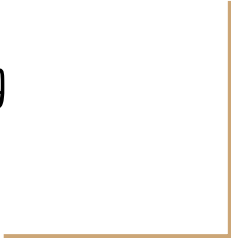




# Game Theory & Networks

(an incredibly brief overview)

**Andrew Smith**  
ECS 253/MAE 289  
May 10th, 2016



**Game theory** can help us answer important questions for scenarios where:

- players/agents (nodes)*** are *autonomous and selfish*, and
- player's connections (edges)*** directly affect their utility.

## Terminology for Games on Networks:

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  - *Pure strategies* correspond to a choice of exactly one action per player (discrete).
  - *Mixed strategies* correspond to a distribution over the action space for each player (continuous).

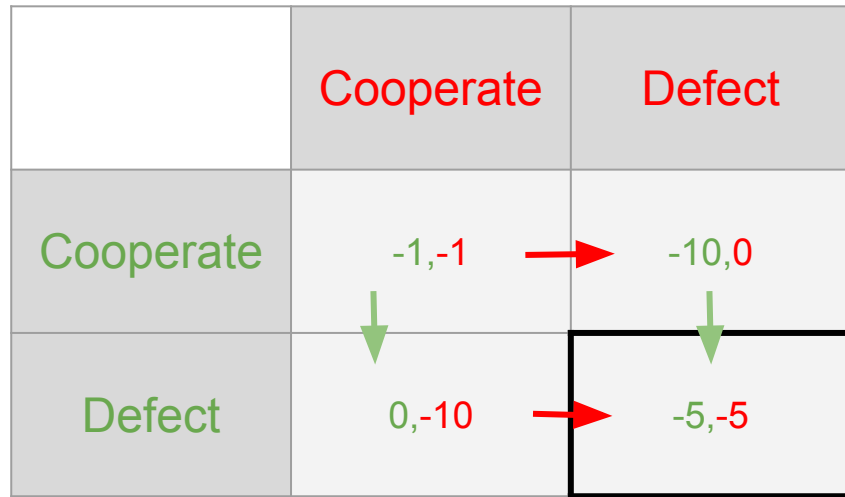
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  - *Pure strategies* correspond to a choice of exactly one action per player (discrete).
  - *Mixed strategies* correspond to a distribution over the action space for each player (continuous).
- **Utility:**  $U_i(S) \forall i \in N$ ; how much benefit a player  $i$  gets from strategy  $S$ .

# Nash Equilibrium

**Pure-strategy Nash equilibrium:** A *pure strategy* for each player, such that, given the strategy of the other players, no player would do better playing a different strategy.

	Cooperate	Defect
Cooperate	-1,-1	-10,0
Defect	0,-10	-5,-5

A 2x2 payoff matrix for the Prisoner's Dilemma. The columns are labeled 'Cooperate' and 'Defect' in red. The rows are labeled 'Cooperate' and 'Defect' in green. The payoffs are: (Cooperate, Cooperate) = -1, -1; (Cooperate, Defect) = -10, 0; (Defect, Cooperate) = 0, -10; (Defect, Defect) = -5, -5. Green arrows point from the top-left cell to the bottom-left cell and from the top-right cell to the bottom-right cell. Red arrows point from the middle-left cell to the middle-right cell and from the bottom-left cell to the bottom-right cell. The bottom-right cell, containing the payoff -5, -5, is highlighted with a thick black border.

**Prisoner's Dilemma**



# Nash Equilibrium

**Mixed-strategy Nash equilibrium:** A *mixed strategy* for each player, such that, given the strategy of the other players, no player would do better by changing their strategy.

	Swerve	Straight
Swerve	0,0	-1,1
Straight	1,-1	-10,-10

Chicken

Two “Unfair” Pure-Strategy Nash Equilibria!

# Mixed-Strategy Nash Equilibrium

- **Player 2** chooses swerve with probability  $p$  and straight with probability  $1-p$ .

	$p$	$1-p$
	Swerve	Straight
Swerve	0,0	-1,1
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Chicken

# Mixed-Strategy Nash Equilibrium

- **Player 2** chooses swerve with probability  $p$  and straight with probability  $1-p$ .
- **Player 2** wishes to make **Player 1** *indifferent* about what strategy to choose

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# Mixed-Strategy Nash Equilibrium

- **Player 2** chooses swerve with probability  $p$  and straight with probability  $1-p$ .
- **Player 2** wishes to make **Player 1** *indifferent* about what strategy to choose

$$u_1(\text{Swerve}) = u_1(\text{Straight})$$

$$0 \cdot p + -1 \cdot (1-p) = 1 \cdot p + -10 \cdot (1-p)$$

$$p - 1 = 11p - 10$$

$$p = 9/10$$

	$p$	$1-p$
	Swerve	Straight
Swerve	0,0	-1,1
Straight	1,-1	-10,-10

Chicken

# Mixed-Strategy Nash Equilibrium

- Now, **Player 1** must also randomize (making **Player 2** indifferent)

$$u_2(\text{Swerve}) = u_2(\text{Straight})$$

$$0 \cdot q + -1 \cdot (1-q) = 1 \cdot q + -10 \cdot (1-q)$$

$$q - 1 = 11q - 10$$

$$q = 9/10$$

		$p=9/10$	$1-p=1/10$
		Swerve	Straight
$q$	Swerve	0,0	-1,1
$1-q$	Straight	1,-1	-10,-10

Chicken

# Mixed-Strategy Nash Equilibrium

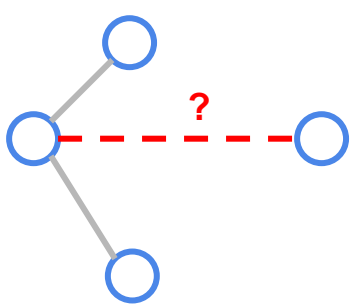
- Now, **Player 1** must also randomize (making **Player 2** indifferent)
- Mixed-strategy Nash equilibria =  $(9/10, 1/10), (9/10, 1/10)$**

		$p=9/10$	$1-p=1/10$
		Swerve	Straight
$q=9/10$	Swerve	0,0	-1,1
$1-q=1/10$	Straight	1,-1	-10,-10

Chicken

# The most well studied network scenarios....

## Network Formation Games

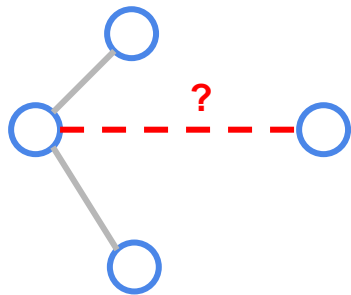


How do networks form  
given selfish, utility-  
driven players?

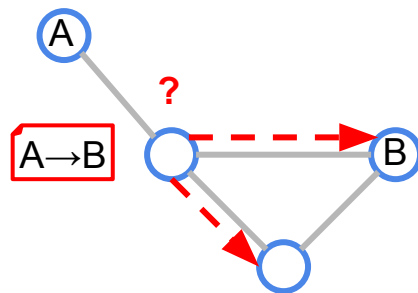
**Social networks, supply  
networks, power grids,  
etc.**

# The most well studied network scenarios....

## Network Formation Games



## Routing Games



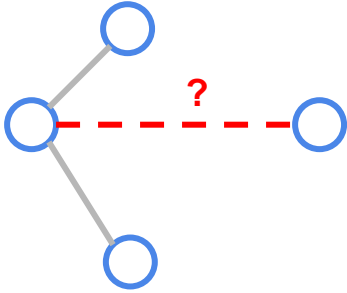
How reliable or efficient is information routing given a network structure (and selfish players)?

**Packet routing, traffic flow, information dissemination**



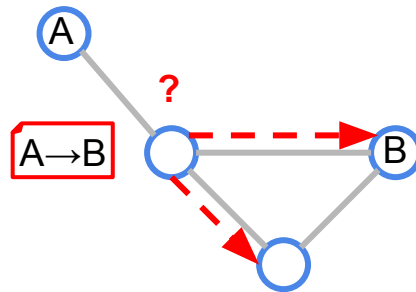
# The most well studied network scenarios....

## Network Formation Games



Equilibria in “Routing Games” can usually be illustrated by Pigou’s Principle

## Routing Games

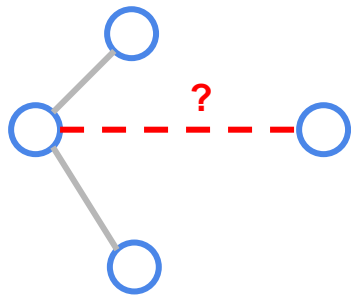


How reliable or efficient is routing flow given a network structure (and selfish players)?

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# The most well studied network scenarios....

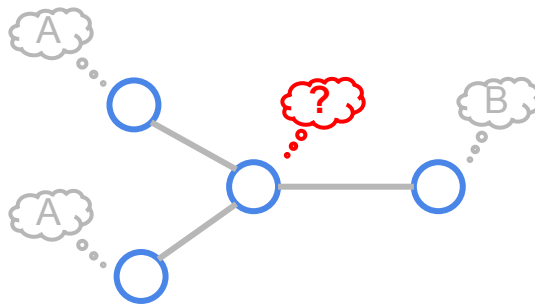
## Network Formation Games



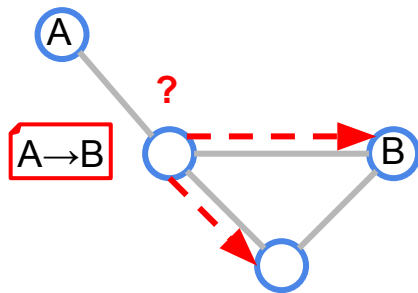
How do opinions/ideas/  
diseases spread in a  
network?

**Epidemic spread,  
voting, technology  
adaptation**

## Opinion Dynamics

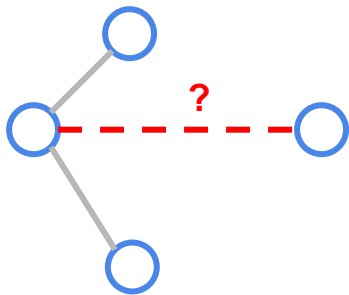


## Routing Games

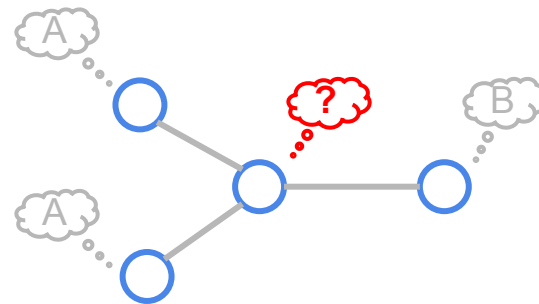


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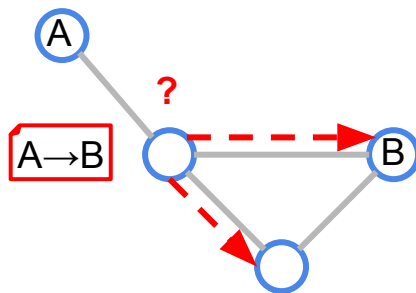
Network Formation Games



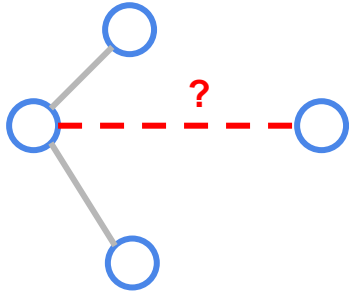
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Routing Games

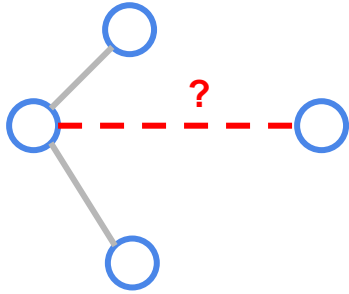


# Network Formation Games



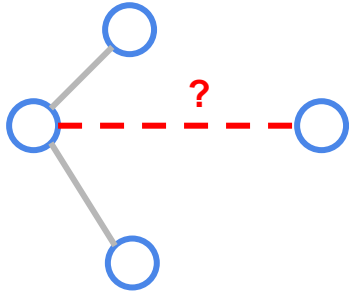
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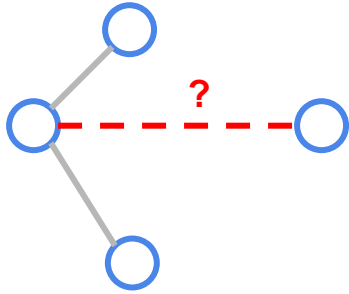


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Actions for player  $i$  (for all  $i$ ):

**{don't build edge, build edge}** <sup>$N$</sup>

# Network Formation Games



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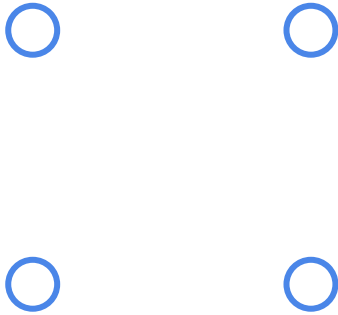
Actions for player  $i$  (for all  $i$ ):

**{don't build edge, build edge}** <sup>$N$</sup>

- **Question:** What networks emerge in Nash equilibria?

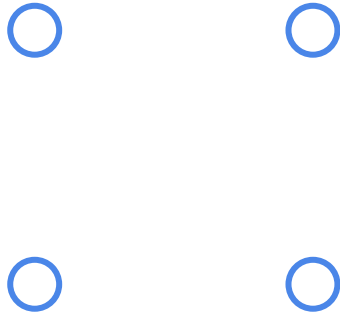
# Network Formation Games

- 4 players/nodes ( $N=4$ );  
empty network



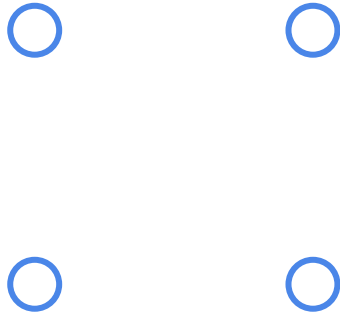


# Network Formation Games



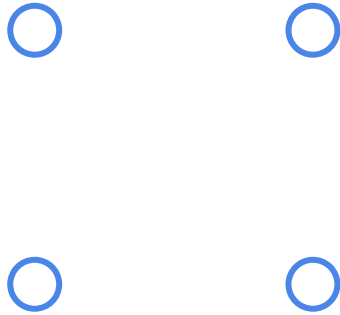
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- Does any *one* player want to deviate from the current **strategy**?

# Network Formation Games



- 4 players/nodes ( $N=4$ ); empty network
- Does any one player want to deviate from the current **strategy**?
  - No! -- They couldn't if they tried.
- **Mutual edge creation** makes Nash equilibria less interesting...

# Network Formation Games



- A network is **pairwise stable** if there is no other network configuration such that:
  - Any two pairs of nodes wishes to add an edge, and...
  - Any one node wishes to remove an edge.
- Now, we care about the *utilities of players*.

# Symmetric Connections Model

## Distance-based utility function

$$u_i = b(\ell_{ij}) - d_i c$$

$b(\ell_{ij})$  = some function on the shortest path between player  $i$  and player  $j$ .

*A game with 4 players/nodes*



# Symmetric Connections Model

## Distance-based utility function

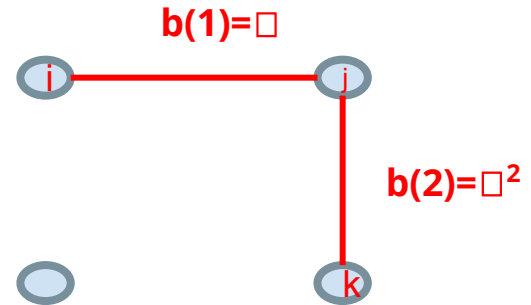
$$u_i = b(\ell_{ij}) - d_i c$$

$b(\ell_{ij})$  = some function on the shortest path between player  $i$  and player  $j$ .

$d_i$  = total degree of player  $i$ .

We will assume  $b(k) = \alpha^k$  (for  $\alpha < 1$ )

A game with 4 players/nodes

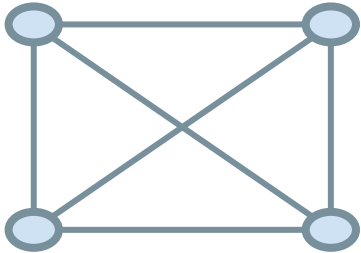


$$u_i = \alpha + \alpha^2 - c$$

$$u_j = \alpha + \alpha - 2c$$

$$u_k = \alpha + \alpha^2 -$$

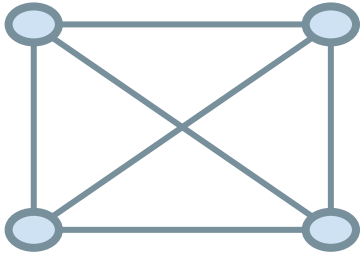
# Pairwise Stability in Symmetric Connections Model



$$c < b(1) - b(2)$$

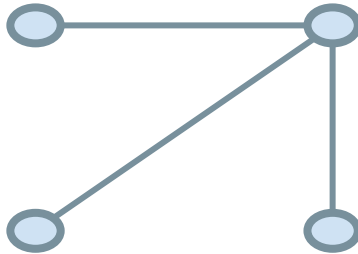
A complete network!

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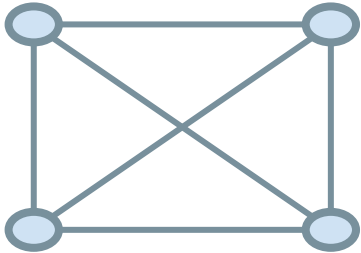
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$$b(1) - b(2) < c < b(1)$$

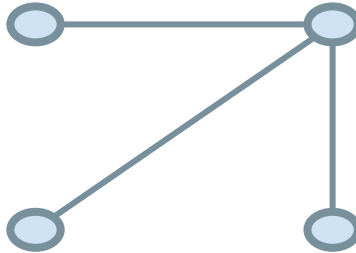
A star! (and possibly others)

# Pairwise Stability in Symmetric Connections Model



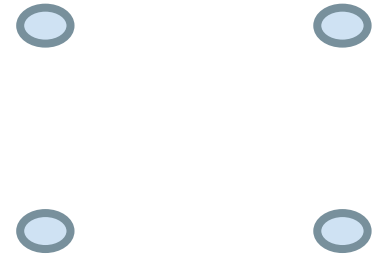
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A complete network!



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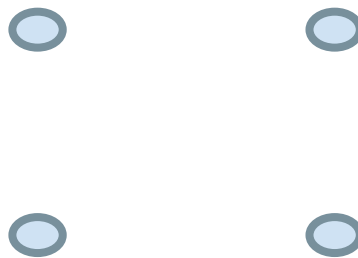
The empty network!



# Efficient Solutions in Symmetric Connections Model

Consider the case when cost is relatively high...

*A game with 4 players/nodes*



$$b(1) < c$$

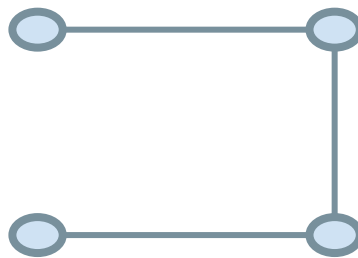
The empty network!  
Each player gets nothing!

# Efficient Solutions in Symmetric Connections Model

**Consider the case when cost is relatively high...**

- A *path* through all nodes is better for everyone!

*A game with 4 players/nodes*



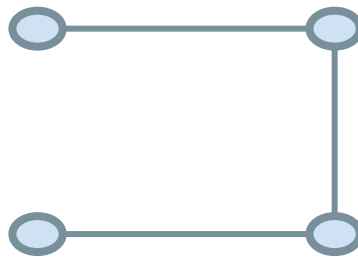
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# Efficient Solutions in Symmetric Connections Model

Consider the case when cost is relatively high...

- A *path* through all nodes is better for everyone!
- **Efficient** solutions maximize the sum of all players' utility

*A game with 4 players/nodes*



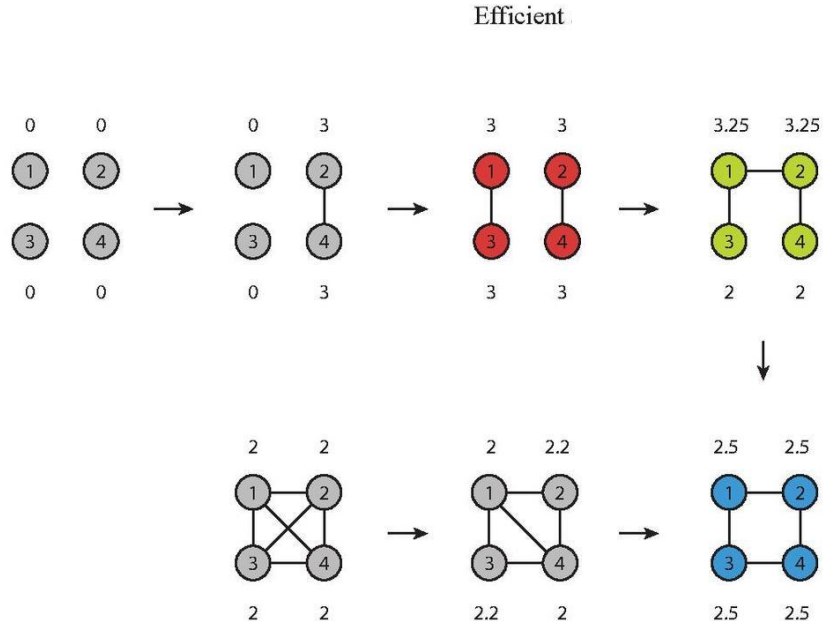
$$b(1) < c$$

Efficient!  
(Given  $c \leq b(1) + b(2)$ )

# Solution concepts in network games

Other solutions (besides NE)  
can also be desired:

- **Efficient strategy**:  
maximizes the sum of  
players' utility



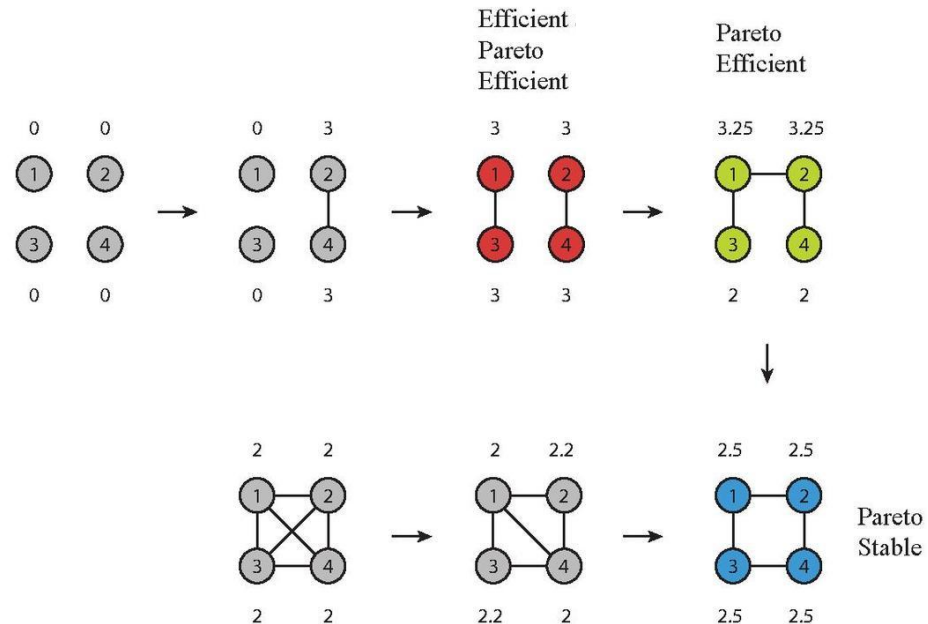
Source:

[https://en.wikipedia.org/wiki/Strategic\\_Network\\_Formation](https://en.wikipedia.org/wiki/Strategic_Network_Formation)

# Solution concepts in network games

Other solutions (besides NE)  
can also be desired:

- **Efficient strategy**:  
maximizes the sum of  
players' utility
- **Pareto optimal** (or  
pareto efficient): network  
such that there **is no  
other network  $g'$**  where:  
 $u_i(g') \geq u_i(g)$  for all  $i$  and  
 $u_i(g') > u_i(g)$  for at least 1  $i$ .



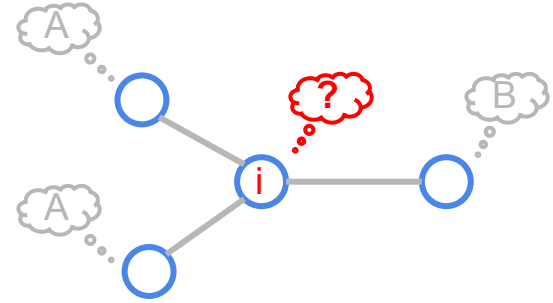
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# Opinion Dynamics via “the Majority Game”

## Majority Game:

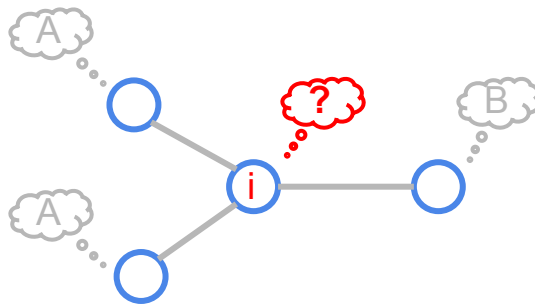
- $N$  players/nodes
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- The set of neighbors of player  $i$  who believe A:  $\mathbf{N}_i(\mathbf{A})$



# Opinion Dynamics via “the Majority Game”

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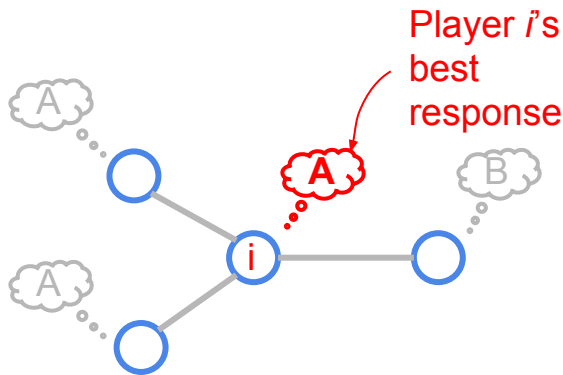
- $N$  players/nodes
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- The set of neighbors of player  $i$  who believe A:  $\mathbf{N}_i(\mathbf{A})$
- **Majority utility function:**
  - If  $|\mathbf{N}_i(\mathbf{A})| > \frac{1}{2} * \deg(i)$ ,  $u_i(A) > u_i(B)$
  - Otherwise,  $u_i(B) > u_i(A)$



# Opinion Dynamics via “the Majority Game”

## Majority Game:

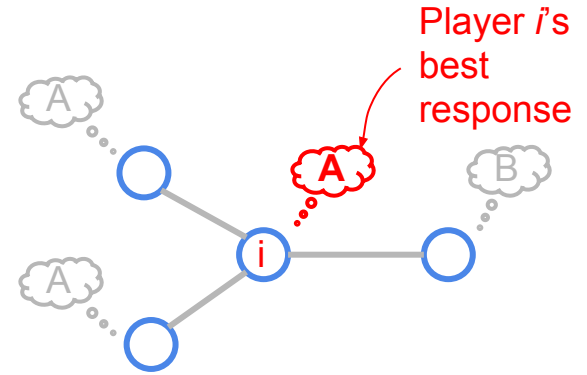
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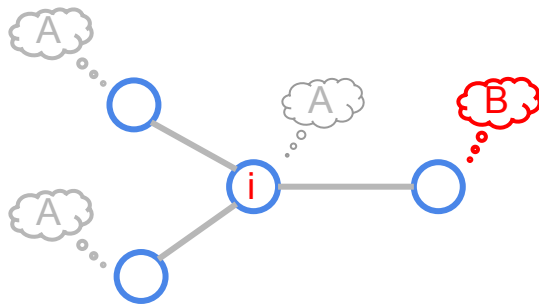
# NE in “the Majority Game”

- Nash equilibrium: When no player wishes to change their belief, given the other players' beliefs.



# NE in “the Majority Game”

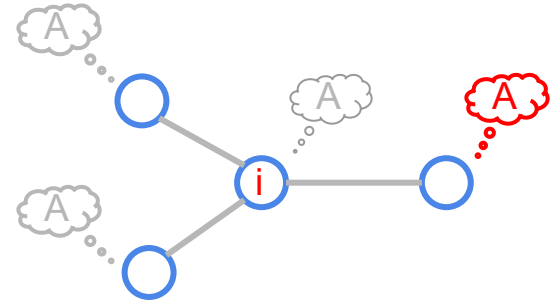
- Nash equilibrium: When no player wishes to change their belief, given the other players' beliefs.



This is not a Nash equilibrium!

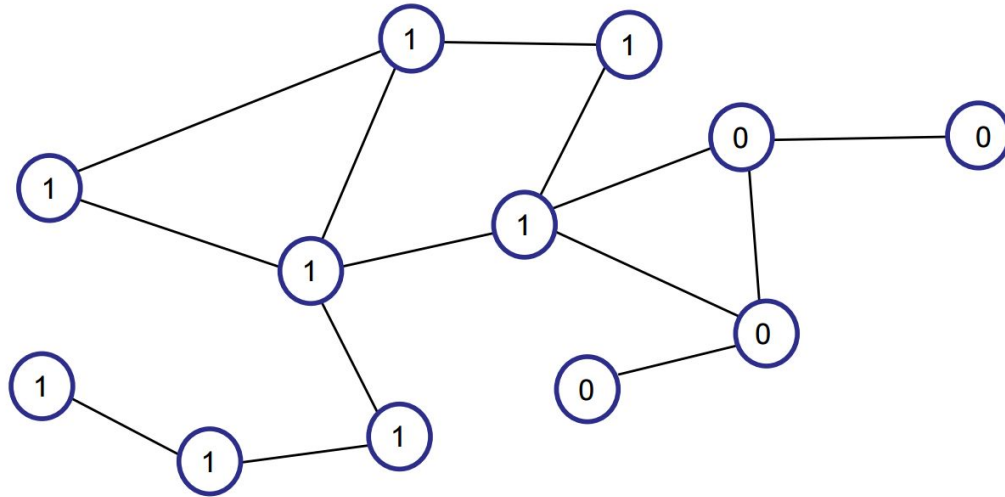
# NE in “the Majority Game”

- Nash equilibrium: When no player wishes to change their belief, given the other players' beliefs.
- Generally, every player choosing A and every player choosing B is a NE.
  - But there can be others...



This IS a Nash equilibrium!

# NE in “the Majority Game”



The initial configuration matters: flipping everyone's opinion is also stable!

(source: Jackson, M., **Games on Networks**, Handbook of Game Theory, Vol. 4, 2014.)

# Extensions of “the Majority Game”

- **Coordination games:** Highest utility is gained by coordinating with neighbors; miscoordination incurs a cost. What thresholds and

	A	B
A	$(b,b)$	$(-c,0)$
B	$(0,-c)$	$(0,0)$

- **Stability analysis of equilibria:** Which equilibria are most stable to a player “changing their mind”?
- **Resources:**
  - Jackson, M.O. and Zenou, Y., 2014. **Games on networks.** *Handbook of game theory*,.
  - Kearns, M., 2007. **Graphical Games.** *Algorithmic Game Theory*.

# Final notes

- Many network-based games can be modeled as evolutionary processes:
  - **Network formation:** Start with an initial network, and add/remove edges until no player wishes to deviate (NE found).
  - **Opinion dynamics:** Seed beliefs randomly (or empirically), and update players' beliefs until no player wishes to change their belief (NE found).
- *Algorithmic Game Theory*, Noam Nisan, Tim Roughgarden et. al
- *Social and Economic Networks*, Matthew Jackson.