

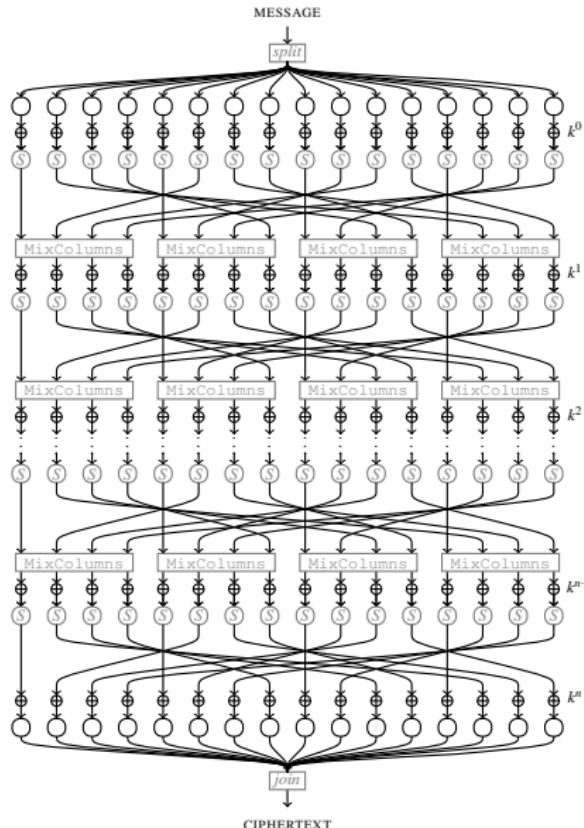
CS 553

CRYPTOGRAPHY

Lecture 11

The Story of AES

Instructor
Dr. Dhiman Saha



- ▶ Winner of AES competition:
Rijndael - SPN Design
- ▶ Designed by **Vincent Rijmen** & **John Daemen**
- ▶ One of the **most widely studied** cryptographic algorithms
- ▶ Used almost everywhere
- ▶ With **microprocessor-level instruction support**

Probably the best intro to AES

Slides from Jeff Moser

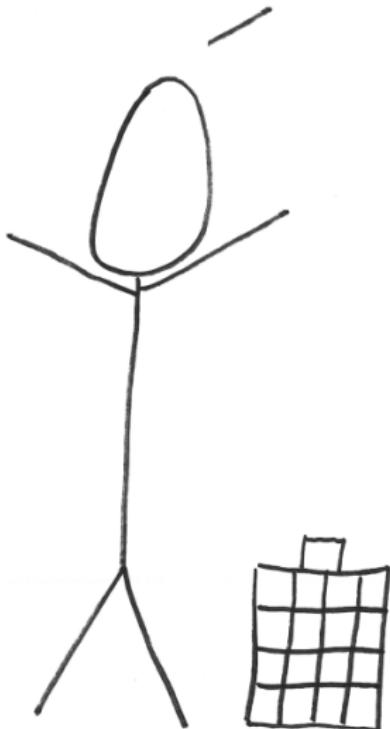
A Stick Figure Guide to the Advanced Encryption Standard (AES)



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<http://www.moserware.com/>

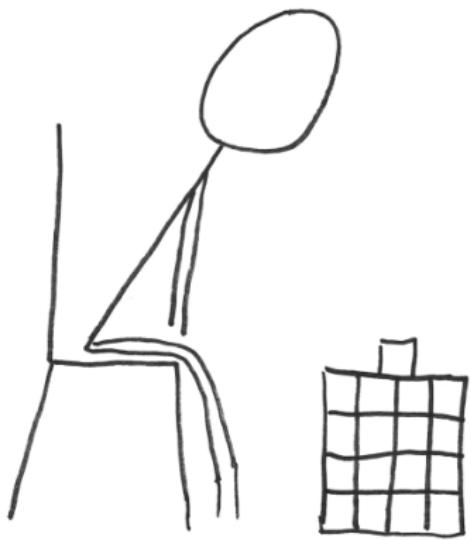
Act 1: Once Upon a Time...

I handle petabytes* of data every day. From encrypting juicy Top Secret intelligence to boring packets bound for your WiFi router, I do it all!

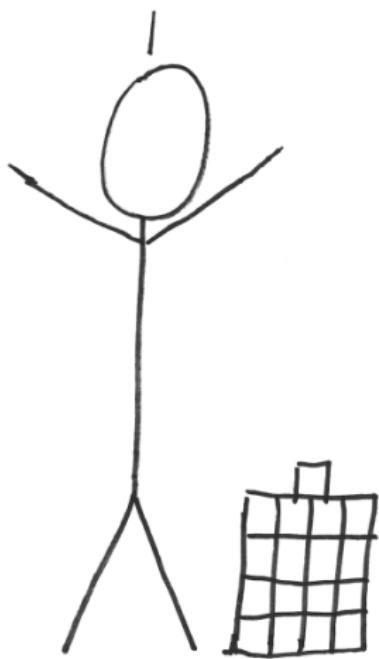


* 1 petabyte ≈ a lot

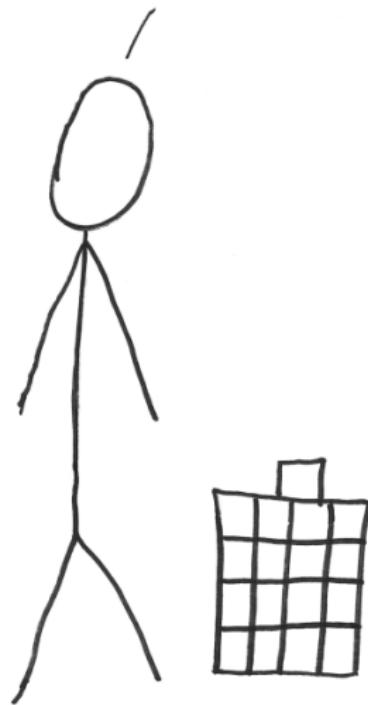
... and still no one seems to care
about me or my story.



I've got a better-than-Cinderella
story as I made my way to become
king of the block cipher world.



Whoa! You're still there. You want
to hear it? Well let's get started...



Once upon a time,* there was no good way for people outside secret agencies to judge good crypto.

EBG13 vf terng!

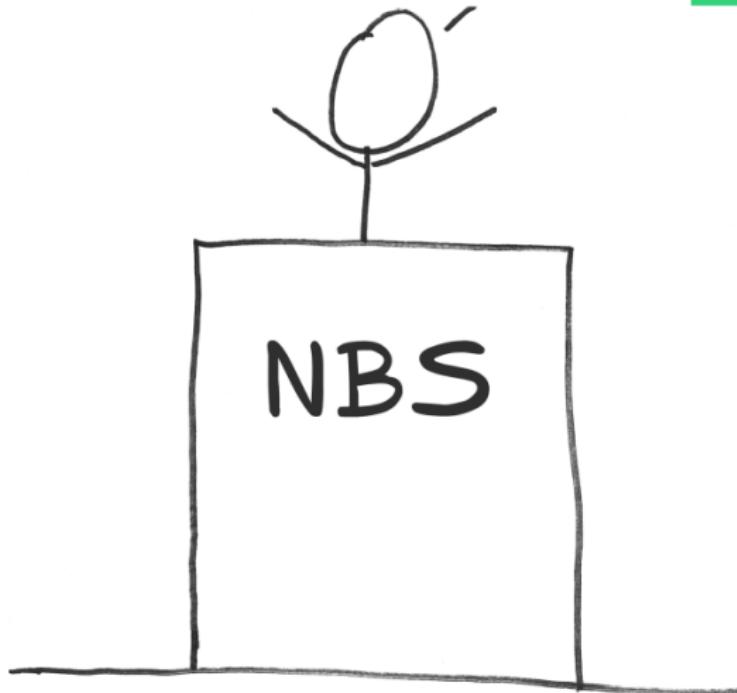
Double ROT13 –
is better!



* ~ pre-1975 for the general public

A decree went throughout the land to find a good, secure, algorithm.

We need a **good cipher!**

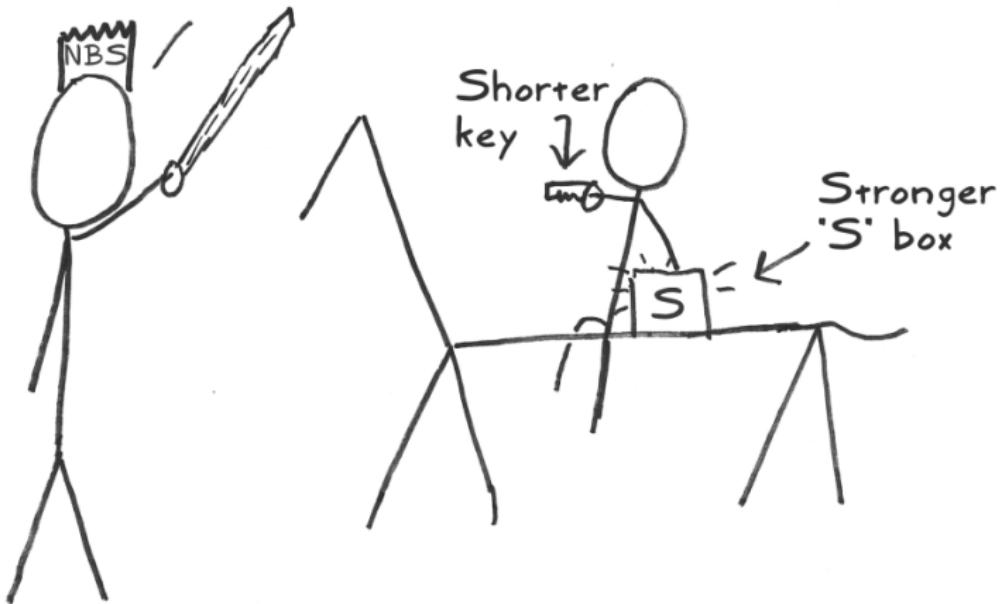


One worthy competitor named Lucifer came forward.



After being modified by the National Security Agency (NSA), he was anointed as the Data Encryption Standard (DES).

I anoint thee as DES!



DES ruled in the land for over 20 years. Academics studied him intently. For the first time, there was something specific to look at. The modern field of cryptography was born.

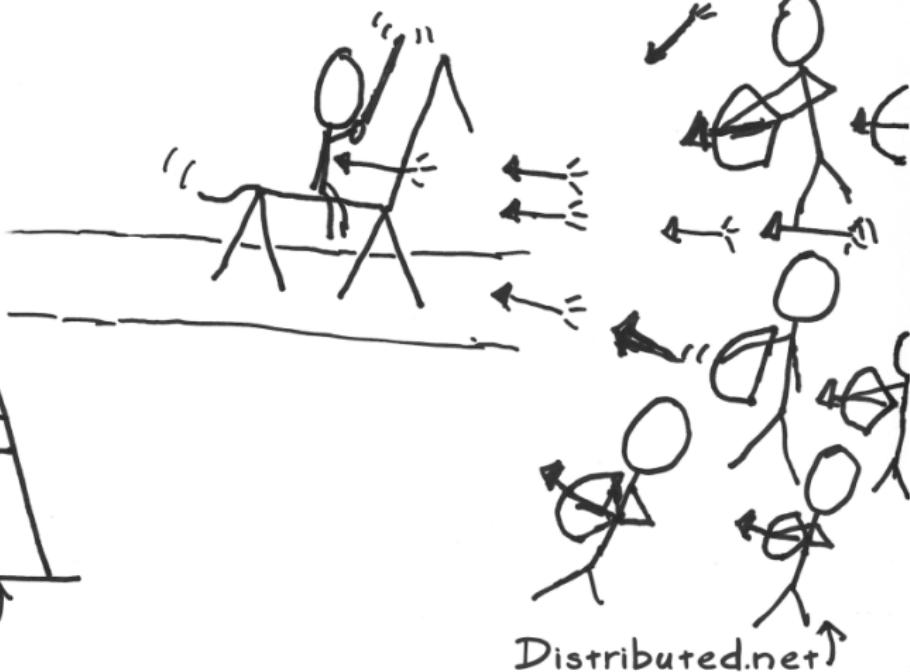
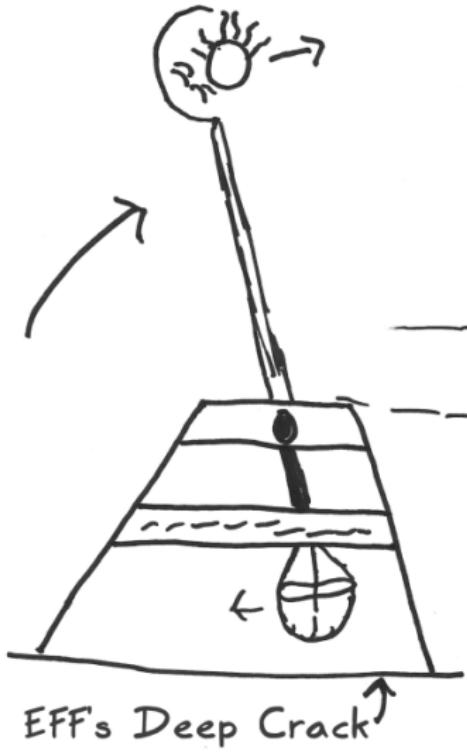
... to the best of our knowledge, DES is free from any statistical or mathematical weakness.



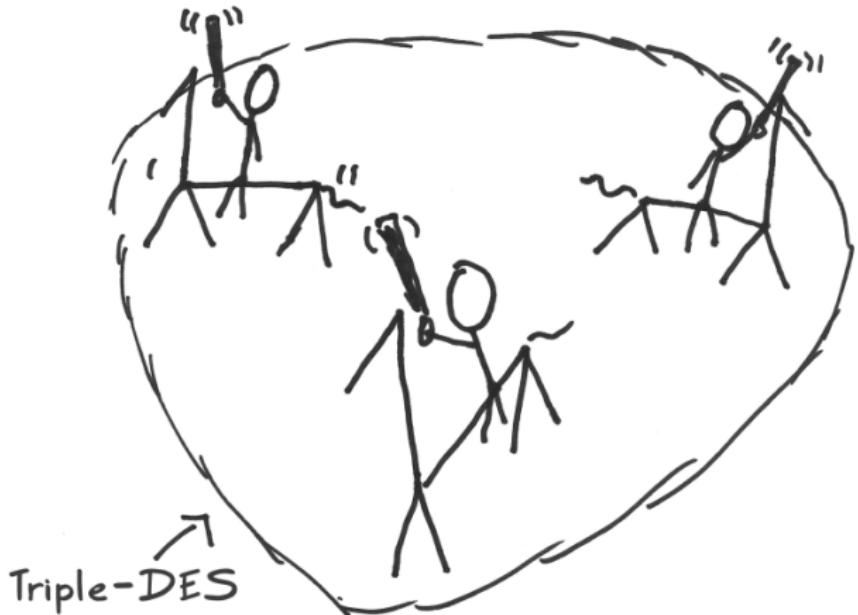
Check out that Feistel network!



Over the years, many attackers challenged DES. He was defeated in several battles.



The only way to stop the attacks was to use DES 3 times in row to form 'Triple-DES'. This worked, but it was awfully slow



Another decree went out*...



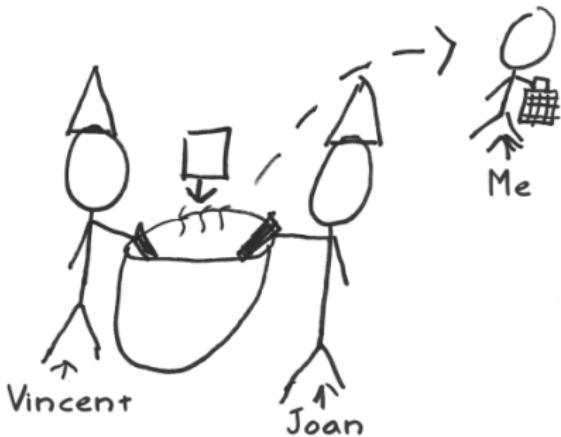
We need something at
least as strong as
Triple-DES, but it has
to be fast and flexible.

* ~ early 1997

This call rallied the crypto wizards
to develop something better.

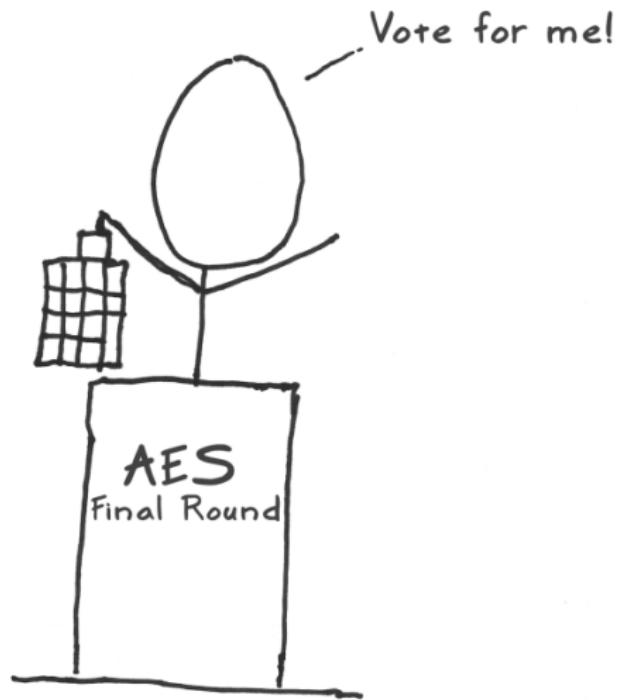


My creators, Vincent Rijmen and Joan Daemen, were among these crypto wizards. They combined their last names to give me my birth name: Rijndael.*



* That's pronounced 'Rhine Dahl' for the non-Belgians out there.

Everyone got together to vote and...

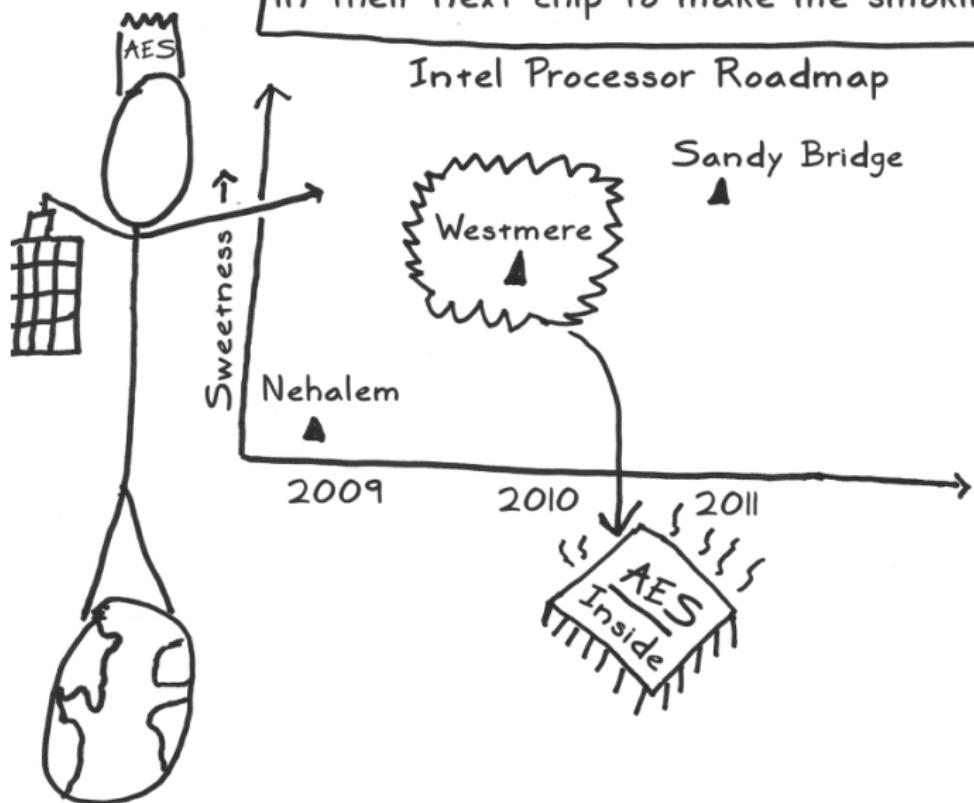


	Rijndael	Serpent	Twofish	MARS	RC6
General Security	2	3	3	3	2
Implementation Difficulty	3	3	2	1	1
Software Performance	3	1	1	2	2
Smart Card Performance	3	3	2	1	1
Hardware Performance	3	3	2	1	2
Design Features	2	1	3	2	1
Total	16	14	13	10	9

I won!!



...and now I'm the new king of the crypto world. You can find me everywhere. Intel is even putting native instructions for me in their next chip to make me smokin' fast!



Any questions?



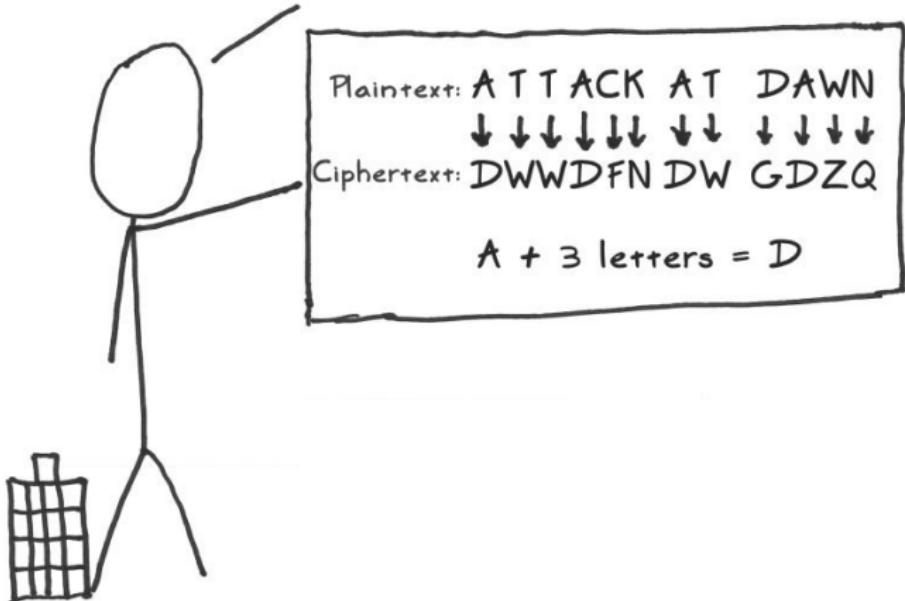
Act 2: Crypto Basics

Great question! You only need to know 3 big ideas to understand crypto.



Big Idea #1: Confusion

It's a good idea to obscure the relationship between your real message and your 'encrypted' message. An example of this 'confusion' is the trusty ol' Caesar Cipher:



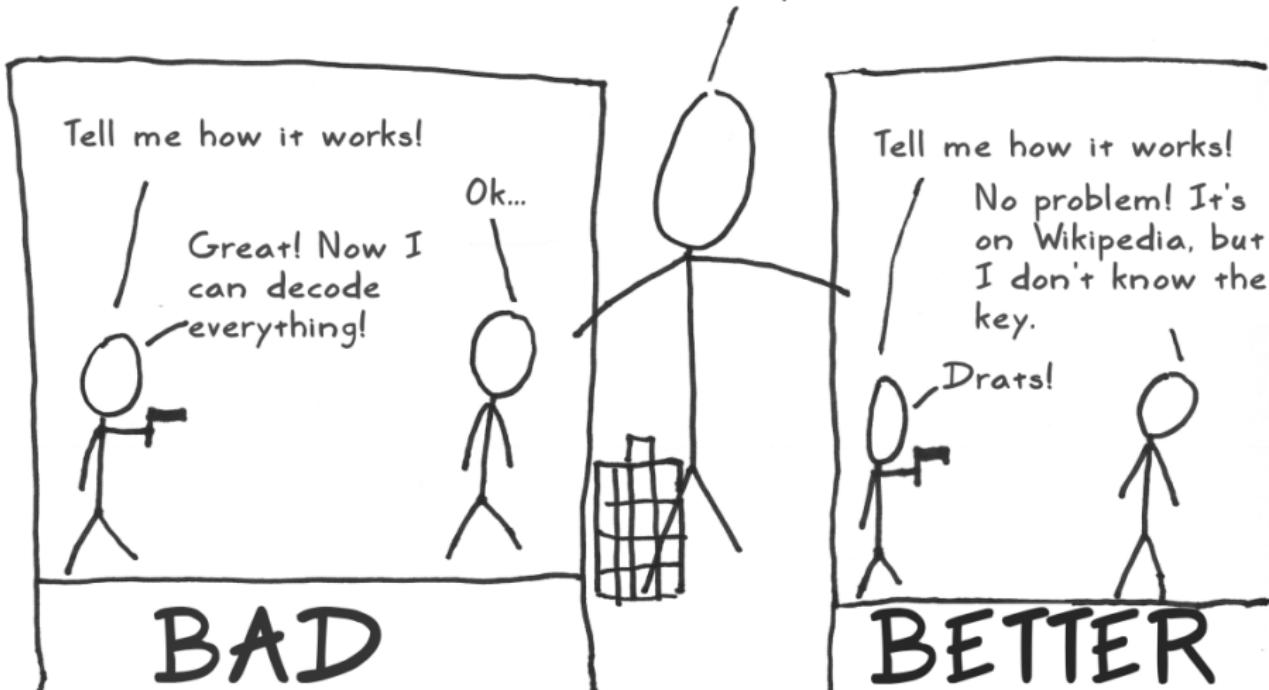
Big Idea #2: Diffusion

It's also a good idea to spread out the message. An example of this "diffusion" is a simple column transposition:

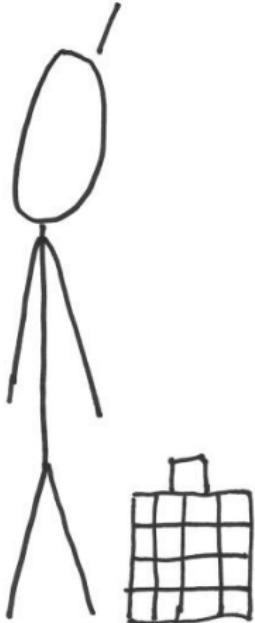


Big Idea #3: Secrecy Only in the Key

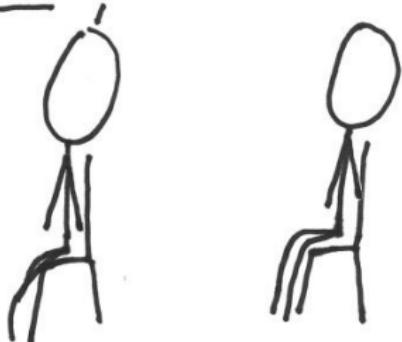
After thousands of years, we learned that it's a bad idea to assume that no one knows how your method works. Someone will eventually find that out.



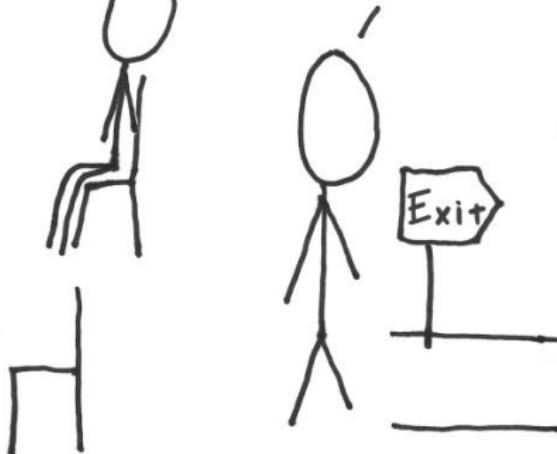
Does that answer
your question?



That helps, but was
too general. How do
you work?

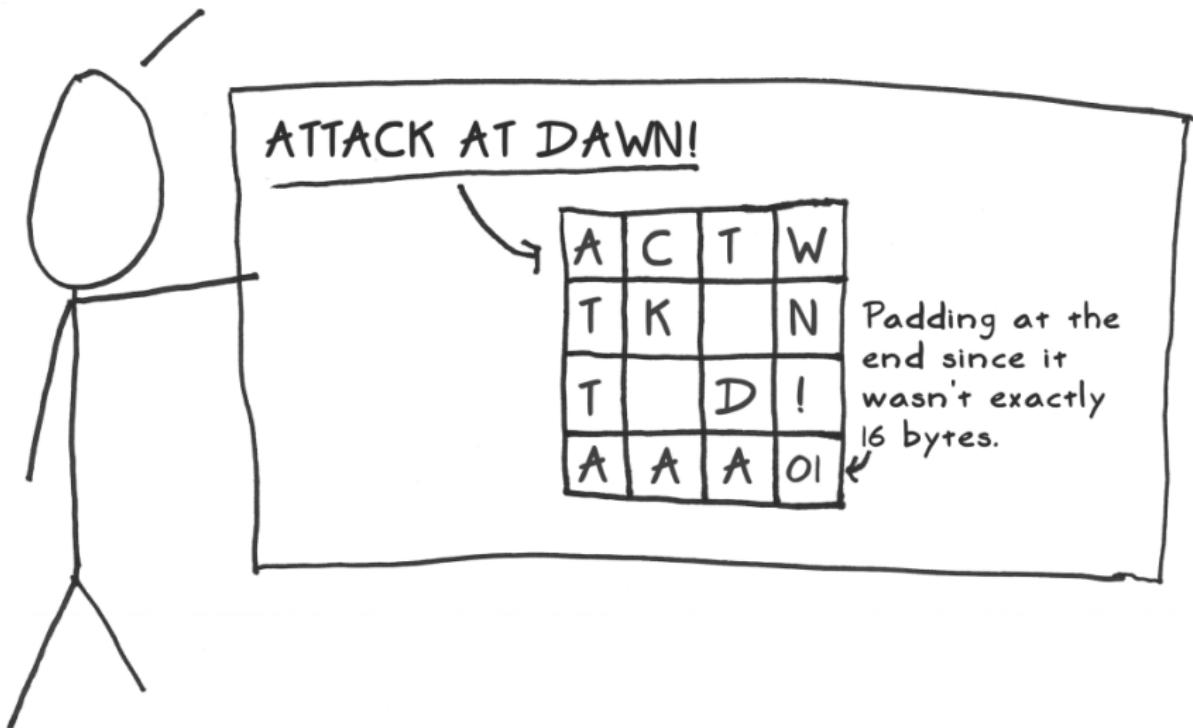


Details? I
can't handle
details!



Act 3: Details

I take your data and load it
into this 4x4 square.*



* This is the 'state matrix' that I carry with me at all times.

The initial round has me xor each input byte with the corresponding byte of the first round key.



A	C	T	W
T	K		N
T		D	!
A	A	A	O

 \oplus

S			
O	I	B	K
M	2	I	E
E	8	T	Y

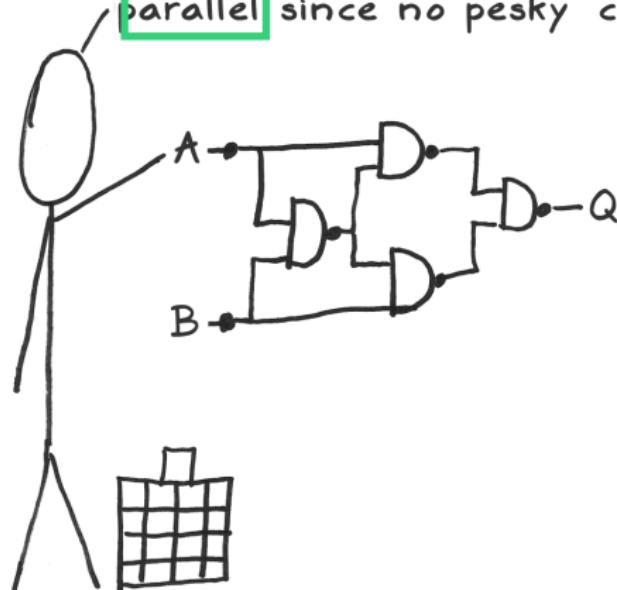
 $=$

12	63	74	77
1b	7a	62	05
19	12	0d	64
04	79	15	58



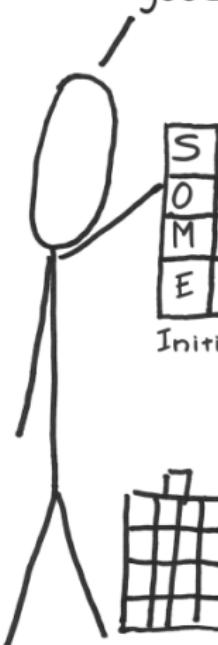
A Tribute to XOR

There's a simple reason why I use xor to apply the key and in other spots: it's fast and cheap - a quick bit flipper. It uses minimal hardware and can be done in parallel since no pesky 'carry' bits are needed.



Key Expansion: Part 1

I need lots of keys for use in later rounds. I derive all of them from the initial key using a simple mixing technique that's really fast. Despite its critics,* it's good enough.



S			
O	1	B	K
M	2	I	E
E	8	T	Y

Initial Key

e1	c1	e1	c1
21	10	52	19
86	b4	fd	b8
f2	ca	9e	c7

#1

...

ae	a6	a0	d4
97	d8	a6	c5
4d	7d	7a	d9
ef	ed	05	06

#9

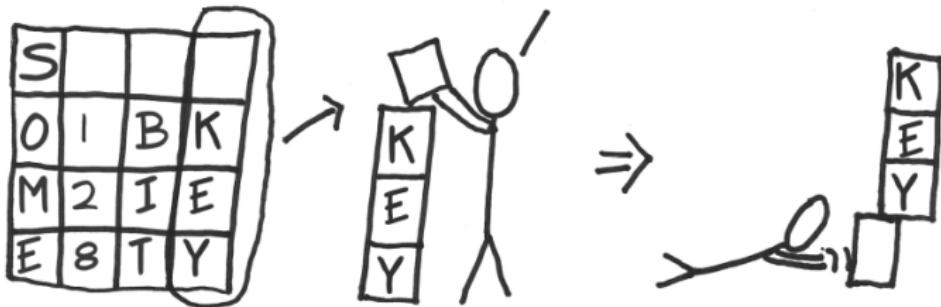
3e	98	38	ec
a2	7a	dc	19
22	5f	25	fc
a7	4a	4f	49

#10

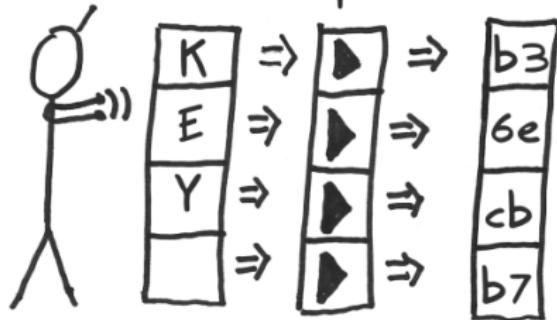
* By far, most complaints against AES's design focus on this simplicity.

Key Expansion: Part 2a

- ① I take the last column of the previous round key and move the top byte to the bottom:

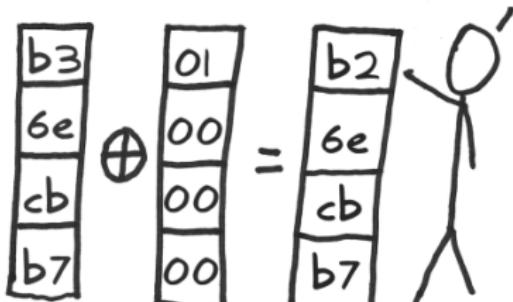


- ② Next, I run each byte through a substitution box that will map it to something else:

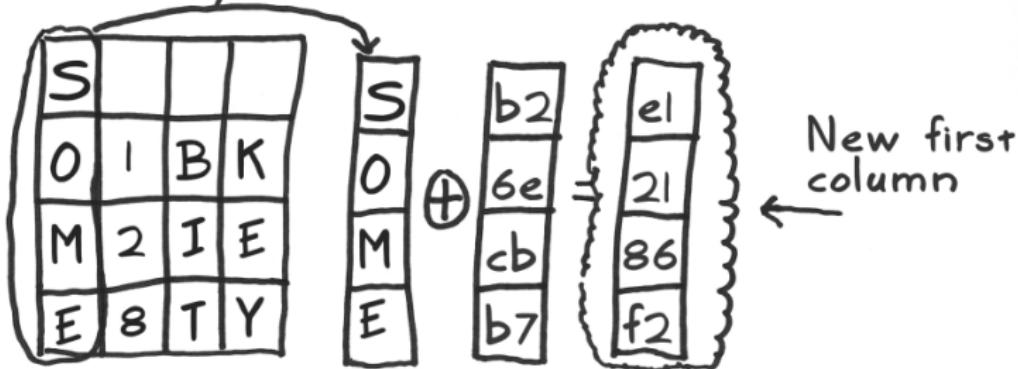


Key Expansion: Part 2b

- ③ I then xor the column with a 'round constant' that is different for each round.

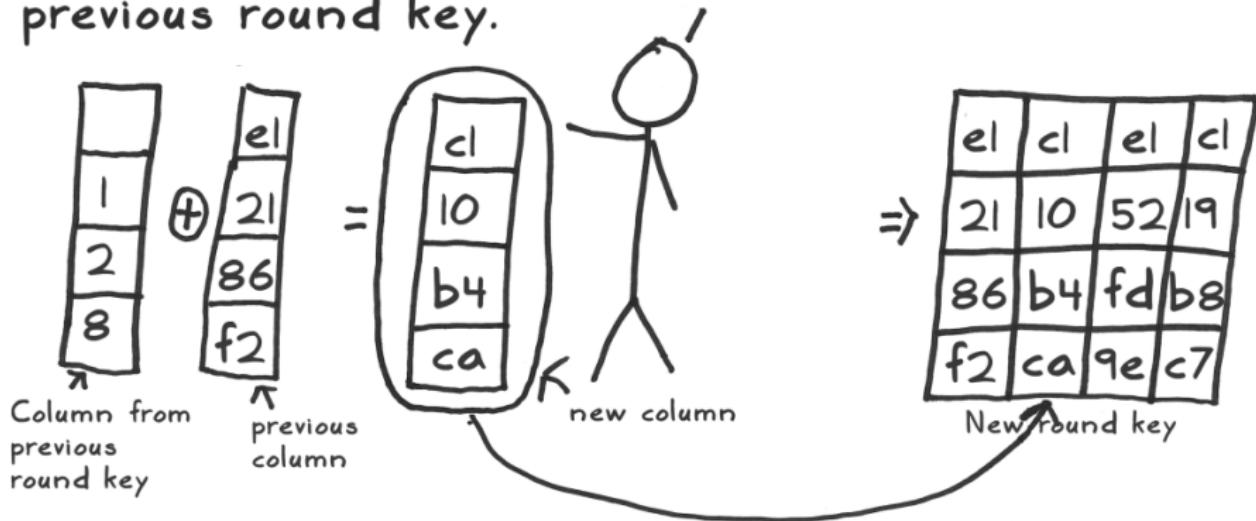


- ④ Finally, I xor it with the first column of the previous round key:



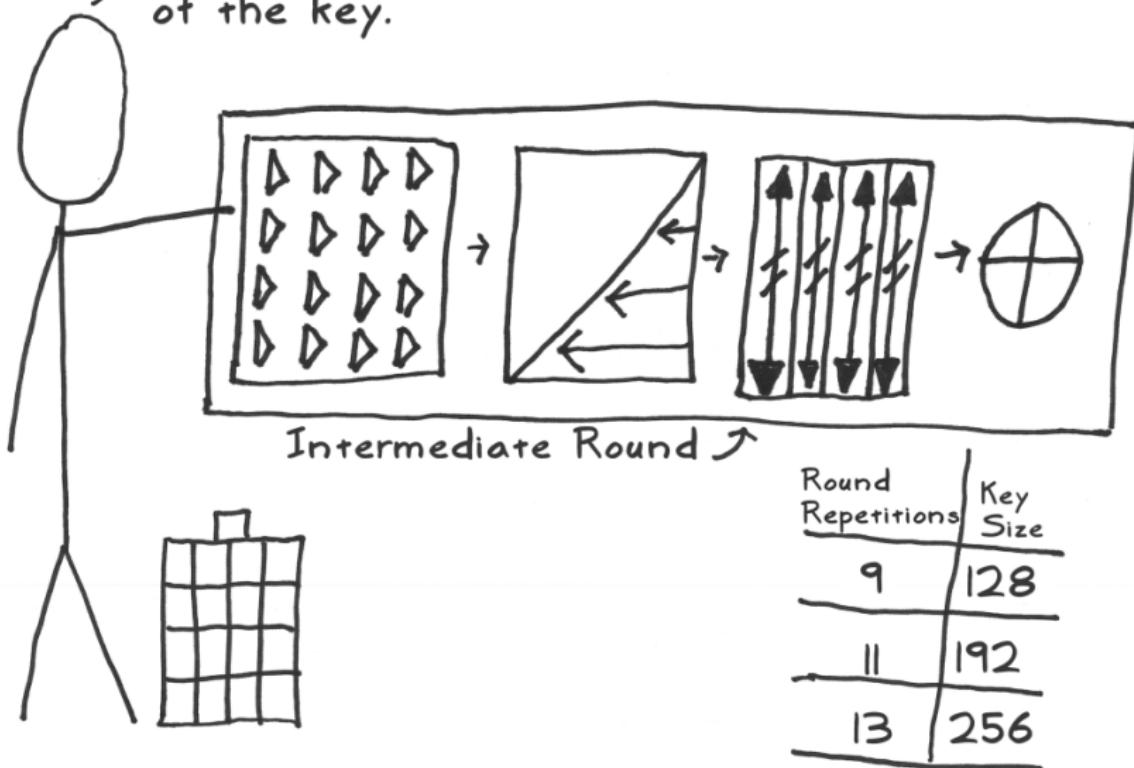
Key Expansion: Part 3

The other columns are super-easy.* I just xor the previous column with the same column of the previous round key.



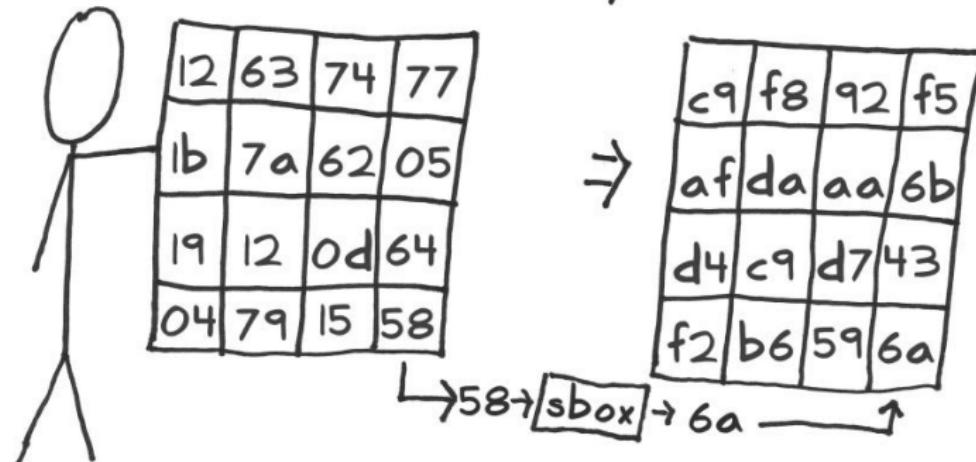
*Note that 256 bit keys are slightly more complicated.

Next, I start the intermediate rounds. A round is just a series of steps I repeat several times. The number of repetitions depends on the size of the key.

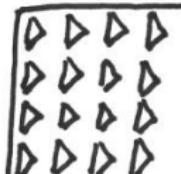


Applying Confusion: Substitute Bytes

I use confusion (Big Idea #1) to obscure the relationship of each byte. I put each byte into a substitution box (sbox), which will map it to a different byte:



Denotes
'confusion'



Applying Diffusion, Part 1: Shift Rows

Next I shift the rows to the left

Hiiiii yaah!

	c9	fb	92	f5
	af	da	aa	6b
d4	c9	d7	43	
f2	b6	59	6a	



c9	fb	92	f5
da	aa	6b	af
d7	43	d4	c9
6a	f2	b6	59

...and then wrap them around the other side

Denotes
permutation



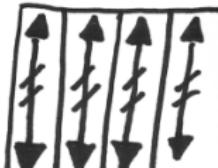
Applying Diffusion, Part 2: Mix Columns

c9	fb	92	f5
da	aa	6b	af
d7	43	d4	ca
6a	f2	b6	59

I take each column and mix up the bits in it.



41	b9	e0	8b
6e	83	95	a9
18	da	8b	38
99	00	65	d0



Applying Key Secrecy: Add Round Key

At the end of each round, I apply the next round key with an xor:

41	b9	e0	8b
6e	83	95	a9
18	da	8b	38
99	00	65	d0



e1	c1	e1	c1
21	10	52	19
86	b4	fd	b8
f2	ca	9e	c7

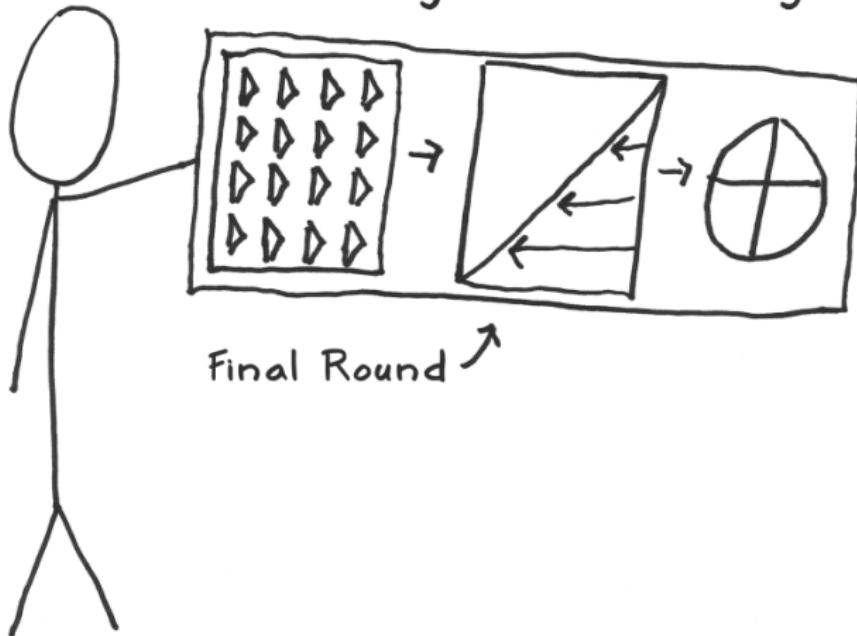


a0	78	01	4a
4f	93	c7	b0
9e	6e	76	80
6b	ca	fb	17

$$d0 \oplus c7 = 17$$



In the final round, I skip the "Mix Columns" step since it wouldn't increase security* and would just slow things down:

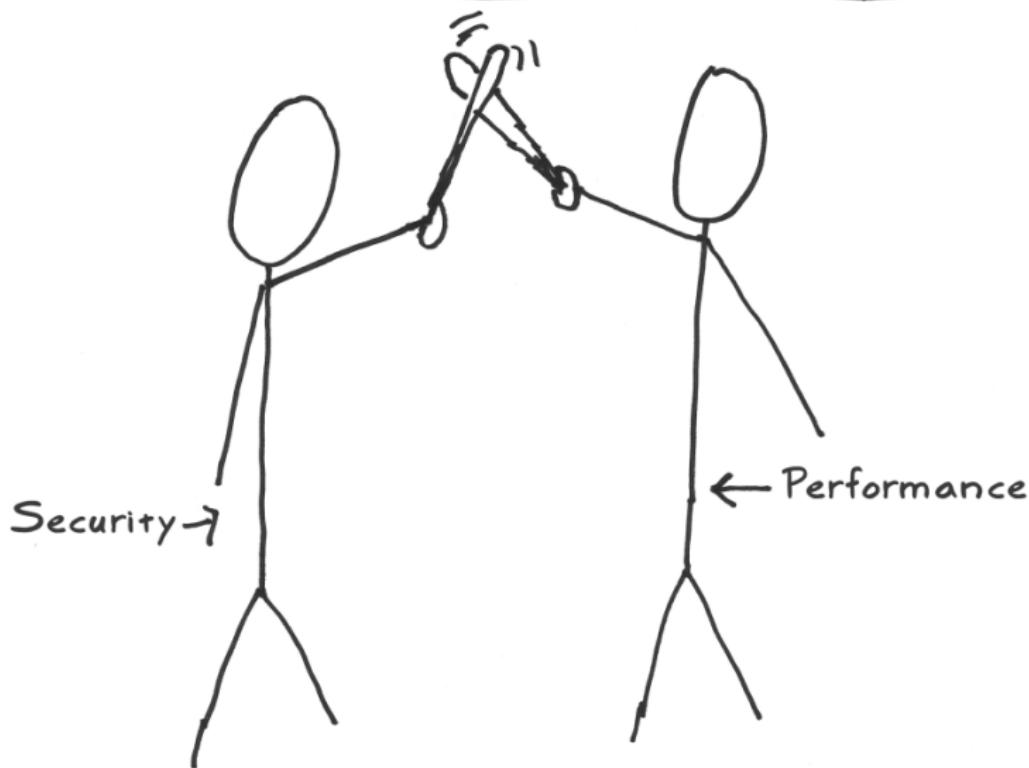


*The diffusion it would provide wouldn't go to the next round.

...and that's it. Each round I do makes the bits more confused and diffused. It also has the key impact them. The more rounds, the merrier!

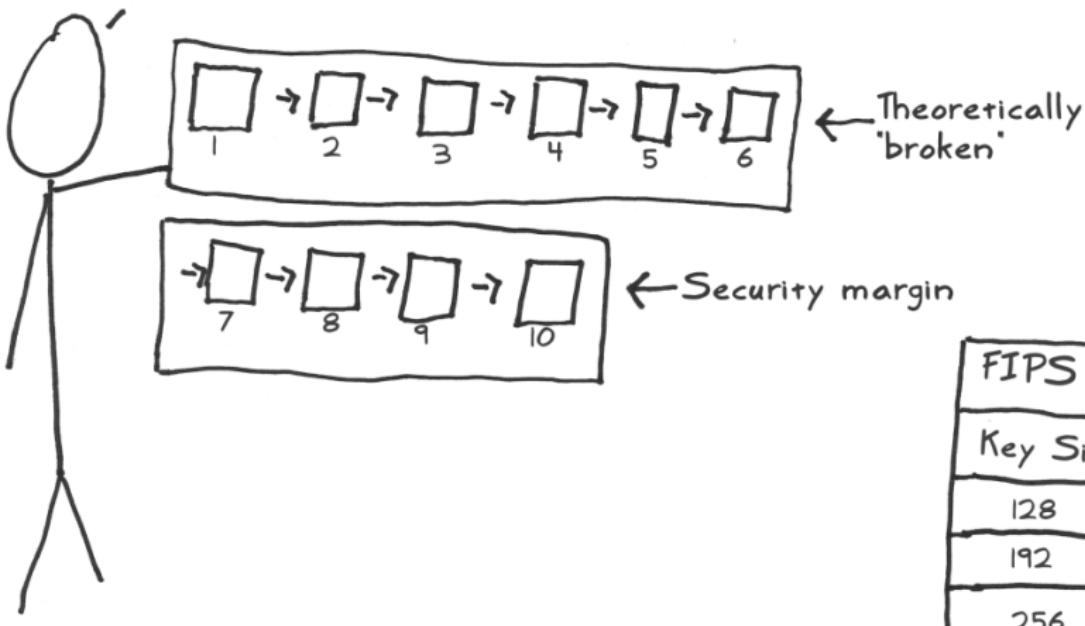


Determining the number of rounds always involves several tradeoffs.



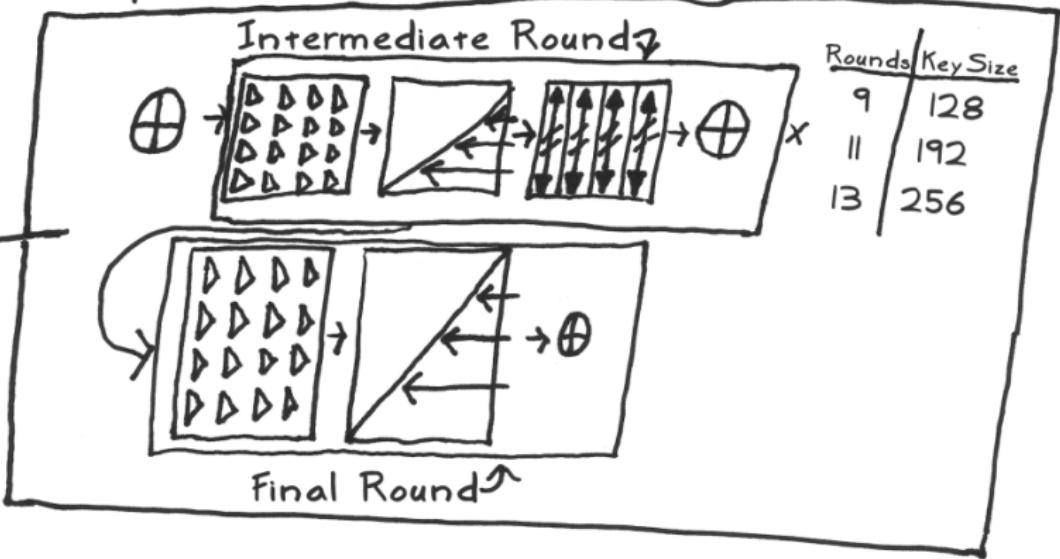
'Security always comes at a cost to performance' - Vincent Rijmen

When I was being developed, a clever guy was able to find a shortcut path through 6 rounds. That's not good! If you look carefully, you'll see that each bit of a round's output depends on every bit from two rounds ago. To increase this diffusion 'avalanche,' I added 4 extra rounds. This is my 'security margin.'



FIPS 197 Spec	
Key Size	Rounds
128	10
192	12
256	14

So in pictures, we have this:



The Story of AES Continues in Next Class...

Check out AES with Python Crypto Library or
openssl