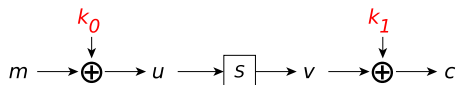


[illegible]

Instructor
Dr. Dhiman Saha

$$\alpha = (1, 0, 0, 1), \beta = (0, 0, 1, 0)$$

Recall (Sypher00A)



x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
S[x]	f	e	b	c	6	d	7	8	0	3	9	a	4	2	1	5
$\alpha \cdot x$	0	1	0	1	0	1	0	1	1	0	1	0	1	0	1	0
$\beta \cdot S[x]$	1	1	1	0	1	0	1	0	0	1	0	1	0	1	0	0

$p = ?$

$$\Pr[\alpha \cdot x = \beta \cdot S[x]] = \frac{2}{16}$$

or

$$\Pr[\alpha \cdot x \oplus 1 = \beta \cdot S[x]] = \frac{14}{16}$$

$$\alpha = 9, \beta = 2,$$

The Linear Approximation Table

► Linear Characteristic:

$$9 \xrightarrow{S} 2$$

► $LAT(9, 2) = -6$

	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
1	-2	.	2	.	-2	4	-2	2	4	2	.	-2	.	2	.
2	2	-2	.	-2	.	.	2	2	4	.	2	4	-2	-2	.
3	4	2	2	-2	2	.	.	.	2	-2	-2	-2	.	4	.
4	.	-2	2	2	-2	.	.	-4	2	2	2	2	.	4	.
5	-2	2	.	2	4	.	2	-2	4	.	-2	.	2	-2	.
6	-2	.	2	.	2	4	2	2	-4	2	.	2	.	-2	.
7	.	.	.	4	.	-4	.	.	.	4	.	4	.	.	.
8	.	-2	2	-4	.	2	2	-4	.	-2	-2	.	2	-2	.
9	-2	-6	.	.	2	-2	.	2	.	-2	-2	.	.	2	.
a	-2	.	-6	-2	.	2	.	-2	.	2	.	-2	.	2	.
b	.	.	.	2	-2	2	-2	.	.	-4	4	2	-2	-2	2
c	.	.	.	-2	-2	-2	-2	.	.	4	-4	2	2	-2	-2
d	-2	.	2	2	.	-2	.	-2	.	2	.	-6	.	-2	.
e	2	-2	.	.	2	2	-4	-2	.	2	-2	.	-4	-2	.
f	-4	2	2	-4	.	-2	-2	.	-2	2	.	.	.	-2	2

► Implication

$$\Pr \left[9 \xrightarrow{S} 2 \right] = \Pr \left[9 \cdot x = 2 \cdot S[x] \right]$$

$$= \left(\frac{-6}{16} + \frac{1}{2} \right) = \frac{1}{8} \text{ 🍷}$$

► So, we take the complement event:

$$9 \cdot x \oplus 1 = 2 \cdot S[x]$$

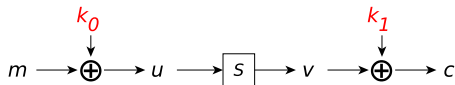
► $\Pr \left[9 \cdot x \oplus 1 = 2 \cdot S[x] \right] = 1 - \frac{1}{8} = \frac{7}{8} \text{ 🍷}$

► For Sypher00A, \Rightarrow

$$\Pr \left[(9 \cdot m) \oplus (2 \cdot c) \oplus 1 = (9 \cdot k_0) \oplus (2 \cdot k_1) \right] = \frac{7}{8} \text{ 🍷}$$

$$\alpha = 9, \beta = 2,$$

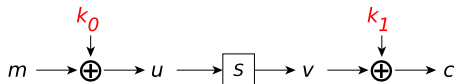
Summarizing LC of Sypher00A



- ▶ $\Pr[9 \cdot u \oplus 1 = 2 \cdot v] = \frac{7}{8}$
- ▶ $\Pr[(9 \cdot m) \oplus (2 \cdot c) \oplus 1 = (9 \cdot k_0) \oplus (2 \cdot k_1)] = \frac{7}{8}$

Procedure

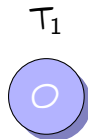
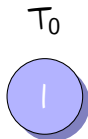
- ▶ Initialize counters T_0 and T_1 to 0
- ▶ Request the encryptions of N known plaintexts.
- ▶ For each plaintext-ciphertext pair, we compute the **left-hand side** of the equation: $(9 \cdot m) \oplus (2 \cdot c) \oplus 1$,
 - ▶ Which is either 0 or 1.
- ▶ Gives an estimate for the value of $(9 \cdot k_0) \oplus (2 \cdot k_1)$
- ▶ T_0++ if LHS evaluates to 0; T_1++ if LHS evaluates to 1

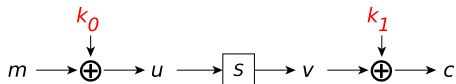


► Message Count #1

$m = 0, c = 6$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 0) \oplus (2 \cdot 6) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

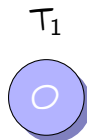
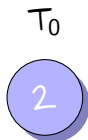


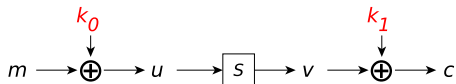


► Message Count #2

$m = 1, c = 0$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 1) \oplus (2 \cdot 0) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

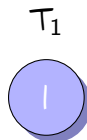
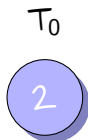


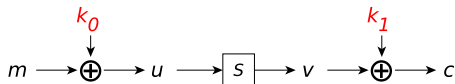


► Message Count #3

$m = 2, c = 1$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 2) \oplus (2 \cdot 1) \oplus 1 \\
 &= 1 \implies T_1 ++
 \end{aligned}$$

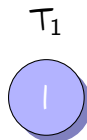
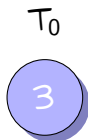


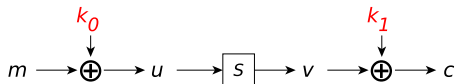


► Message Count #4

$m = 3, c = 5$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 3) \oplus (2 \cdot 5) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

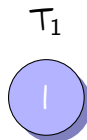
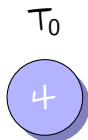


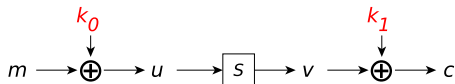


► Message Count #5

$m = 4, c = 7$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 4) \oplus (2 \cdot 7) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

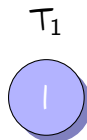
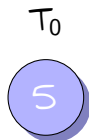


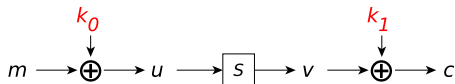


► Message Count #6

$m = 5, c = 4$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 5) \oplus (2 \cdot 4) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

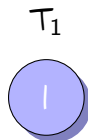
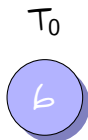


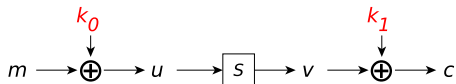


► Message Count #1

$m = 6, c = 14$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 6) \oplus (2 \cdot 14) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

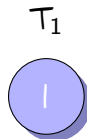
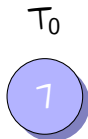


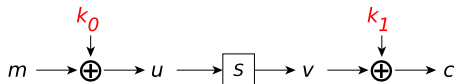


► Message Count #8

$m = 7, c = 13$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 7) \oplus (2 \cdot 13) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

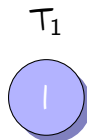
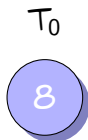


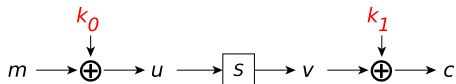


► Message Count #9

$m = 8, c = 9$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 8) \oplus (2 \cdot 9) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

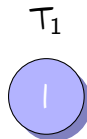
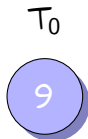


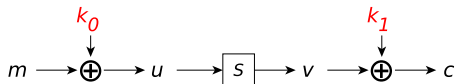


► Message Count #10

$m = 9, c = 2$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 9) \oplus (2 \cdot 2) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

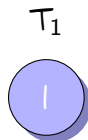
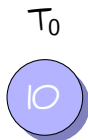


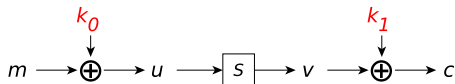


► Message Count #11

$m = 10, c = 12$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 10) \oplus (2 \cdot 12) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

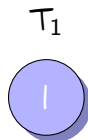
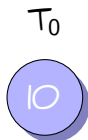


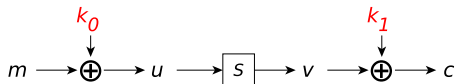


► Message Count #12

$m = 11, c = 3$

$$\begin{aligned} LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\ &= (9 \cdot 11) \oplus (2 \cdot 3) \oplus 1 \\ &= 0 \implies T_0++ \end{aligned}$$

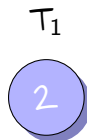
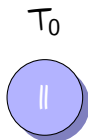


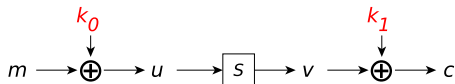


► Message Count #13

$m = 12, c = 10$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 12) \oplus (2 \cdot 10) \oplus 1 \\
 &= 1 \implies T_1++
 \end{aligned}$$

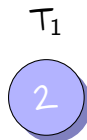
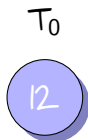


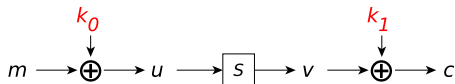


► Message Count #14

$m = 13, c = 11$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 13) \oplus (2 \cdot 11) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

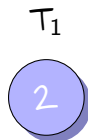
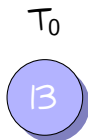


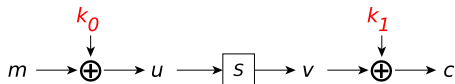


► Message Count #15

$m = 14, c = 8$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 14) \oplus (2 \cdot 8) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$

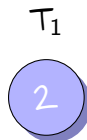
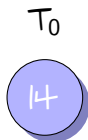




► Message Count #16

$m = 15, c = 15$

$$\begin{aligned}
 LHS &= (\alpha \cdot m) \oplus (\beta \cdot c) \oplus 1 \\
 &= (9 \cdot 15) \oplus (2 \cdot 15) \oplus 1 \\
 &= 0 \implies T_0++
 \end{aligned}$$



- ▶ $RHS = (\alpha \cdot k_0) \oplus (\beta \cdot k_1) \stackrel{?}{=} 0/1$
- ▶ Key-bit estimation correct with prob. $\frac{14}{16}$
- ▶ What to expect at T_0/T_1 after N KP encryptions

If $(\alpha \cdot k_0) \oplus (\beta \cdot k_1) = 1$

$$T_0 \leftarrow \frac{2N}{16}$$

$$T_1 \leftarrow \frac{14N}{16}$$

If $(\alpha \cdot k_0) \oplus (\beta \cdot k_1) = 0$

$$T_0 \leftarrow \frac{14N}{16}$$

$$T_1 \leftarrow \frac{2N}{16}$$

Here, $N = 16$

T_0



T_1



- ▶ Verifying any one counter say, T_0
 - ▶ Reveals one bit $\rightarrow (\alpha \cdot k_0) \oplus (\beta \cdot k_1)$
 - ▶ **Attack Outcome** $\rightarrow (9 \cdot k_0) \oplus (2 \cdot k_1) = 0$

m	c	$(9 \cdot m) \oplus (2 \cdot c) \oplus 1$	T_0	T_1	Remarks
0	6	0	1	0	T_0++
1	0	0	2	0	T_0++
2	1	1	2	1	T_1++
3	5	0	3	1	T_0++
4	7	0	4	1	T_0++
5	4	0	5	1	T_0++
6	14	0	6	1	T_0++
7	13	0	7	1	T_0++
8	9	0	8	1	T_0++
9	2	0	9	1	T_0++
10	12	0	10	1	T_0++
11	3	0	10	1	T_0++
12	10	1	11	2	T_1++
13	11	0	12	2	T_0++
14	8	0	13	2	T_0++
15	15	0	14	2	T_0++

- ▶ Every group needs to generate the Hawk-Eye Table from last slide using their own oracles for (m, c) pairs and submit in the **notebook**.
- ▶ You are free to choose any of the masks you used for the In-Class assignment.