ECS 253 / MAE 253, Network Theory Spring 2016 Problem Set # 1, Solutions

Problem 1: The Erdős-Rényi random graph – analyzing the phase transition

(a) For each j that not equals i: (i) The probability of that i is not linked to j is 1-p. (ii) The probability that j is not a part of giant component is u and the probability that i is linked to j is p. Since these two events are independent, the probability that i is linked to j and j is not a part of giant component is pu. (i) and (ii) are mutually exclusive therefore the probability of (i) or (ii) is (1-p+pu). There are N-1 choice for j and all of them must satisfy (i) or (ii) and they are independent too. Thus:

$$u = (1 - p + pu)^{N-1}. (1)$$

(b)

$$\ln u = \ln(1 - p + pu)^{N-1}$$

$$= (N-1)\ln(1 - p + pu)$$

$$= (N-1)\ln(1 - p(1-u))$$

$$= (N-1)\ln(1 - \frac{\langle k \rangle}{N-1}(1-u))$$

$$\approx (N-1)\frac{-\langle k \rangle}{N-1}(1-u)$$

$$= -\langle k \rangle(1-u).$$
(2)

(c)

$$u = e^{-\langle k \rangle S}$$

$$(1 - u) = S = 1 - e^{-\langle k \rangle S}$$
(3)

- (d) See figure below.
- (e) Take derivative of both side w.r.t S we have:

$$1 = \langle k \rangle e^{-\langle k \rangle S} \tag{4}$$

Plug S = 0 into it we get:

$$k_c = \langle k \rangle = 1 \tag{5}$$

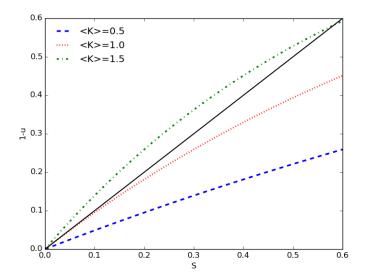
(f)

$$I_G = N(1-p)^{N_G}$$

$$\approx N(1-p)^N$$

$$= N(1 - \frac{Np}{N})^N$$

$$\approx Ne^{-Np}$$
(6)



Setting $I_G = 1$ gives:

$$1 = Ne^{-Np}$$

Therefore:

$$e - Np = \frac{1}{N}$$
$$p = \frac{\ln N}{N}$$

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Problem 2: The Erdős-Rényi random graph – cluster size distribution

(a) During the edge adding process, c_k can be changed in two ways. The first one is joining two smaller components and form a new component of size k. The other one is adding edges to a component of size k and this will decrease the number of components of size k.

$$\frac{dc_k}{dt} = \left(\sum_{i+j=k} \frac{1}{2} (ic_i)(jc_j)\right) - kc_k$$

(b) Let k = 1, we have:

$$\frac{dc_1}{dt} = -c_1$$

Solve this differential equation we have:

$$c_1 = e^{-t}$$

Similarly, we can get:

$$c_2 = \frac{1}{2}te^{-2t}$$

$$c_3 = \frac{1}{2}t^2e^{-3t}$$

(c) We probably would guess that:

$$c_k = \alpha_k t^{k-1} e^{-kt}$$

(d) Using Stirling's formula to approximate factorial:

$$k! \approx \sqrt{2\pi k} \left(\frac{k}{e}\right)^k$$

When t = 1:

$$c_{k} = \frac{k^{k-2}}{k!} e^{-k}$$

$$\approx \frac{1}{\sqrt{2\pi k}} k^{-k} e^{k} k^{k-2} e^{-k}$$

$$= \frac{1}{\sqrt{2\pi k}} k^{-2}$$

$$= \frac{1}{\sqrt{2\pi}} k^{-2.5}$$
(7)