

# NYU Tandon School of Engineering

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## Homework Assignment 3

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**Homework Assignment 3** [released Monday September 26<sup>th</sup> 2022] [due Wednesday October 5<sup>th</sup> by 11:59PM]

You *are allowed* to discuss HW assignments with anyone. You are *not allowed* to share your solutions with other colleagues in the class. Please feel free to reach out to the Course Assistants or the Instructor during office hours or by appointment if you need any help with the HW. Please enter your responses in this Word document after you download it from NYU Classes. *Please use the Brightspace portal to upload your completed HW.*

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- *In RISC V, only load and store instructions access memory locations*
- *These instructions must follow a 'format' to access memory*
- *Assume a 32-bit machine in all problems unless asked to assume otherwise*

### **Problem 1:**

Assume address in memory of 'A[0]', 'B[0]' and 'C[0]') are stored in Registers **x27**, **x30**, **x31**. Assume values of variables **f**, **g**, **h**, **i**, and **j** are assigned to registers **x5**, **x6**, **x7**, **x28**, **x29** respectively

Write down RISC V Instruction(s) to

- (a) Load Register **x5** with content of **A[10]**
- (b) Store contents of Register **x5** into **A[17]**
- (c) add 2 operands: one in **x5** - a register, the other in in Register **x6**. Assume result of operation to be stored in register **x7**
- (d) copy contents at one memory location to another: **C[g] = A[i+j+31]**
- (e) implement in RISC V these line of code in C:
  - (i) **f = g - A[B[9]]**
  - (ii) **f = g - A[C[8] + B[4]]**
  - (iii) **A[i] = B[2i+1], C[i] = B[2i]**

(iv)  $A[i] = 4B[i-1] + 4C[i+1]$

(v)  $f = g - A[C[4] + B[12]]$

### **Problem 2:**

Assume the following register contents:

$x5 = 0x00000000AAAAAAAA$ ,  $x6 = 0x1234567812345678$

a. For the register values shown above, what is the value of x7 for the following sequence of instructions?

```
srli x7, x5, 16
addi x7, x7, -128
srai x7, x7, 2
and x7, x7, x6
```

b. For the register values shown above, what is the value of x7 for the following sequence of instructions?

```
slli x7, x6, 4
```

c. For the register values shown above, what is the value of x7 for the following sequence of instructions?

```
srli x7, x5, 3
andi x7, x7, 0xFEF
```

### **Problem 3:**

For each RISC-V instruction below, identify the instruction format and show, wherever applicable, the value of the opcode (**op**), source register (**rs1**), source register (**rs2**), destination register (**rd**), immediate (**imm**), **func3**, **func7** fields. Also provide the 8 hex char (or 32 bit) instruction for each of the instructions below

```
add x5, x6, x7
```

```
addi x8, x5, 512
```

```
ld x3, 128(x27)
```

```
sd x3, 256(x28)
```

```
beq x5, x6 ELSE #ELSE is the label of an instruction 16 bytes larger
#than the current content of PC
```

```
add x3, x0, x0
auipc x3, FFEFA
jal x3 ELSE
```

**Problem 4:**

(a) For the following C statement, write a minimal sequence of RISC-V assembly instructions that performs the identical operation. Assume `x5 = A`, and `x11` is the base address of `C`.

```
A = C[0] << 16;
```

(b) Find the shortest sequence of RISC-V instructions that extracts `bits 12 down to 7` from register `x3` and uses the value of this field to replace bits `28 down to 23` in register `x4` without changing the other bits of registers `x3` or `x4`. (Be sure to test your code using `x3 = 0` and `x4 = 0xffffffffffffffff`. Doing so may reveal a common oversight.)

(c) Provide a minimal set of RISC-V instructions that may be used to implement the following pseudoinstruction:

```
not x5, x6 // bit-wise invert
```

*[Hint: note that there is no ‘not’ instruction in RISC-V. However, an XOR immediate instruction could be used]*

**Problem 5:**

Suppose the program counter (PC) is set to `0x60000000hex`.

- a. What range of addresses can be reached using the RISC-V *jump-and-link* (`jal`) instruction? (In other words, what is the set of possible values for the PC after the jump instruction executes?)
- b. What range of addresses can be reached using the RISC-V *branch if equal* (`beq`) instruction? (In other words, what is the set of possible values for the PC after the branch instruction executes?)

**Problem 6:**

Assume that the register `x6` is initialized to the value 10. What is the final value in register `x5` assuming the `x5` is initially zero?

```
LOOP:    beq x6, x0, DONE
         addi x6, x6, -1
         addi x5, x5, 2
         jal x0, LOOP
DONE:
```

- a. For the loop above, write the equivalent C code. Assume that the registers `x5` and `x6` are integers `acc` and `i`, respectively.

- b. For the loop written in RISC-V assembly above, assume that the register **x6** is initialized to the value **N**. How many RISC-V instructions are executed?
- c. For the loop written in RISC-V assembly above, replace the instruction “**beq x6, x0, DONE**” with the instruction “**blt x6, x0, DONE**” and write the equivalent C code.

**Problem 7:**

- a. Translate the following C code to RISC-V assembly code. Use a minimum number of instructions. Assume that the values of **a**, **b**, **i**, and **j** are in registers **x5**, **x6**, **x7**, and **x29**, respectively. Also, assume that register **x10** holds the base address of the array **D**.

```
for(i=0; i<a; i++)
    for(j=0; j<b; j++)
        D[4*j] = i + j;
```

- b. How many RISC-V instructions does it take to implement the C code from 7a. above? If the variables **a** and **b** are initialized to **10** and **1** and all elements of **D** are initially 0, what is the total number of RISC-V instructions executed to complete the loop?

**Problem 8:**

Consider the following code:

```
lb x6, 0(x7)
sd x6, 8(x7)
```

Assume that the register **x7** contains the address **0x10000000** and the data at address is **0x1122334455667788**.

- a. What value is stored in **0x10000007** on a bigendian machine?
- b. What value is stored in **0x10000007** on a littleendian machine?

**Problem 9:**

Write the RISC-V assembly code that creates the 64-bit constant **0x1234567812345678<sub>hex</sub>** and stores that value to register **x10**.

**Problem 10:** Assume that **x5** holds the value **128<sub>10</sub>**.

- a. For the instruction **add x30, x5, x6**, what is the range(s) of values for **x6** that would result in overflow?
- b. For the instruction **sub x30, x5, x6**, what is the range(s) of values for **x6** that would result in overflow?
- c. For the instruction **sub x30, x6, x5**, what is the range(s) of values for **x6** that would result in overflow?