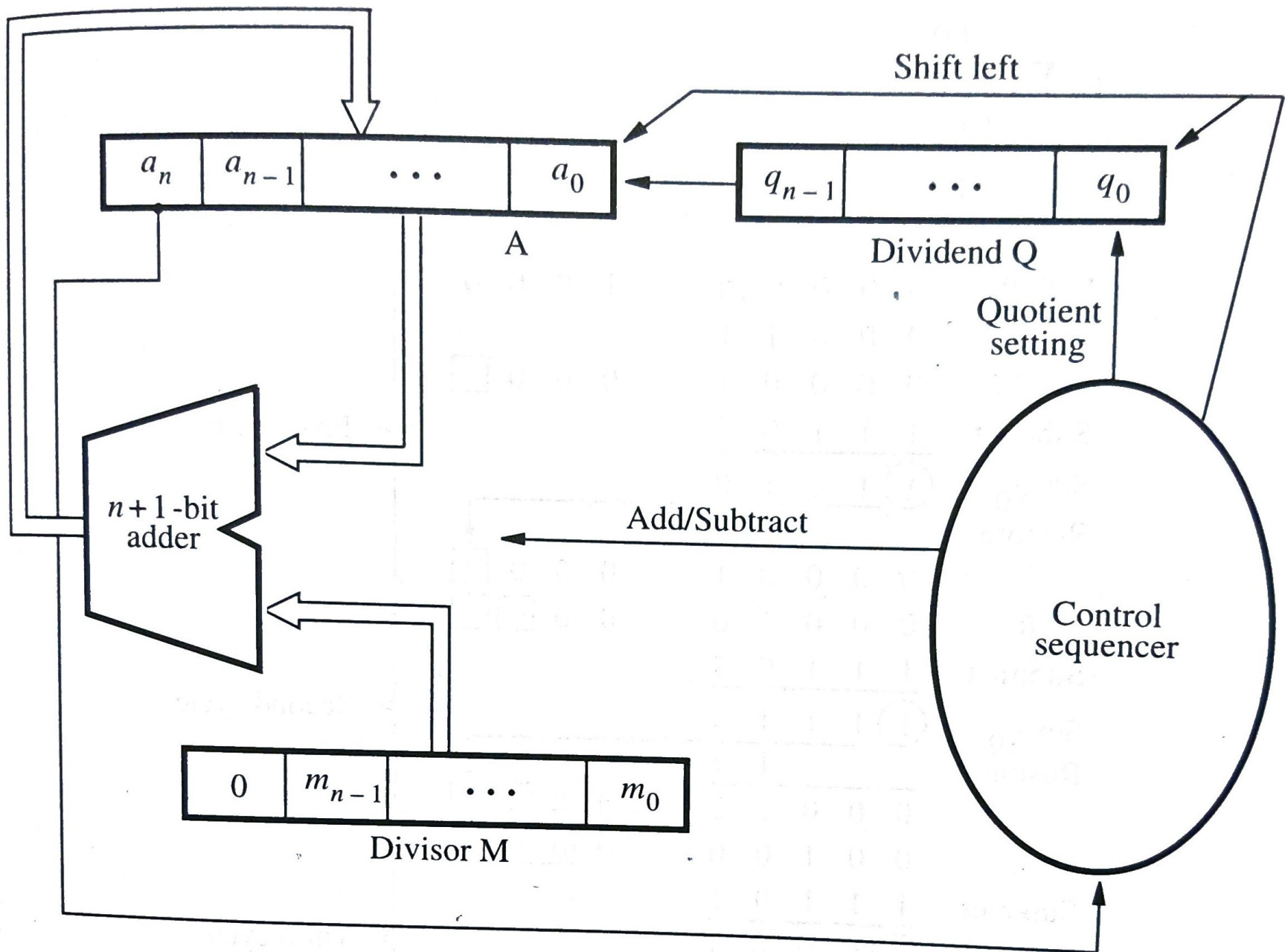


Figure 6.21 Circuit arrangement for binary division.

Restoring Division

- Algorithm:

Do the following steps n times

1. Shift A and Q left one binary position
2. Subtract M from A , and place the answer back in A
3. If the sign of A is 1
 1. Set q_0 to 0
 2. Add M back to A (i.e. restore A)
4. Otherwise, set q_0 to 1

- 4-bit division:

- 8/3
 - Quotient: 2
 - Remainder: 2

$$\begin{array}{r} 11 \overline{) 1000} 10 \\ \underline{11} \\ 10 \end{array}$$

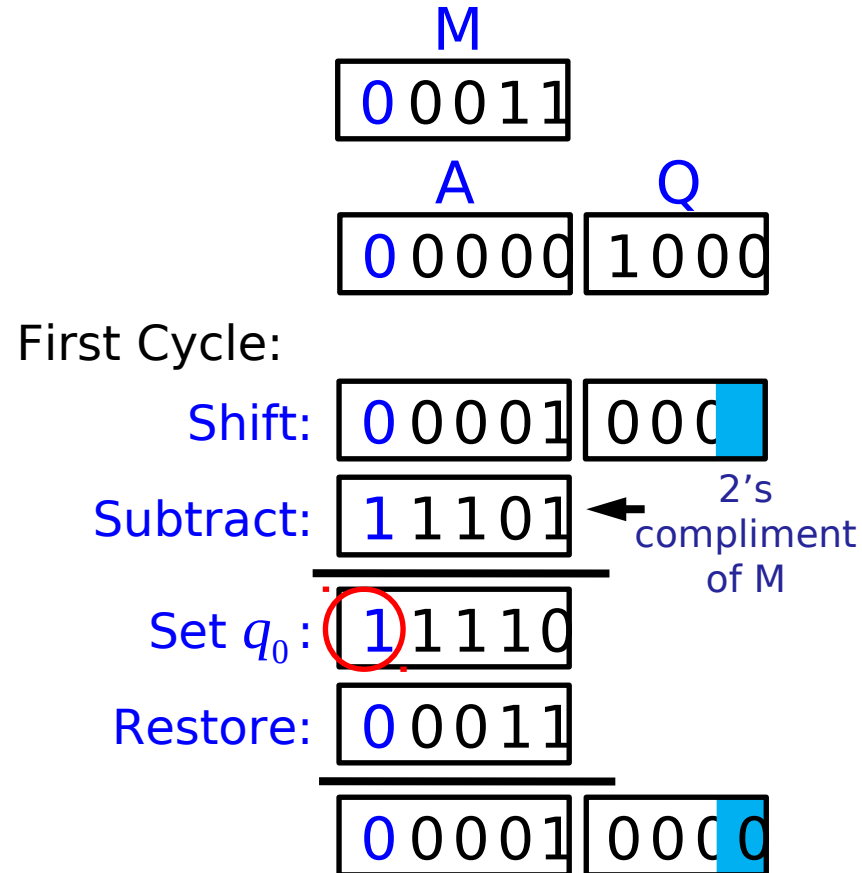
Restoring Division

- Algorithm:

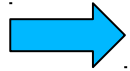
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4. Otherwise, set q_0 to 1

- 4-bit division:



Restoring Division

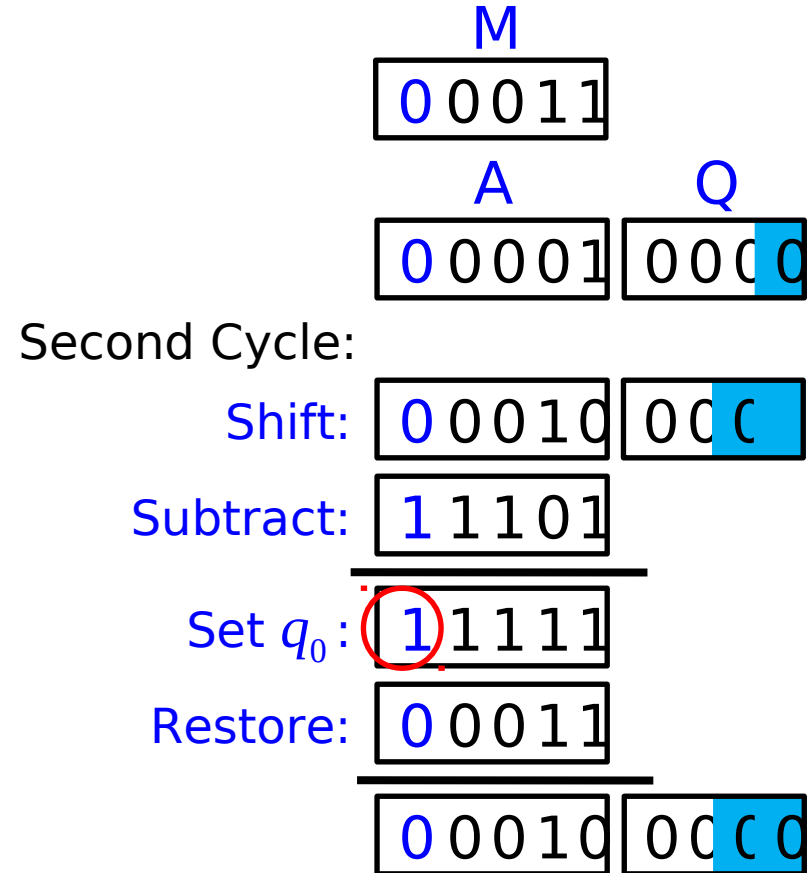


- Algorithm:

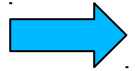
Do the following steps n times

1. Shift A and Q left one binary position
2. Subtract M from A , and place the answer back in A
3. If the sign of A is 1
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4. Otherwise, set q_0 to 1

- 4-bit division:



Restoring Division

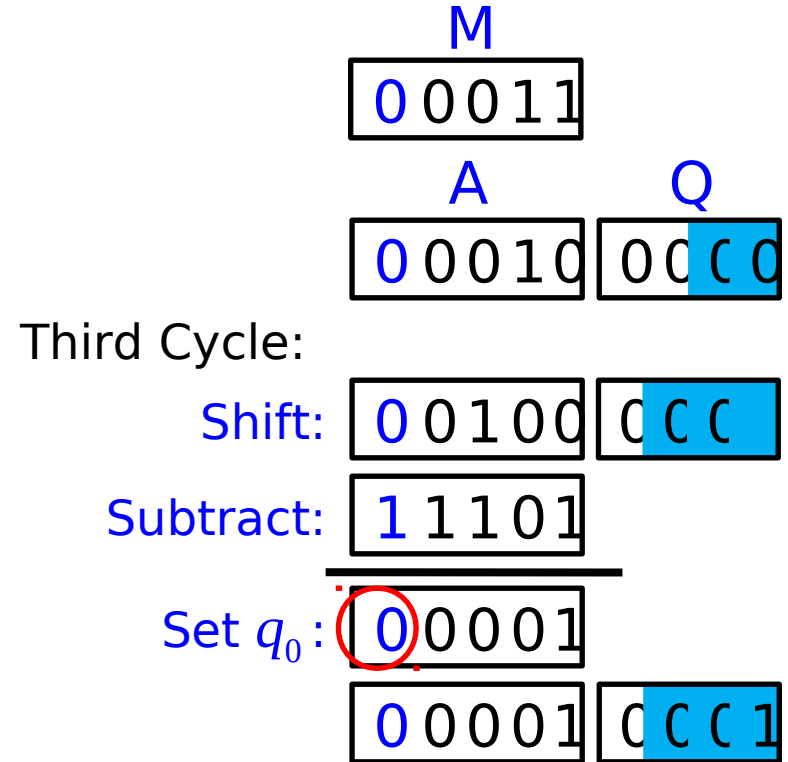


- Algorithm:

Do the following steps n times

1. Shift A and Q left one binary position
2. Subtract M from A , and place the answer back in A
3. If the sign of A is 1
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 2. Add M back to A (i.e. restore A)
4. Otherwise, set q_0 to 1

- 4-bit division:



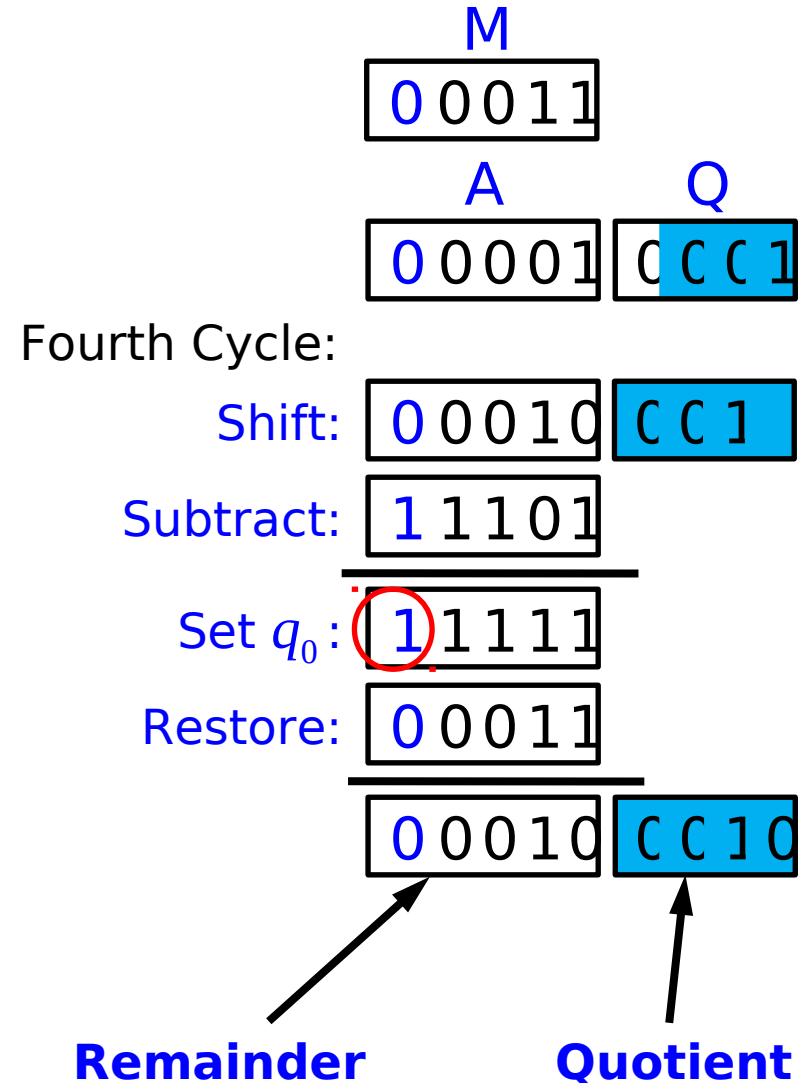
Restoring Division

- Algorithm:

Do the following steps n times

1. Shift A and Q left one binary position
2. Subtract M from A , and place the answer back in A
3. If the sign of A is 1
 1. Set q_0 to 0
 2. Add M back to A (i.e. restore A)
4. Otherwise, set q_0 to 1

- 4-bit division:



Restoring Division

- The restoring division can be improved by **avoiding the need for restoring A** after unsuccessful subtraction
- Subtraction is said to be unsuccessful if the result is negative

Non-Restoring Division

- Algorithm:

Step-1:

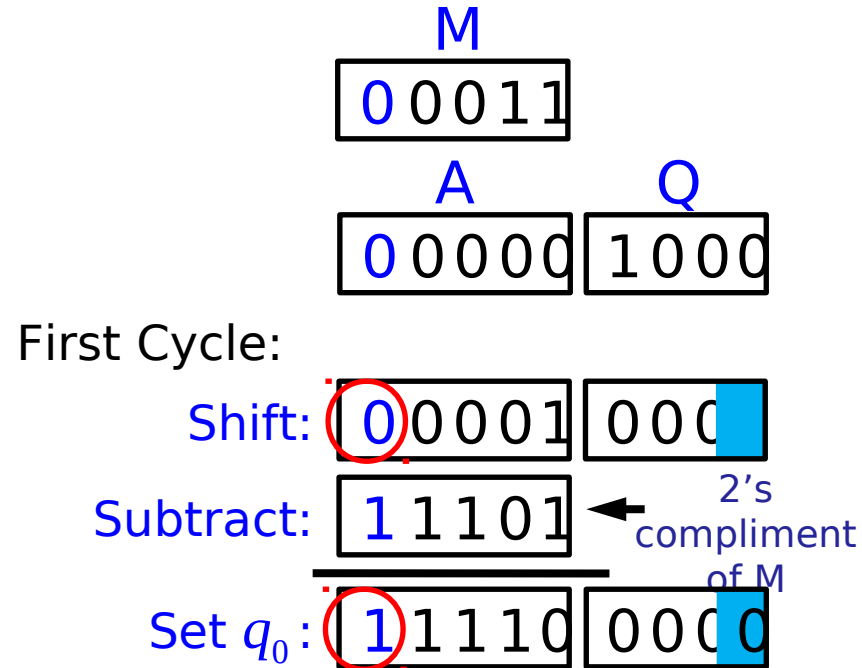
Do the following steps n times

1. Shift A and Q left one binary position
2. If the sign of A is 0
 then, subtract M from A
 else, add M to A
3. If the sign of resulting A is 0
 then set q_0 to 1
 Otherwise, set q_0 to 0

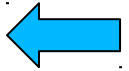
Step-2:

If the sign of A is 1, add M to A

- 4-bit division:



Non-Restoring Division



- Algorithm:

Step-1:

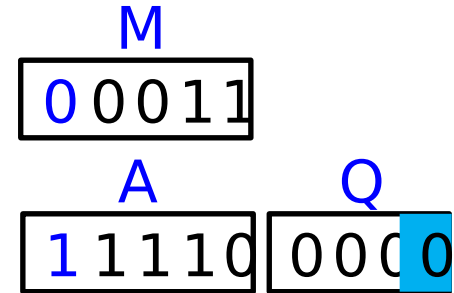
Do the following steps n times

1. Shift A and Q left one binary position
2. If the sign of A is 0
then, subtract M from A
else, add M to A
3. If the sign of resulting A is 0
then set q_0 to 1
Otherwise, set q_0 to 0

Step-2:

If the sign of A is 1, add M to A

- 4-bit division:



Second Cycle:

Shift:

1	1 1 0 0	0 0	C
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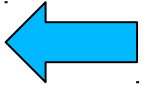
Add:

0	0 0 1 1
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Set q_0 :

1	1 1 1 1	0 0	C 0
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Non-Restoring Division



- Algorithm:

Step-1:

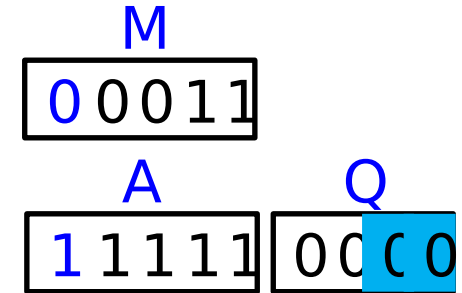
Do the following steps n times

1. Shift A and Q left one binary position
2. If the sign of A is 0
then, subtract M from A
else, add M to A
3. If the sign of resulting A is 0
then set q_0 to 1
Otherwise, set q_0 to 0

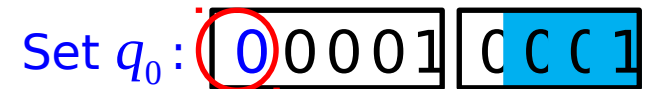
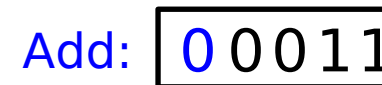
Step-2:

If the sign of A is 1, add M to A

- 4-bit division:



Third Cycle:



Non-Restoring Division

- Algorithm:

Step-1:

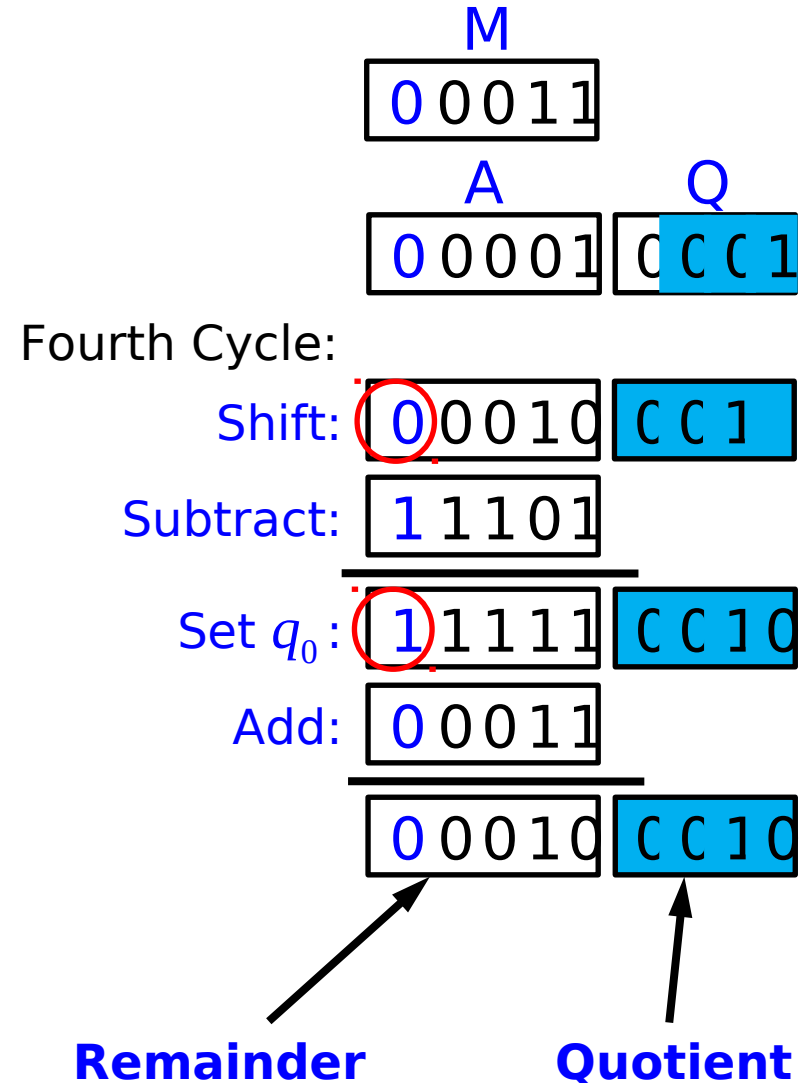
Do the following steps n times

1. Shift A and Q left one binary position
2. If the sign of A is 0
then, subtract M from A
else, add M to A
3. If the sign of resulting A is 0
then set q_0 to 1
Otherwise, set q_0 to 0

Step-2:

If the sign of A is 1, add M to A

- 4-bit division:

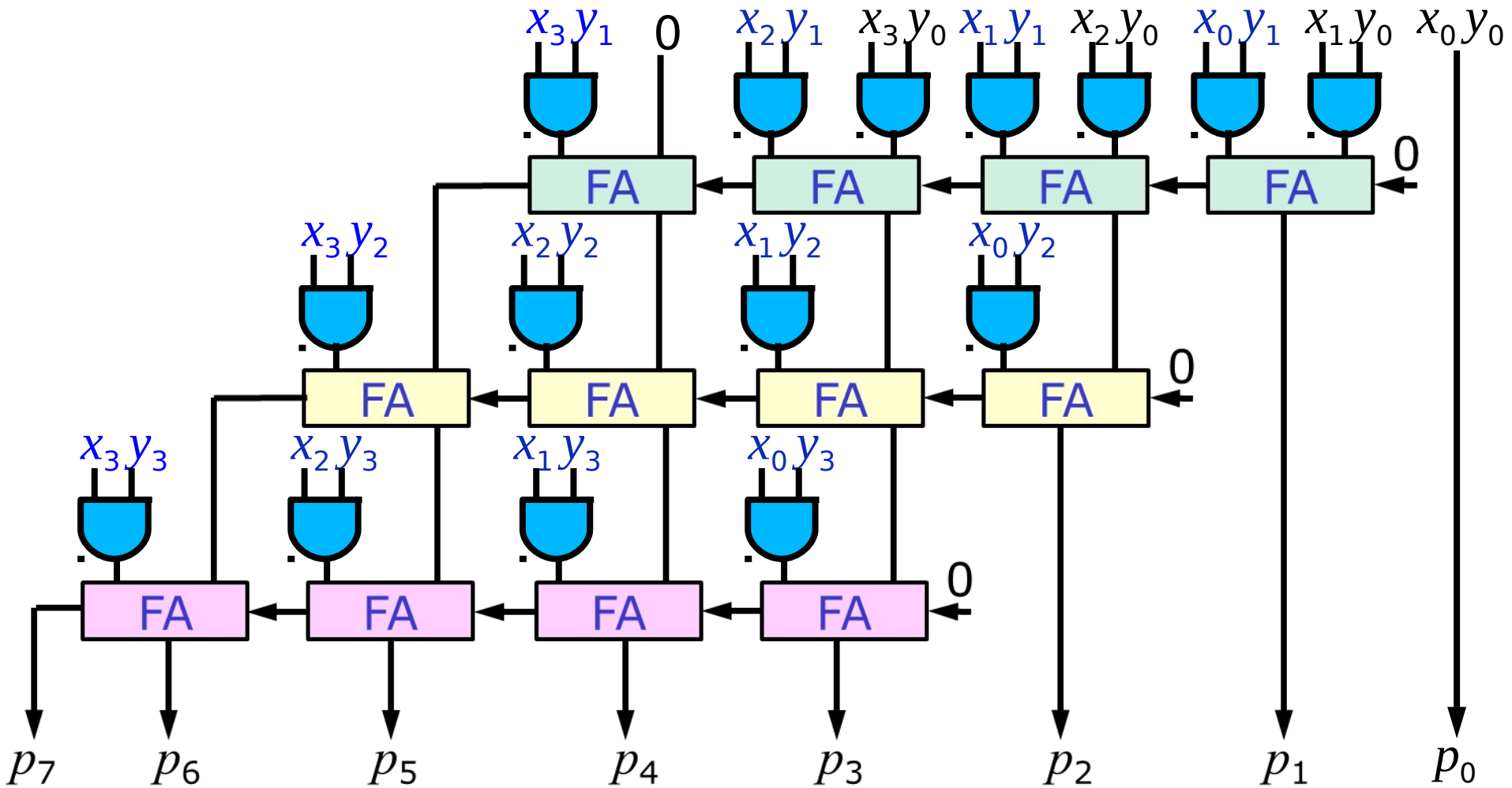


Signed Integer Division

- There is no simple algorithms for directly performing division on signed operands
- In signed division, the negative operands can be pre-processed to transform them into positive values
- After using restoring or non-restoring division method, the results are transformed to the correct signed values

Combinational Array Division Circuit

Combinational Array Multiplier Circuit



- Uses ripple carry adder

Combinational Array (CA) Division Circuit

- Analogous to the array multiplier
- Array divider can be realized by implementing the behaviour of each division step on a **row of basic cells**
- **Basic cell** in array divider depends on the **specific division algorithm (restored or non-restored)** to be implemented
- **Non-Restoring Algorithm:**
 - A row of cells accepts the intermediate remainder and divisor as input
 - Addition or Subtraction should be implemented i.e. **1-bit full adder/subtractor** to be implemented based on the sign

Non-Restoring Division

- Algorithm:

Step-1:

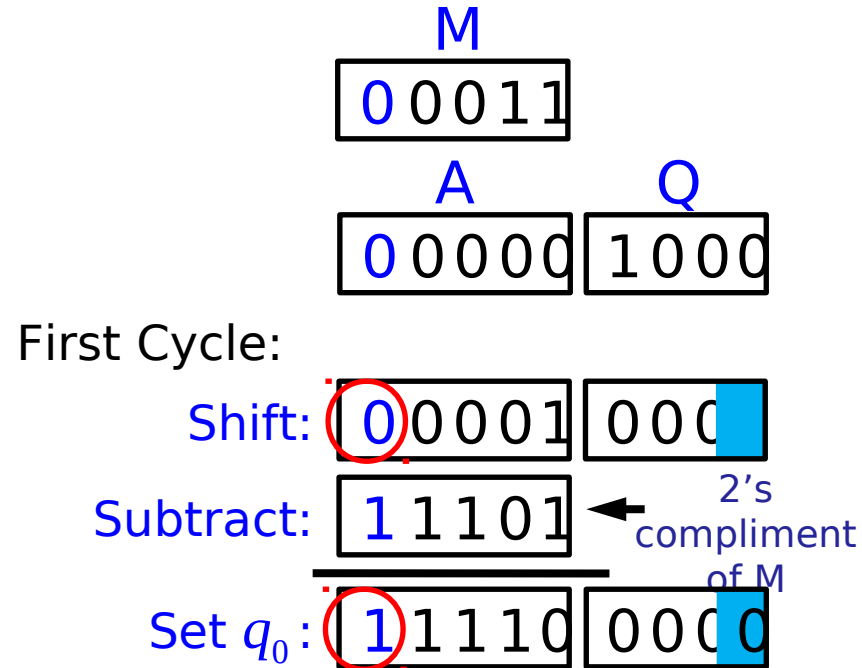
Do the following steps n times

1. Shift A and Q left one binary position
2. If the sign of A is 0
 then, subtract M from A
 else, add M to A
3. If the sign of resulting A is 0
 then set q_0 to 1
 Otherwise, set q_0 to 0

Step-2:

If the sign of A is 1, add M to A

- 4-bit division:



Basic Cell of a CA Division - Non-restoring Algorithm

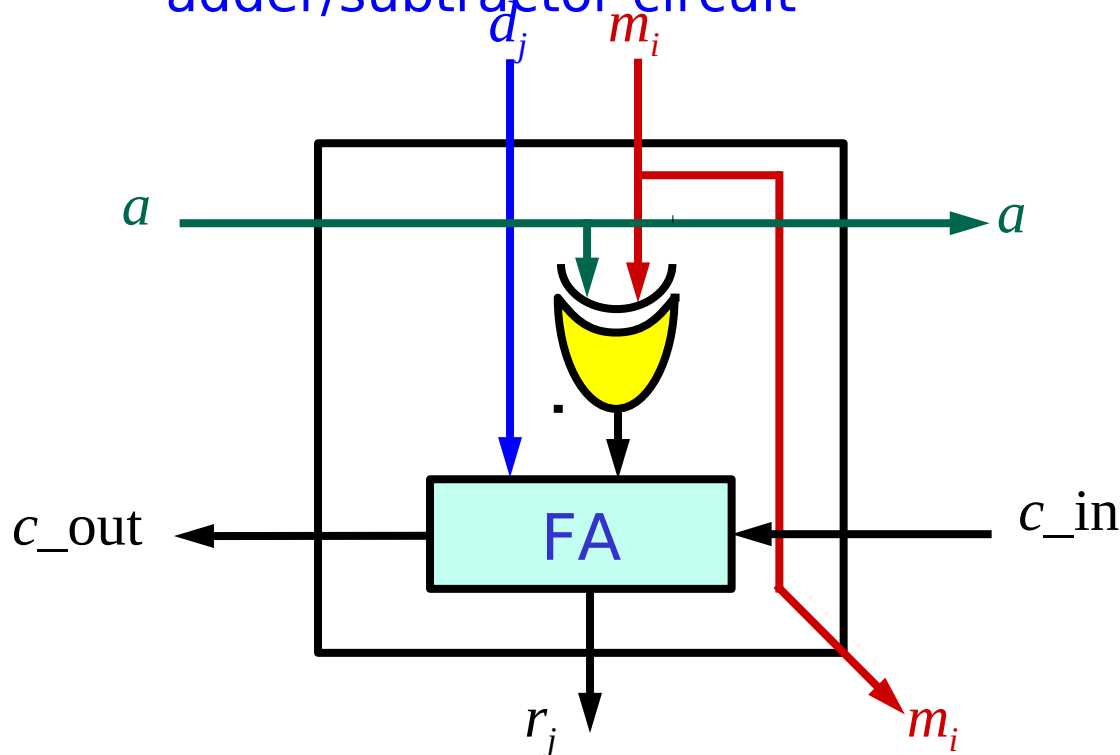
- Dividend is $2n$ -bit length

$$D = d_{2n-1} d_{2n-2} \dots d_1 d_0$$

- Divisor is n -bit length

$$M = m_{n-1} m_{n-2} \dots m_1 m_0$$

- Basic cell is a full adder (FA) that acts as an adder/subtractor circuit



c_{in} : Carry in

c_{out} : Carry out

r_j : Remainder bit

when,

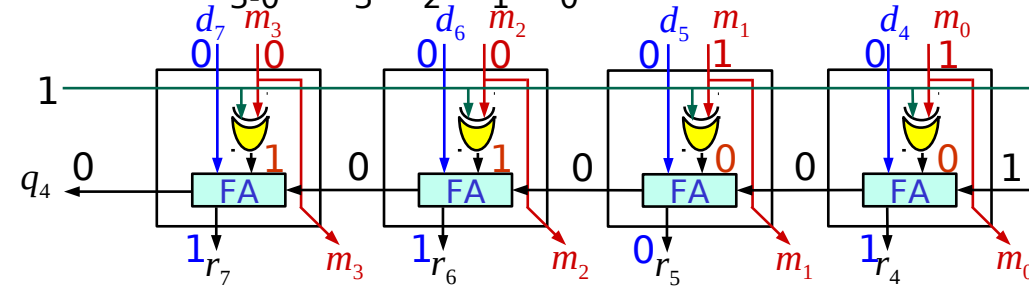
$a=1$, m_i is 1's complemented

$a=0$, m_i is unchanged

4-bit CA Division: Non-restore Algorithm

• $D_{7:0}: d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

• $M_{3:0}: m_3 m_2 m_1 m_0$ Divisor: 0011

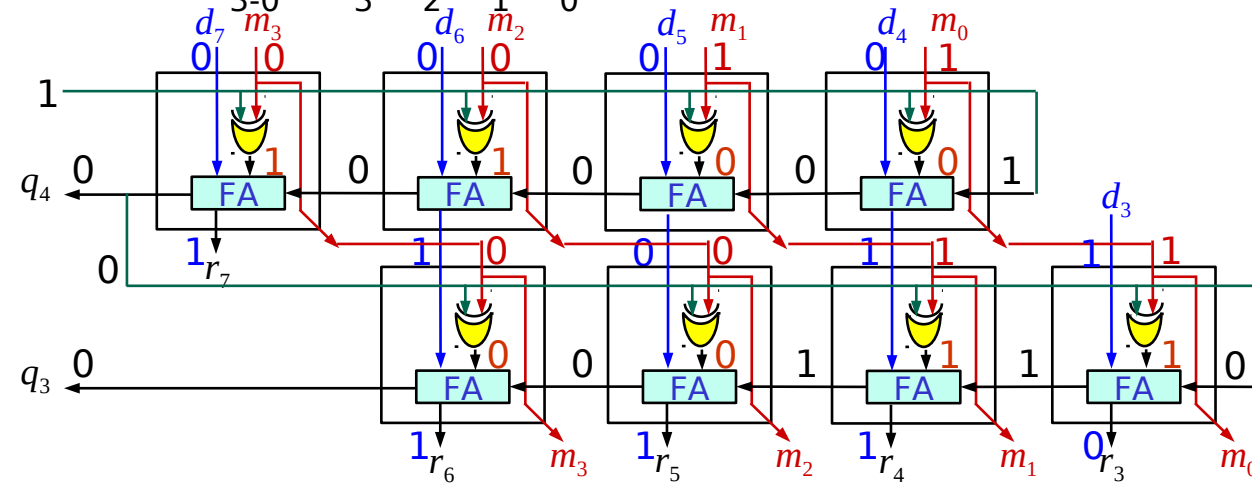


Subtraction

4-bit CA Division: Non-restore Algorithm

• $D_{7:0}$: $d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

• $M_{3:0}$: $m_3 m_2 m_1 m_0$ Divisor: 0011

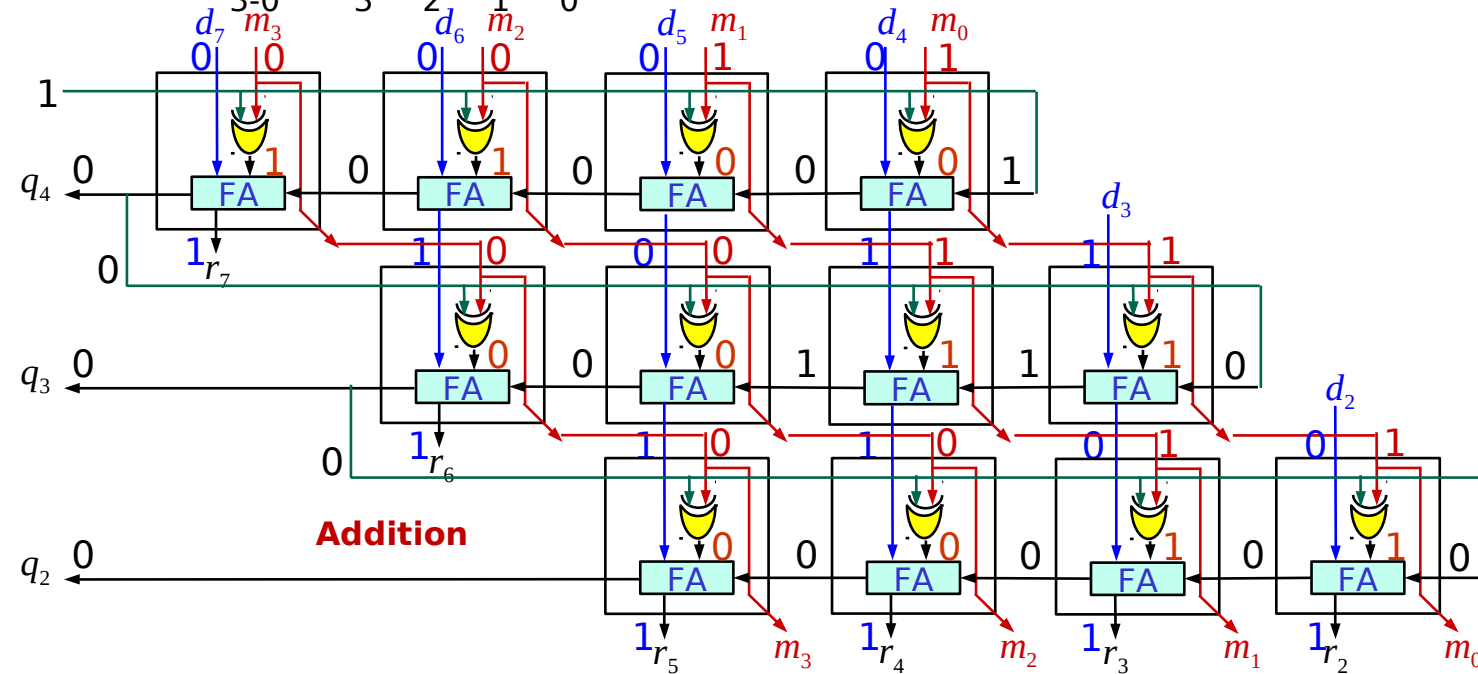


Addition

4-bit CA Division: Non-restore Algorithm

• $D_{7:0}$: $d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

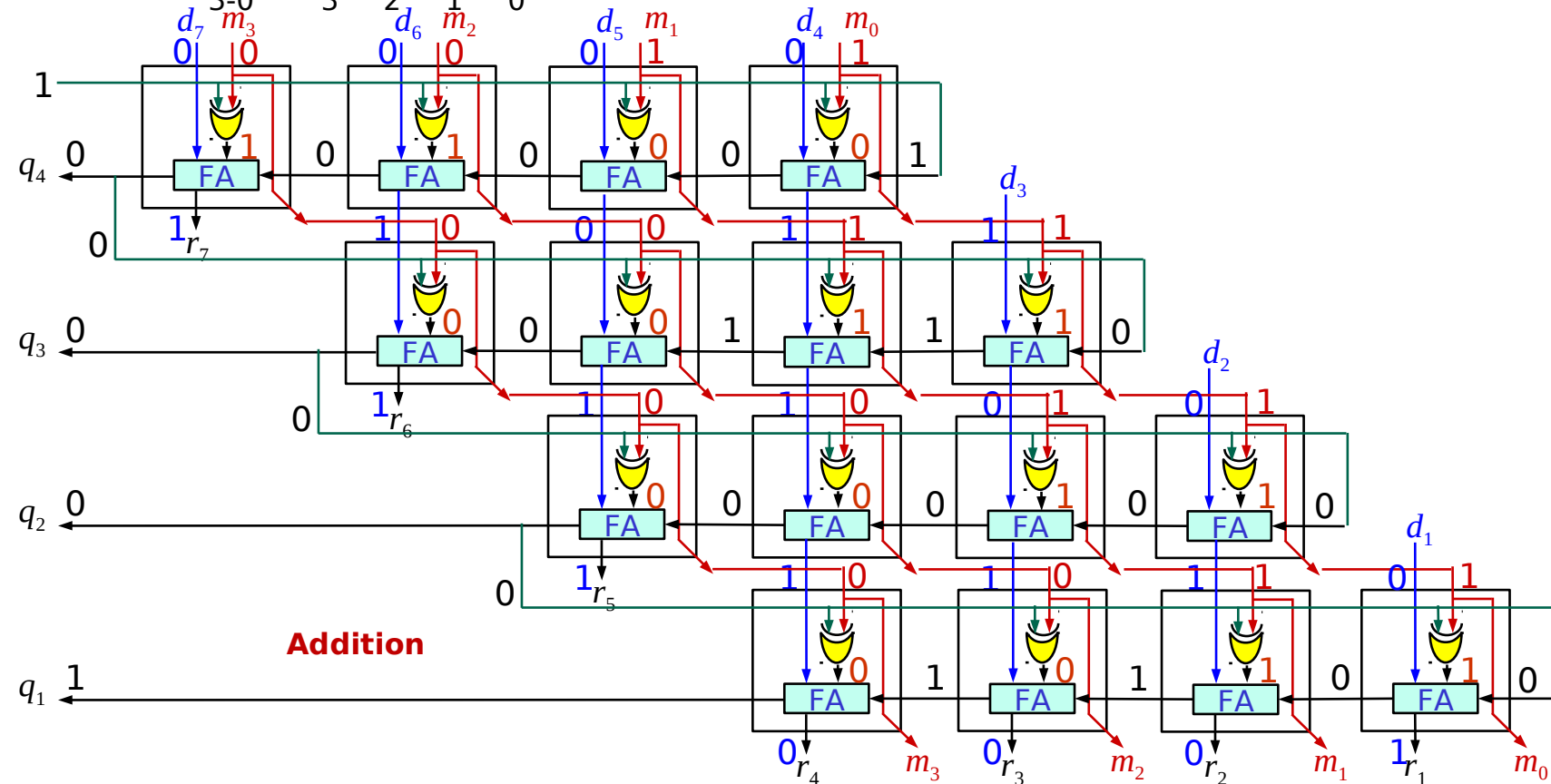
• $M_{3:0}$: $m_3 m_2 m_1 m_0$ Divisor: 0011



4-bit CA Division: Non-restore Algorithm

• $D_{7:0}$: $d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

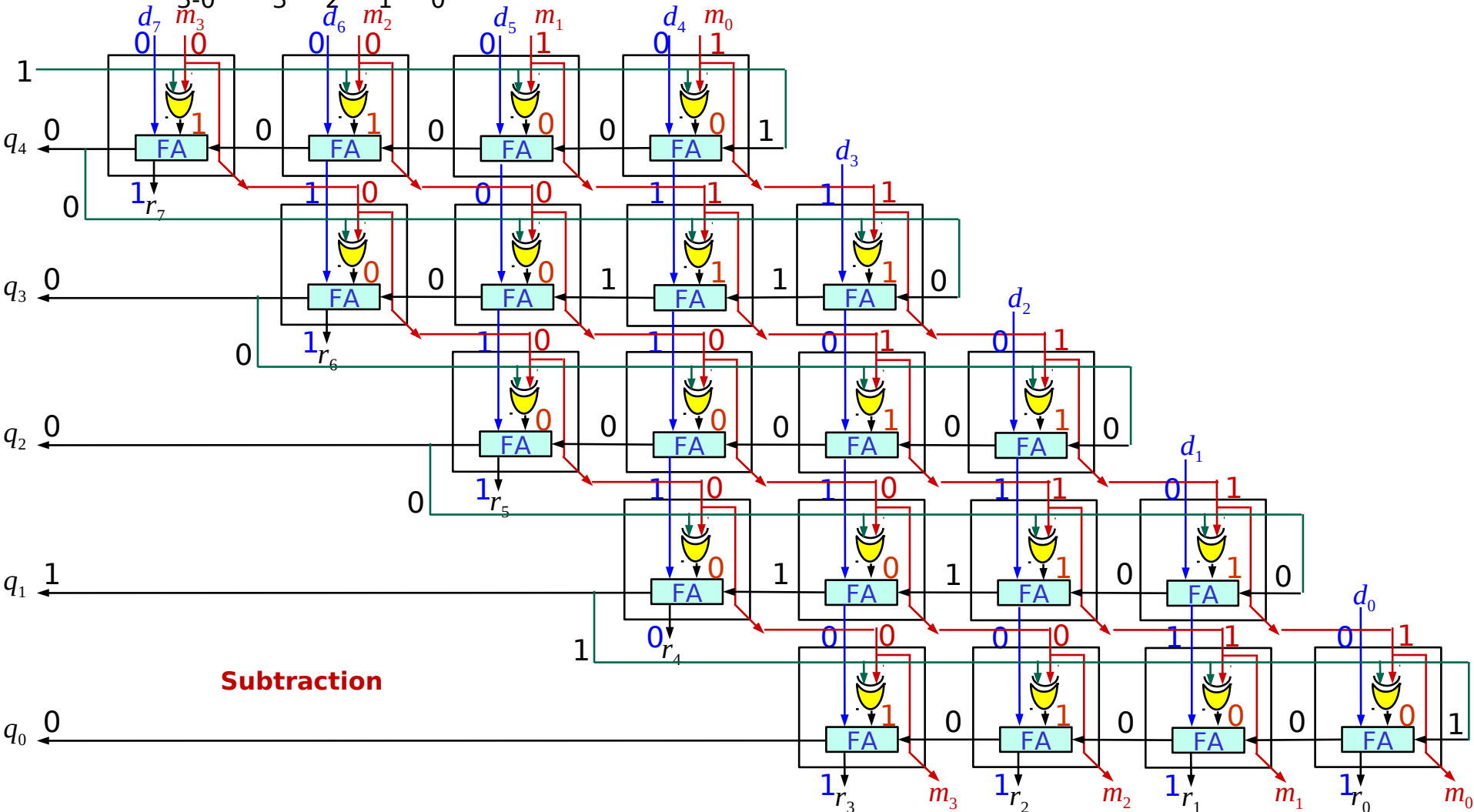
• $M_{3:0}$: $m_3 m_2 m_1 m_0$ Divisor: 0011



4-bit CA Division: Non-restore Algorithm

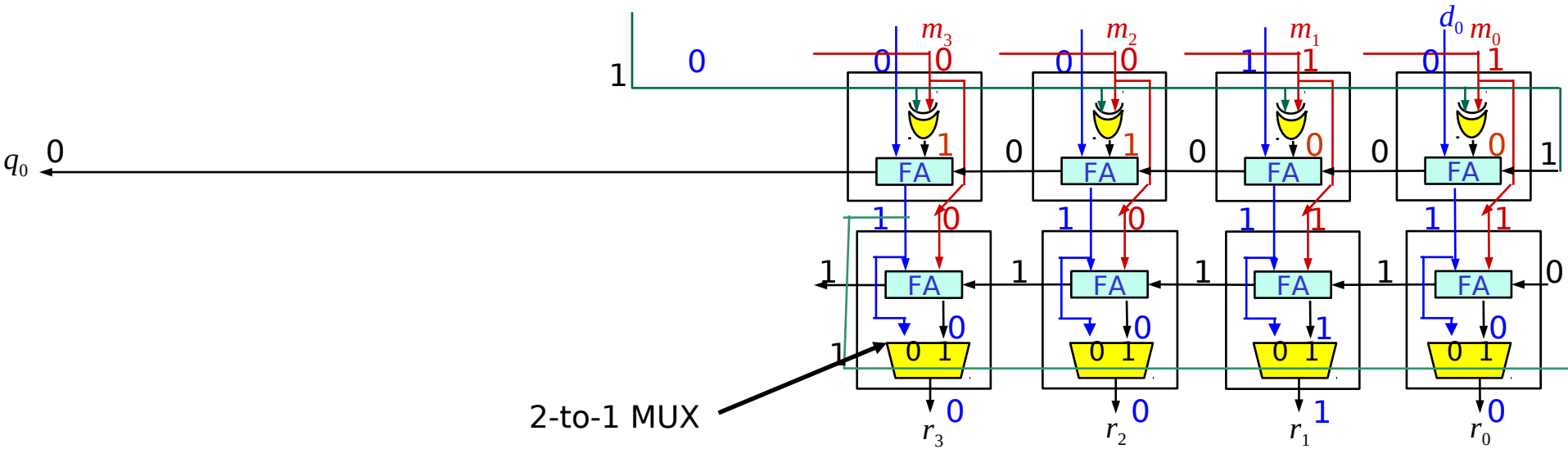
• D_{7-0} : $d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

• M_{3-0} : $m_3 m_2 m_1 m_0$ Divisor: 0011



4-bit CA Division: Non-restore Algorithm

- $D_{7:0}: d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000
- $M_{3:0}: m_3 m_2 m_1 m_0$ Divisor: 0011



CA Division Circuit - Restoring Algorithm

- A row of cells accepts the intermediate remainder and divisor as input
- Subtraction should be implemented i.e. 1-bit full subtractor to be implemented
- Depending upon the outcome of subtraction
 - the row output can be **restored intermediate remainder** input to the row or
 - **the result of the subtraction**

Basic Cell of a CA Division - Restoring Algorithm

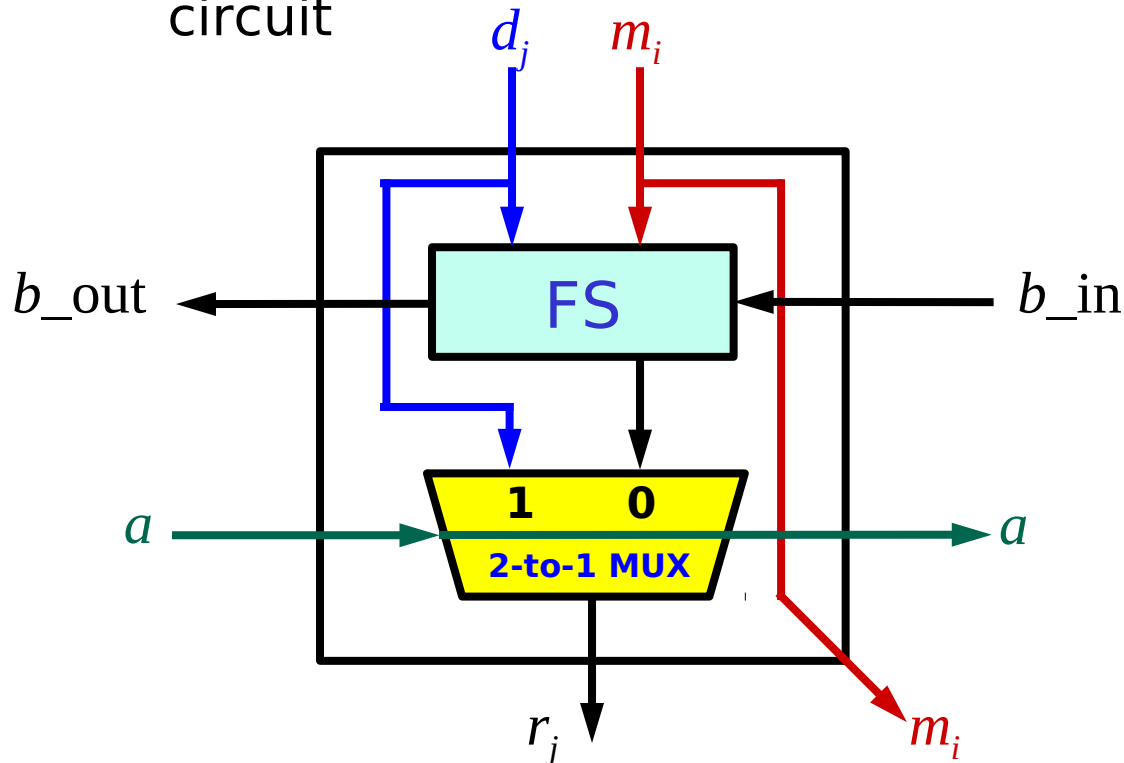
- Dividend is $2n$ -bit length

$$D = d_{2n-1} d_{2n-2} \dots d_1 d_0$$

- Divisor is n -bit length

$$M = m_{n-1} m_{n-2} \dots m_1 m_0$$

- Basic cell is a full subtractor (FS) along with extra logic circuit



b_{in} : Barrow in

b_{out} : Barrow out

r_j : Remainder bit

when,

$$a = 1, r_j = d_j$$

$$a = 0, r_j = d_j - m_i - b_{in}$$

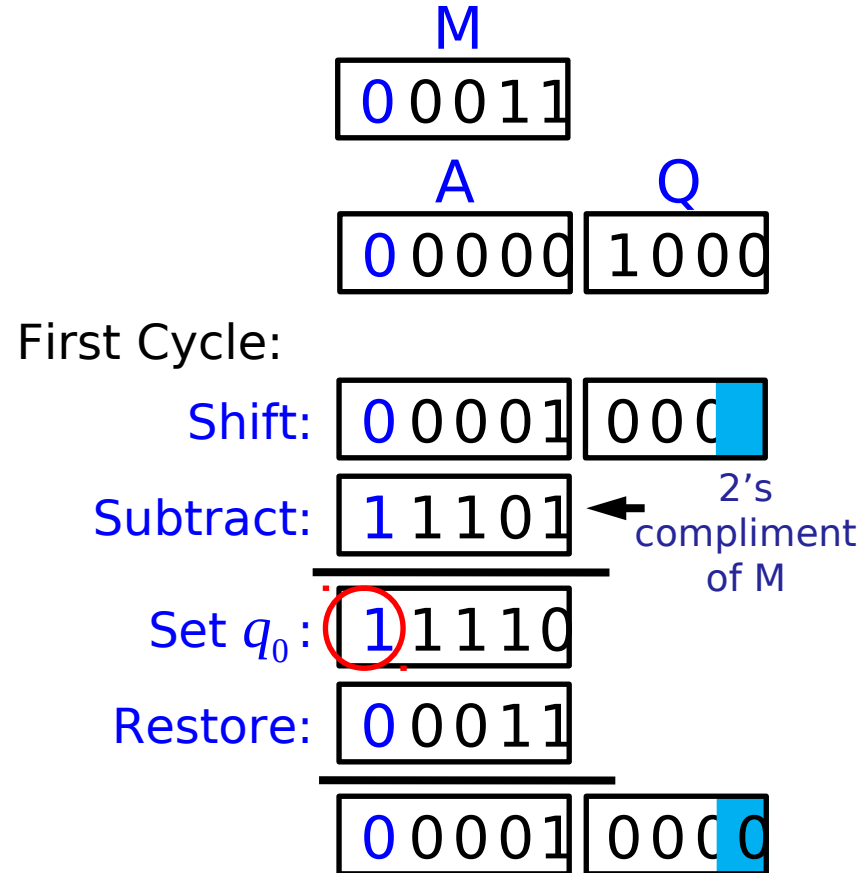
Restoring Division

- Algorithm:

Do the following steps n times

1. Shift A and Q left one binary position
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 1. Set q_0 to 0
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4. Otherwise, set q_0 to 1

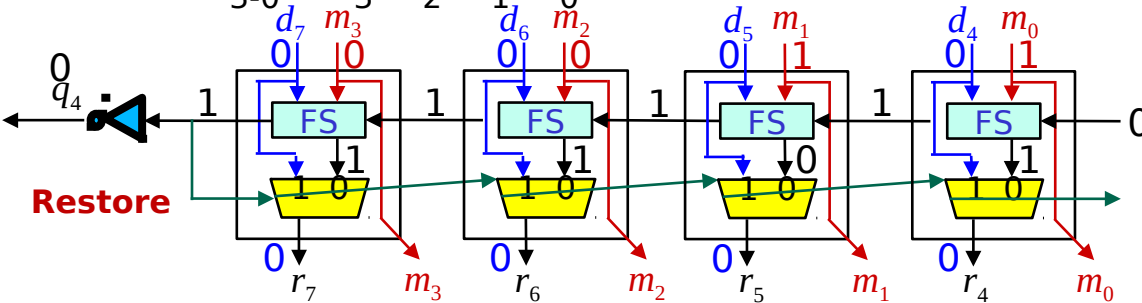
- 4-bit division:



4-bit CA Division: Restore Algorithm

• $D_{7:0}: d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

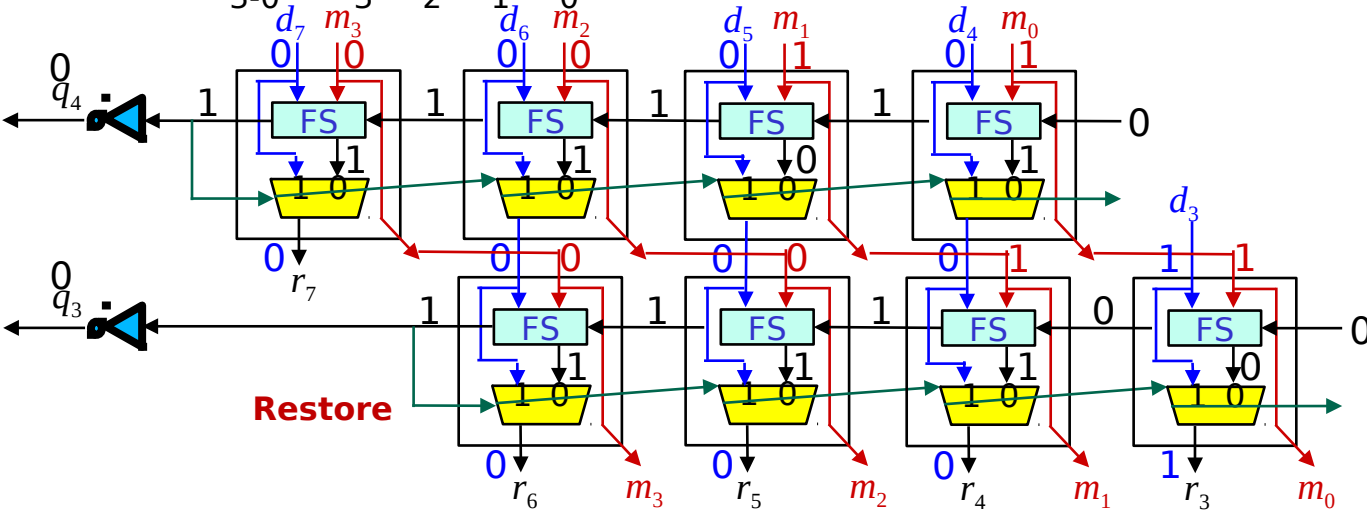
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4-bit CA Division: Restore Algorithm

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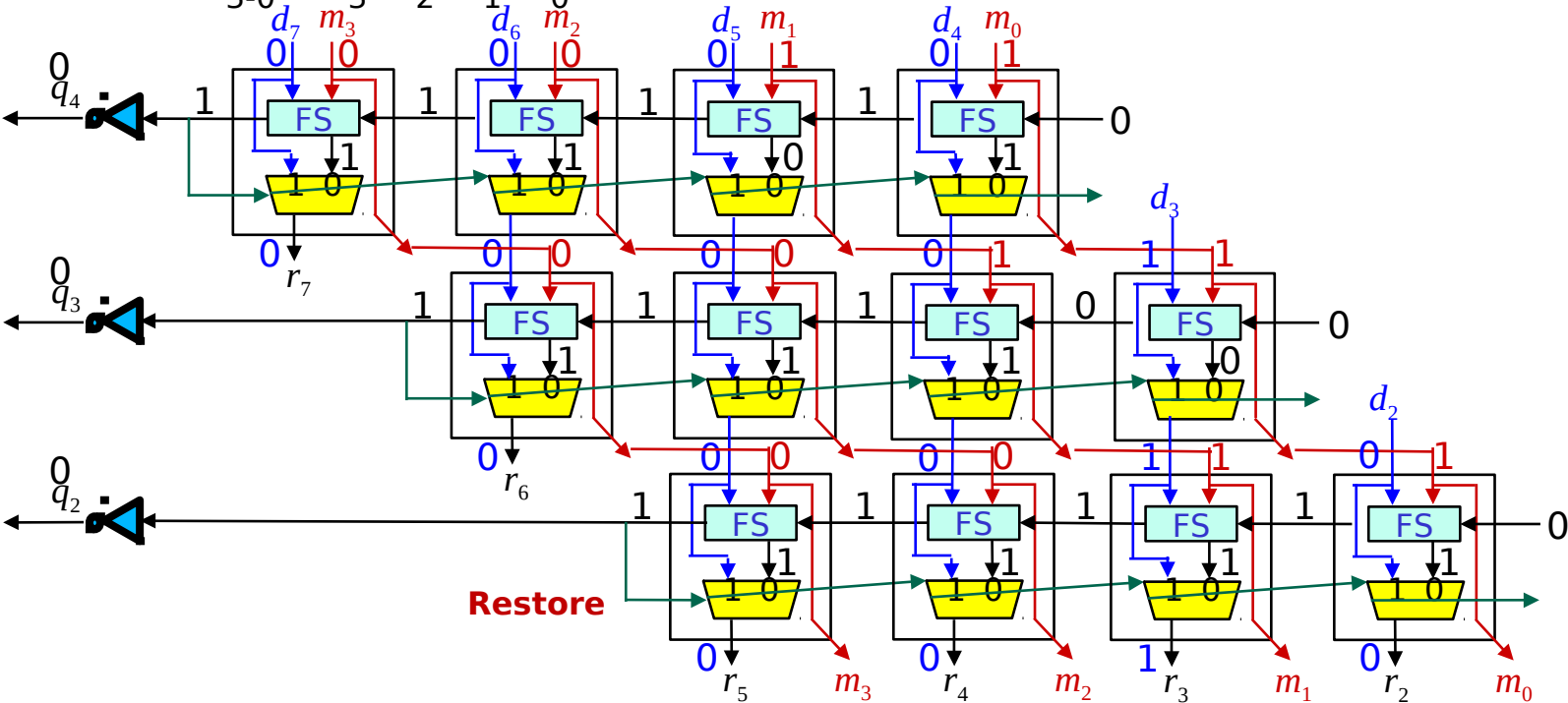
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4-bit CA Division: Restore Algorithm

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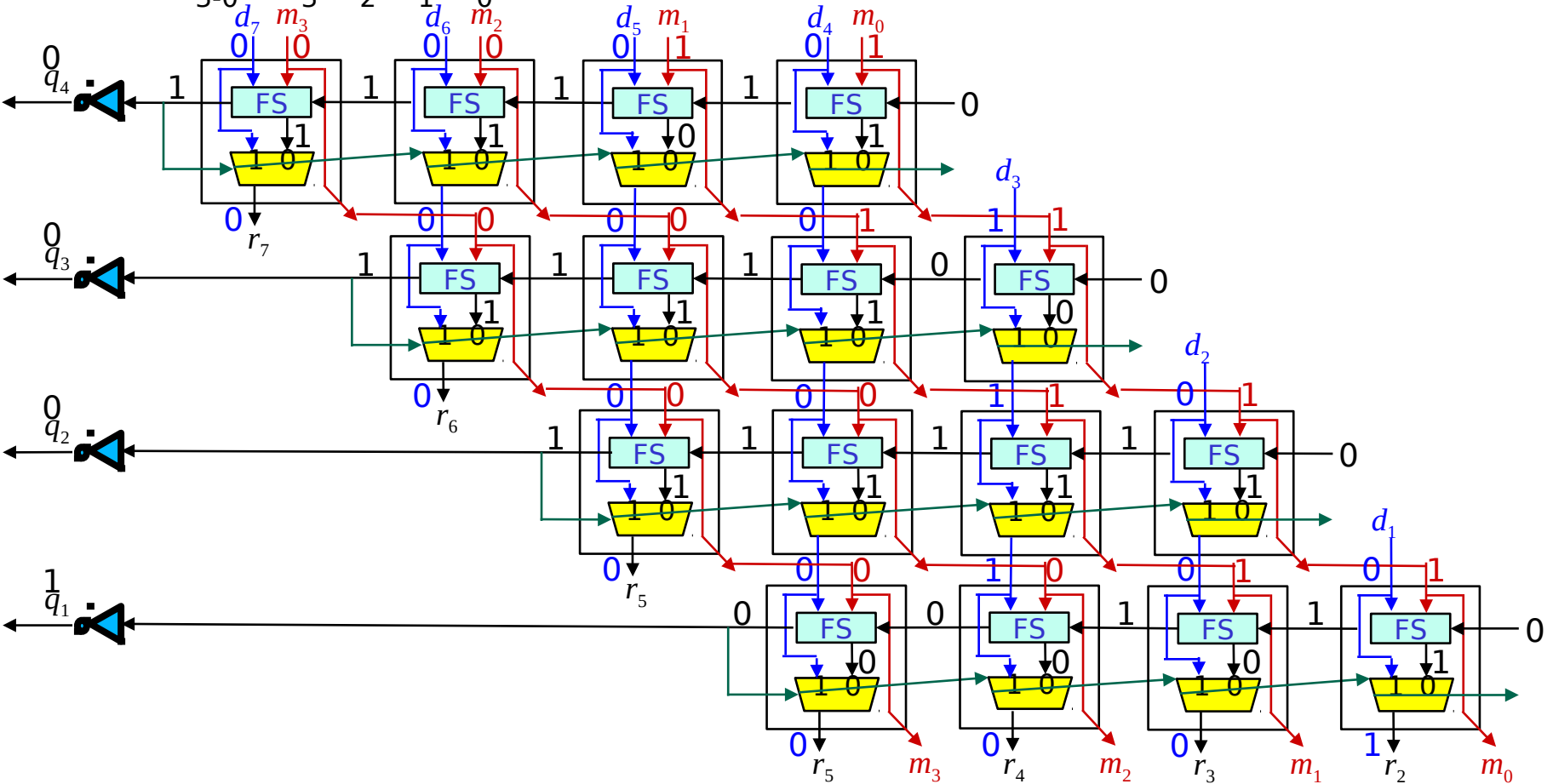
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4-bit CA Division: Restore Algorithm

• $D_{7:0}$: $d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

• $M_{3:0}$: $m_3 m_2 m_1 m_0$ Divisor: 0011



4-bit CA Division: Restore Algorithm

• $D_{7:0}: d_7 d_6 d_5 d_4 d_3 d_2 d_1 d_0$ Dividend: 00001000

• $M_{3:0}: m_3 m_2 m_1 m_0$ Divisor: 0011

