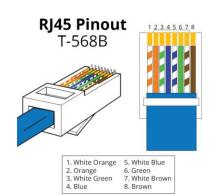
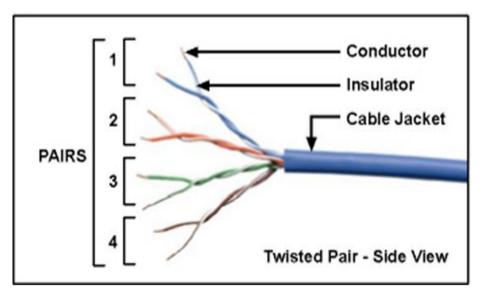
# Crimping LAN/Ethernet cables



Cable/Network Crimping tool







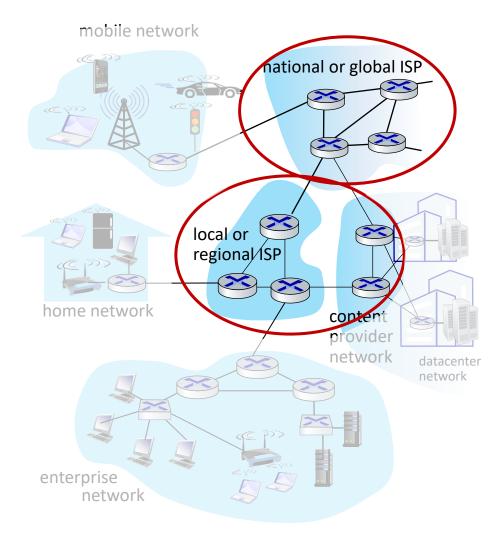
# Chapter 1: roadmap

- What is the Internet?
- What is a protocol?
- Network edge: hosts, access network, physical media, Internet structure
- Network core: packet/circuit switching
- Performance: loss, delay, throughput
- Security
- Protocol layers, service models
- History

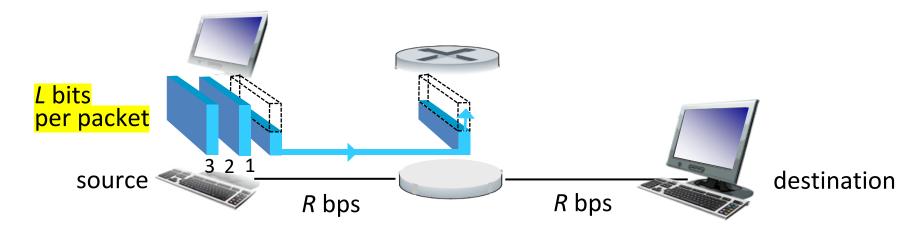


### The network core

- mesh of interconnected routers
- packet-switching: hosts break application-layer messages into packets
  - forward packets from one router to the next, across links on path from source to destination
  - each packet transmitted at full link capacity



# Packet-switching: store-and-forward

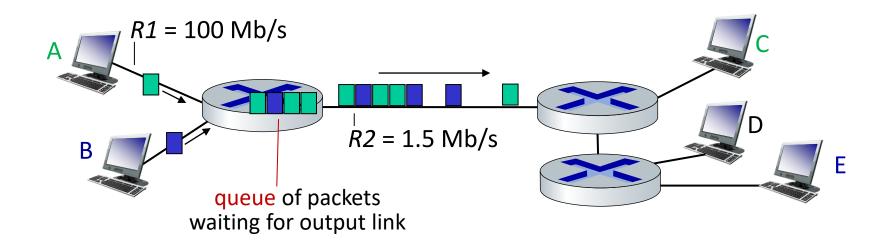


- Transmission delay: takes L/R seconds to transmit (push out) L-bit packet into link at R bps
- Store and forward: entire packet must arrive at router before it can be transmitted on next link
- End-end delay: 2L/R (above), assuming zero propagation delay (more on delay shortly)

### One-hop numerical example:

- *L* = 10 Kbits
- *R* = 100 Mbps
- one-hop transmission delay= 0.1 msec

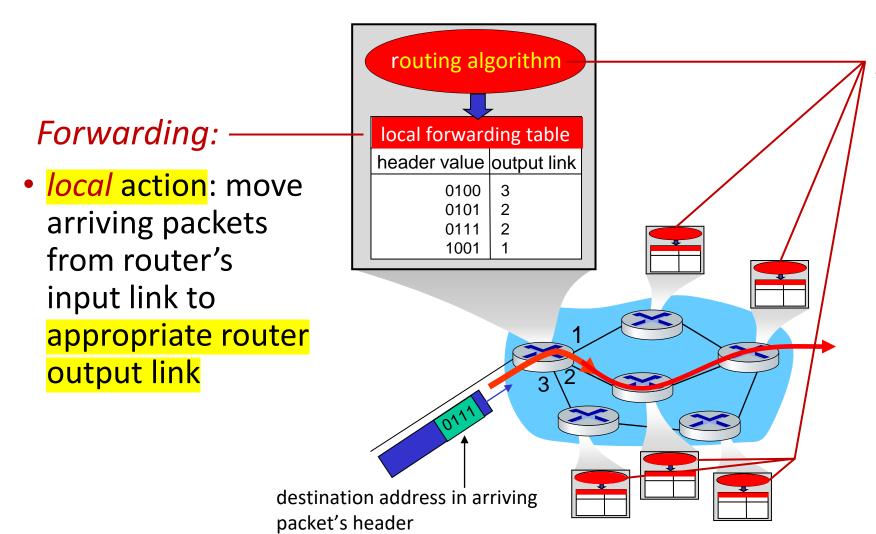
# Packet-switching: queueing delay, loss



Packet queuing and loss: if arrival rate (in bps) to link exceeds transmission rate (bps) of link for a period of time:

- packets will queue, waiting to be transmitted on output link
- packets can be dropped (lost) if memory (buffer) in router fills up

# Two key network-core functions



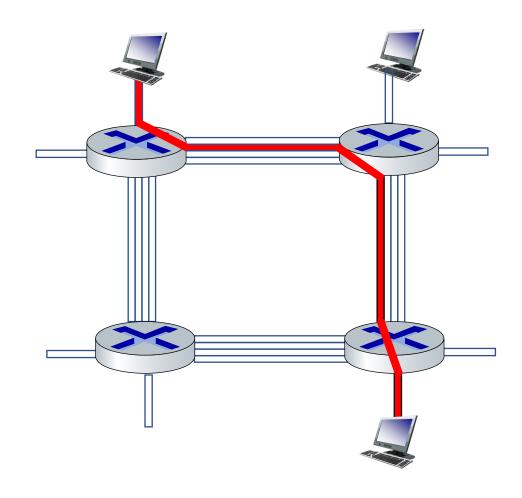
### Routing:

- global action:
   determine source destination paths
   taken by packets
- routing algorithms

### Alternative to packet switching: circuit switching

end-end resources allocated to, reserved for "call" between source and destination

- in diagram, each link has four circuits.
  - call gets 2<sup>nd</sup> circuit in top link and 1<sup>st</sup> circuit in right link.
- dedicated resources: no sharing
  - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (no sharing)
- commonly used in traditional telephone networks



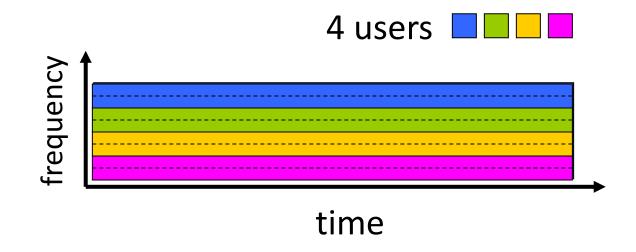
# Circuit switching: FDM and TDM

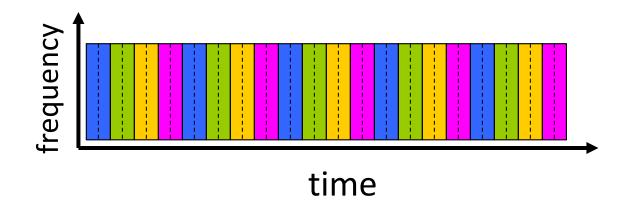
# Frequency Division Multiplexing (FDM)

- optical, electromagnetic frequencies divided into (narrow) frequency bands
- each call allocated its own band, can transmit at max rate of that narrow band

### Time Division Multiplexing (TDM)

- time divided into slots
- each call allocated periodic slot(s), can transmit at maximum rate of (wider) frequency band, but only during its time slot(s)



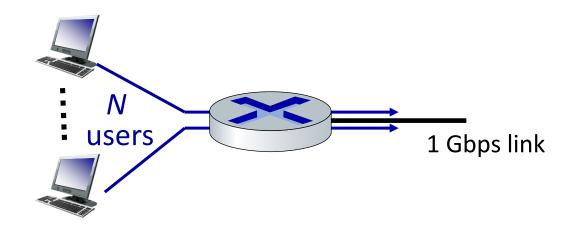


# Packet switching versus circuit switching

packet switching allows more users to use network!

#### Example:

- 1 Gb/s link
- each user:
  - 100 Mb/s when "active"
  - active 10% of time
- circuit-switching: 10 users
- packet switching: with 35 users, probability > 10 active at same time is less than .0004 \*



Q: how did we get value 0.0004?

Q: what happens if > 35 users?

<sup>\*</sup> Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose\_ross/interactive

• N = 35 users

P (# Active Users > 10) = 1 - P(# Active Users = 10) - P(# Active Users = 9)
- P(# Active Users = 8) ....P(# Active Users = 1)

• P(# Active Users =10) =C(35, 10) \* 0.1^10 \* 0.9^25

Note: Binomial Distribution

Probability of k out of n ways:

$$P(k \text{ out of } n) = \frac{n!}{k!(n-k)!} p^{k}(1-p)^{(n-k)}$$

The General Binomial Probability Formula

# Packet switching versus circuit switching

#### Is packet switching a "winner"?

- great for "bursty" data sometimes has data to send, but at other times not
  - resource sharing
  - simpler, no call setup
- excessive congestion possible: packet delay and loss due to buffer overflow
  - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
  - bandwidth guarantees traditionally used for audio/video applications