

System Programming

Assembly

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Topics

Intel Assembly

- x86 Processors
- Instructions
- Directives
- System Calls

Assembly and C

- Subroutines
- Calling Conventions
- C from Assembly
- Assembly from C

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x86 Processors

- ▶ Intel family of processors: x86 (32 bit), x64 (64 bit)
- ▶ very similar from the programming standpoint
- ▶ 8086: 16 bit processor, real mode
- ▶ 80386: 32 bit processor, protected mode (virtual memory)

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Segments

- ▶ programs are divided into **segments**
- ▶ **code** segment: instructions
- ▶ **data** segment: initialized data
- ▶ **bss** segment: uninitialized data
- ▶ **stack** segment

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8086 Registers

- ▶ 4 general purpose data registers
- ▶ 2 index registers
- ▶ 2 pointer registers
- ▶ 4 segment registers
- ▶ 2 control registers

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Data Registers

- ▶ AX: accumulator register
- ▶ BX: base register
 - ▶ used to address data in memory
- ▶ CX: counter register
 - ▶ used as repetition counter in loop operations
- ▶ DX: data register
 - ▶ used in multiplication and division operations
- ▶ high and low halves can be accessed as 8-bit registers:
AH-AL, BH-BL, CH-CL, DH-DL

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Index and Pointer Registers

- ▶ index registers:
 - ▶ DI: data index
 - ▶ SI: stack index
 - ▶ they can be used like general purpose registers
- ▶ pointer registers:
 - ▶ BP: base pointer
 - ▶ SP: stack pointer

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Segment Registers

- ▶ CS: code segment register
- ▶ DS: data segment register
- ▶ SS: stack segment register
- ▶ ES: extra segment register

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Control Registers

- ▶ IP: instruction pointer
 - ▶ CS + IP: address of next instruction
- ▶ FLAGS: status conditions
 - ▶ ZF (zero), OF (overflow), SF (sign), CF (carry), PF (parity)

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80386

- ▶ in 80386, registers are extended to 32 bits:
EAX EBX ECX EDX ESI EDI EBP ESP
EIP
- ▶ AX, BX, ..., BP, SP are still valid (lower 16 bits)
- ▶ AH, AL, ..., DH, DL are still valid

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Operand Types

- ▶ register
- ▶ memory: offset from beginning of segment
- ▶ immediate: listed in the instruction itself
- ▶ implied: not explicitly specified

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Basic Instructions

<code>mov</code>	<code>dest, src</code>	move src to dest
<code>add</code>	<code>dest, src</code>	add src to dest
<code>adc</code>	<code>dest, src</code>	add src to dest with carry
<code>sub</code>	<code>dest, src</code>	subtract src from dest
<code>sbb</code>	<code>dest, src</code>	subtract src from dest with borrow
<code>inc</code>	<code>dest</code>	increment dest
<code>dec</code>	<code>dest</code>	decrement dest
<code>mul</code>	<code>src</code>	multiply eax with src, result in edx:eax
<code>div</code>	<code>src</code>	divide edx:eax by src, result in eax and edx

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Bitwise Instructions

<code>not</code>	dest	bitwise not (one's complement)
<code>and</code>	dest, src	bitwise and
<code>or</code>	dest, src	bitwise or
<code>xor</code>	dest, src	bitwise xor
<code>neg</code>	dest	negate (two's complement)
<code>shl</code>	dest, amount	logical shift left
<code>shr</code>	dest, amount	logical shift right
<code>asl</code>	dest, amount	arithmetic shift left
<code>asr</code>	dest, amount	arithmetic shift right
<code>rol</code>	dest, amount	rotate left
<code>ror</code>	dest, amount	rotate right
<code>rcl</code>	dest, amount	rotate left with carry
<code>rcr</code>	dest, amount	rotate right with carry

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Branching Instructions

<code>jmp</code>	unconditional
<code>jz</code>	if ZF is set
<code>jnz</code>	if ZF is unset
<code>jo</code>	if OF is set
<code>jno</code>	if OF is unset
<code>js</code>	if SF is set
<code>jns</code>	if SF is unset
<code>jc</code>	if CF is set
<code>jnc</code>	if CF is unset
<code>jp</code>	if PF is set
<code>jnp</code>	if PF is unset

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Branching Instructions

- `cmp vleft, vright`: compare vleft and vright

condition	signed	unsigned
vleft = vright	<code>je</code>	<code>je</code>
vleft ≠ vright	<code>jne</code>	<code>jne</code>
vleft < vright	<code>jl</code>	<code>jb</code>
vleft <= vright	<code>jnl</code>	<code>jnb</code>
vleft ≤ vright	<code>jle</code>	<code>jbe</code>
vleft ≥ vright	<code>jnle</code>	<code>jnbe</code>
vleft > vright	<code>jg</code>	<code>ja</code>
vleft >= vright	<code>jng</code>	<code>jna</code>
vleft ≥ vright	<code>jge</code>	<code>jae</code>
vleft ≠ vright	<code>jnge</code>	<code>jnae</code>

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Directives

- assembler needs extra info: *directives*
- not part of the instruction set
- labels: mark points in code and data
 - entry labels have to be marked `global`
- segments
- data definition
- named constants: `equ`
 - no memory allocated

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Code Template

```
segment .data
; initialized data definitions

segment .bss
; uninitialized data definitions

segment .text
global _start

_start:
; entry point
```

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Data Definition

type	initialized	uninitialized
byte	db	resb
word	dw	resw
dword	dd	resd
qword	dq	resq
tword	dt	rest

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Data Definition Examples

```
L1 db 0
L2 dw 1000
L3 dd 1A92h
L4 db 0, 1, 2, 3
L5 db "w", "o", "r", "d", 0
L6 db 'word', 0
L7 times 100 db 0
L8 resb 1
L9 resw 100
```

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Dereferencing

► plain label:
address of memory

Example

```
mov eax, L1
```

► label in brackets:
contents of memory

Example

```
mov eax, [L1]
```

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System Calls

- ▶ system calls are implemented using software interrupt 80h

system call setup

eax ← system call number

ebx ← first argument

ecx ← second argument

edx ← third argument

int 80h

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System Call Examples

- ▶ exit system call number: 1
- ▶ arg. 1: return status
 - ▶ 0: success, 1: failure
- ▶ read system call number: 3
- ▶ arg. 1: input descriptor
 - ▶ 0: stdin, 1: stdout, 2: stderr
- ▶ arg. 2: start of input buffer
- ▶ arg. 3: length of input
- ▶ write system call number: 4
- ▶ arg. 1: output descriptor
 - ▶ 0: stdin, 1: stdout, 2: stderr
- ▶ arg. 2: start of output buffer
- ▶ arg. 3: length of output

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Example: Hello, world!

```
segment .data
msg db "Hello, world!", 10
len equ 14
```

```
segment .text
global _start

_start:
    mov eax, 4
    mov ebx, 1
    mov ecx, msg
    mov edx, len
    int 80h

    mov eax, 1
    mov ebx, 0
    int 80h
```

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References

Required Reading: Carter

- ▶ Chapter 1: Introduction
 - ▶ 1.2. Computer Organization
 - ▶ 1.3. Assembly Language

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Stack

- ▶ the stack is accessed in 4-byte units

push operand

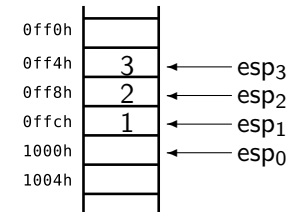
- ▶ subtract 4 from esp
- ▶ store operand to address [esp]

pop register

- ▶ store operand at address [esp] to register
- ▶ add 4 to esp

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Stack Example



```
push dword 1
push dword 2
push dword 3
pop  eax
pop  ebx
pop  ecx
```

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Subroutine Call

call target

- ▶ push address of next instruction
- ▶ jump to target

ret

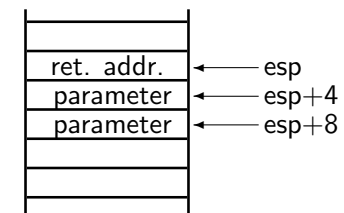
- ▶ pop return address
- ▶ jump to return address

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Stack Parameters

- ▶ called subroutine does not pop parameters
- ▶ accesses parameters on the stack

stack layout

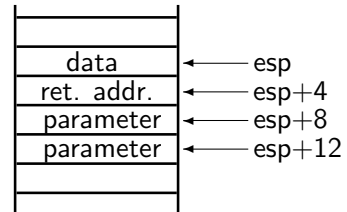


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Accessing Parameters

- ▶ offsets from esp may change

example: after a push



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Accessing Parameters

- ▶ use ebp (frame pointer)

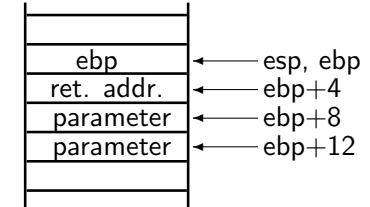
subroutine template

```
push ebp
mov  ebp, esp

...

pop  ebp
ret
```

stack layout



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Example: Factorial

```
segment .bss
f    resd 1

segment .text

fact:
    push ebp
    mov  ebp, esp

    mov  dword [f], 1
    mov  ecx, [ebp+8]
```

```
back:
    mov  eax, [f]
    mul  ecx
    mov  [f], eax
    dec  ecx
    cmp  ecx, 1
    jne  back

    pop  ebp
    ret
```

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Example: Calling Factorial

```
segment .data
k    dd  5

segment .bss
f    resd 1

segment .text
global _start

fact:
    ...
```

```
_start:
    push ebp
    mov  ebp, esp

    push dword [k]
    call fact
    add  esp, 4

    pop  ebp
    ret
```

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Calling Conventions

- ▶ how will parameters be passed?
- ▶ if using stack:
 - ▶ in what order will the parameters be pushed?
 - ▶ who will remove parameters from the stack?
- ▶ how will the result be returned?
- ▶ which registers should remain unchanged?

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C Calling Conventions

- ▶ parameters are passed via the stack
 - ▶ caller pushes parameters in reverse order
 - ▶ caller removes parameters from the stack
- ▶ result is returned over eax
- ▶ ebx, esi, edi, ebp, cs, ds, ss, es should remain unchanged

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Calling C from Assembly

- ▶ to call a C function from Assembly:
- ▶ declare function as `extern`
- ▶ push arguments in reverse order
- ▶ call function
- ▶ adjust esp

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Example: printf

```
segment .data
k    dd    5
intf db    "%d", 10, 0
```

```
segment .bss
f    resd 1
```

```
segment .text
global main
extern printf
```

```
fact:
...
```

```
main:
```

```
...
```

```
push dword [k]
call fact
add esp, 4
```

```
push dword [f]
push intf
call printf
add esp, 8
```

```
...
```

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C Variables

- ▶ global: in fixed memory locations
- ▶ static: same as global, only scope is different
- ▶ automatic: on stack
- ▶ register: in a register (if possible)
- ▶ volatile: do not optimize

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Automatic Variables

- ▶ allocation is done by subtracting from esp

subroutine template

```
push ebp
mov  ebp, esp
sub  esp, N_BYTES

...

mov  esp, ebp
pop  ebp
ret
```

stack layout

var. 2	← esp, ebp-8
var. 1	← ebp-4
ebp	← ebp
ret. addr.	← ebp+4
param. 1	← ebp+8
param. 2	← ebp+12

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Example: Factorial (C)

```
int y;

void fact(int k)
{
    register int i;

    y = 1;
    for (i = k; i > 1; i--)
        y = y * i;
}
```

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Example: Factorial (C)

```
int fact(int k)
{
    int y;
    register int i;

    y = 1;
    for (i = k; i > 1; i--)
        y = y * i;
    return y;
}
```

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Example: Factorial

```
segment .text
global fact
```

```
fact:
```

```
    push ebp
    mov  ebp, esp
    sub  esp, 4
```

```
    mov  dword [ebp-4], 1
    mov  ecx, [ebp+8]
```

```
back:
```

```
    mov  eax, [ebp-4]
    mul  ecx
    mov  [ebp-4], eax
    dec  ecx
    cmp  ecx, 1
    jne  back
```

```
    mov  eax, [ebp-4]
    mov  esp, ebp
    pop  ebp
    ret
```

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Calling Assembly from C

- ▶ to call an Assembly function from C:
- ▶ in Assembly file: declare function as `global`
- ▶ in C file: declare the prototype

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Example: Factorial

```
int fact(int k);

int main(void)
{
    int x, y;

    ...
    y = fact(x);
    ...
}
```

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References

Required Reading: Carter

- ▶ Chapter 4: `Subprograms`

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