CS251: Introduction to Language Processing

Machine Independent Optimizations

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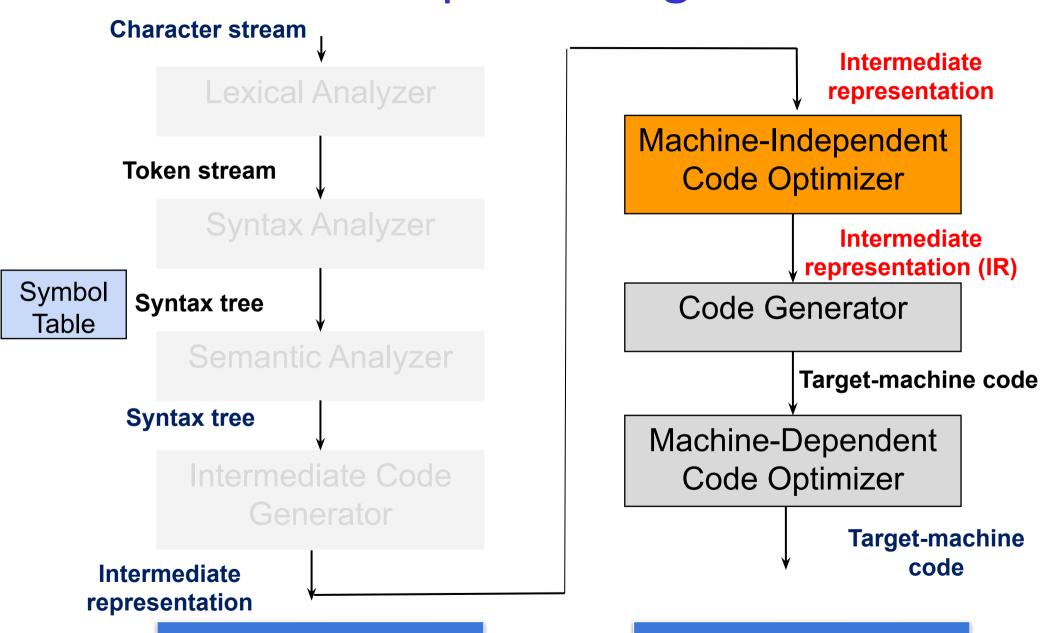
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Acknowledgement

- References for today's slides
 - Stanford University
 https://web.stanford.edu/class/archive/cs/cs143/cs143.1128/
 - Prof. Y. N Srikant, IISc Bangalore
 https://iith.ac.in/~ramakrishna/Compilers-Aug14/slides/
 - http://sei.pku.edu.cn/~yaoguo/ACT11/slides/lect2-opt.
 ppt
 - Course textbook

Compiler Design



Front End

Optimizations

- Global common subexpression elimination
- Copy propagation
- Constant propagation and constant folding
- Loop invariant code motion
- Induction variable elimination and strength reduction

Loop Invariant Code Motion Example

```
t1 = 202
    i = 1
L1: t2 = i > 100
    if t2 goto L2
    t1 = t1-2
    t3 = addr(a)
    t4 = t3 - 4
    t5 = 4*i
    t6 = t4 + t5
    *t6 = t1
    i = i + 1
     goto L1
L2:
```

Before LIV code motion

After LIV code motion

Optimizations

- Global common subexpression elimination
- Copy propagation
- Constant propagation and constant folding
- Loop invariant code motion
- Induction variable elimination and strength reduction

Strength Reduction

Before strength reduction for t5

After strength reduction for t5 and copy propagation

Induction Variable Elimination

```
t1 = 202
    i = 1
    t3 = addr(a)
    t4 = t3 - 4
    t7 = 4
L1: t2 = i > 100
    if t2 goto L2
    t1 = t1-2
    t6 = t4 + t7
    *t6 = t1
    i = i+1
    t7 = t7 + 4
    goto L1
L2:
```

t1 = 202t3 = addr(a)t4 = t3 - 4t7 = 4L1: t2 = t7 > 400if t2 goto L2 t1 = t1-2t6 = t4 + t7*t6 = t1t7 = t7 + 4goto L1 L2:

After eliminating i and replacing it with t7

Before induction variable elimination (i)

Summary

- Machine Independent Optimizations
 - Improve the quality of code: performance, memory, and energy efficiency
 - Still hot area of research
- Formalize:
 - Basic blocks
 - Control flow graph
- Optimizations:
 - Examples