

Dynamic - Programming (DP)

- DP is an algorithm design technique
- "Programming" = use of tables

Toy Example

What is the sum of

$$20 + 21 + \dots + 30 = ?$$

Toy Example

What is the sum of

$$20 + 21 + \dots + 30 + 8 + 2$$

- DP solves Problems by Combining solns to subproblems.
- Storing soln to a subproblem the first time it is solved.
- Looking up the solution when subproblem is encountered again.

DP

- The Programming refers to the use of tables (arrays) to construct a solution.
- In dynamic programming we usually reduce time by increasing the amount of space.
- We solve the problem by solving sub-problems of increasing size and saving each optimal solution in a table (usually).
- The table is then used for finding the optimal solution to larger problems.
- Time is saved since each sub-problem is solved only once.

Fibonacci Numbers

0, 1, 1, 2, 3, 5, 8, - - - -

Fib(n)

If $n \leq 1$

return 1

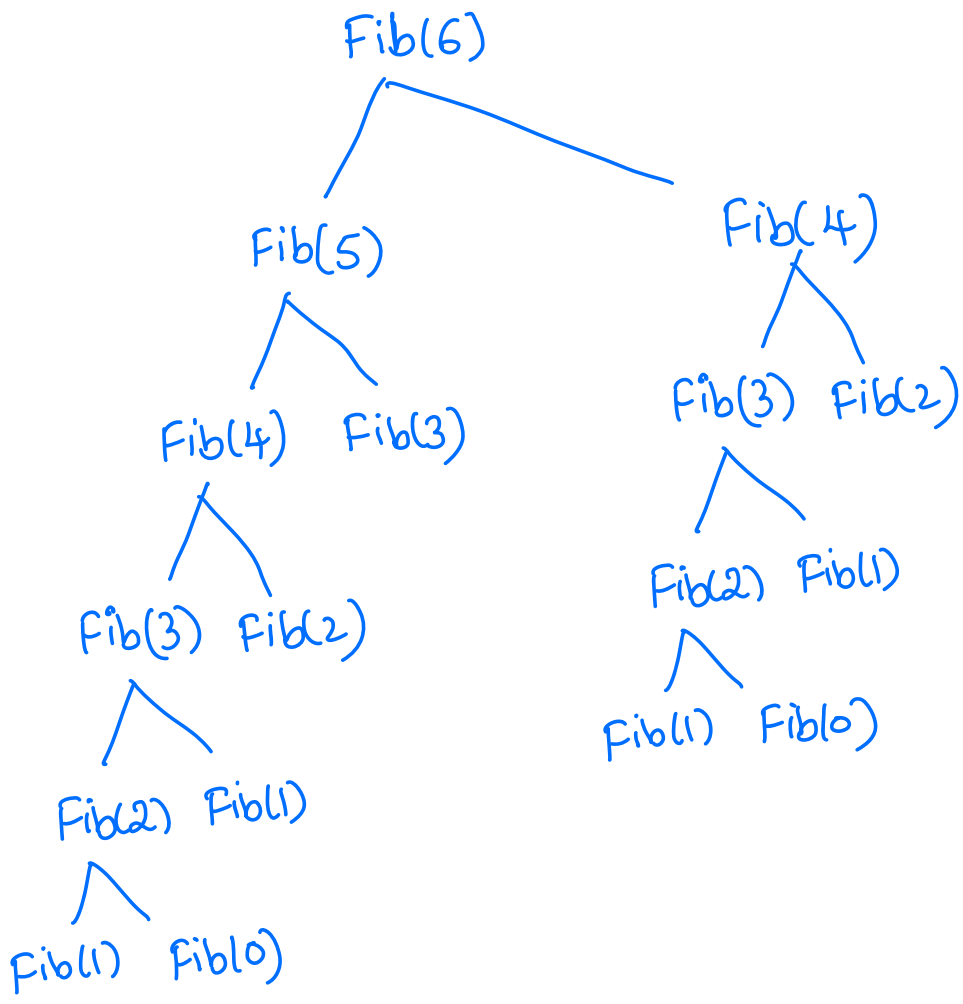
else

return Fib(n-1) + Fib(n-2)

Running time:

$$T(n) = T(n-1) + T(n-2)$$

$$T(n) = O(2^n) \quad [\text{Exponential}]$$



There is a Problem with the
above approach ?

Algorithm - I

def fib(n):

```
1 memo = {}  
2 for k in range(1, n+1):  
3     if k <= 2: f = 1  
4     else: f = memo[k-1] + memo[k-2]  
5     memo[k] = f  
6 return memo[n]
```

Running time : $O(n)$

Algorithm - II

$A \leftarrow$ Array of size $n+1$
with all values set to zero

I/P: n, A

$F(n, A)$

If $n \leq 1$

return n

If $A[n] \neq 0$ then

return $A[n]$

else
 $A[n] = F(n-1, A) + F(n-2, A)$

return $A[n]$.

There are usually two equivalent ways to implement a dynamic-programming approach.

top-down ~~with~~ (memoization)

bottom-up ~~method~~ (Tabulation)

We say that the recursive procedure has been memoized; it “remembers” what results it has computed previously.

Top-down vs Bottom-up