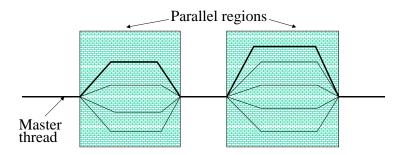
Lecture 22 – OpenMP Shared-Memory Parallel Computing

- Now let's continue our discussion on how we can use shared-memory systems to perform parallel computing
- Remember that in shared-memory systems, all of the processors see all of the data

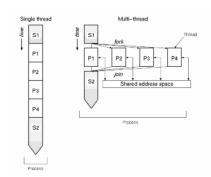
Execution Model

- Master thread spawns a team of threads as needed
 - Concept of threads is also used for GPU parallel computing
- Parallelism is added incrementally



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Execution Model



- OpenMP programs start as single thread
- In parallel regions, additional threads are created
- · Master thread is part of the team
- Outside parallel regions, additional threads go away (or sleep)

Directives or pragmas

Fortran (fixed format)

!\$omp ... c\$omp ... *\$omp ...

If the character in column 6 is different from a space or a 0 it is a continuation line

- Fortran (free form)
 - A line that begins with !\$omp is an OpenMP directive
 - A directive that needs to be continued on the next line:!\$OMP &
- C:

#pragma omp

OpenMP Constructs

- OpenMP's constructs (directives) fall into 5 categories:
 - Parallel regions
 - Worksharing
 - Synchronization
 - Data environment
 - Runtime functions/environment variables

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Parallel regions

Threads are created in OpenMP with the omp parallel directive

FORTRAN

!\$OMP PARALLEL

Code to be executed by each thread
!\$OMP END PARALLEL

<u>C</u>

#pragma omp parallel {Code to be executed by each thread}

Parallel Regions

- Dynamic mode (the default):
 - The number of threads used can vary from one parallel region to another
 - Setting the number of threads only sets the maximum number of threads
- · Static mode:
 - The number of threads is fixed and controlled by the programmer:
 - Use library routine: call omp_set_num_threads(16)
 - Environmental variable: setenv OMP_NUM_THREADS 16

OpenMP Constructs

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C Example: Work Sharing: for

Serial for $(i=0;i< N,i++){a[i]=b[i]+c[i];}$

}

OpenMP

Parallel region

#pragma omp parallel
{ int myid,i,nthreads,mystart,myend;
 myid=omp_get_thread_num();
 nthreads=omp_get_num_threads();
 mystart=myid*N/nthreads;
 myend=(Myid+1)*N/nthreads;
 for (i=mystart;i<myend,i++){ a[i]=b[i]+c[i];}</pre>

OpenMP Parallel region and work-sharing #pragma omp parallel
#pragma omp for schedule(static)
for (i=0;i<N,i++){ a[i]=b[i]+c[i];}</pre>

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Fortran Example: Work Sharing: for

Serial

do i = 0,N a(i) = b(i)+c(i) enddo

OpenMP Parallel region !\$omp parallel

myid= omp_get_thread_num()

nthreads= omp_get_num_threads()
mystart = myid*N/nthreads

myend = (myid+1)*N/nthreads

do I = mystart, myend a(i) = b(i)+c(i) enddo

(No dependencies within loop)

!\$omp end parallel

OpenMP
Parallel region
and work-sharing

!\$omp parallel !\$omp do schedule(static) do i =0,N a(i) = b(i)+c(i)

enddo !\$omp end schedule !\$omp end parallel These directives could be combined to: !\$omp parallel do schedule(static)

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Notice that operation within

loop is index-independent!!

Allows the programmer to

break up loop into nthreads

The loop is automatically broken up into nthreads as evenly as possible

Work Sharing: schedule

• Schedule keyword specifies how loop is split:

- schedule(static,[chunk]):
 - Each thread gets a block of size chunk
- schedule(dynamic,[chunk]):
 - Each thread gets a block of size chunk from a queue until all iterations are scheduled
- schedule(guided,[chunk]):
 - Dynamic schedule that shrinks down to size "chunk" as the calculation proceeds.
- schedule(runtime):
 - Schedule and chunk size taken from OMP_SCHEDULE environment variable

Work Sharing: sections

• The "sections" construct is used to give different workload to different threads:

 Assuming that subroutines (tasks) sub_x, sub_y, sub_z, etc. are independent of each other

!\$OMP PARALLEL !\$OMP SECTIONS Call sub_x(...) !\$OMP SECTION Call sub_y(...) !\$OMP SECTION Call sub_z(...) !\$OMP END SECTIONS !\$OMP END PARALLEL

OpenMP Constructs

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Synchronization

- There are different OpenMP constructs that allow the processors to synchronize:
 - Mutual exclusion
 - Atomic update
 - Barrier synchronization
 - Master section

Mutual Exclusion

· CRITICAL section:

- Can be named, but the name should not conflict with subroutine or module
- Used to ensure that only one thread (first to arrive) is within a "critical" section (i.e.serial execution)
- Can not jump into or out of critical sections

!SOMP PARALLEL

•••••

......

!\$OMP END PARALLEL

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Atomic Update

• Optimization of critical directive for a certain case

- Operation should not be broken into sub-parts (single-thread operation)
- Update to a single memory location should be performed as a single-threaded operation

!\$OMP PARALLEL

........

!SOMP ATOMIC

a(i)=a(i)+alocal

!\$OMP END ATOMIC

.....

!\$OMP END PARALLEL

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Barrier Synchronization

BARRIER directive

- Threads wait until all the threads reach this point:
- An implicit barrier exists at the end of each parallel region

!\$OMP PARALLEL

.....

!\$OMP BARRIER

•••••

!\$OMP END PARALLEL

Master section

Only the master thread executes the section

- Rest of the team skip execution
- No barrier at the end of the master section
- A specialization of the "single" directive

!SOMP PARALLEL

.....

!\$OMP MASTER

print *,"I am the master"

!SOMP END MASTER

......

!\$OMP END PARALLEL

OpenMP Constructs

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• OpenMP's constructs fall in 5 categories:

- Parallel regions
- Synchronization We discussed these in the last lecture.
- Worksharing
- Data environment
- Runtime functions/environment variables

Default Storage Attributes

- Global variables are SHARED among threads
 - FORTRAN:
 - SAVF variables
 - MODULE variables
 - C: File scope variables, static
- Stack variables in sub-programs called from parallel regions are PRIVATE

Default Storage Attributes

subroutine work program sort use input use input real dummy(10) integer index(10) integer count call input **!SOMP PARALLEL** save count call work **!SOMP END PARALLEL** module input print *,index(1) dimension a(10) end module input

a(10), index(10), count are shared dummy(10) is private

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Storage Attributes

- Storage attributes can be changed using the following clauses of an OMP directive:
 - SHARED
 - PRIVATE
 - FIRSTPRIVATE
 - THREADPRIVATE
- The value of private data inside a parallel loop can be transmitted to a global value outside the loop with:
 - LASTPRIVATE
- The default status can be modified with:
 - DEFAULT (PRIVATE | SHARED | NONE)

PRIVATE Clause

- PRIVATE(var) creates a local copy of var for each thread
 - The value is initialized
 - Private copy is not storage associated with the original

```
program wrong
is=0

!$OMP PARALLEL DO PRIVATE(IS)
do j=1,1000
is=is+j !is is not initialized
end do

!$OMP END PARALLEL DO
print *,is !regardless of initialization, "is" is undefined
stop
end
```

FIRSTPRIVATE Clause

 FIRSTPRIVATE(var) creates a local copy of var for each thread with the corresponding value from the master thread

```
program wrong2
is=0

!$OMP PARALLEL DO FIRSTPRIVATE(IS)
do j=1,1000
is=is+j !is is now initialized to 0
end do

!$OMP END PARALLEL DO
print *,is !regardless of initialization, "is" is undefined because it
! was not lastprivate
stop
end
```

LASTPRIVATE Clause

 LASTPRIVATE(var) passes the value of a private from the last iteration to a global variable

```
program wrong3
is=0
!$OMP PARALLEL DO FIRSTPRIVATE(IS) LASTPRIVATE(IS)
do j=1,1000
is=is+j !is is now initialized to 0
end do
!$OMP END PARALLEL DO
print *,is
stop
end
```

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Example

Serial output program wrong3 Is=0 ~> a.out !\$OMP PARALLEL DO FIRSTPRIVATE(IS) 500500 **!\$OMP LASTPRIVATE(IS)** do j=1,1000Parallel output: is=is+i !is starts at 0 ~> setenv OMP NUM THREADS 2 end do ~> a.out **!SOMP END PARALLEL DO** 375250 print *,is stop ~> setenv OMP_NUM_THREADS 4 end ~> a.out 218875

Why do we get different answers?

Answer: Because 'is' is the final value for each thread, not the global sum. A reduction is needed!

Note: The loop will be broken up into NUM_THREADS even if it is greater than the number of available CPUs (in which case, a warning will be issued). ²⁷

REDUCTION Clause

- REDUCTION(op: list) :local copies are reduced into a single global copy at the end of the construct
 - The variables in "list" must be shared.
 - Inside a parallel or worksharing construct a local variable is made and properly initialized.
 - "op" includes +, *, -, .AND., .OR., .EQV., .NEQV., or the intrinsics MAX, MIN, IAND, IOR, or IEOR

Example

program correct3
is=0
call omp_set_num_threads(2)
!\$OMP PARALLEL DO REDUCTION(+:IS)
do j=1,1000
is=is+j !is is now initialized to 0
end do
!\$OMP END PARALLEL DO
print *,is
stop
end

Parallel Output:

./correct3 500500 FORTRAN STOP

"Variables that appear in a REDUCTION clause must be SHARED in the enclosing context. A private copy of each variable in *list* is created for each thread as if the PRIVATE clause had been used. The private copy is initialized according to the operator.

At the end of the REDUCTION, the shared variable is updated to reflect the result of combining the original value of the (shared) reduction variable with the final value of each of the private copies using the operator specified."

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THREADPRIVATE Clause

- Make global data private to a thread:
 - FORTRAN: MODULE blocks
 - C: file scope and static variables
- Each thread has its own copy of the module block
 - Persistent across multiple parallel regions provided that the number of threads remains the same

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OpenMP's Constructs

- OpenMP's constructs fall in 5 categories:
 - Parallel regions
 - Synchronization
 - Worksharing
 - Data environment
 - Runtime functions/environment variables
 - · We'll discuss this during the next lecture

Project 5

- Now that you have "completed" Projects 1 4 (hopefully), you are ready to parallelize a multi-block code using MPI
- Take your multi-block code that was developed under Project 3 and parallelize it using MPI across the node/CPU partitions created in Project 4

Project 5

• Due Monday, December 9 (Start on this NOW!)

- Overall description of heat conduction problem including governing equations, discretization scheme (recursion formulas), domain decomposition, processor mapping, parallel processing strategy, etc.
- Demonstration that the parallel heat conduction code works on Wopr with:
 - For 501x501 grid using 10x10 decomposition with 4 processors, figures of
 - Grid
 - Converged solution
 - Convergence history (log(residual) vs iteration) comparison with single-block and multi-block serial codes. Plot 3 code results on single plot.
 - Plots of computational speed-up and efficiency vs number of processors for up to 10 processors for 501x501 grid using 10x10 decomposition
 - Show along with comparison with ideal (linear) and optimal (based on your load balance numbers in Project 4)
 - » Make sure you are the only one running on freed nodes when you are doing timing studies by using the qsub queuing system. Use MPI_Wtime for wall clock times in Project-5.
- Listing of your parallel heat conduction code

