

CS621/CSL611

Quantum Computing For Computer Scientists

Quantum Search

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Quantum Search

Index View of Amplitude Vectors

Adapted from Bernstein's Invited Talk at Indocrypt 2021

- Do **not** worry about **normalization**!

Non-Quantum Computer

- Data stored in 3 bits
- List of 3 elements $\in \{0, 1\}$
- e.g (0,0,0)
- e.g (0,1,1)
- Data stored in 64 bits
- List of 64 elements $\in \{0, 1\}$
- e.g.: (1; 1; 1; 1; 1; 0; 0; 0; 1; 0; 0; 0; 0; 0; 0; 1; 1; 0; 0; 0; 0; 1; 0; 0; 1; 0; 0; 0; 0; 1; 1; 0; 1; 0; 0; 0; 1; 0; 0; 0; 1; 0; 0; 1; 1; 1; 0; 0; 1; 0; 0; 0; 1; 1; 0; 1; 1; 0; 0; 1; 0; 0; 1):

Quantum Computer

- Data stored in 3 qubits
- List of 8 numbers,
- **Not** all zero
- e.g.: [3; 1; 4; 1; 5; 9; 2; 6].
- e.g.: [-2; 7; -1; 8; 1; -8; -2; 8].
- e.g.: [0; 0; 0; 0; 0; 1; 0; 0].
- Data stored in 64 qubits:
- List of 2^{64} numbers
- Not all zero.

Adapted from Bernstein's Invited Talk at Indocrypt 2021

Interpreting The Quantum State Vector

- If n qubits have state $[a_0; a_1; \dots; a_q; \dots; a_{2^n-1}]$ then measurement produces q with probability

$$\frac{|a_q|^2}{\sum_r |a_r|^2}$$

- Recall measuring n qubits
 - produces n bits and
 - collapses the state.
- Collapse \implies New state is all zeros except 1 at position q .

$$[a_0; a_1; \dots; a_q; \dots; a_{2^n-1}] \xrightarrow{\text{Measure}} [0; 0; \dots; \underbrace{1}_{\text{Position}-q}; \dots; 0]$$

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- e.g.: Say 3 qubits have state

$$[1; 1; 1; 1; 1; 1; 1; 1] \rightarrow \sum_r |a_r|^2 = 8$$

- Measurement produces q with probability $\frac{|a_q|^2}{\sum_r |a_r|^2}$
 - $000 = 0$ with probability $\frac{1}{8}$
 - $001 = 1$ with probability $\frac{1}{8}$
 - $010 = 2$ with probability $\frac{1}{8}$
 - $011 = 3$ with probability $\frac{1}{8}$
 - $100 = 4$ with probability $\frac{1}{8}$
 - $101 = 5$ with probability $\frac{1}{8}$
 - $110 = 6$ with probability $\frac{1}{8}$
 - $111 = 7$ with probability $\frac{1}{8}$

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- e.g.: Say 3 qubits have state

$$[3; 1; 4; 1; 5; 9; 2; 6] \rightarrow \sum_r |a_r|^2 = 173$$

- Measurement produces q with probability $\frac{|a_q|^2}{\sum_r |a_r|^2}$
 - $000 = 0$ with probability $\frac{9}{173}$
 - $001 = 1$ with probability $\frac{1}{173}$
 - $010 = 2$ with probability $\frac{4}{173}$
 - $011 = 3$ with probability $\frac{1}{173}$
 - $100 = 4$ with probability $\frac{25}{173}$
 - $101 = 5$ with probability $\frac{81}{173}$
 - $110 = 6$ with probability $\frac{4}{173}$
 - $111 = 7$ with probability $\frac{36}{173}$

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- e.g.: Say 3 qubits have state

$$[0; 0; 0; 0; 0; 0; 1; 0; 0] \rightarrow \sum_r |a_r|^2 = 1$$

- Measurement produces q with probability $\frac{|a_q|^2}{\sum_r |a_r|^2}$
 - $000 = 0$ with probability 0
 - $001 = 1$ with probability 0
 - $010 = 2$ with probability 0
 - $011 = 3$ with probability 0
 - $100 = 4$ with probability 0
 - $101 = 5$ with probability 1 \rightarrow guaranteed outcome
 - $110 = 6$ with probability 0
 - $111 = 7$ with probability 0

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- NOT₀ gate on 3 qubits:

$$\begin{array}{c}
 \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \\
 \downarrow \text{Flipping qubit 0} \\
 \left[\underbrace{3}_{001}; \underbrace{1}_{000}; \underbrace{4}_{011}; \underbrace{1}_{010}; \underbrace{5}_{101}; \underbrace{9}_{100}; \underbrace{2}_{111}; \underbrace{6}_{110} \right] \\
 \downarrow \text{Rearranging Indices} \\
 \left[\underbrace{1}_{000}; \underbrace{3}_{001}; \underbrace{1}_{010}; \underbrace{4}_{011}; \underbrace{9}_{100}; \underbrace{5}_{101}; \underbrace{6}_{110}; \underbrace{2}_{111} \right]
 \end{array}$$

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- NOT₀ gate on 3 qubits:

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right]$$

$$\left[\underbrace{1}_{000}; \underbrace{3}_{001}; \underbrace{1}_{010}; \underbrace{4}_{011}; \underbrace{9}_{100}; \underbrace{5}_{101}; \underbrace{6}_{110}; \underbrace{2}_{111} \right]$$

- **Note:** Adjacent values in state vector have been swapped

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- NOT₀ gate on 4 qubits:

$$[3; 1; 4; 1; 5; 9; 2; 6; 5; 3; 5; 8; 9; 7; 9; 3] \rightarrow [1; 3; 1; 4; 9; 5; 6; 2; 3; 5; 8; 5; 7; 9; 3; 9]$$

- NOT₁ gate on 3 qubits:

$$[3; 1; 4; 1; 5; 9; 2; 6] \rightarrow [4; 1; 3; 1; 2; 6; 5; 9]$$

- NOT₂ gate on 3 qubits:

$$[3; 1; 4; 1; 5; 9; 2; 6] \rightarrow [5; 9; 2; 6; 3; 1; 4; 1].$$

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Interpreting NOT_0 w.r.t Measurement

state	measurement
[1, 0, 0, 0, 0, 0, 0, 0]	000 ←
[0, 1, 0, 0, 0, 0, 0, 0]	001 ←
[0, 0, 1, 0, 0, 0, 0, 0]	010 ←
[0, 0, 0, 1, 0, 0, 0, 0]	011 ←
[0, 0, 0, 0, 1, 0, 0, 0]	100 ←
[0, 0, 0, 0, 0, 1, 0, 0]	101 ←
[0, 0, 0, 0, 0, 0, 1, 0]	110 ←
[0, 0, 0, 0, 0, 0, 0, 1]	111 ←

- Operation on quantum state: NOT_0 , swapping pairs.
- Operation after measurement: flipping bit 0 of result.

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Revisiting Controlled-NOT (CNOT) Gates

- e.g. $C_1\text{NOT}_0$:

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right]$$

↓ Flipping qubit 0 based on qubit 1

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{011}; \underbrace{1}_{010}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{111}; \underbrace{6}_{110} \right]$$

↓ Rearranging Indices

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{1}_{010}; \underbrace{4}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{6}_{110}; \underbrace{2}_{111} \right]$$

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Revisiting Controlled-NOT (CNOT) Gates

- e.g. $C_2\text{NOT}_0$:

$$\begin{aligned} & \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \\ & \quad \downarrow \text{Flipping qubit 0 based on qubit 2} \\ & \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{101}; \underbrace{9}_{100}; \underbrace{2}_{111}; \underbrace{6}_{110} \right] \\ & \quad \downarrow \text{Rearranging Indices} \\ & \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{9}_{100}; \underbrace{5}_{101}; \underbrace{6}_{110}; \underbrace{2}_{111} \right] \end{aligned}$$

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- Compute $C_0\text{NOT}_2$:

$$\begin{aligned}
 & \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \\
 & \quad \downarrow \text{Flipping qubit 2 based on qubit 0} \\
 & \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \\
 & \quad \downarrow \text{Rearranging Indices} \\
 & \left[\underbrace{}_{000}; \underbrace{}_{001}; \underbrace{}_{010}; \underbrace{}_{011}; \underbrace{}_{100}; \underbrace{}_{101}; \underbrace{}_{110}; \underbrace{}_{111} \right]
 \end{aligned}$$

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$$(q_2, q_1, q_0) \rightarrow (q_2, q_1, q_0 \oplus q_1 q_2)$$

Revisiting Toffoli gates

- e.g. $C_2C_1\text{NOT}_0$:

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \rightarrow \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{6}_{110}; \underbrace{2}_{111} \right]$$

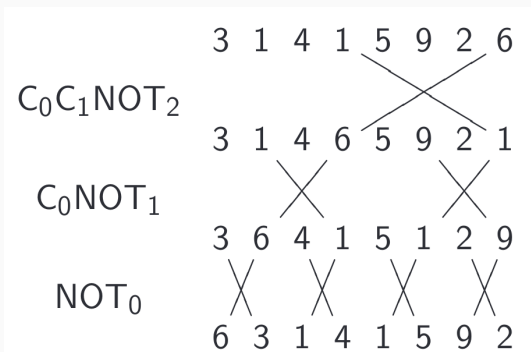
- e.g. $C_0C_1\text{NOT}_2$:

$$\left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{1}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{6}_{111} \right] \rightarrow \left[\underbrace{3}_{000}; \underbrace{1}_{001}; \underbrace{4}_{010}; \underbrace{6}_{011}; \underbrace{5}_{100}; \underbrace{9}_{101}; \underbrace{2}_{110}; \underbrace{1}_{111} \right]$$

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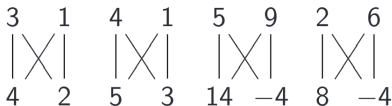
Building Other Permutations

- Combine NOT, CNOT, Toffoli to build other permutations.
- e.g. series of gates to rotate 8 positions by distance 1:



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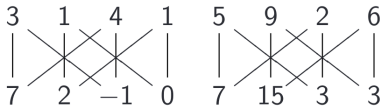
$$[a, b] \mapsto [a + b, a - b].$$



Hadamard₁:

$$[a, b, c, d] \mapsto$$

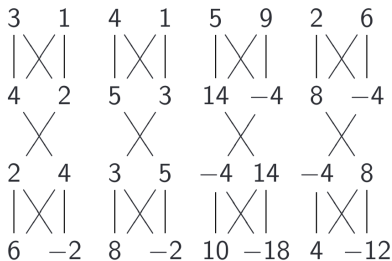
$$[a + c, b + d, a - c, b - d].$$



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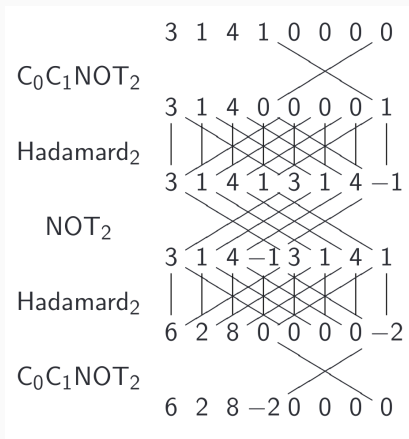
Some uses of Hadamard gates

Hadamard₀, NOT₀, Hadamard₀:



- “Multiplied each amplitude by 2.”
- This is not physically observable.
- What other change has happened?
- “Negated amplitude if q_0 is set.” No effect on measuring now.

- “Negate amplitude if q_0q_1 is set.”
- Assumes $q_2 = 0$: “ancilla” qubit.



- “Negate amplitude around its average.”

$$[3, 1, 4, 1] \rightarrow [1.5, 3.5, 0.5, 3.5]$$

