

Algorithms and Data Structures 2 CS 1501



Fall 2022

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Announcements

- Upcoming Deadlines
 - Homework 7: next Friday @ 11:59 pm
 - Lab 6: Monday 10/31 @ 11:59 pm
 - Nothing due this week
- Midterm Exam
 - Wednesday 10/19 (MW Section) and Thursday 10/20 (TuTh Section)
 - in-person, closed-book
- Weekly Live QA Session on Piazza
 - Friday 4:30-5:30 pm

Previous lecture

- ADT Priority Queue (PQ)
 - Heap implementation
- Heap Sort
- Indexable PQ

This Lecture

- Muddiest Points
- Introduction to ADT Graph

- Q: what is going to be on the midterm?
- Up to and including material covered on Monday 10/10 and Tuesday 10/11
- Please check the study guide on Canvas
 - practice test on GradeScope
 - old exam
 - will try to post the answer as soon as possible

- Q: I'm confused about entropy and the equations for that. How is it useful?
- It helps us determine the "information content" in a file.
- In lossless compression, a file cannot be compressed to less than its entropy

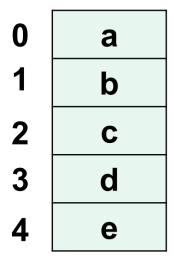
- Q: How do we determine the entropy of a given file we're trying to compress?
- Given that
 - the file has *n* total characters and *K* unique characters,
 - *f*(*c*) is the frequency of character *c* in the file
- Shannon's entropy: H(file) = -1 * $\sum_{c=0}^{K} \frac{f(c)}{n} log_2(\frac{f(c)}{n})$ bits/character
- Underlying assumption:
 - the file has been generated by a source that produces independent characters
 - Huffman Compression is optimal under that assumption
 - This assumption may be wrong though!
 - repeated long strings → use LZW
 - long sequences of identical characters → use RLE (Run Length Encoding)
 - We may need to try different compression algorithms on the file

- Q: What might be some examples in where we might pick one compression type over the other?
- Depends on the structure of the input file
- long strings of identical values
 - → RLE
- repeated long strings
 - → LZW
- frequently occurring values
 - → Huffman
- few different characters
 - → fixed-length codewords with <8 bits

- Q: Does entropy play a role in the implementation of the compression algorithm itself? Or is it only used for selecting the best algorithm?
- Some compression algorithms attempt to reach Shannon's Entropy lower bound on the file size
 - e.g., Huffman Encoding and Arithmetic Encoding

Q: Not exactly clear on move to front encoding

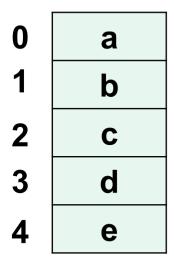






Q: Not exactly clear on move to front encoding

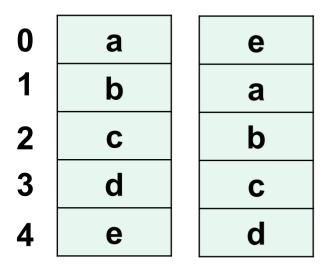






Q: Not exactly clear on move to front encoding





Output:

4

Q: Not exactly clear on move to front encoding



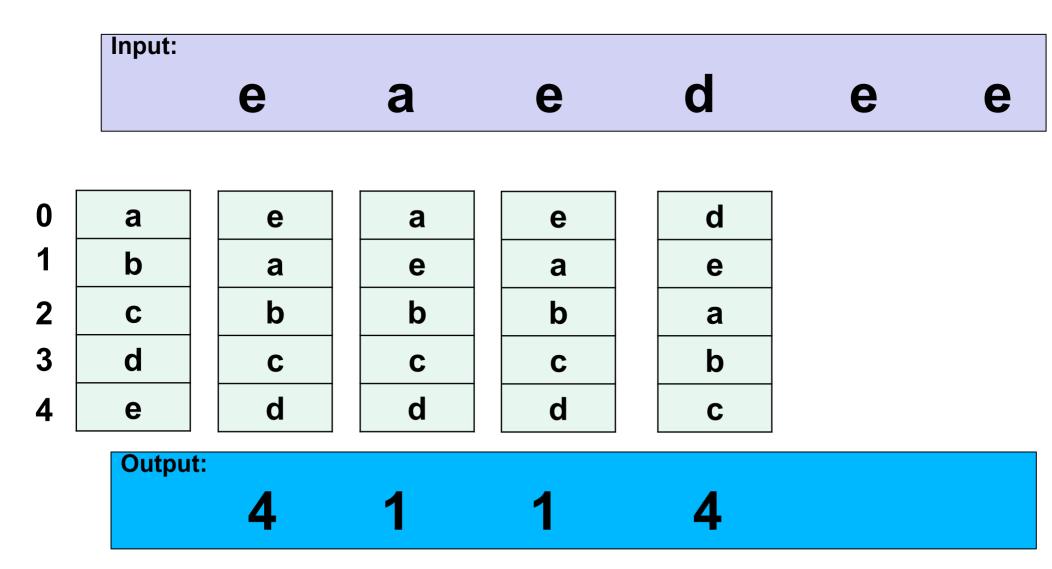
0 a a e b a 2 b b C 3 d 4 d d e

Output: 4 1

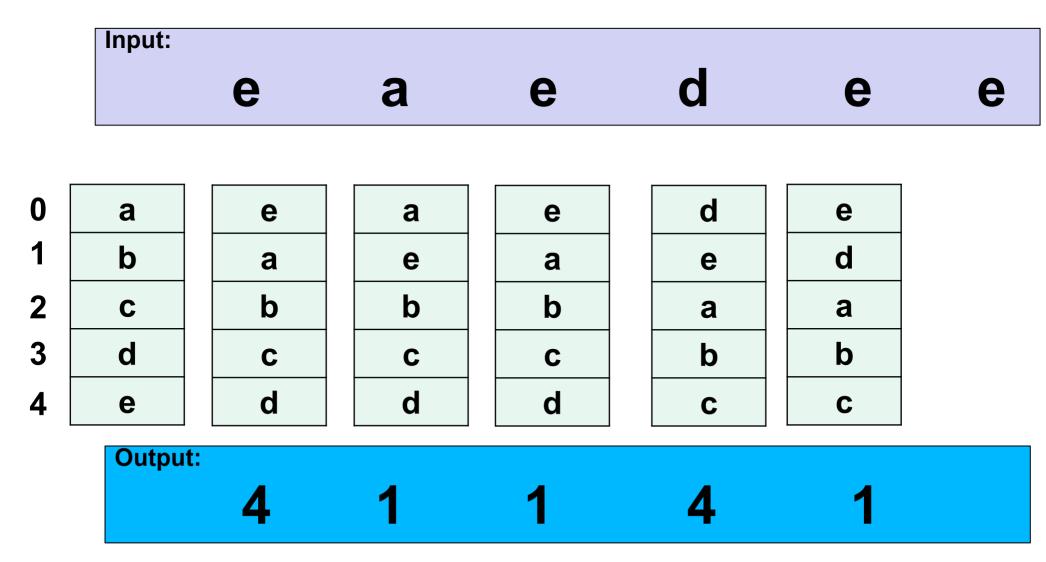
Q: Not exactly clear on move to front encoding

Input: 0 a a e e b a a 2 b b b C 3 d 4 d d d e **Output:**

Q: Not exactly clear on move to front encoding



Q: Not exactly clear on move to front encoding

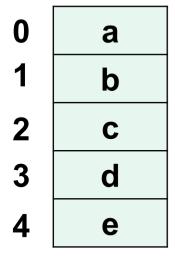


Q: Not exactly clear on move to front encoding

Input: 0 d a e a e e e b d d a e a 2 b b b C a a a 3 d b b b 4 d d d e C C C **Output:**

- Q: Not exactly clear on move to front encoding
- Decoding

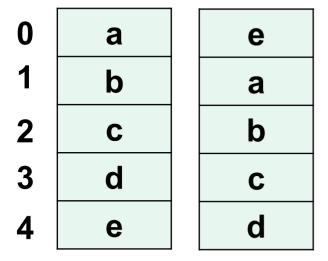






- Q: Not exactly clear on move to front encoding
- **Decoding**

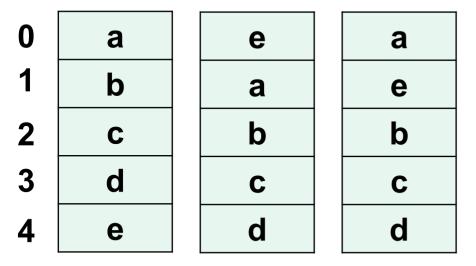






- Q: Not exactly clear on move to front encoding
- **Decoding**

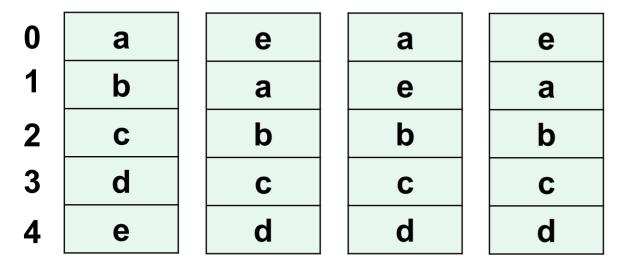






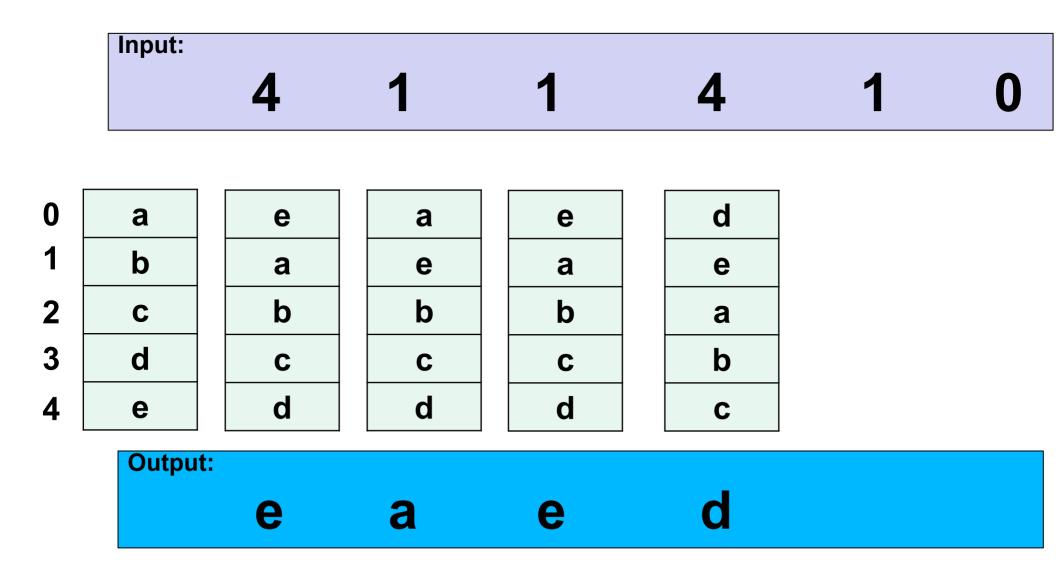
- Q: Not exactly clear on move to front encoding
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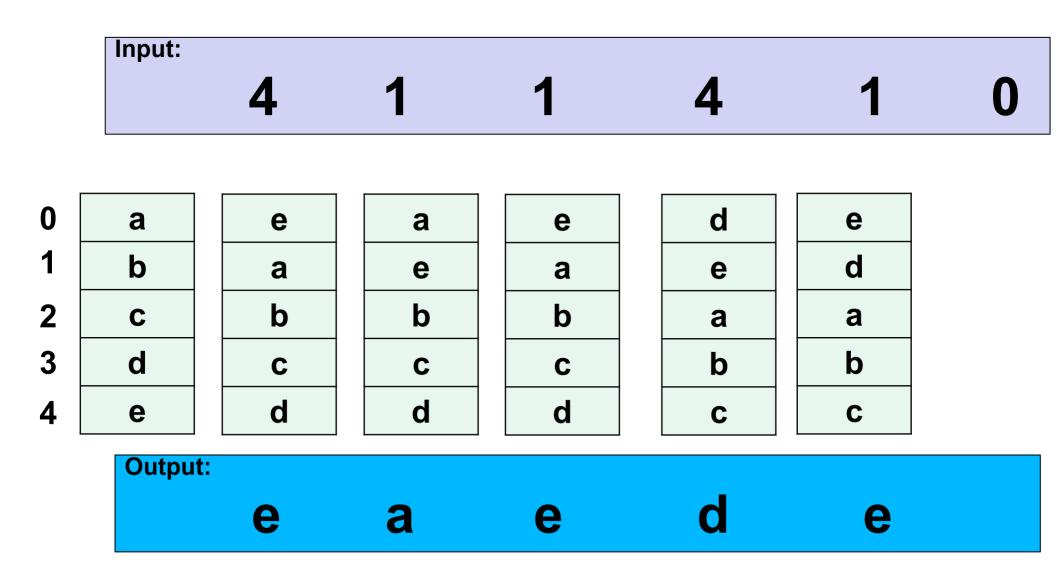




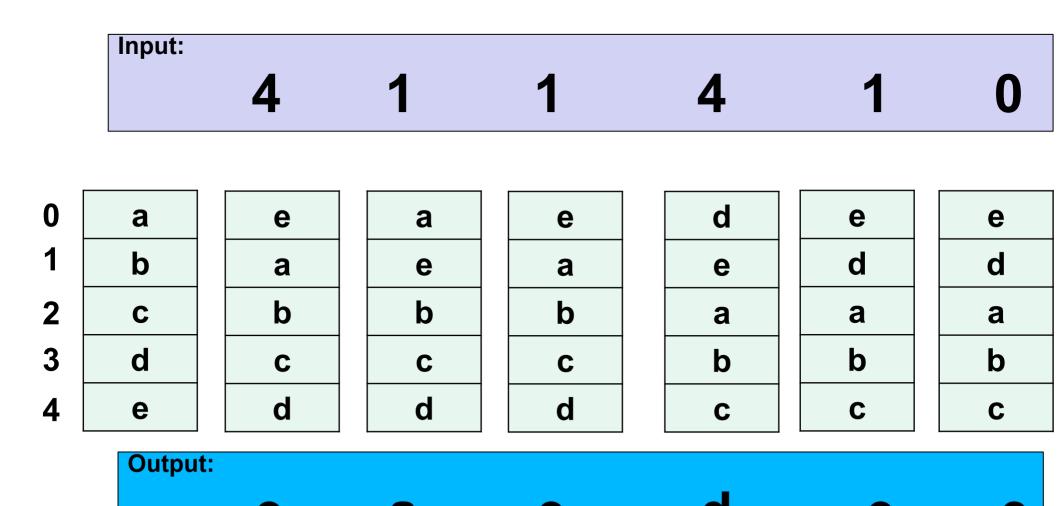
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- Q: Not exactly clear on move to front encoding
- **Decoding**



- Q: LZW doesn't seem like prefix-free compression. Do we use delimiters between the codes or is there some way to tell the numbers apart. Even just 255 contains multiple interpretations of ASCII codes (2, 55 or 25, 5)
- Great Question!
- Each integer is the same number of bits (e.g., 12 bits)
- So, the expansion program reads 12 bits at a time
- No need for delimiters nor prefix-free encoding of the integers

- Q: calculating the bits for the lzw compression (last question on the tophat), we didn't go over this during the example in class
- The second column represents the output of LZW compression, i.e., the compressed file
- Each integer is 12-bit codeword (per the question)
- Total compressed file size = # codewords * 12

- Q: Corner case of lzw expansion
- The tricky (corner) case happens when the longest match in compressions happens to be the string that was just added in the previous step
- Expansion sees that as a codeword that is not (yet) in its codebook
- Remember that expansion builds the same codebook as compression but is one step behind
- Handling the tricky case:
 - output: previous output + first character of previous output
 - add the same string to the codebook

LZW corner case example

Compress, using 12 bit codewords: AAAAAA

Cur (Outpu	t Add
Α	65	AA:256
AA	256	AAA:257
AAA	257	

Expansion:

Cur	Output	Add
65	Α	
256	AA	256:AA
257	AAA	257:AAA

- Q: How can the uncompressed file have more entropy than compressed if the entropy is the average number of bits to represent a word?
- In lossless compression,
 - entropy of compressed file >= entropy of uncompressed file
- Since compressed file has fewer characters than uncompressed
 - entropy/char of compressed file is > entropy/char of uncompressed file

- Q: I was wondering why LZW choose 12bits instead of any other number
- 12 bits was used just as an example
- Actual implementation use an adaptive codeword size
 - start with 9 bits
 - when codebook full, change to 10 bits
 - when codebook full, change to 11 bits
 - •
 - when codebook full and codeword is 16 bits.
 - either stop adding to codebook or reset it
 - depends on the compression ratio

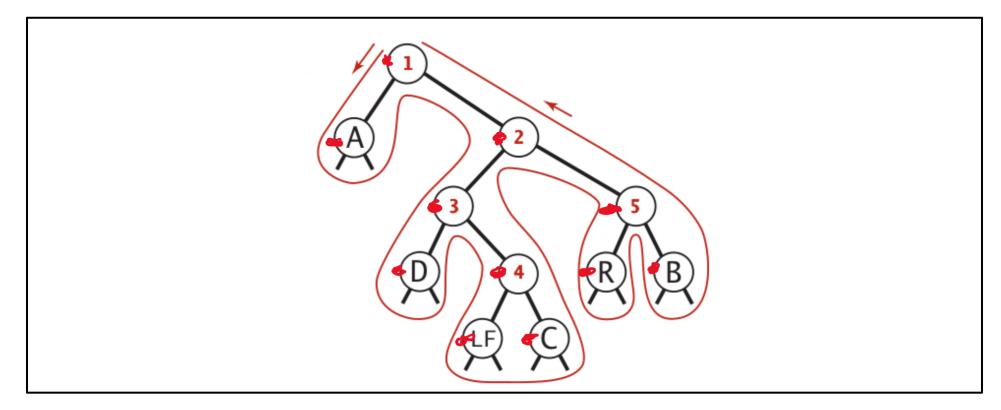
- Q: please make assignments easier
- Yep!

- Q: Is there a best way to handle ties when doing Huffman compression or is it a purely arbitrary choice?
- Q: what are the rules for drawing out the tree in the huffman approach?
- Ties are arbitrarily handled
- The compressed file size is the same no matter how ties are handled

- Q: the bits needed for the compression post huffman encoding
- Huffman compression may store the trie in the compressed file
- The trie is encoded using preorder traversal of the nodes
 - encoding internal nodes with 0
 - leaf nodes with 1 followed by the ASCII code of the character inside the leaf

Representing tries as bitstrings

Preorder traversal

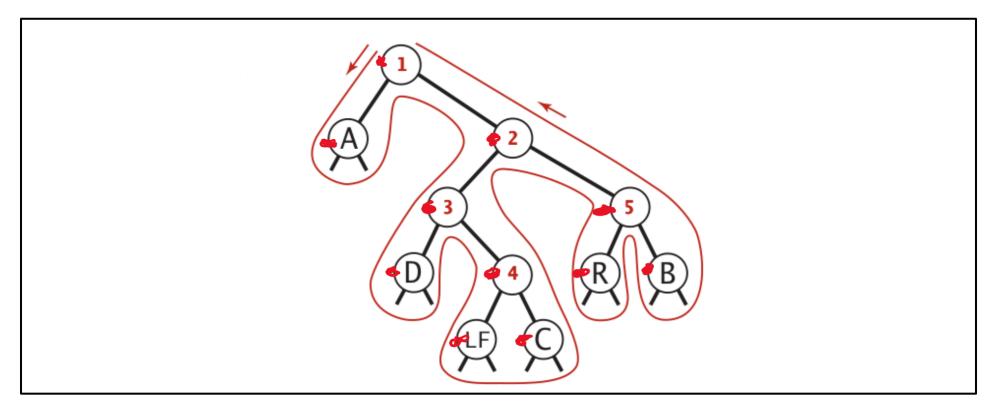


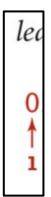
internal node → 0

leaf node -> 1 followed by ASCII code of char inside

Representing tries as bitstrings

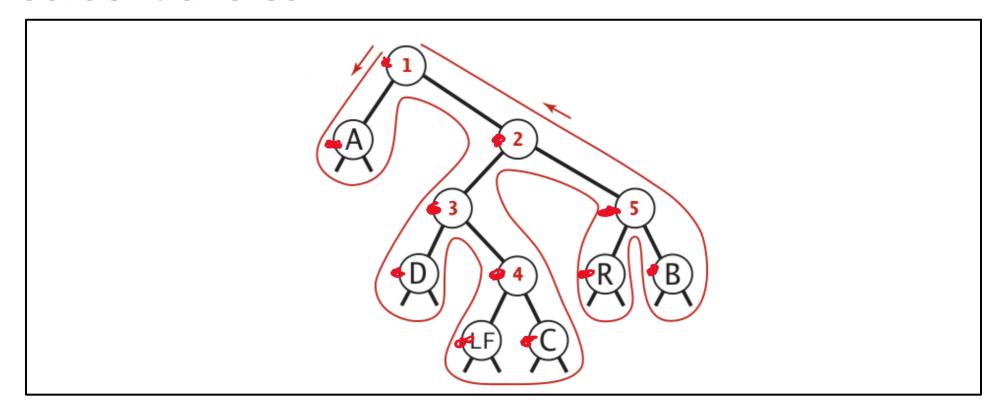
Preorder traversal

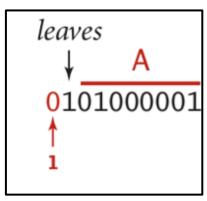


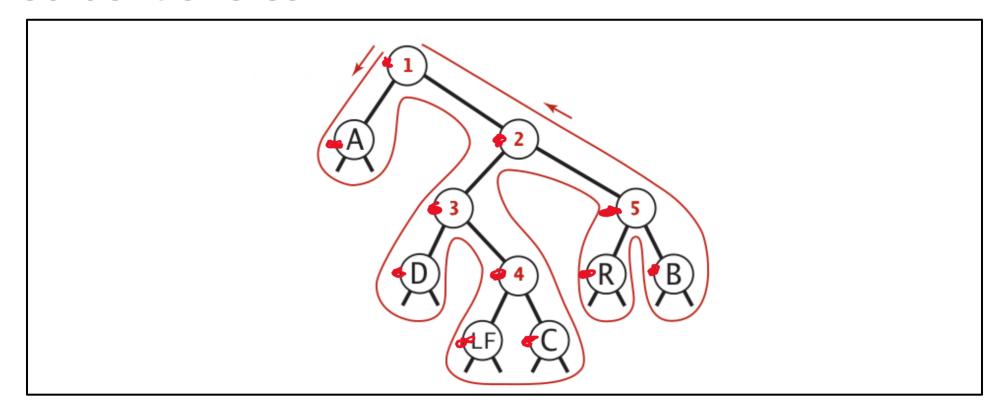


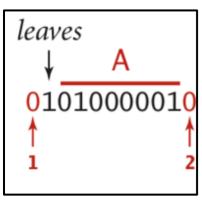
Representing tries as bitstrings

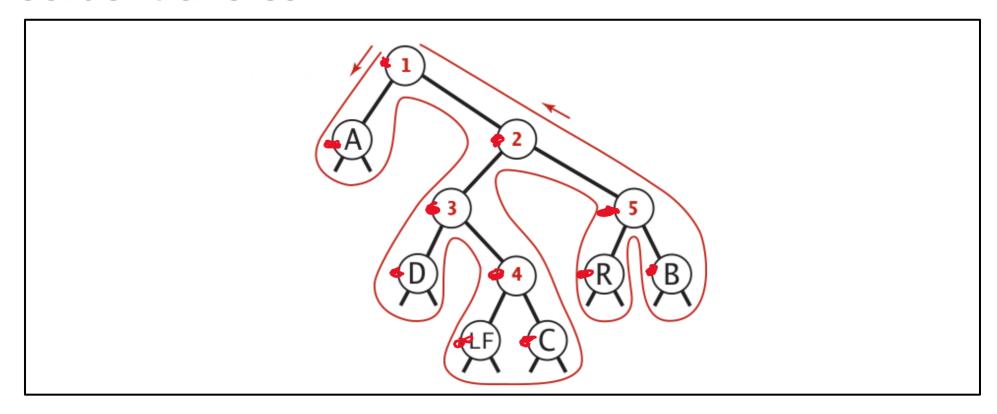
Preorder traversal

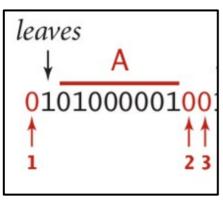


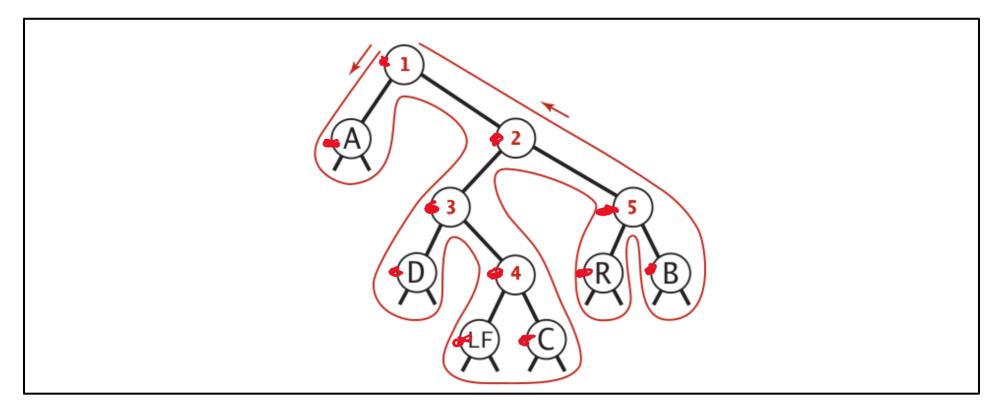


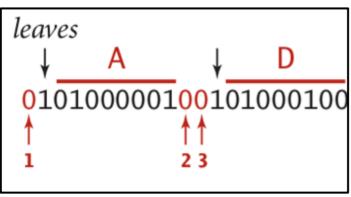


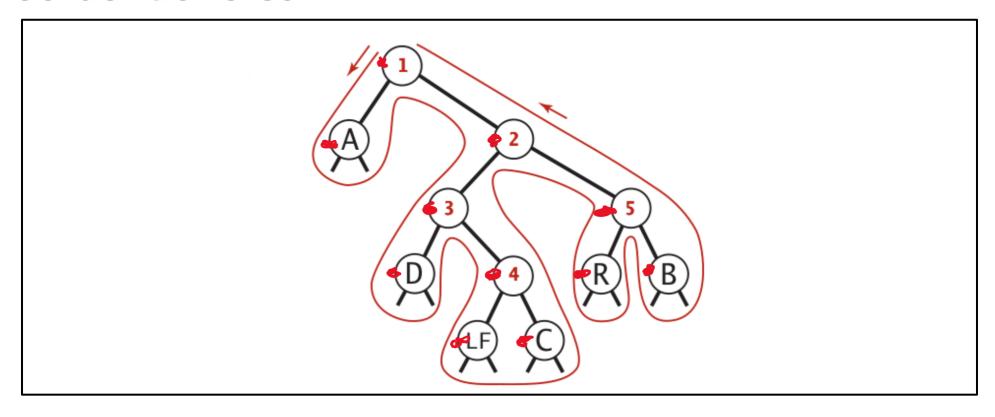


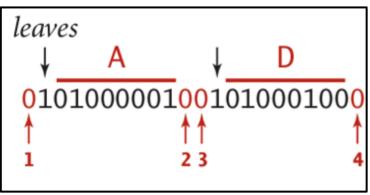


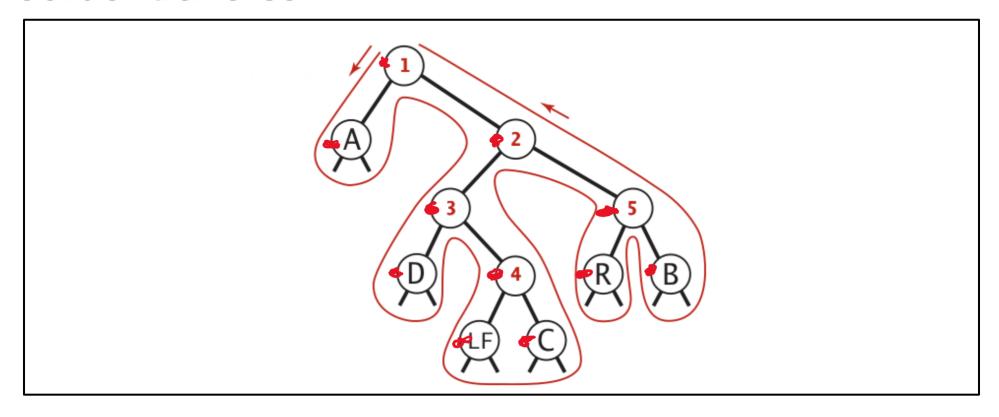


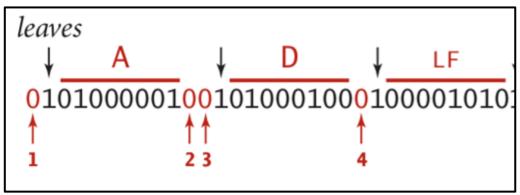


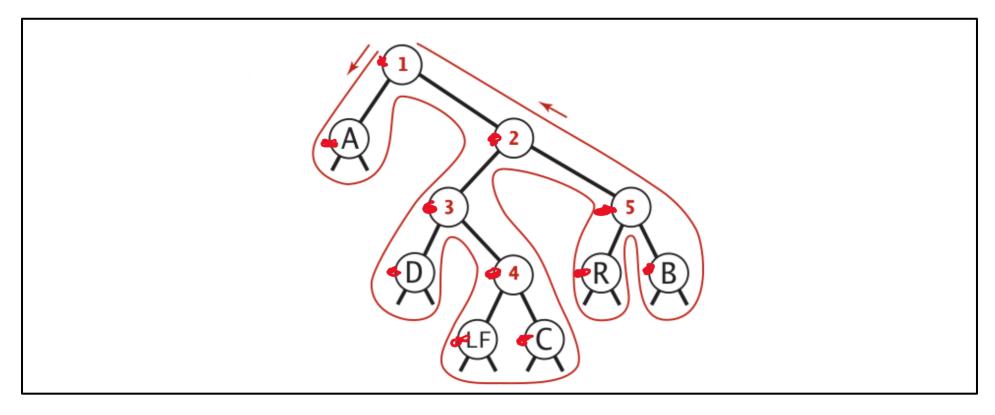


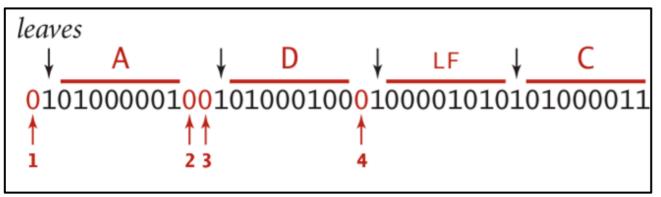


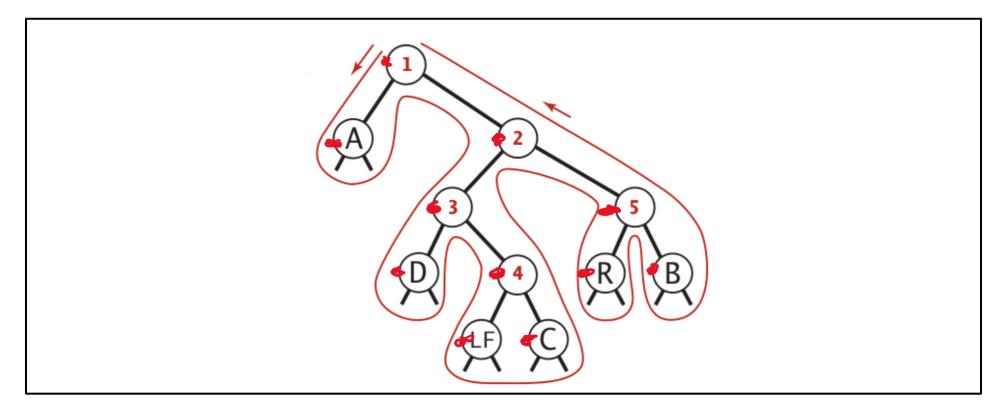


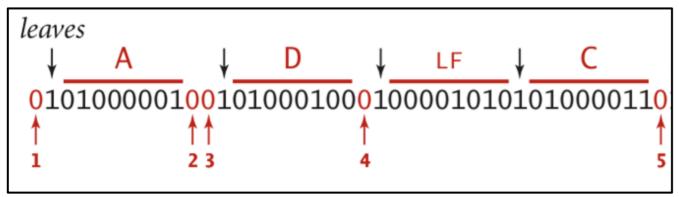


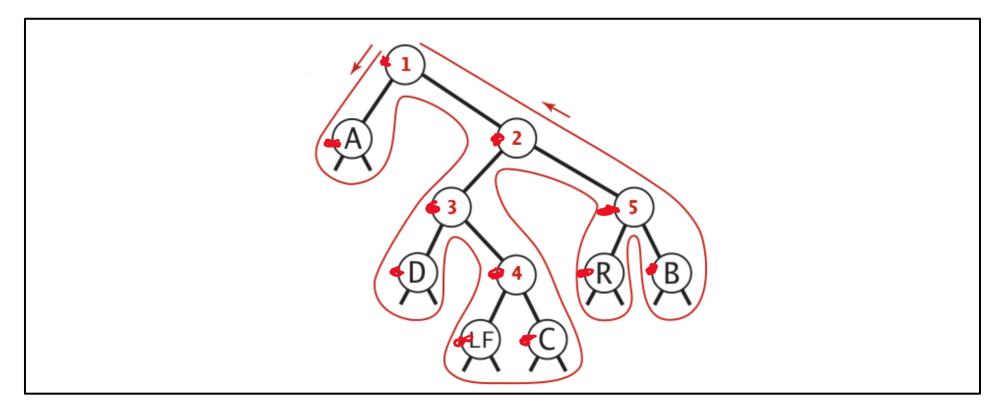


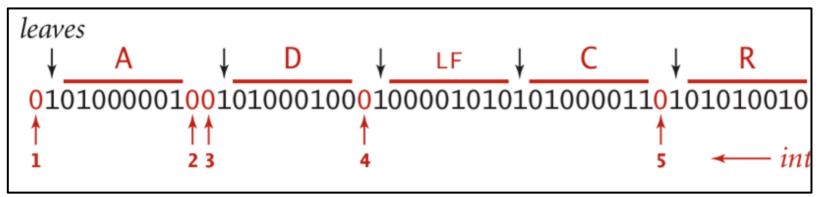


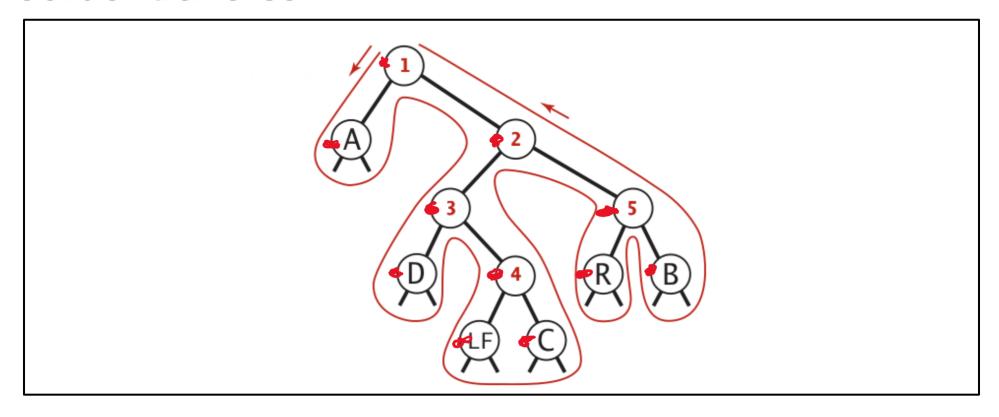


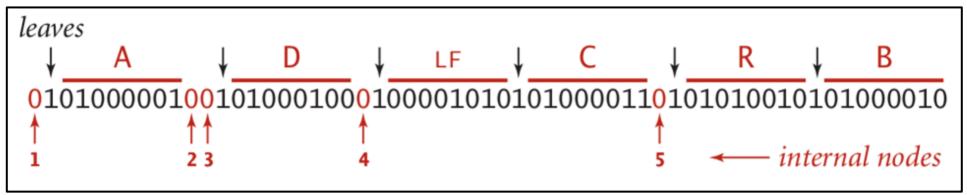












- Q: After writing out the Trie as a bitstring, how can you tell where each character starts?
- First, we are writing bits not characters!
- Each bit represents a node
 - except for leaves, their 1 bit is followed by 8 bits (ASCII code of char inside)
 - but that's also fixed length

- Q: how compression and tries combine. Are you supposed to recreate the original trie while decompressing
- The compressed file contains:
 - the trie representation in bits
 - the number of characters in the original file
 - the Huffman encoding of the file characters
- The original trie is reconstructed from the trie representation in the compressed file

- Q: Questions are not getting responded to on Piazza in a timely manner. Students had asked several questions pertaining to items on Homework 4 last week but the questions were not addressed until AFTER Homework 4 was already graded. The questions I got marked wrong this week were the exact questions I had asked for clarification on but didn't get a response to (until after the homework was graded and returned...by that point it was too late).
- I will hold a Live QA Session on Piazza every Friday 4:30-5:30 pm

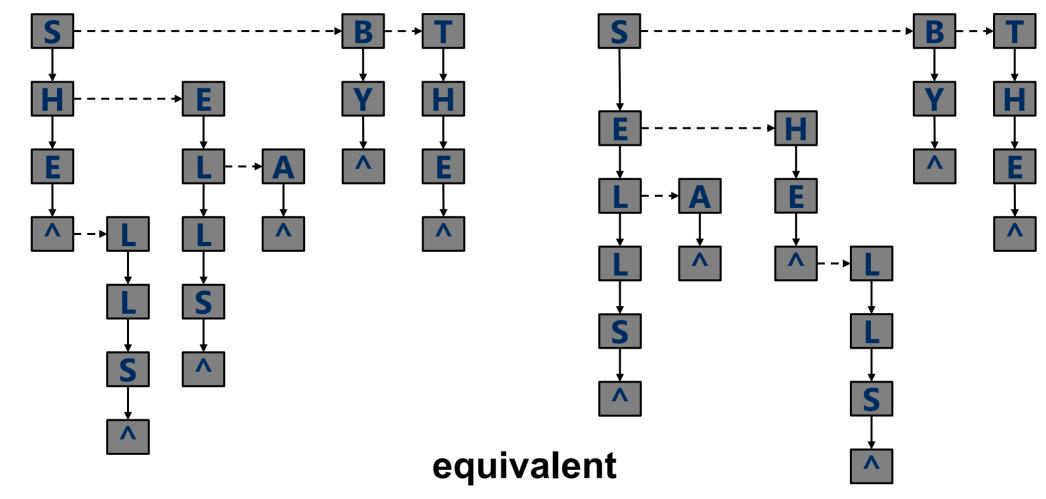
- Q: the overall concepts of code blocks/code words, like how do they fit into everything?
- The input file is divided into code blocks
- Each code block is replaced by a codeword
- For Huffman:
 - code blocks are single characters
 - codewords are variable-length bit strings
- For LZW:
 - code blocks are the longest-match strings (variable length)
 - codewords are fixed-length integers (e.g., 12 bits)
- For RLE:
 - code block are long strings with identical characters (variable length)
 - codewords are fixed-length integer followed by fixed-length ASCII of the character

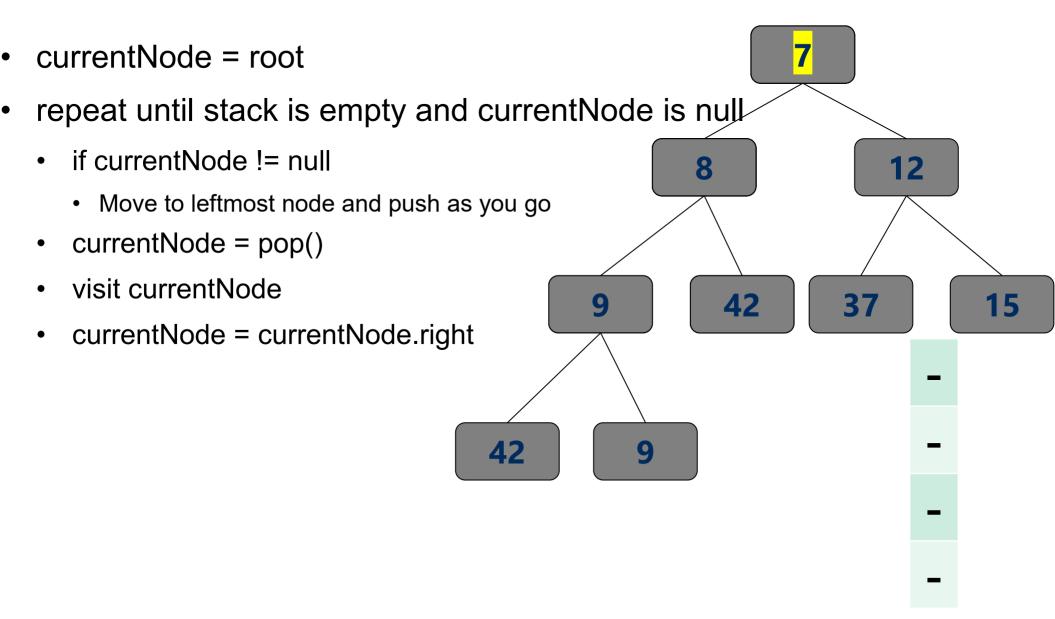
- Q: Just making sure I got the order of the compression framework: We go from file to code block to code word via compression then back to code block via expansion
- Yep!

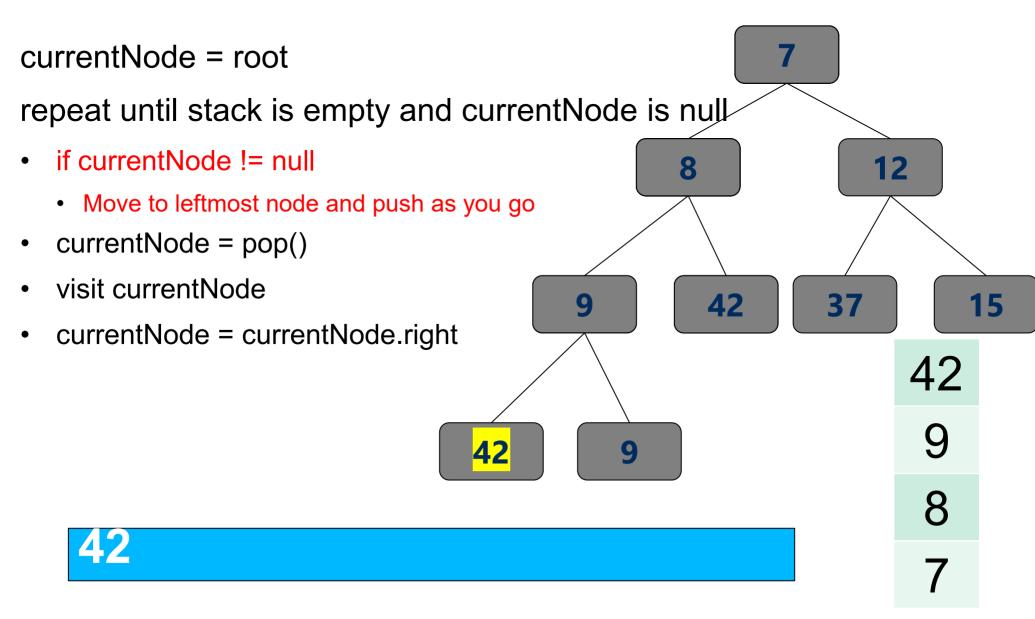
- Q: Can we get partial credit on labs and projects?
- We assign partial credit on projects only and only when autograder score <=
 40%
 - partial credit has a ceiling of 60% of the autograder score
- Unless you think you lost points in the autograder because of an autograder error

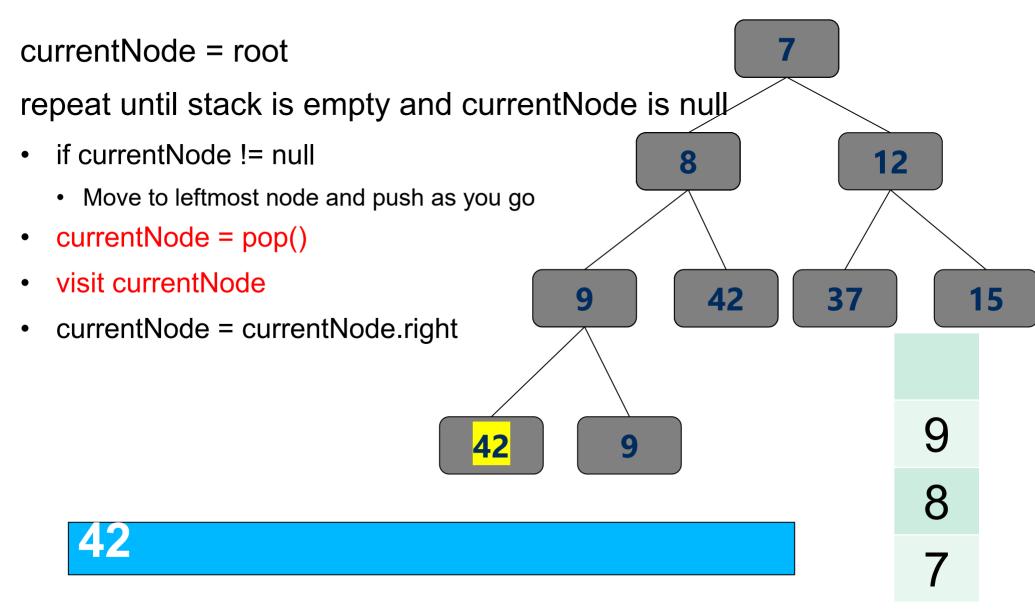
- Q:How are you able to create unique codes without leading to another prefix
- Characters are leaves in the Huffman tree
- No two characters share the same path from root
- Codewords encode the root-to-leaf path
- No two codewords share a prefix

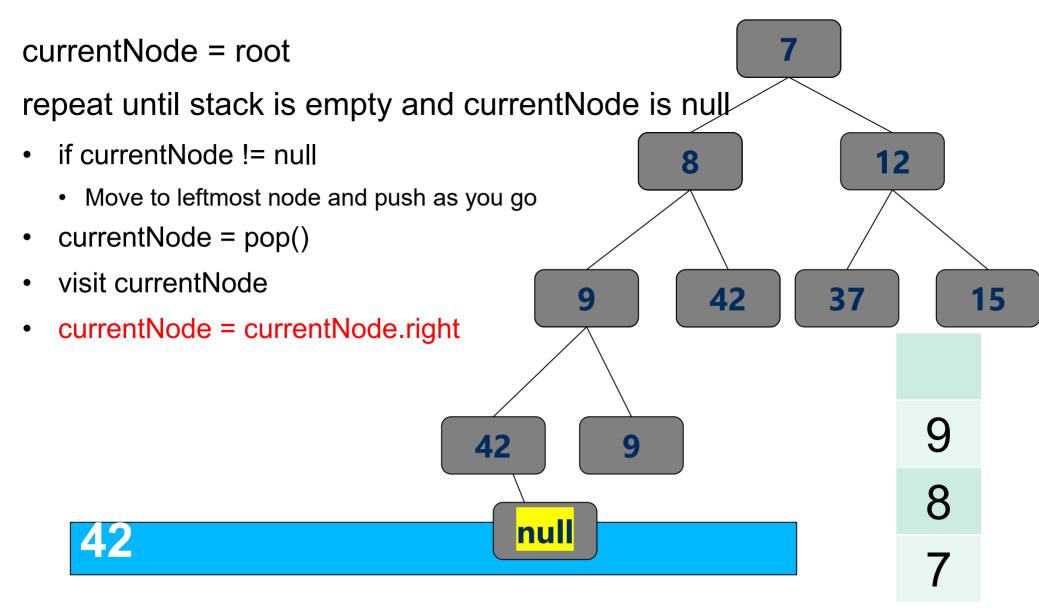
- Q: In DLB tries, can you interchange a parent node's child with any one of the child's siblings? I think you can but you would have to change the sibling links, is that correct?
- Correct!

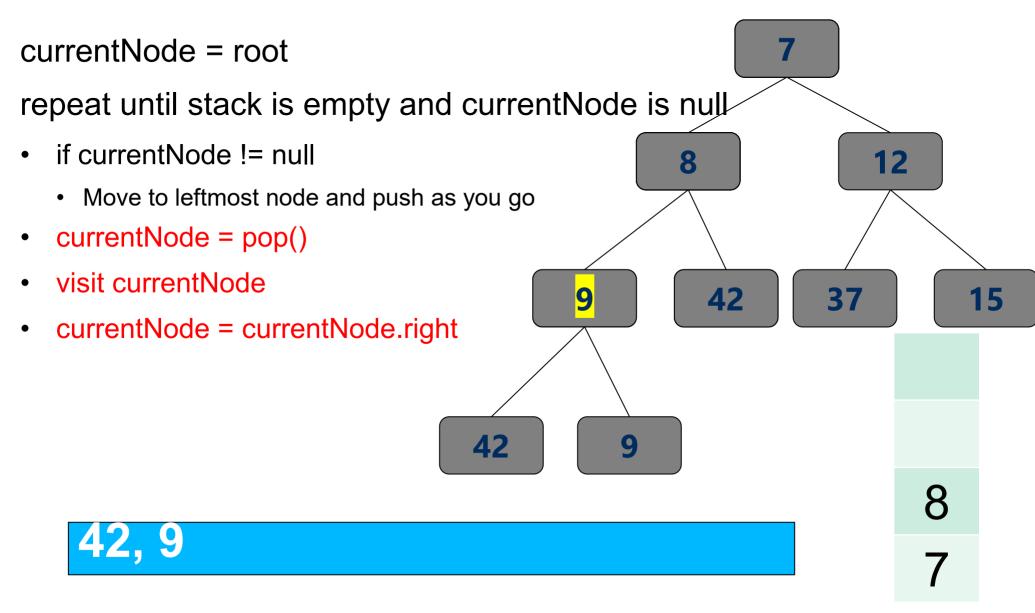


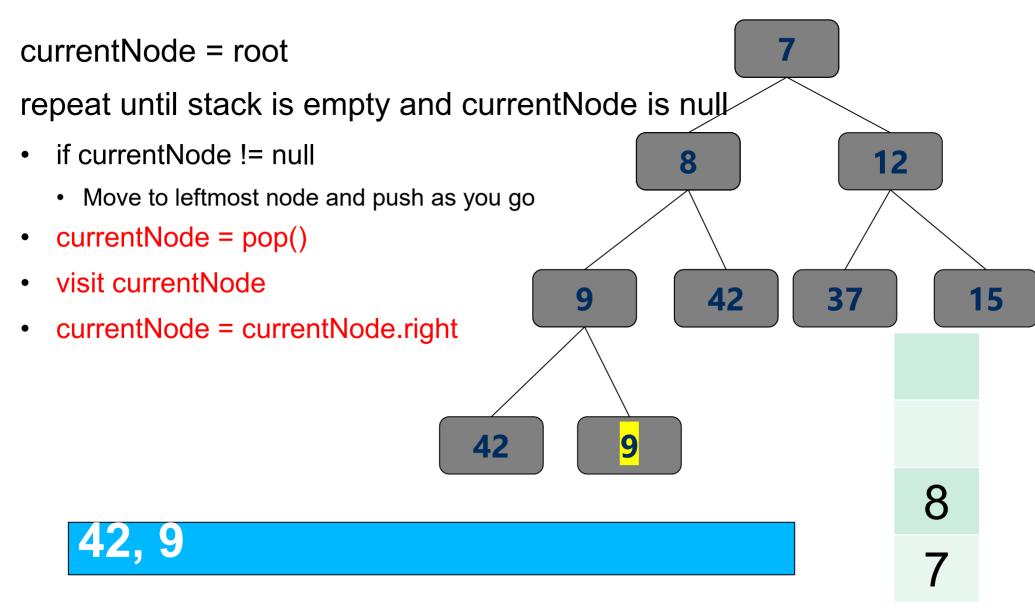


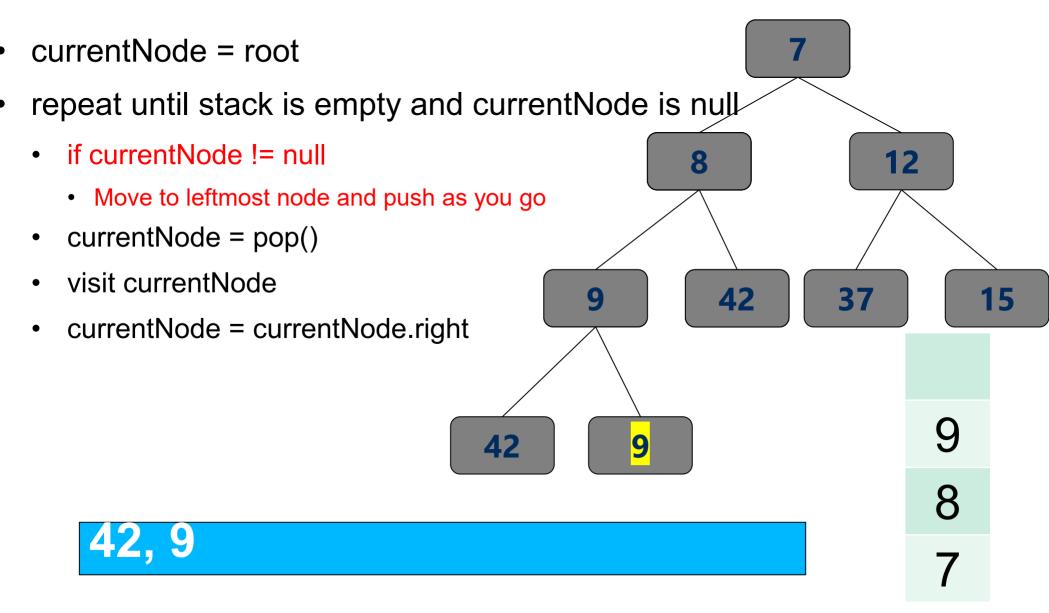


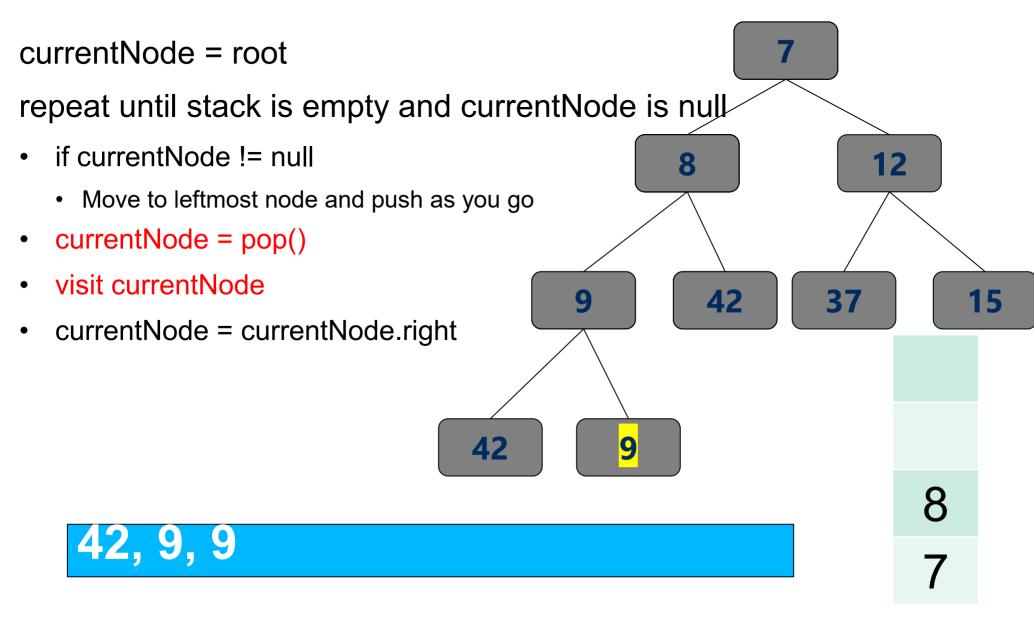


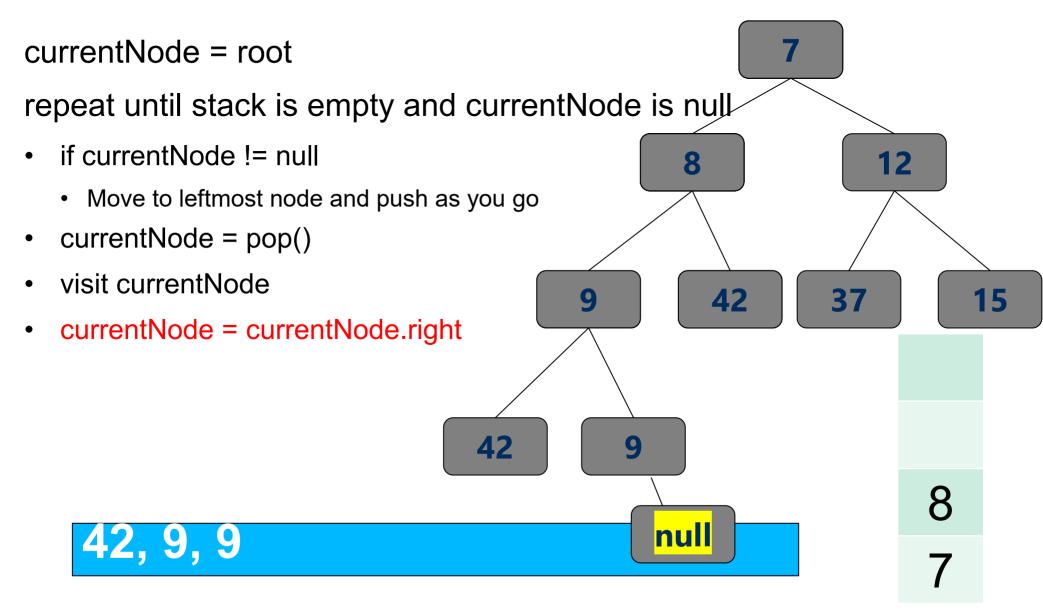


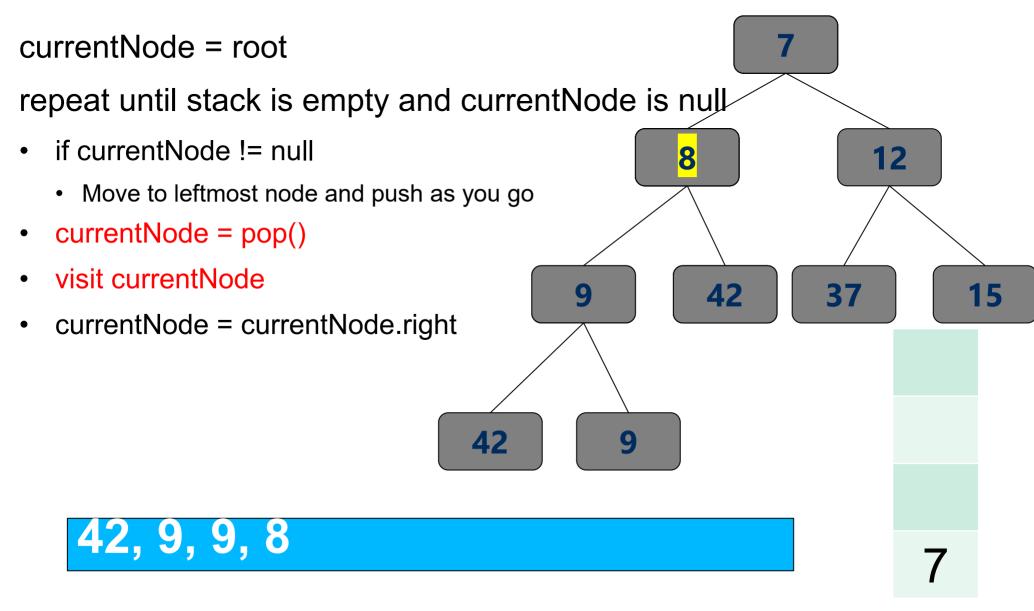


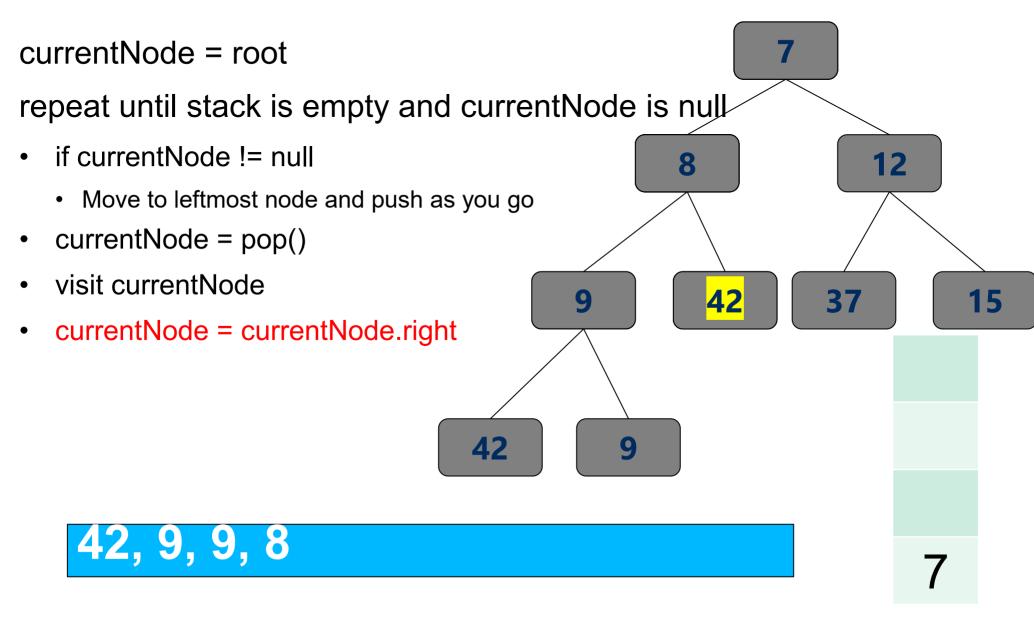


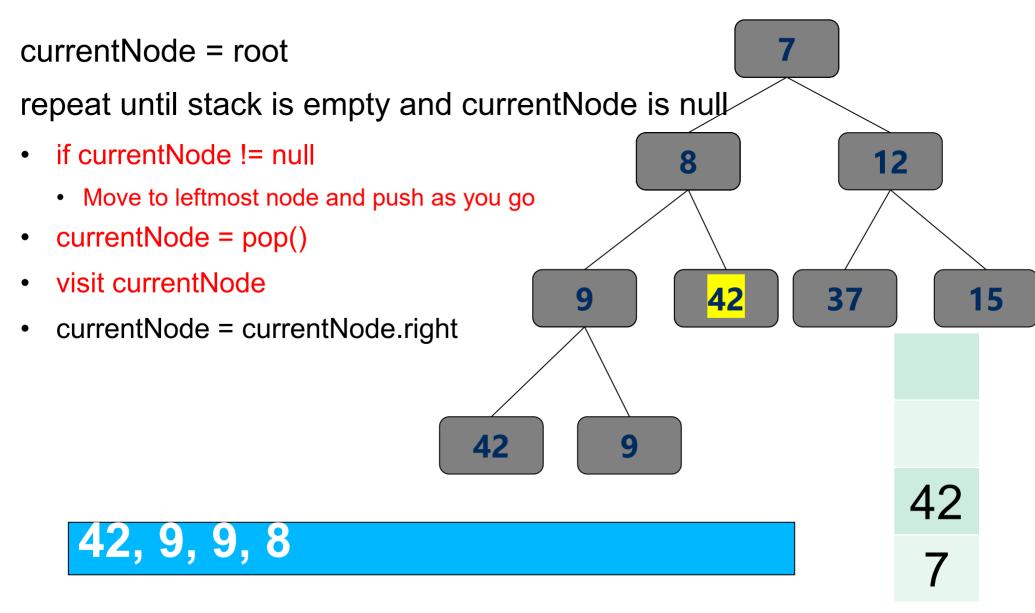


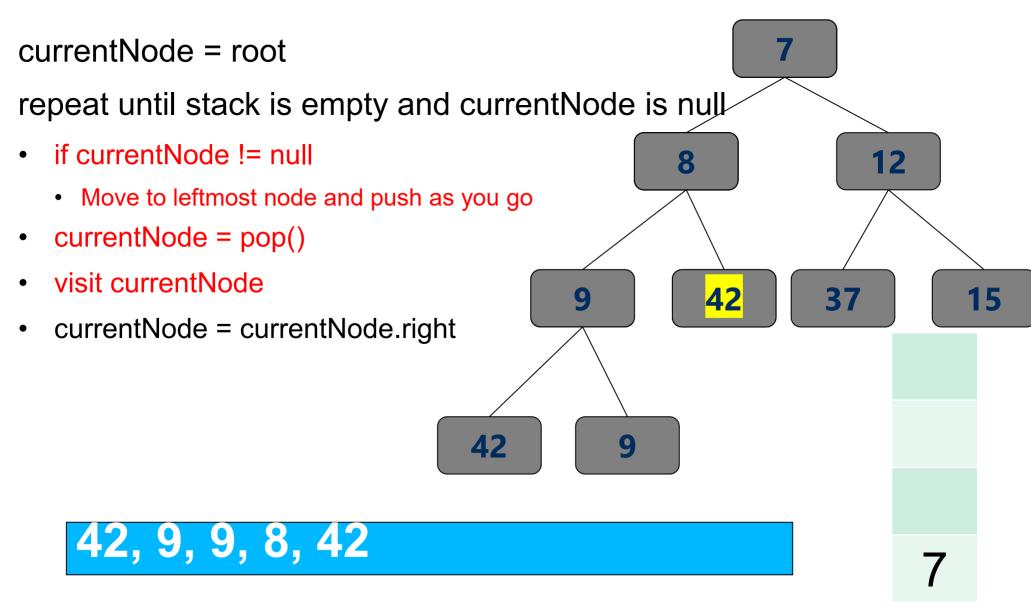


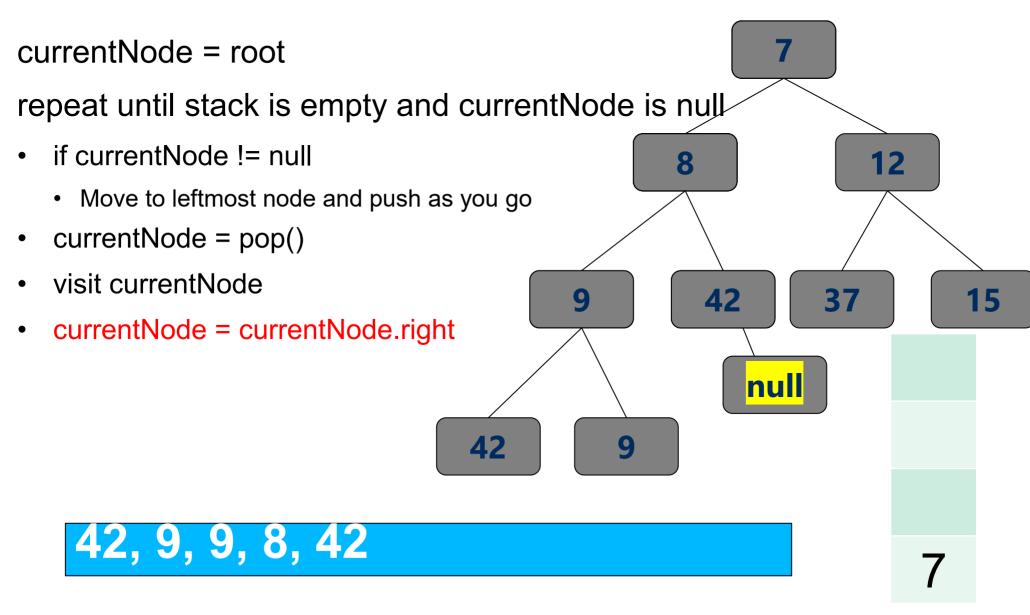


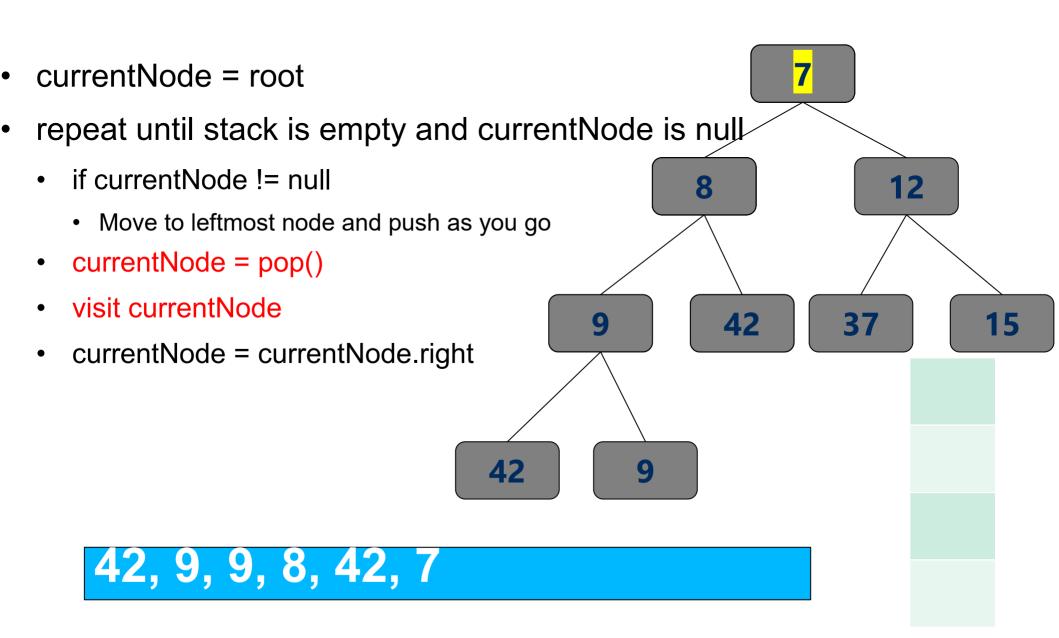


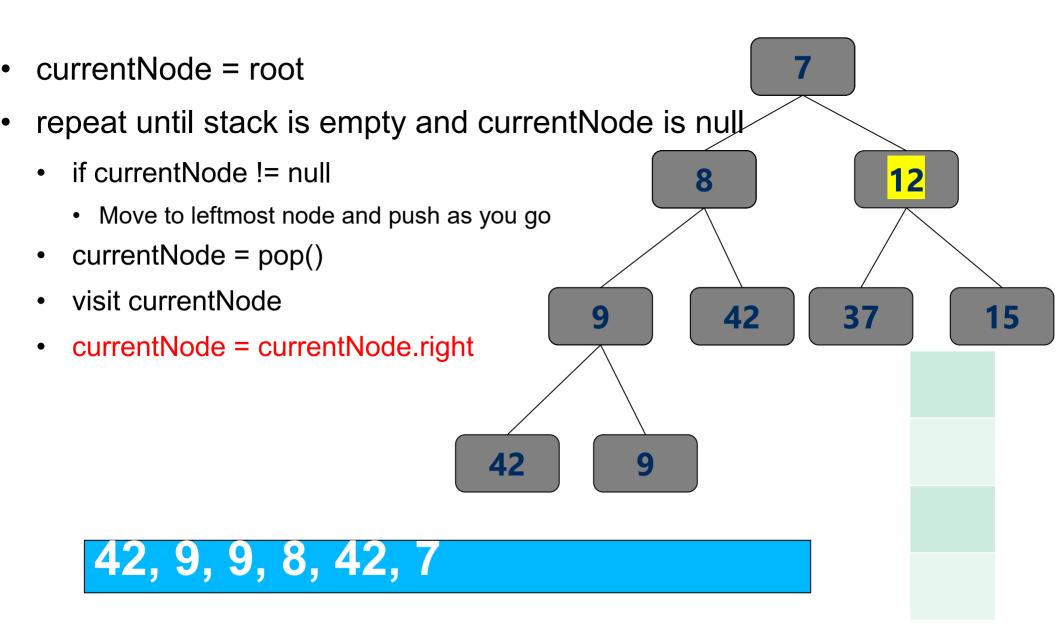


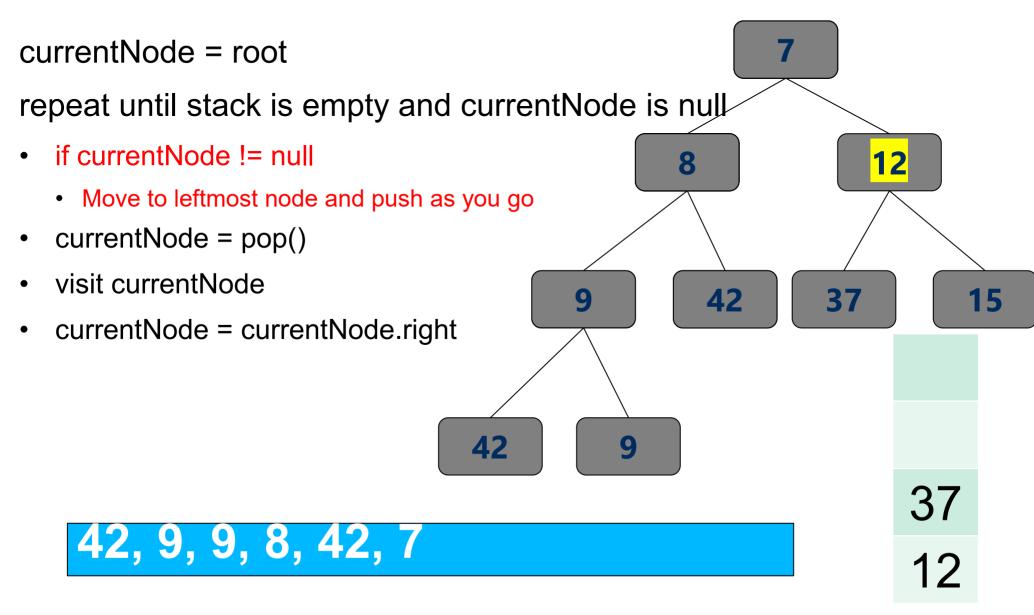


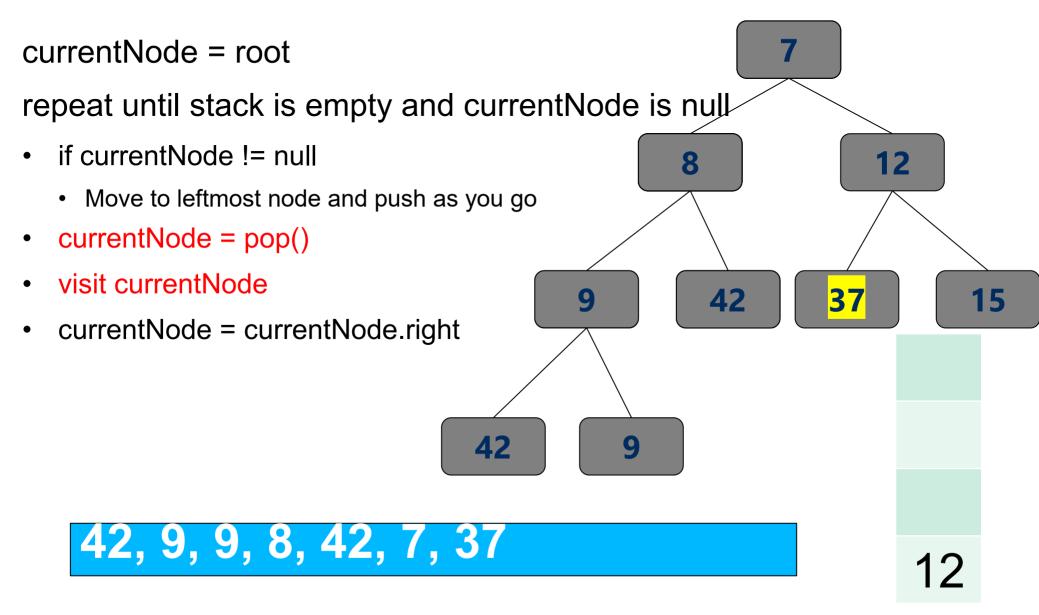


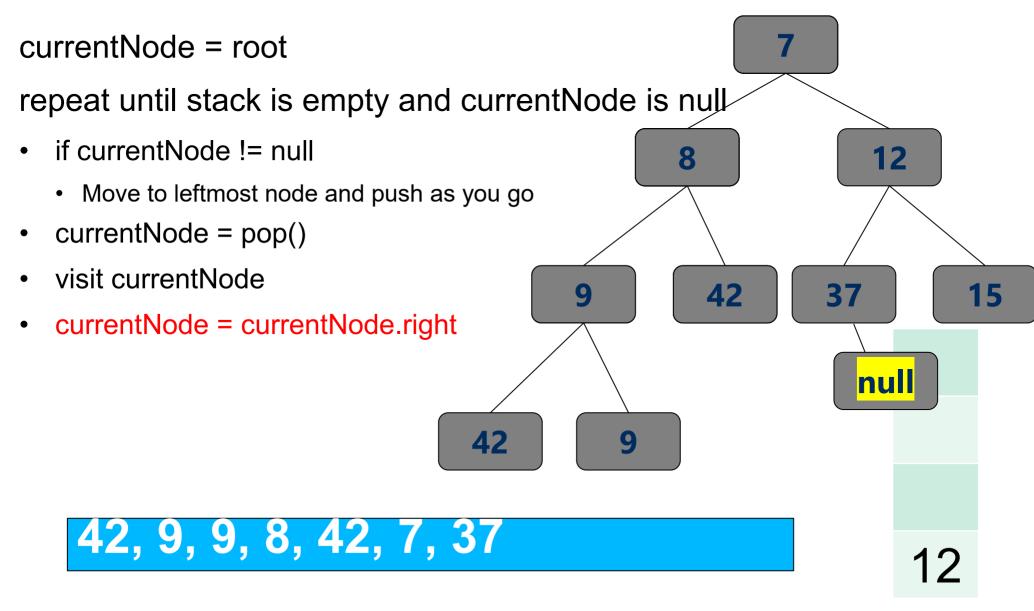


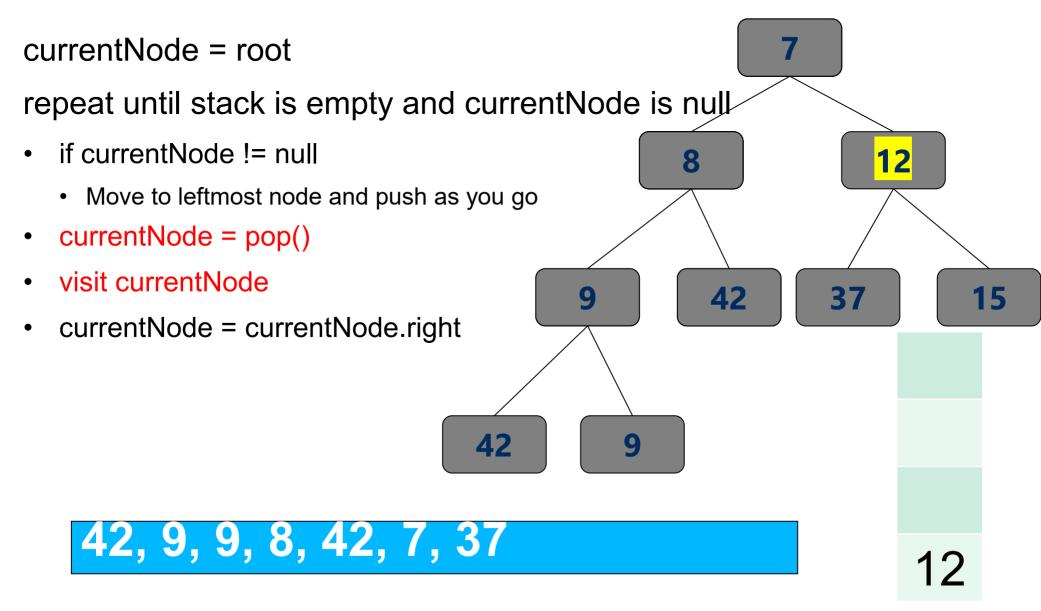


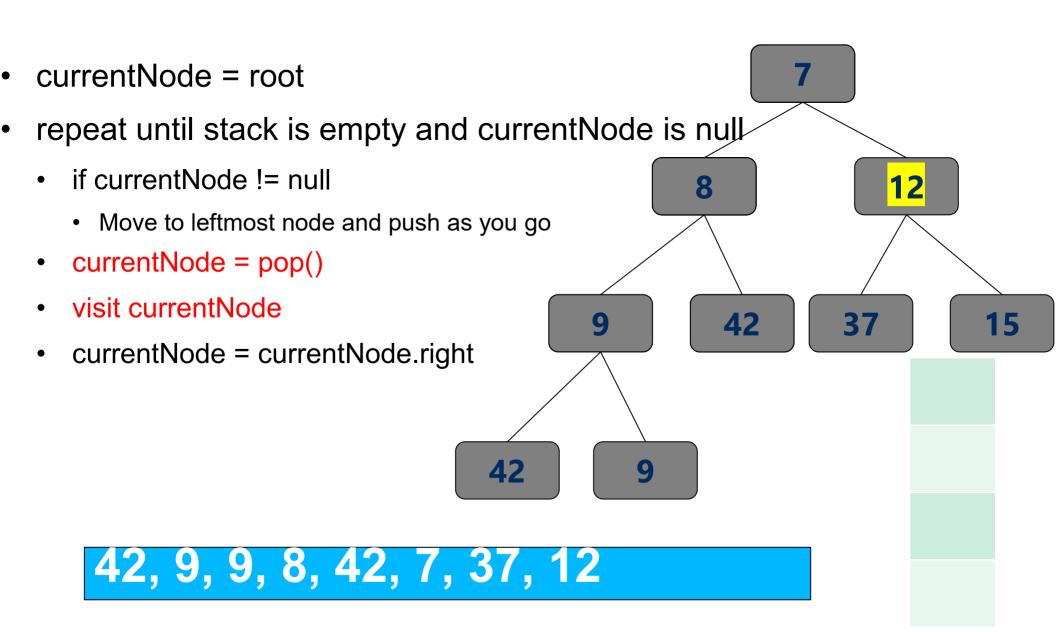


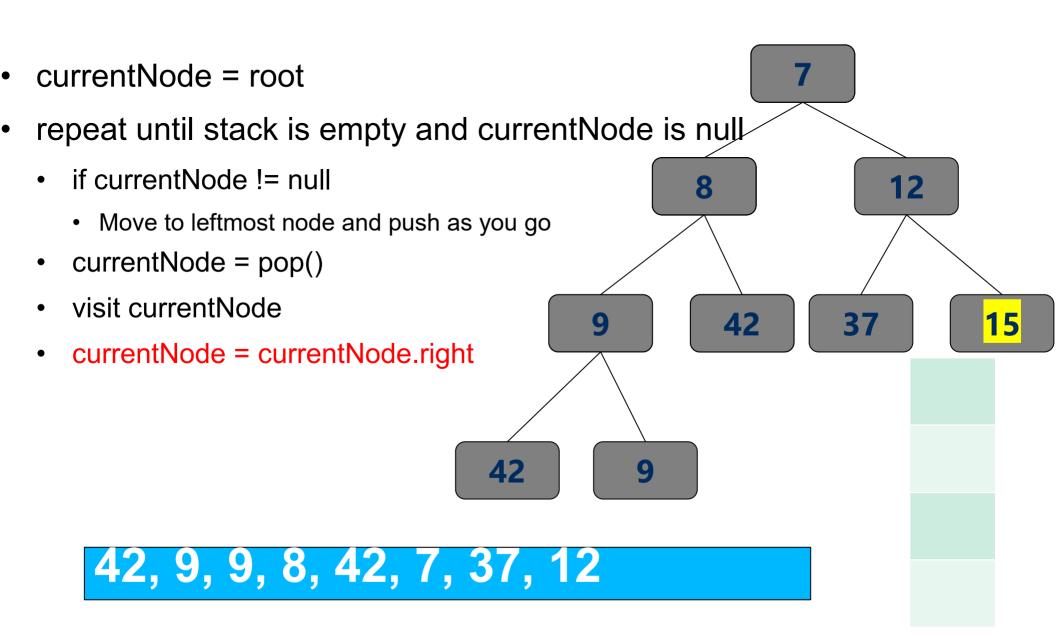


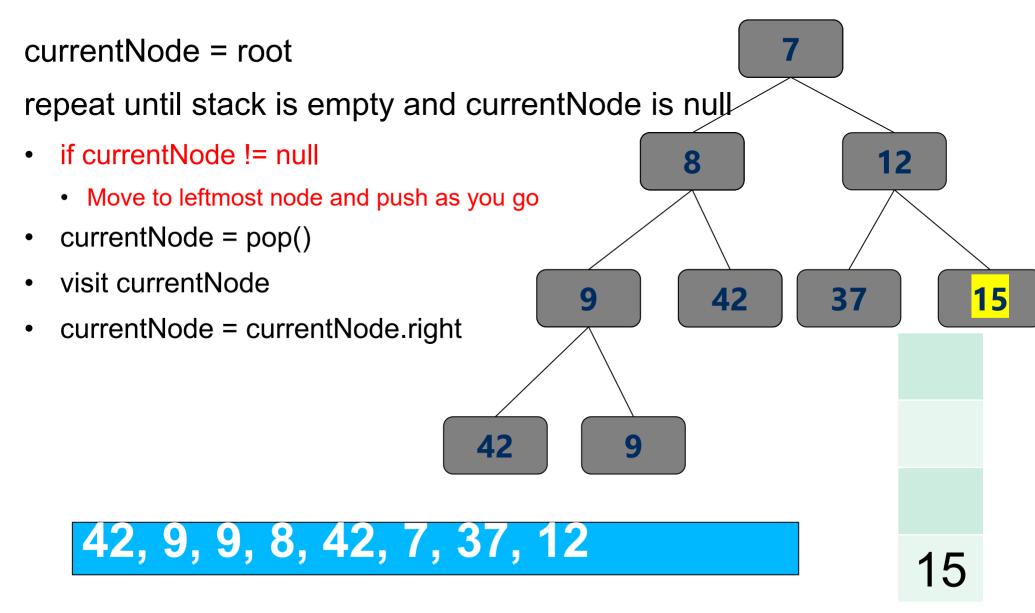


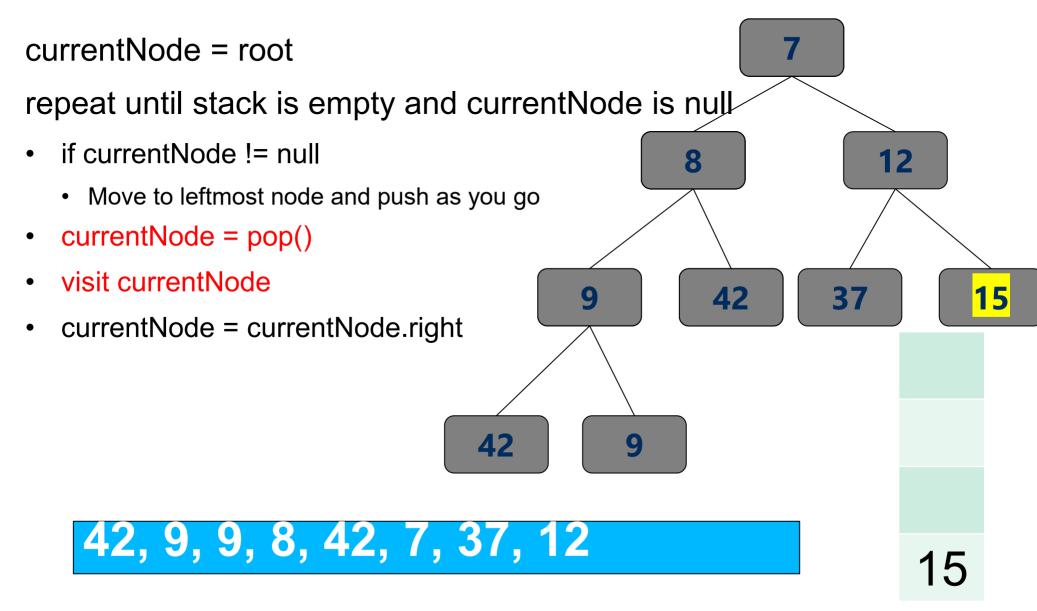


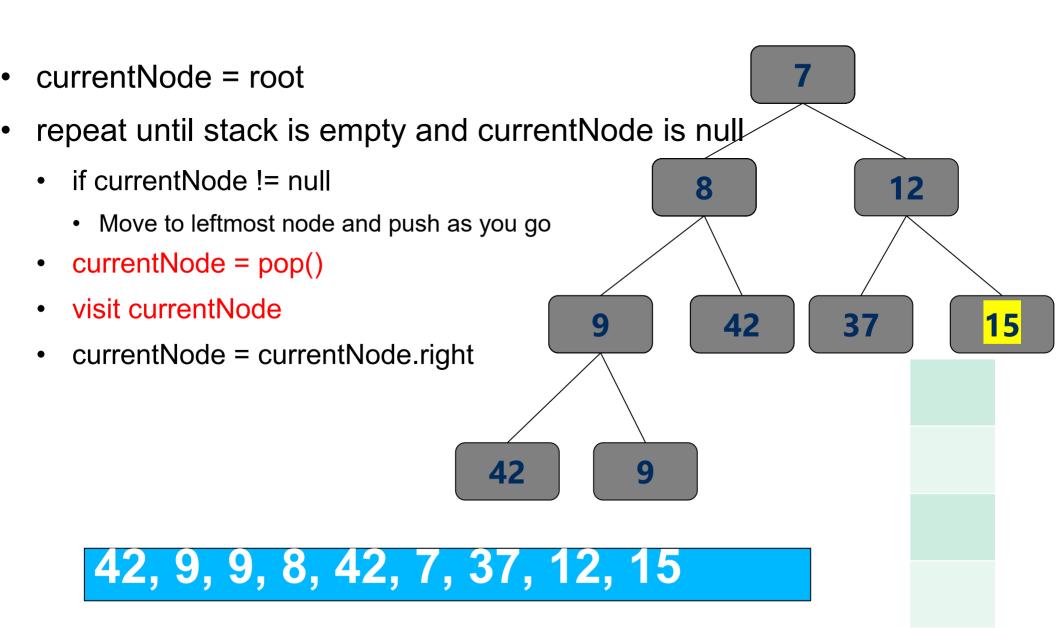


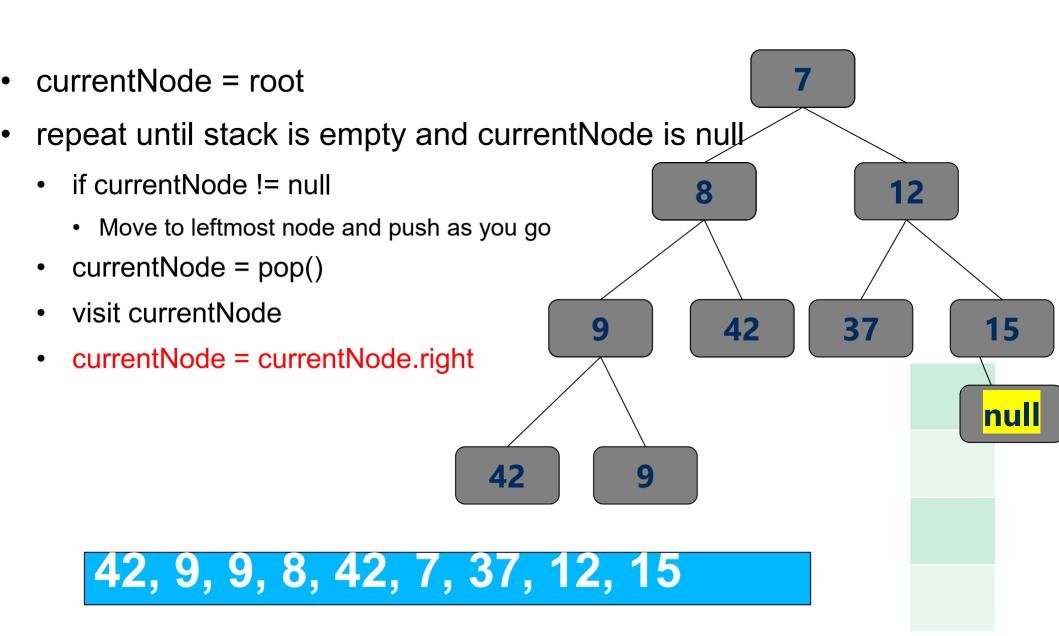


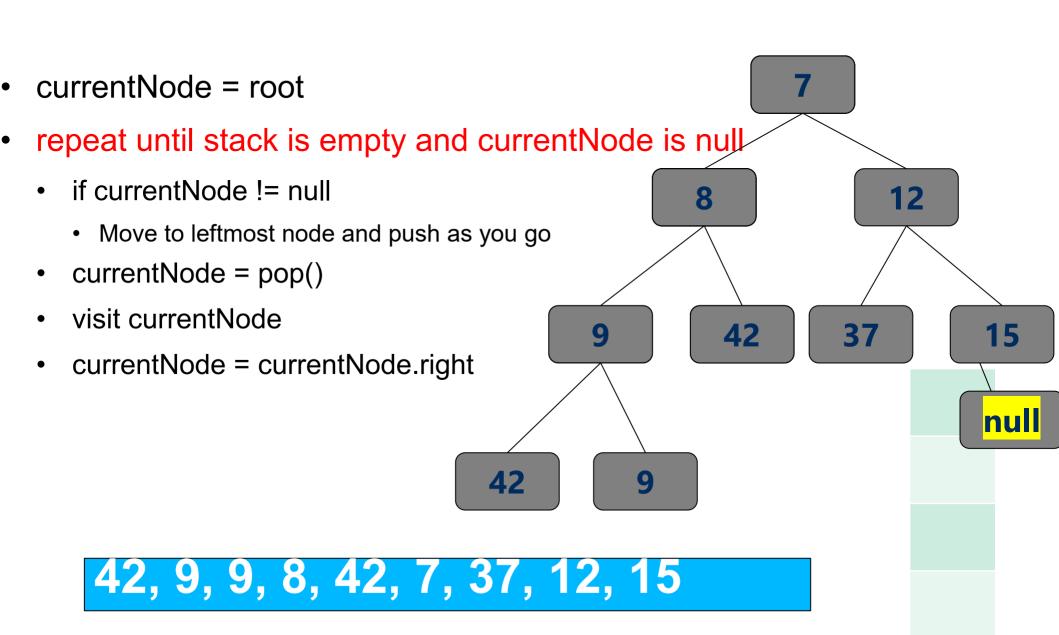








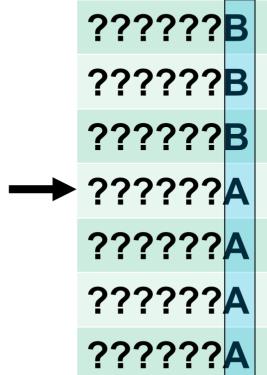




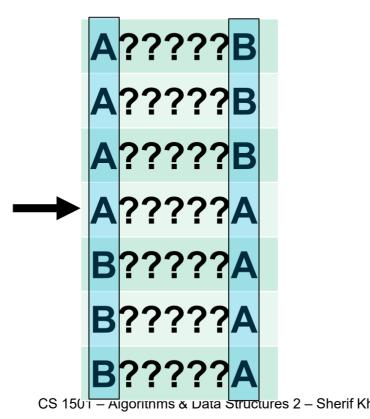
Burrows-Wheeler Transform Example

- ABABABA
- Step 1: Build an array of 7 strings, each a circular rotation of the original by one character
- ABABABA
- BABABAA sorted array original array **ABABAAB ARARARA** AABABAB BABAABA **BABABAA** ARAARAR **ABABAAB** ABABAAB **ABAABAB BABAABA** ABABABA BAABABA BAABABA **ABAABAB AABABAB BAABABA** BABAABA Step 2: Sort the array alphabetic **AABA** Notice that the first column of the some characters as BBBAAAA and 3 the last column
 - all columns have the same set of letters
- Step 3: Output the last column of the sorted array and the index of the input string in the sorted array

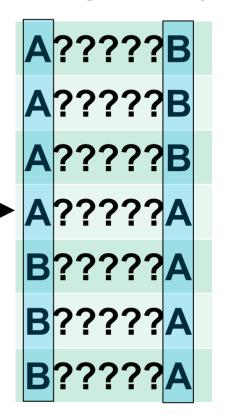
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 1: Sort the encoded string
 - BBBAAAA → AAAABBB
 - The first column of the sorted array has the same characters as the last column
 - but in sorted order



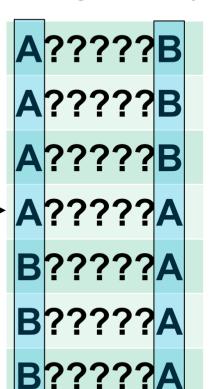
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 1: Sort the encoded string
 - BBBAAAA → AAAABBB
 - This gives us the first column of the sorted array



- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 2: Fill an array next[]
 - defined for each entry in the sorted array
 - tells us the index in sorted array of the next string in the original array
 - Scan through the first column
 - for each row *i* holding character *c*
 - next[i] = first unassigned index of c in the last column



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next -

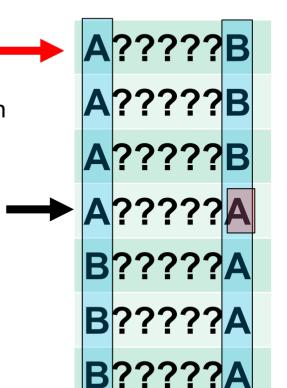
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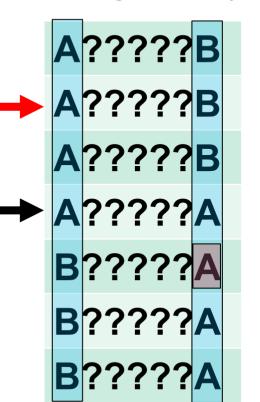
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 - defined for each entry in the sorted array
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 - Scan through the first column
 - for each row *i* holding character *c*
 - next[i] = first unassigned index of c in the last column



next

3

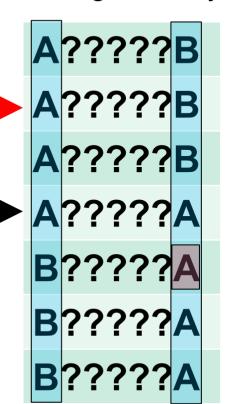
4

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- Output of BWT:
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next

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4

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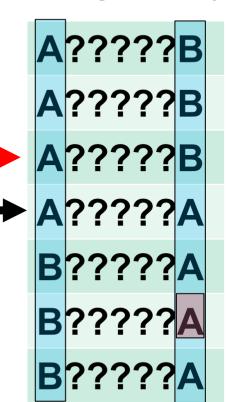
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next

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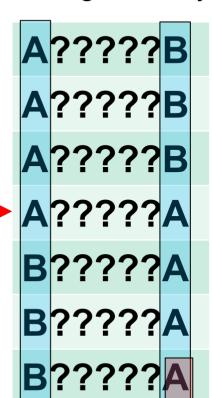
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next

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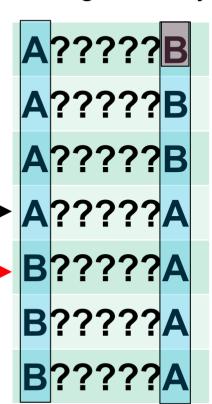
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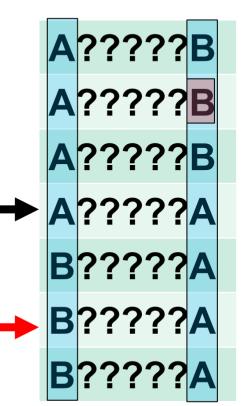
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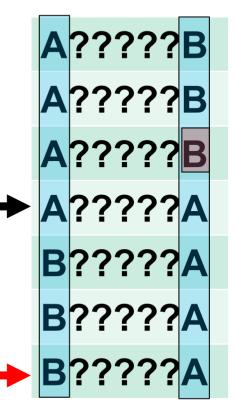
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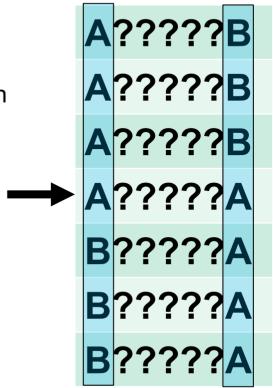
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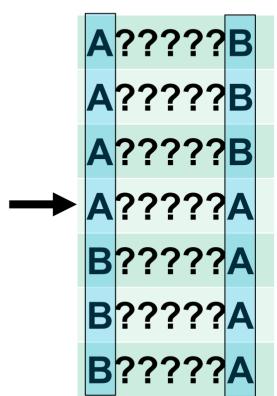


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- Step 2: Fill an array next[]
 - defined for each entry in the sorted array
 - tells us the index in sorted array of the next string in the original array
 - Scan through the first column
 - for each row i holding character c
 - next[i] = first unassigned index of *c* in the last column
- Why does that work?
 - first character of a string becomes the last character in the next string in the original order



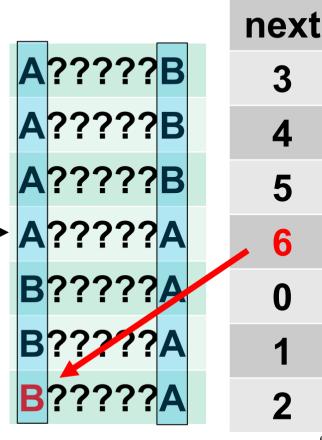
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?

A??????



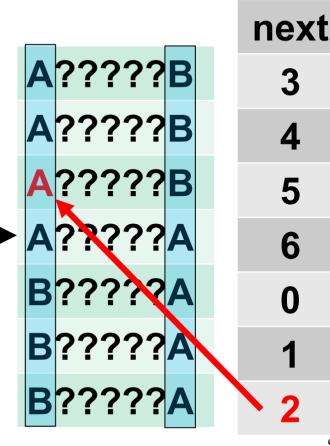
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?
- The next character is the first character of the next string in the original order
 - first character in string at next[3]

AB?????



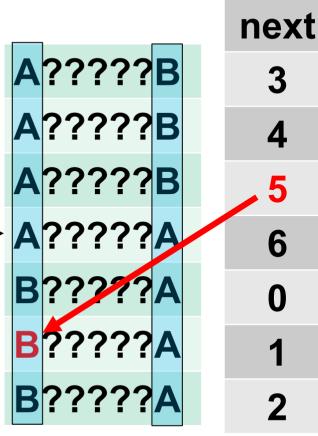
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?
- The next character is the first character of the next string in the original order
 - first character in string at next[6]

ABA????



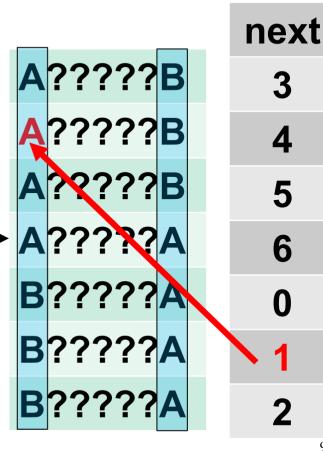
- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?
- The next character is the first character of the next string in the original order
 - first character in string at next[2]

ABAB???



- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?
- The next character is the first character of the next string in the original order
 - first character in string at next[5]

ABABA??



- Output of BWT:
 - BBBAAAA and 3
- How can we recover ABABABA?
- Step 3: Recover the input string using the next[] array
- We can conclude that A is the first character in the input string
 - why?
- The next character is the first character of the next string in the original order
 - first character in string at next[5]



