



Computer Architecture

Introduction and Review

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DIGITAL LOGIC CIRCUITS

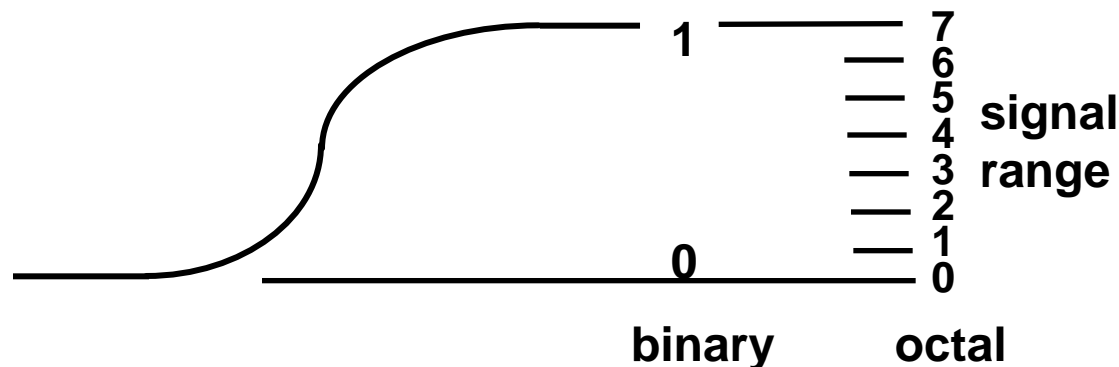
- **Logic Gates**
- **Boolean Algebra**
- **Map Specification**
- **Combinational Circuits**
- **Flip-Flops**
- **Sequential Circuits**
- **Memory Components**

LOGIC GATES

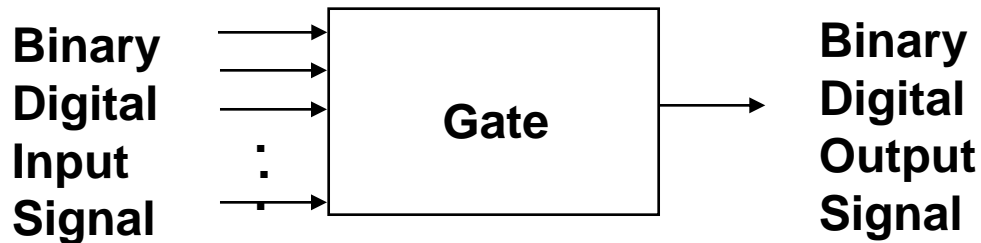
Digital Computers

- Imply that the computer deals with digital information, i.e., it deals with the information that is represented by binary digits
- Why *BINARY*? instead of Decimal or other number system ?

* Consider electronic signal



BASIC LOGIC BLOCK - GATE -



Types of Basic Logic Blocks

- **Combinational Logic Block**

Logic Blocks whose output logic value depends only on the input logic values

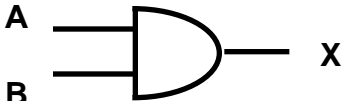


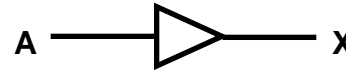
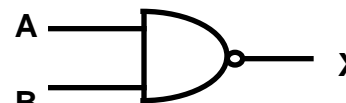
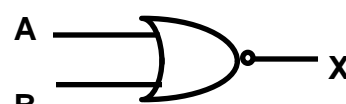
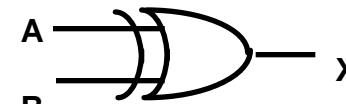
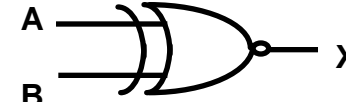
- **Sequential Logic Block**

Logic Blocks whose output logic value depends on the input values and the state (stored information) of the blocks

Functions of Gates can be described by

- **Truth Table**
- **Boolean Function**
- **Karnaugh Map**

COMBINATIONAL GATES

Name	Symbol	Function	Truth Table															
AND		$X = A \cdot B$ or $X = AB$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table>	A	B	X	0	0	0	0	1	0	1	0	0	1	1	1
A	B	X																
0	0	0																
0	1	0																
1	0	0																
1	1	1																
OR		$X = A + B$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table>	A	B	X	0	0	0	0	1	1	1	0	1	1	1	1
A	B	X																
0	0	0																
0	1	1																
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I		$X = A$	<table><tr><th>A</th><th>X</th></tr><tr><td>0</td><td>1</td></tr><tr><td>1</td><td>0</td></tr></table>	A	X	0	1	1	0									
A	X																	
0	1																	
1	0																	
Buffer		$X = A$	<table><tr><th>A</th><th>X</th></tr><tr><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td></tr></table>	A	X	0	0	1	1									
A	X																	
0	0																	
1	1																	
NAND		$X = (AB)'$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table>	A	B	X	0	0	1	0	1	1	1	0	1	1	1	0
A	B	X																
0	0	1																
0	1	1																
1	0	1																
1	1	0																
NOR		$X = (A + B)'$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table>	A	B	X	0	0	1	0	1	0	1	0	0	1	1	0
A	B	X																
0	0	1																
0	1	0																
1	0	0																
1	1	0																
XOR Exclusive OR		$X = A \oplus B$ or $X = A'B + AB'$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table>	A	B	X	0	0	0	0	1	1	1	0	1	1	1	0
A	B	X																
0	0	0																
0	1	1																
1	0	1																
1	1	0																
XNOR Exclusive NOR or Equivalence		$X = (A \oplus B)'$ or $X = A'B' + AB$	<table><tr><th>A</th><th>B</th><th>X</th></tr><tr><td>0</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table>	A	B	X	0	0	1	0	1	0	1	0	0	1	1	1
A	B	X																
0	0	1																
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1	1	1																

BOOLEAN ALGEBRA

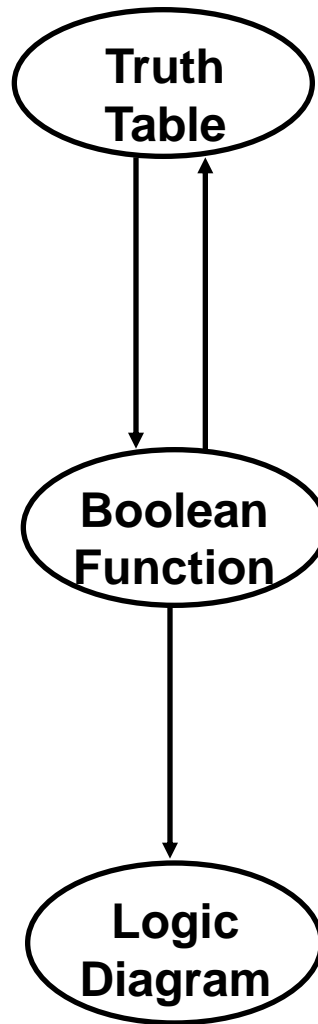
Boolean Algebra

- * Algebra with Binary(Boolean) Variable and Logic Operations
- * Boolean Algebra is useful in Analysis and Synthesis of Digital Logic Circuits
 - Input and Output signals can be represented by Boolean Variables, and
 - Function of the Digital Logic Circuits can be represented by Logic Operations, i.e., Boolean Function(s)
 - From a Boolean function, a logic diagram can be constructed using AND, OR, and I

Truth Table

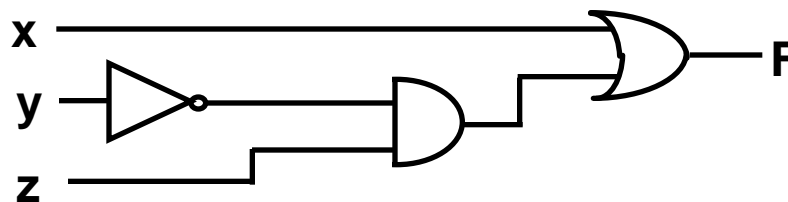
- * The most elementary specification of the function of a Digital Logic Circuit is the Truth Table
 - Table that describes the Output Values for all the combinations of the Input Values, called **MINTERMS**
 - n input variables $\rightarrow 2^n$ minterms

LOGIC CIRCUIT DESIGN



x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$F = x + y'z$$



BASIC IDENTITIES OF BOOLEAN ALGEBRA

[1] $x + 0 = x$	[2] $x \cdot 0 = 0$
[3] $x + 1 = 1$	[4] $x \cdot 1 = x$
[5] $x + x = x$	[6] $x \cdot x = x$
[7] $x + x' = 1$	[8] $x \cdot x' = 0$
[9] $x + y = y + x$	[10] $xy = yx$
[11] $x + (y + z) = (x + y) + z$	[12] $x(yz) = (xy)z$
[13] $x(y + z) = xy + xz$	[14] $x + yz = (x + y)(x + z)$
[15] $(x + y)' = x'y'$	[16] $(xy)' = x' + y'$
[17] $(x')' = x$	

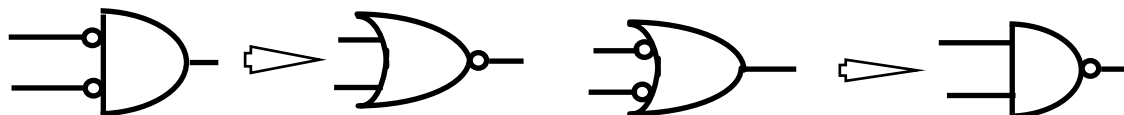
Usefulness of this Table

[15] and [16] : De Morgan's Theorem

- Simplification of the Boolean function
 - Derivation of equivalent Boolean functions to obtain logic diagrams utilizing different logic gates
 - Ordinarily ANDs, ORs, and Inverters
 - But a certain different form of Boolean function may be convenient to obtain circuits with NANDs or NORs
- Applications of De Morgans Theorem

$x'y' = (x + y)'$
I, AND → NOR

$x' + y' = (xy)'$
I, OR → NAND



EQUIVALENT CIRCUITS

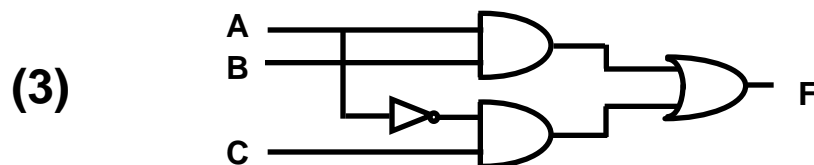
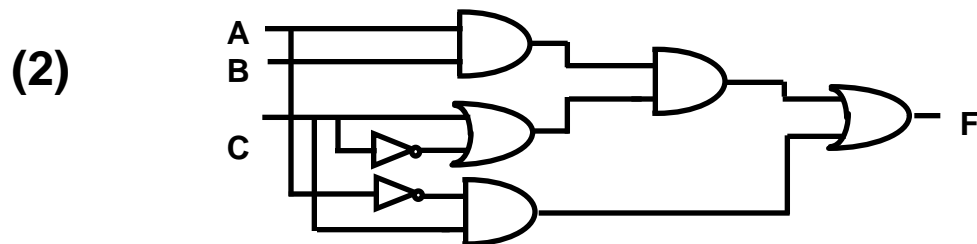
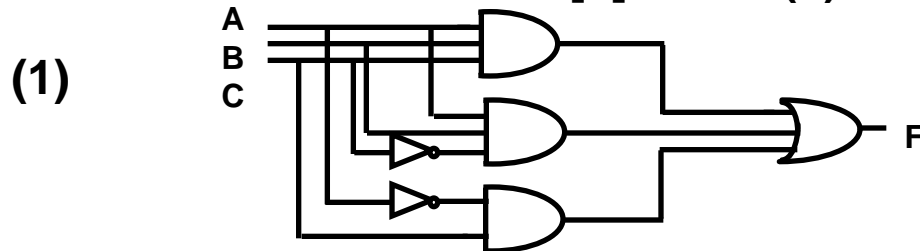
Many different logic diagrams are possible for a given Function

$$F = ABC + ABC' + A'C \quad \dots\dots\dots (1)$$

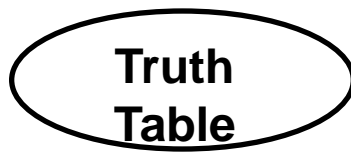
$$= AB(C + C') + A'C \quad [13] \dots\dots (2)$$

$$= AB \cdot 1 + A'C \quad [7]$$

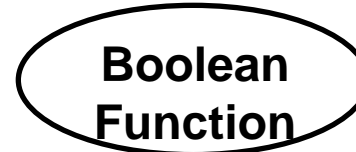
$$= AB + A'C \quad [4] \dots\dots (3)$$



SIMPLIFICATION



Unique

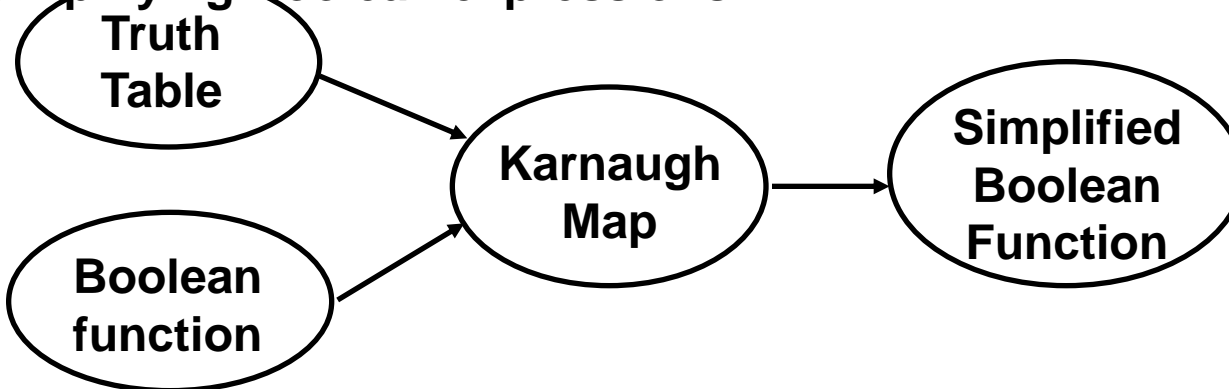


Many different expressions exist

Simplification from Boolean function

- Finding an equivalent expression that is least expensive to implement
- For a simple function, it is possible to obtain a simple expression for low cost implementation
- But, with complex functions, it is a very difficult task

Karnaugh Map (K-map) is a simple procedure for simplifying Boolean expressions.



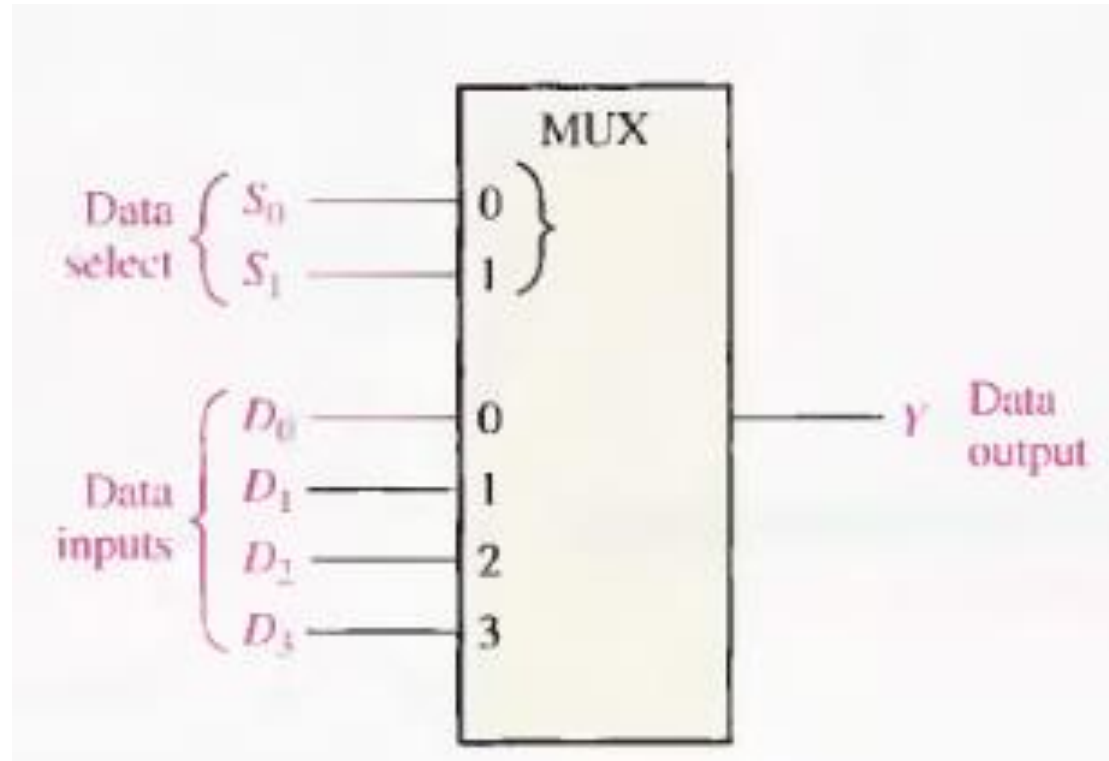
COMBINATIONAL LOGIC CIRCUITS

- **Multiplexer**
- **Encoder**
- **Decoder**
- **Parity Checker**
- **Parity Generator**
- **etc**

Selecting

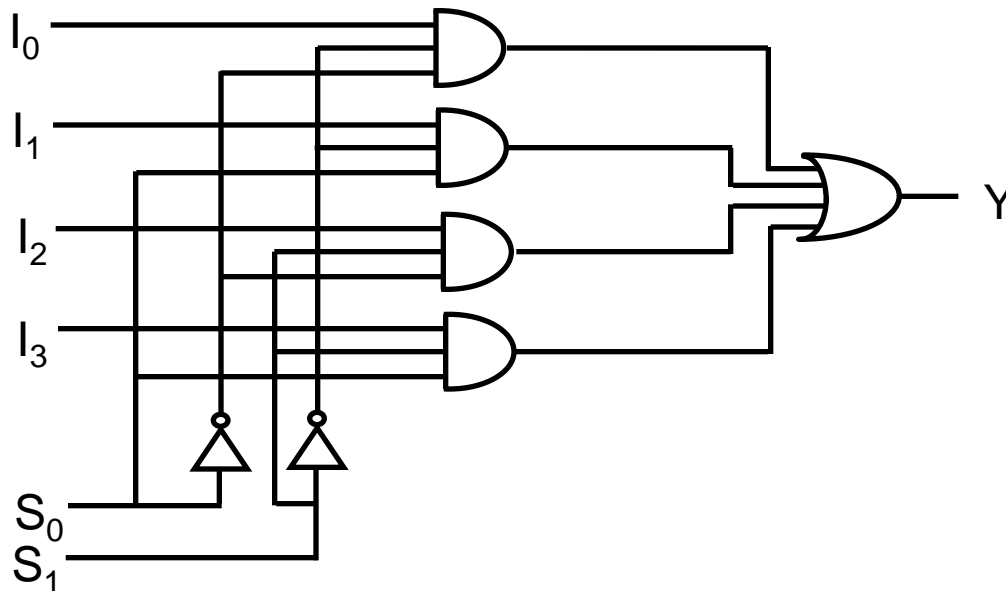
- Selecting of data or information is a critical function in digital systems and computers
- Circuits that perform selecting have:
 - A set of information inputs from which the selection is made
 - A single output
 - A set of control lines for making the selection
- Logic circuits that perform selecting are called *multiplexers*

4-to-1-line Multiplexer



4-to-1-line Multiplexer

Select		Output
S_1	S_0	Y
0	0	I_0
0	1	I_1
1	0	I_2
1	1	I_3

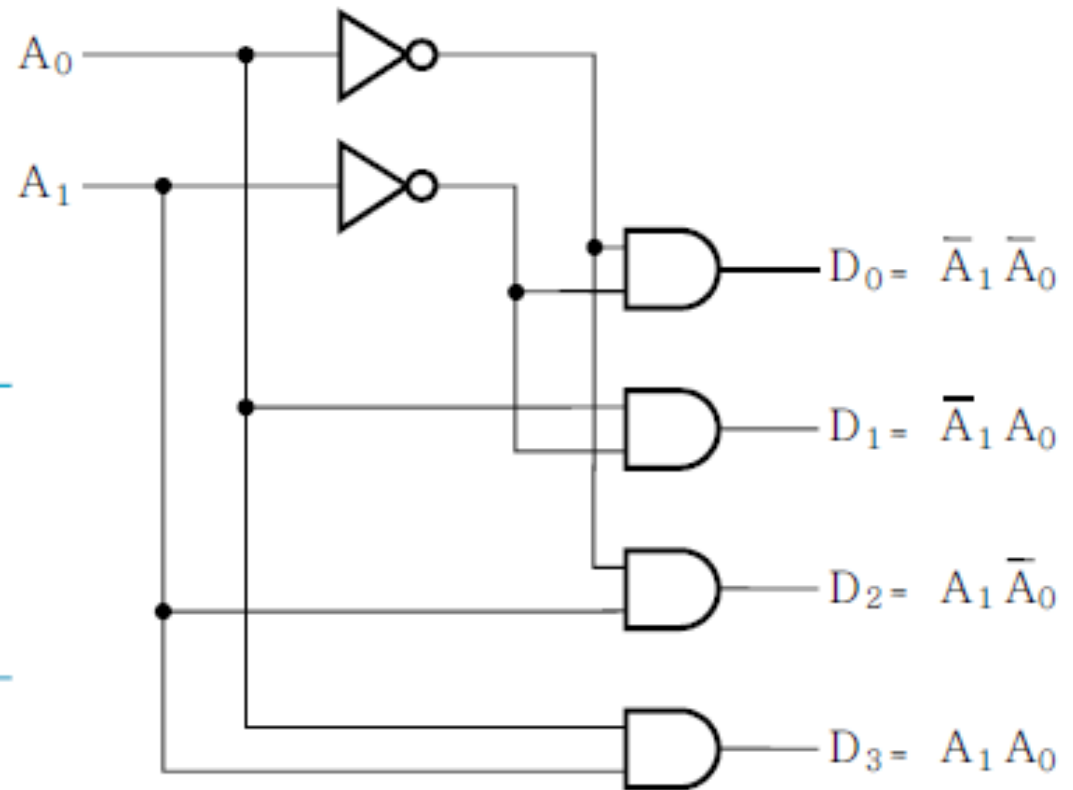


Decoding

- Decoding - the conversion of an n -bit input code to an m -bit output code with $n \leq m \leq 2^n$ such that each valid code word produces a unique output code
- Circuits that perform decoding are called *decoders*

2-to-4 Decoder

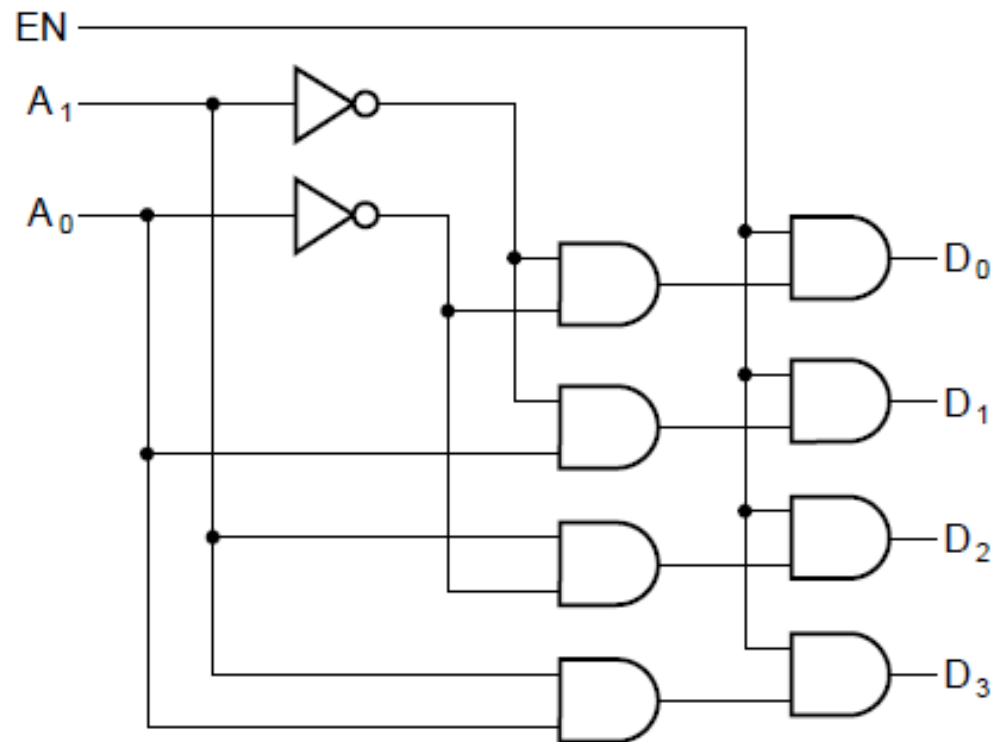
A_1	A_0	D_0	D_1	D_2	D_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1



2-to-4 Decoder with enable

EN	A ₁	A ₀	D ₀	D ₁	D ₂	D ₃
0	X	X	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

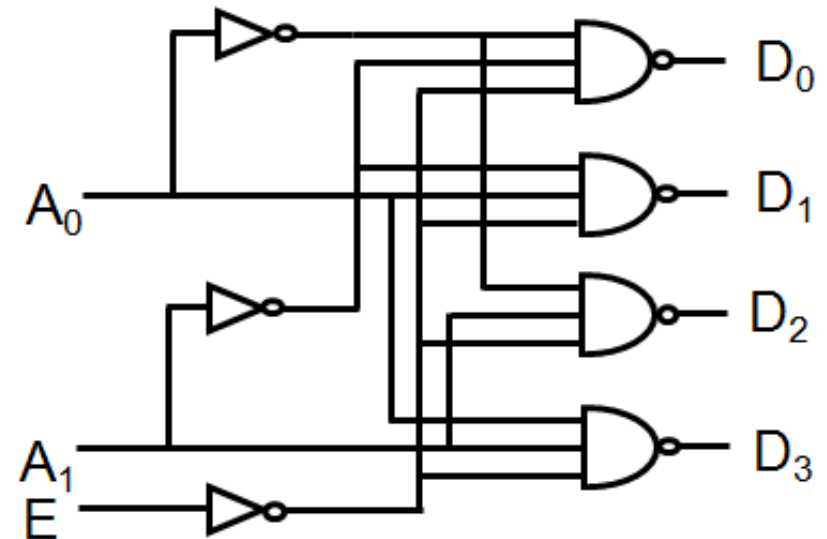
(a)



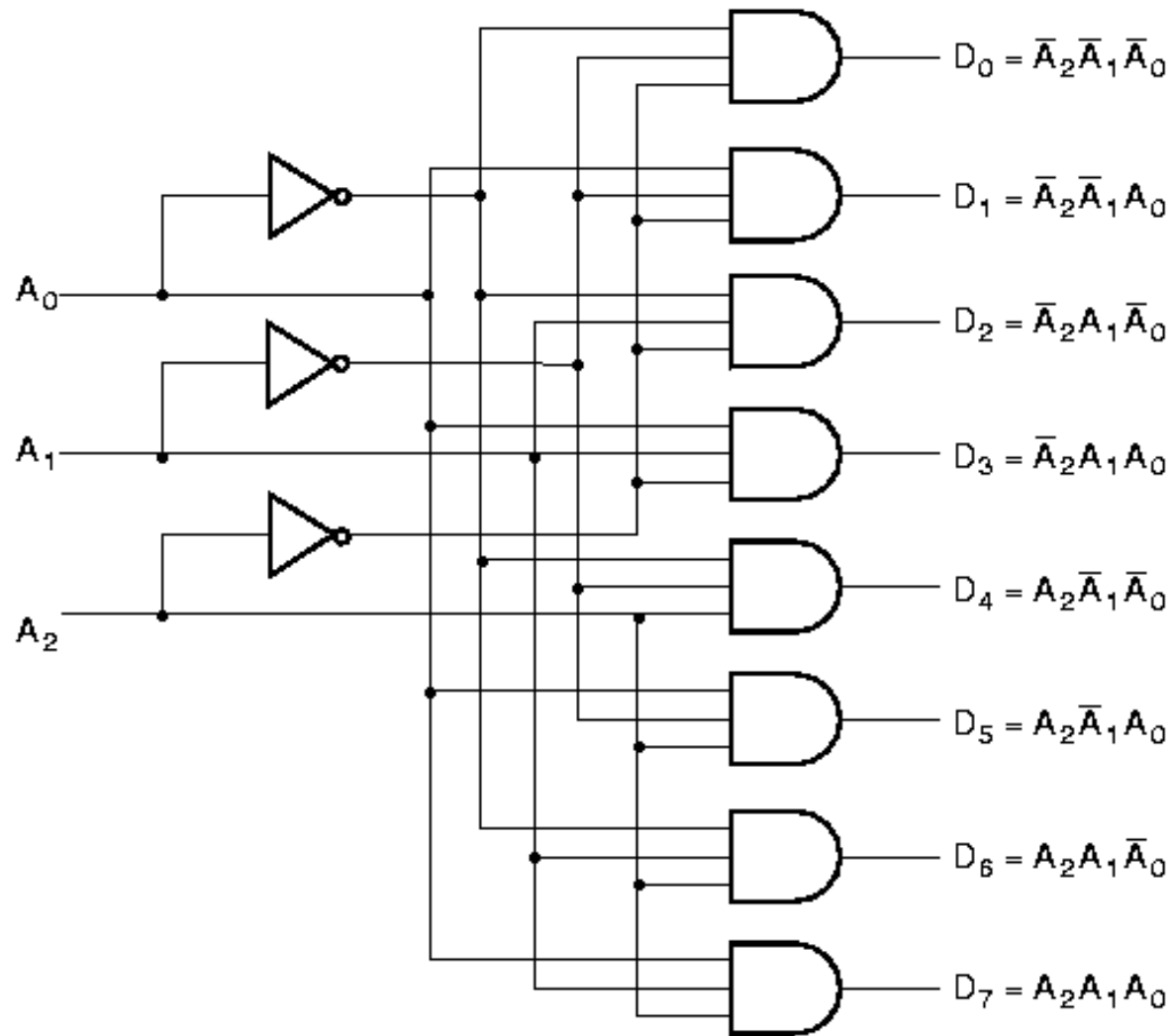
(b)

2-to-4 Decoder with enable

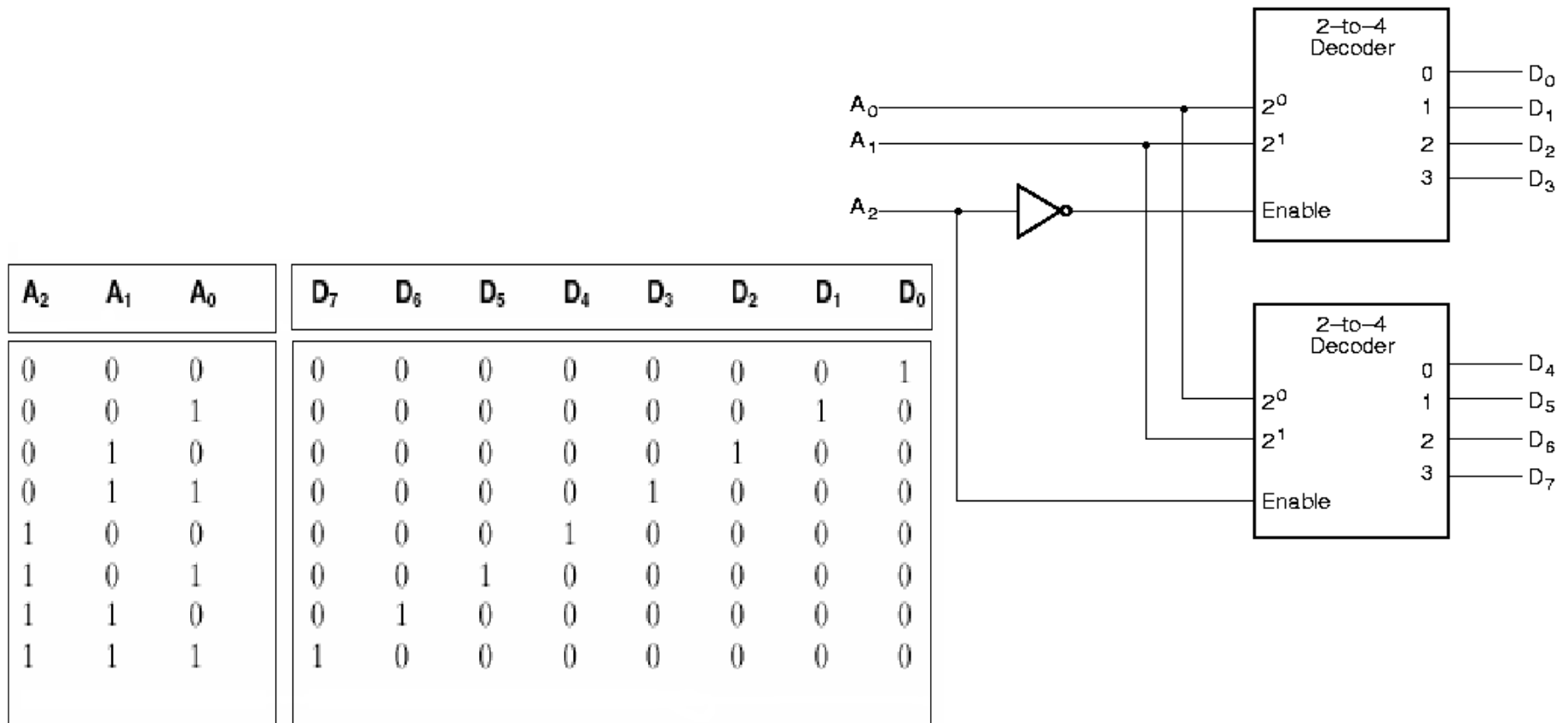
E	A ₁	A ₀	D ₀	D ₁	D ₂	D ₃
0	0	0	0	1	1	1
0	0	1	1	0	1	1
0	1	0	1	1	0	1
0	1	1	1	1	1	0
1	d	d	1	1	1	1



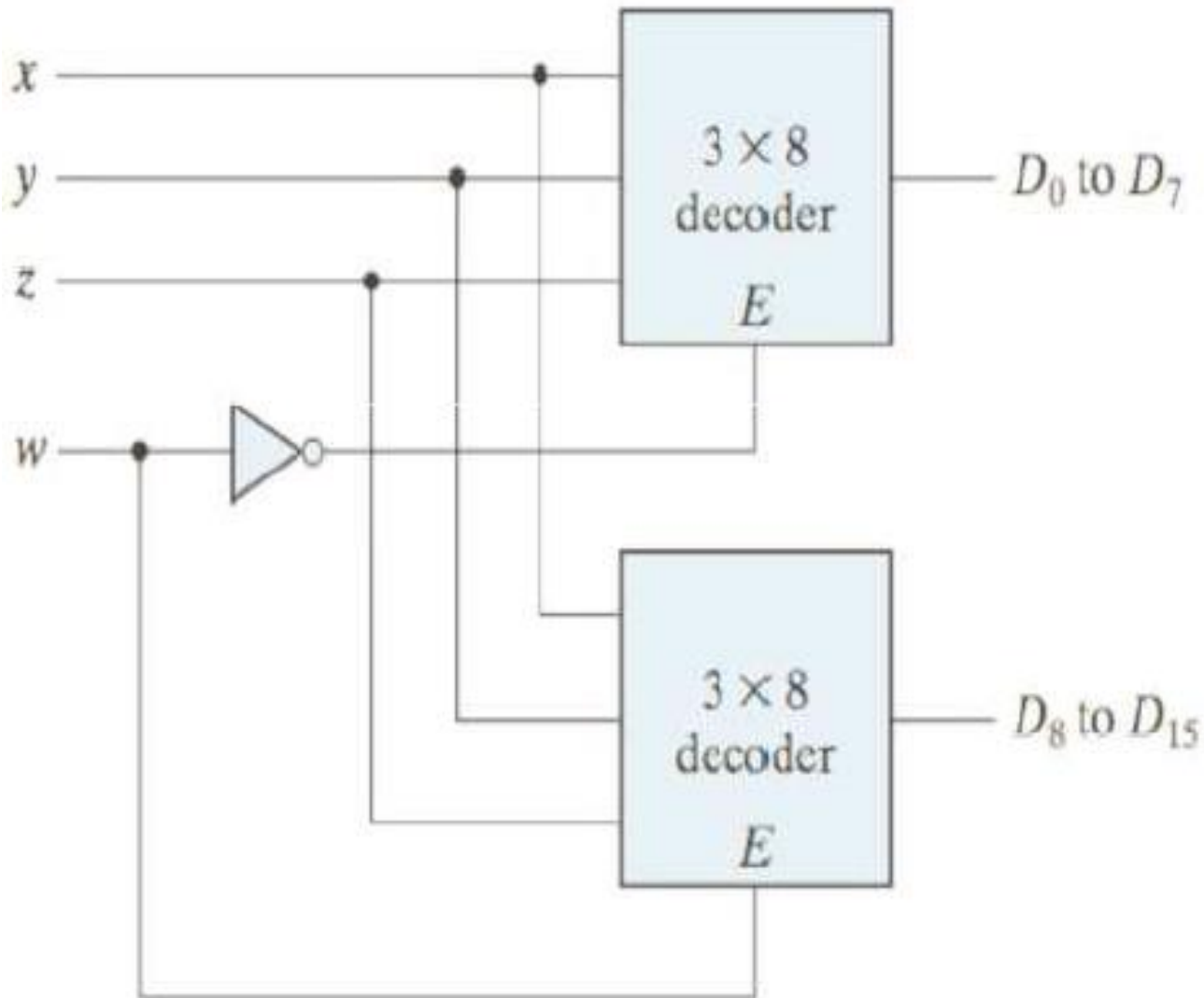
3-to-8-line Decoder



Decoder Expansion



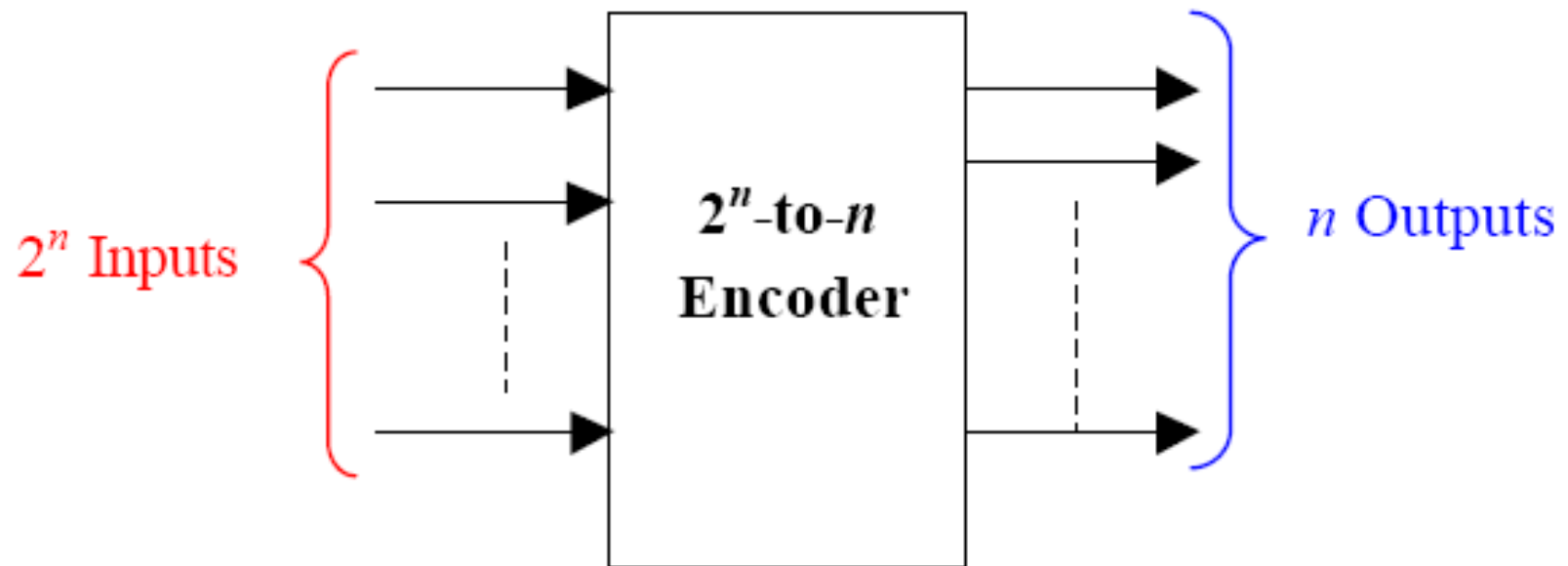
Decoder Expansion



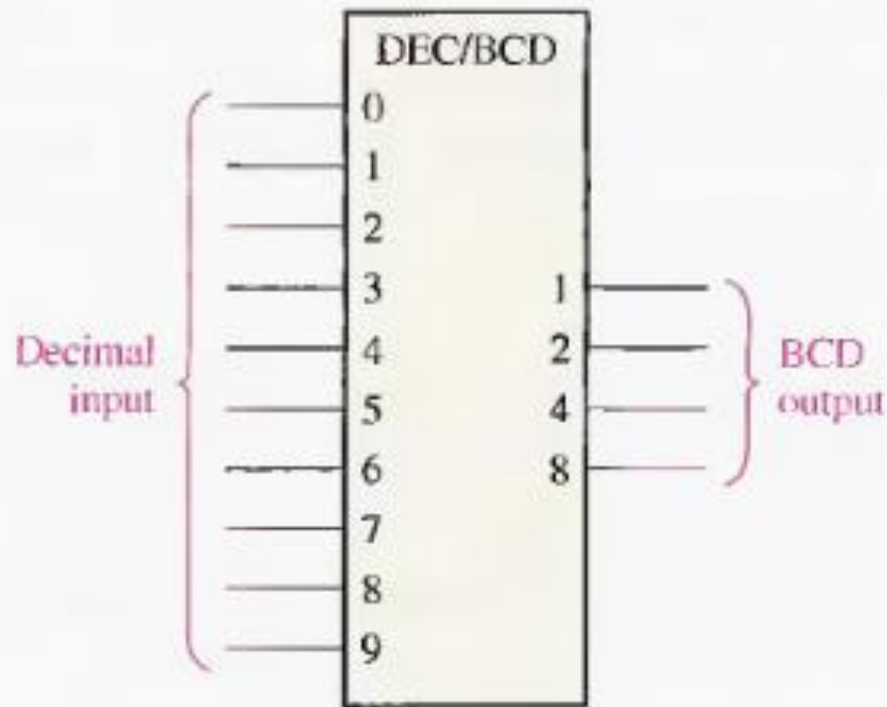
Encoding

- Encoding - the opposite of decoding - the conversion of an m -bit input code to a n -bit output code with $n \leq m \leq 2^n$ such that each valid code word produces a unique output code
- Circuits that perform encoding are called *encoders*
- An encoder has 2^n (or fewer) input lines and n output lines which generate the binary code corresponding to the input values

Encoder

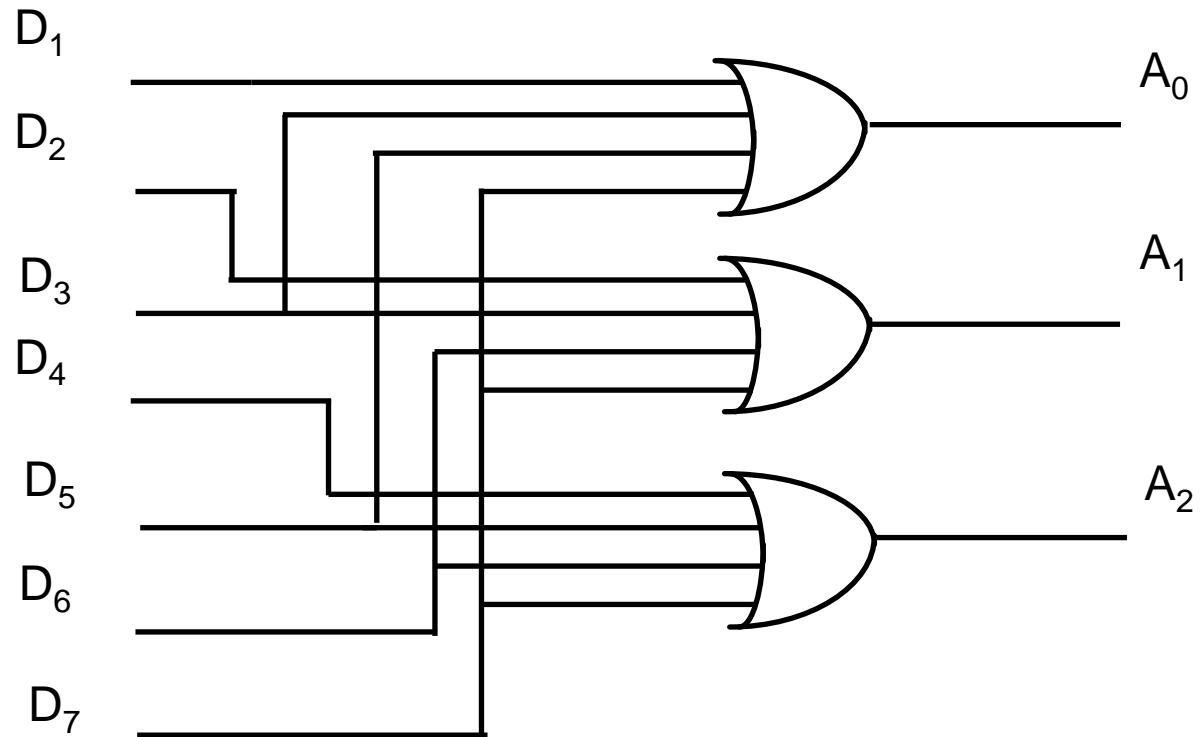


Decimal-to-BCD Encoder



DECIMAL DIGIT	BCD CODE			
	A ₃	A ₂	A ₁	A ₀
0	0	0	0	0
1	0	0	0	1
2	0	0	1	0
3	0	0	1	1
4	0	1	0	0
5	0	1	0	1
6	0	1	1	0
7	0	1	1	1
8	1	0	0	0
9	1	0	0	1

Octal-to-Binary Encoder



Octal-to-Binary Encoder

Inputs								Outputs		
D_0	D_1	D_2	D_3	D_4	D_5	D_6	D_7	x	y	z
1	0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1
0	0	1	0	0	0	0	0	0	1	0
0	0	0	1	0	0	0	0	0	1	1
0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	1	0	0	1	0	1
0	0	0	0	0	0	1	0	1	1	0
0	0	0	0	0	0	0	1	1	1	1

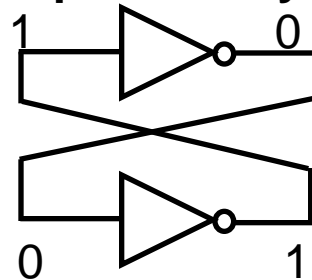
$$x = D_4 + D_5 + D_6 + D_7 \quad y = D_2 + D_3 + D_6 + D_7$$

$$z = D_1 + D_3 + D_5 + D_7$$

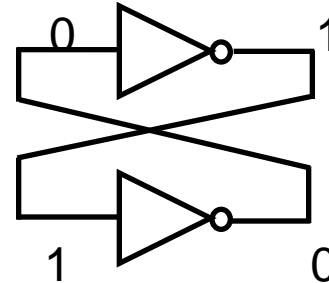
FLIP FLOPS

Characteristics

- 2 stable states
- Memory capability
- Operation is specified by a Characteristic Table

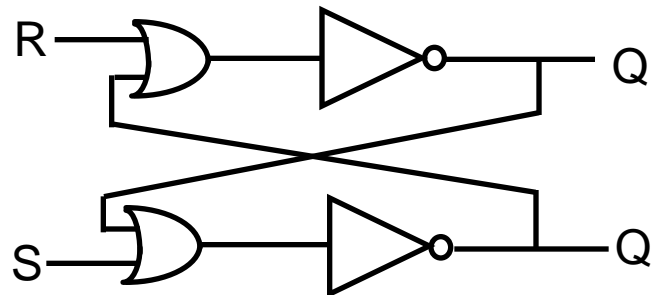


0-state



1-state

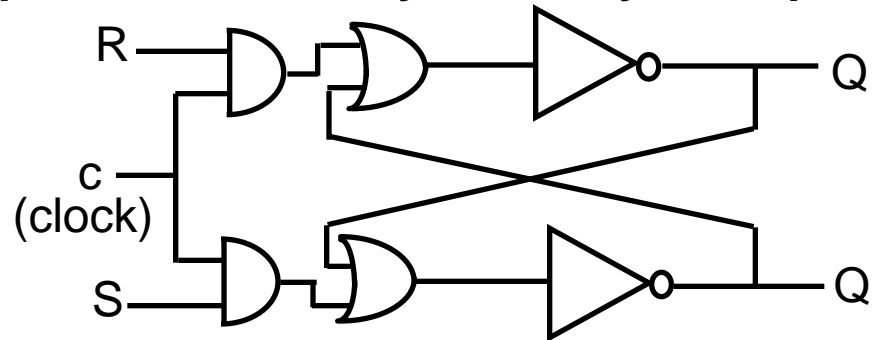
In order to be used in the computer circuits, state of the flip flop should have input terminals and output terminals so that it can be set to a certain state, and its state can be read externally.



S	R	Q(t+1)
0	0	Q(t)
0	1	0
1	0	1
1	1	indeterminate (forbidden)

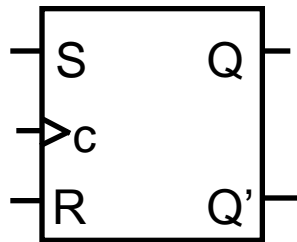
CLOCKED FLIP FLOPS

In a large digital system with many flip flops, operations of individual flip flops are required to be synchronized to a clock pulse. Otherwise, the operations of the system may be unpredictable.

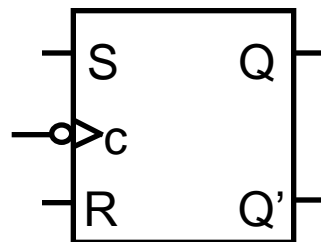


Clock pulse allows the flip flop to change state only when there is a clock pulse appearing at the c terminal.

We call above flip flop a Clocked RS Latch, and symbolically as

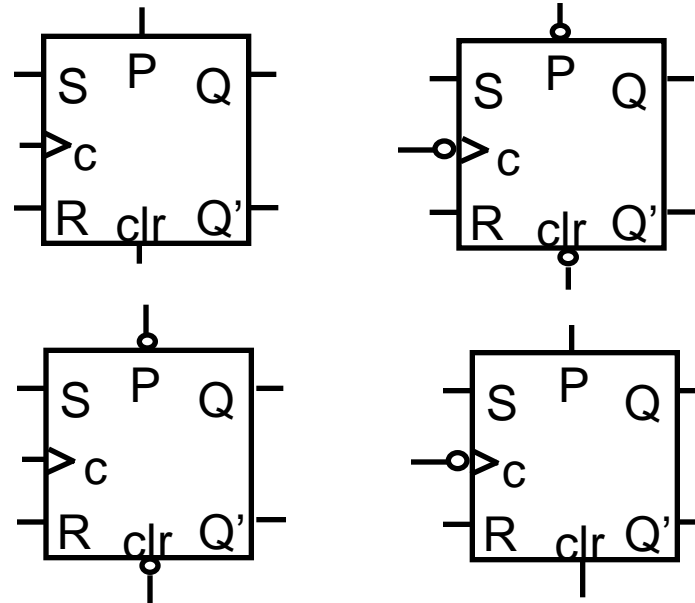
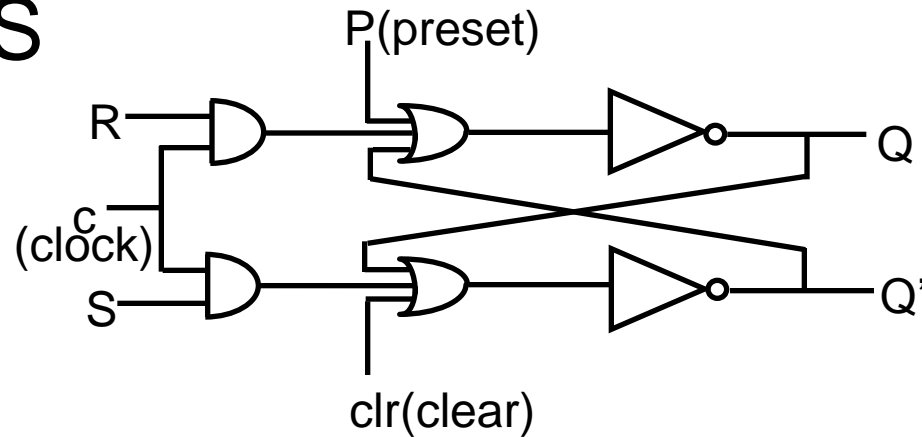


operates when
clock is high



operates when
clock is low

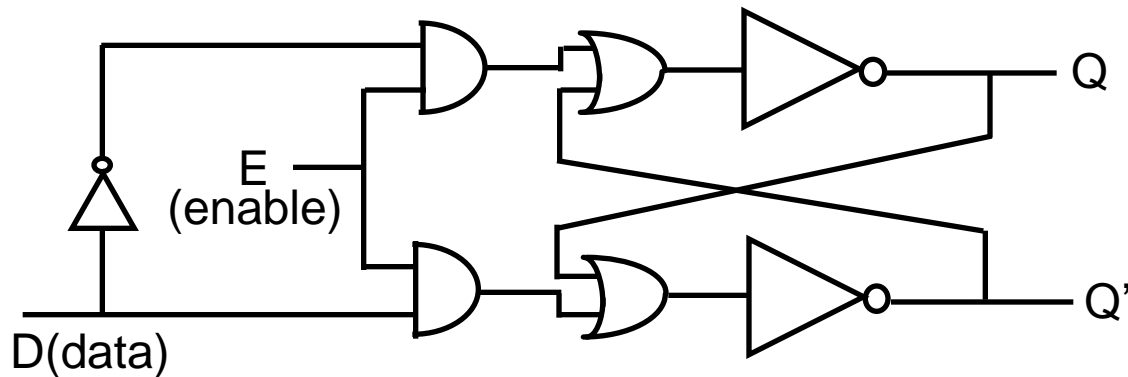
RS-LATCH WITH PRESET AND CLEAR INPUTS



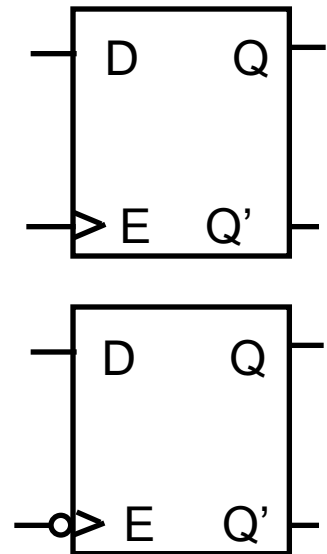
D-LATCH

D-Latch

Forbidden input values are forced not to occur by using an inverter between the inputs



D	Q(t+1)
0	0
1	1

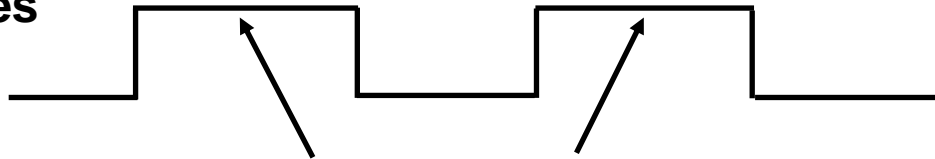


EDGE-TRIGGERED FLIP FLOPS

Characteristics

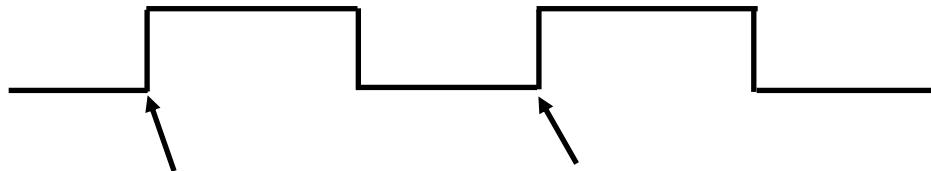
- State transition occurs at the rising edge or falling edge of the clock pulse

Latches



respond to the input only during these periods

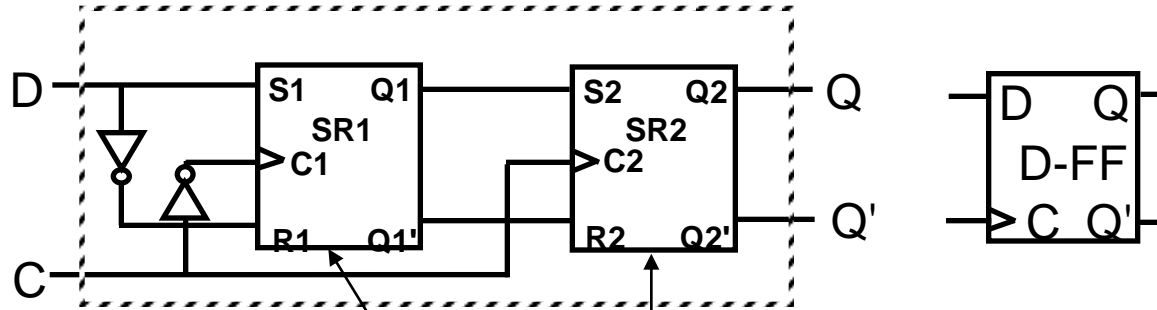
Edge-triggered Flip Flops (positive)



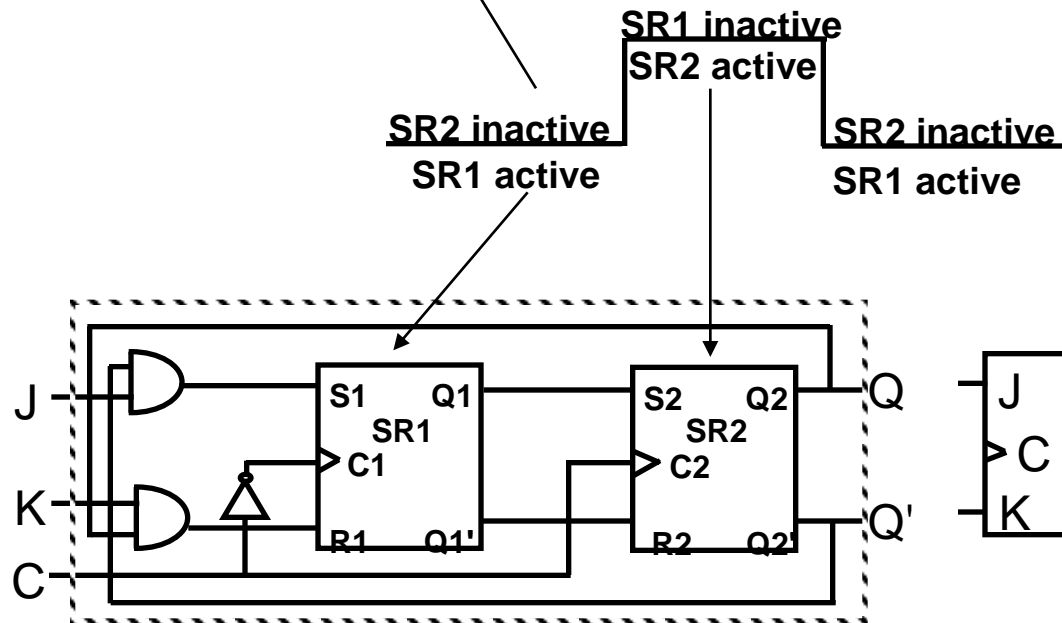
respond to the input only at this time

POSITIVE EDGE-TRIGGERED

D-Flip Flop



JK-Flip Flop

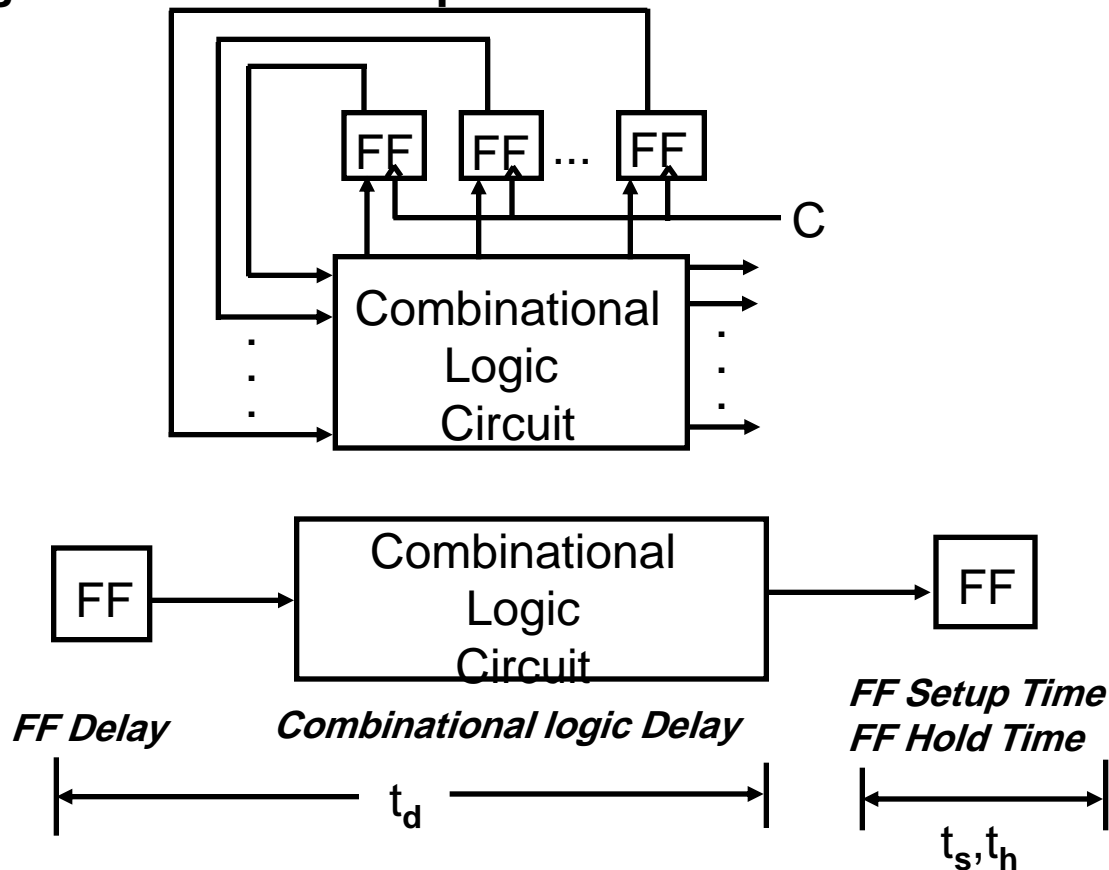


T-Flip Flop: JK-Flip Flop whose J and K inputs are tied together to make T input. Toggles whenever there is a pulse on T input.

CLOCK PERIOD

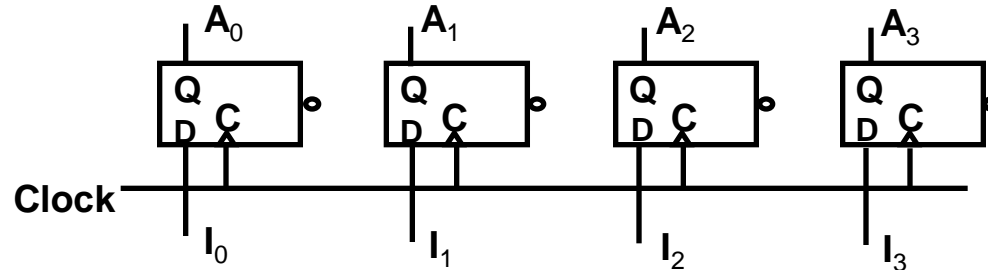
Clock period determines how fast the digital circuit operates.
How can we determine the clock period ?

Usually, digital circuits are sequential circuits which has some flip flops

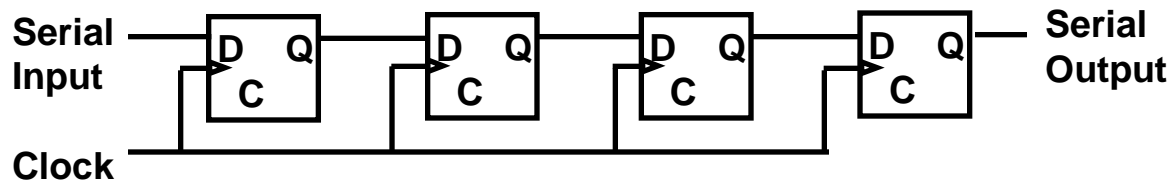


$$\text{clock period } T = t_d + t_s + t_h$$

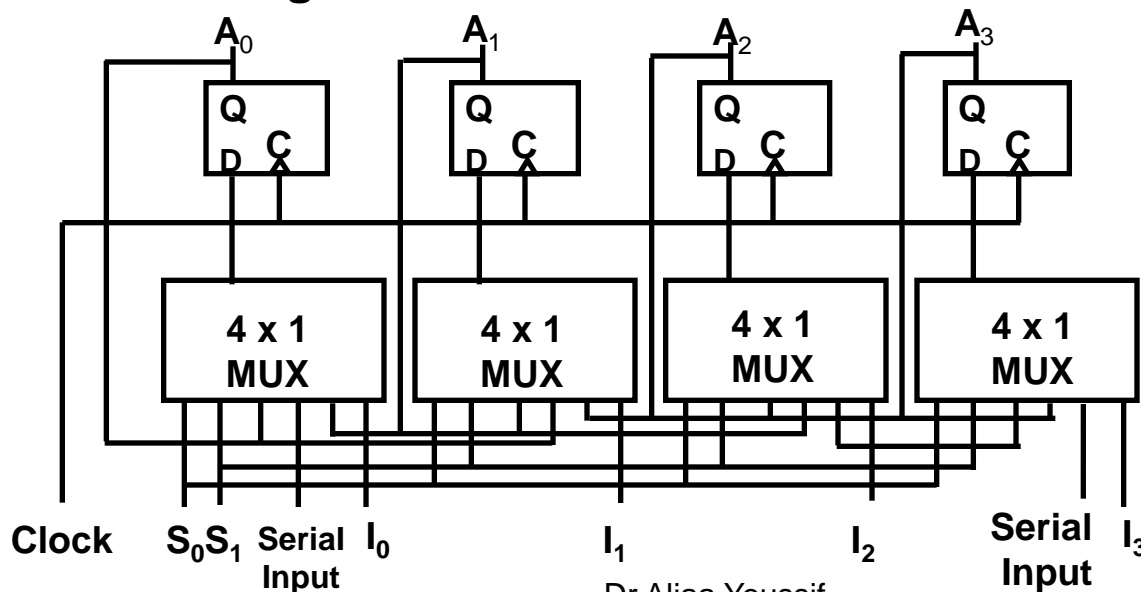
SEQUENTIAL CIRCUITS - Registers



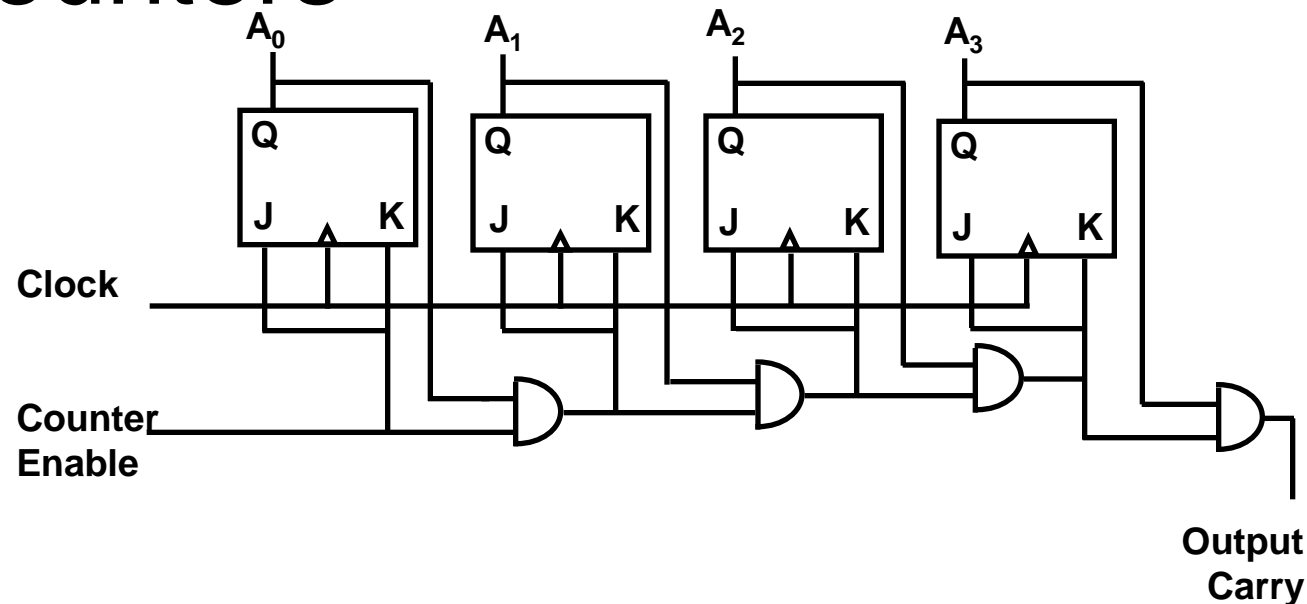
Shift Registers



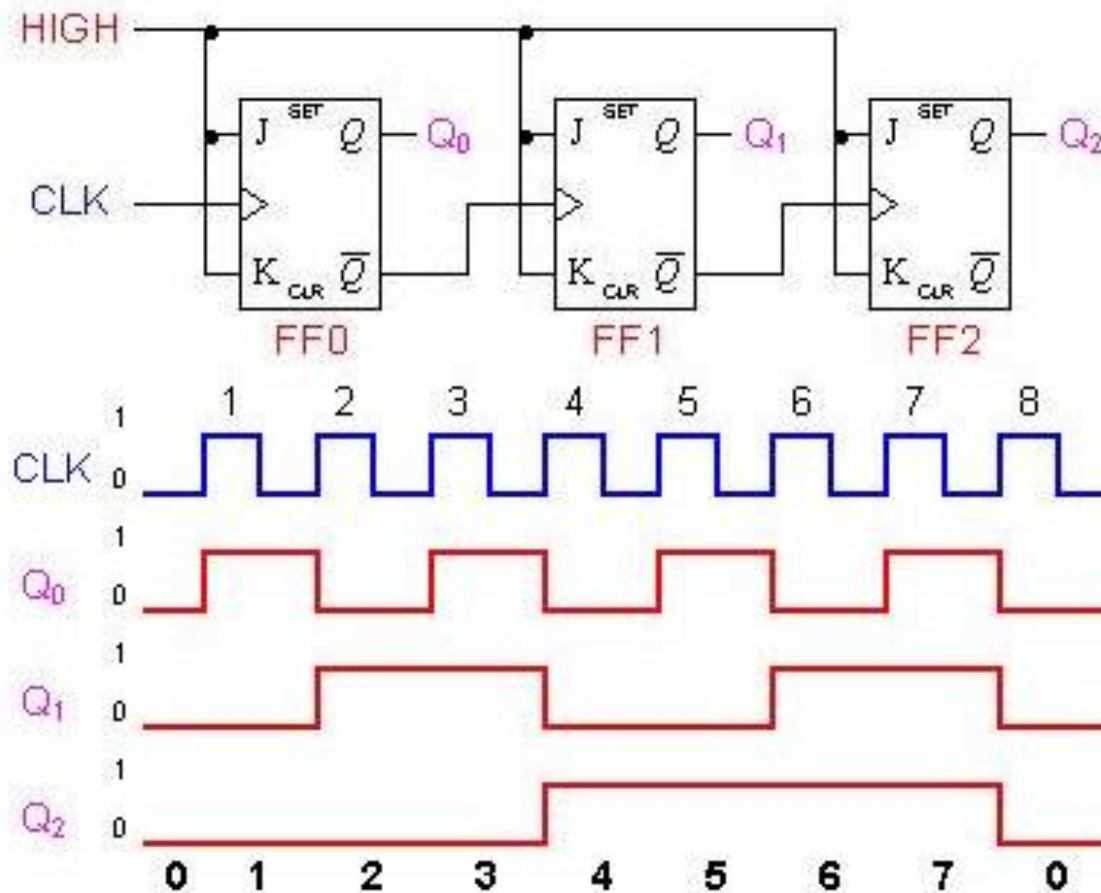
Bidirectional Shift Register with Parallel Load



SEQUENTIAL CIRCUITS - Counters



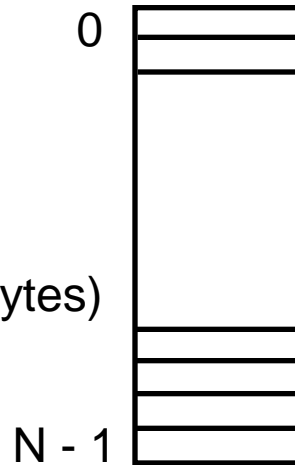
SEQUENTIAL CIRCUITS - Counters



MEMORY COMPONENTS

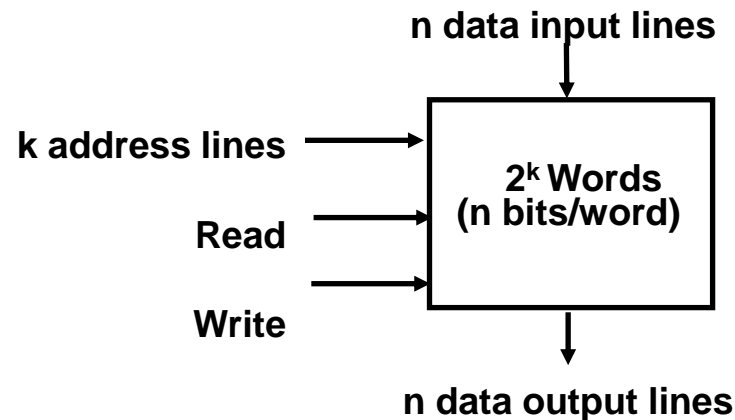
Logical Organization

words
(byte, or n bytes)

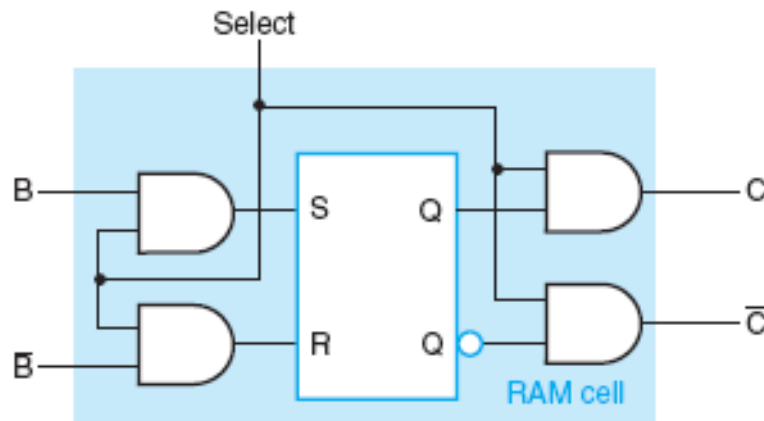


Random Access Memory

- Each word has a unique address
- Access to a word requires the same time independent of the location of the word
- Organization



MEMORY COMPONENTS

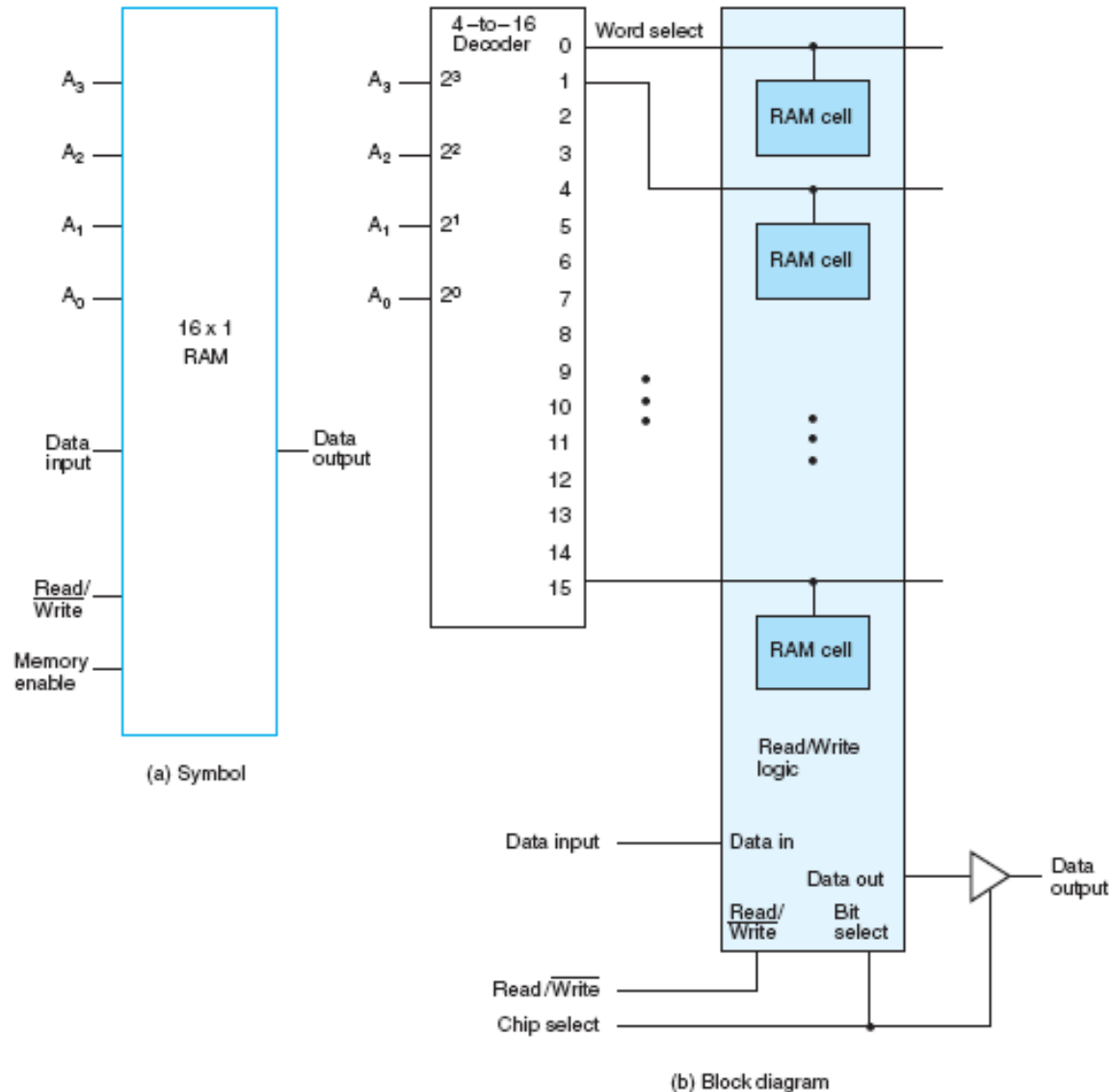


Static RAM Cell

Memory address		Memory contents
Binary	Decimal	
0000000000	0	10110101 01011100
0000000001	1	10101011 10001001
0000000010	2	00001101 01000110
.	.	.
.	.	.
.	.	.
.	.	.
.	.	.
1111111101	1021	10011101 00010101
1111111110	1022	00001101 00011110
1111111111	1023	11011110 00100100

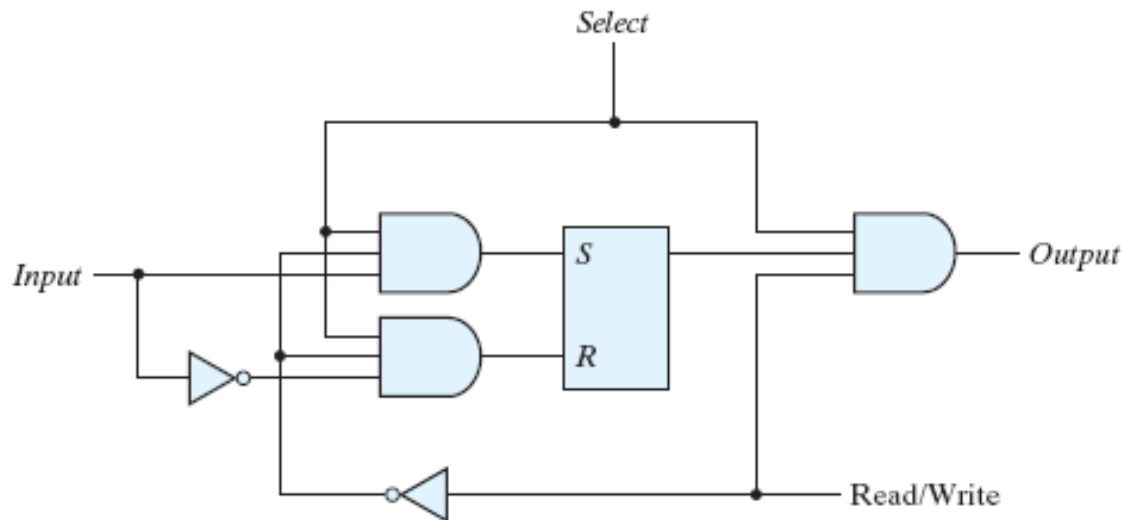
Contents of a 1024×16 Memory

MEMORY COMPONENTS

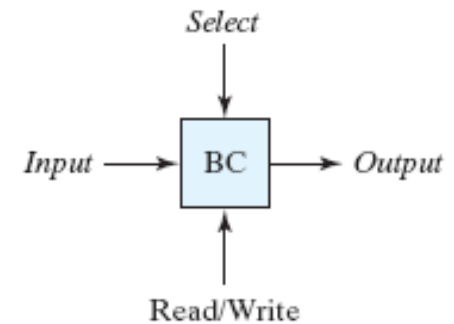


16-Word by 1-Bit RAM Chip

MEMORY COMPONENTS

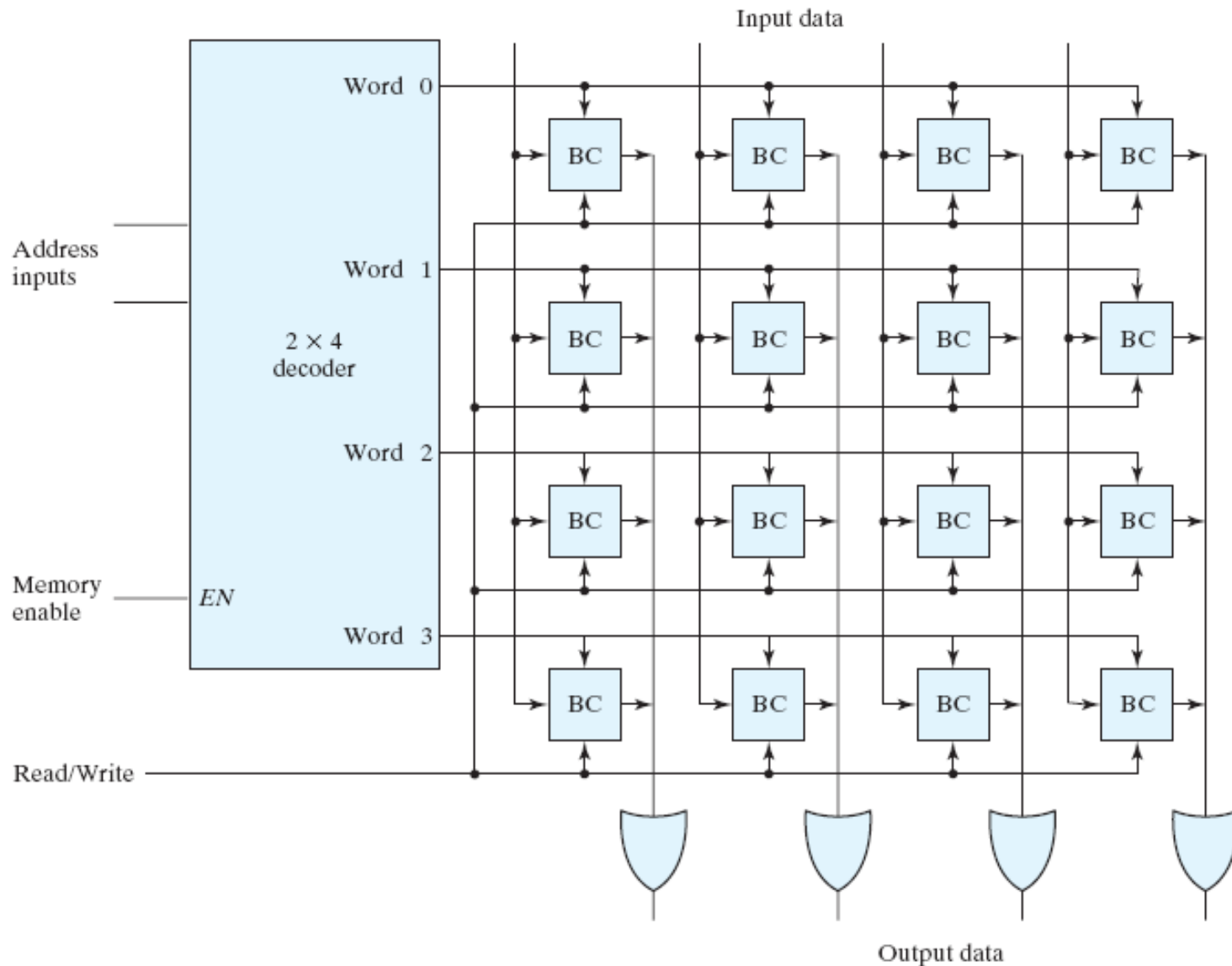


(a) Logic diagram



(b) Block diagram

MEMORY COMPONENTS

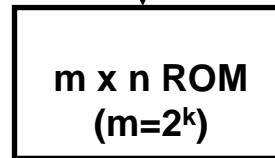


READ ONLY MEMORY(ROM)

Characteristics

- Perform read operation only, write operation is not possible
- Information stored in a ROM is made permanent during production, and cannot be changed
- Organization

k address input lines



n data output lines

Information on the data output line depends only on the information on the address input lines.

--> Combinational Logic Circuit

$$\begin{aligned} X_0 &= A'B' + B'C \\ X_1 &= A'B'C + A'BC' \\ X_2 &= BC + AB'C' \\ X_3 &= A'BC' + AB' \\ X_4 &= AB \end{aligned}$$

Canonical minterms

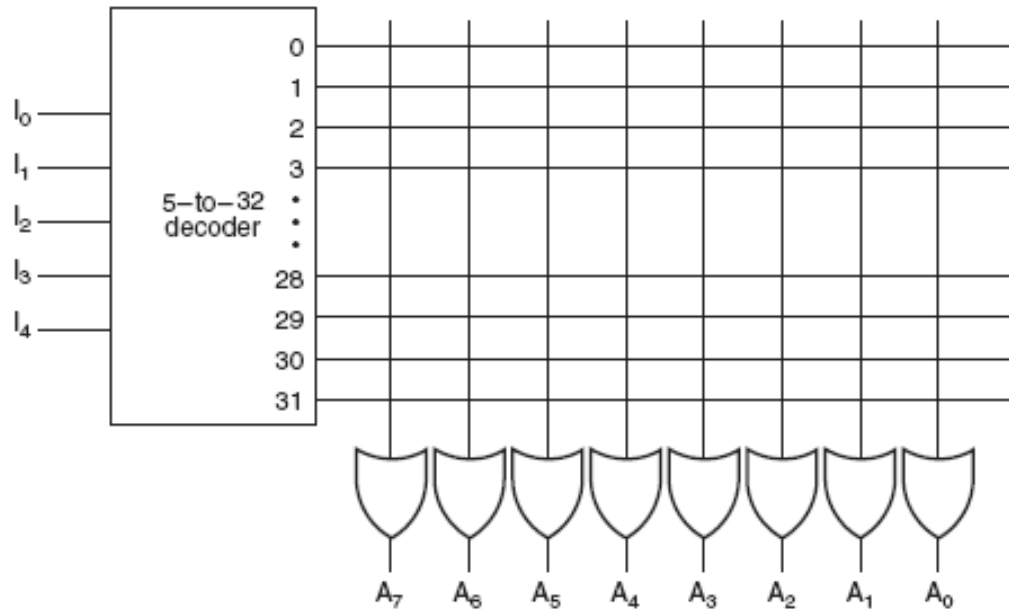
$$\begin{aligned} X_0 &= A'B'C' + A'B'C + AB'C \\ X_1 &= A'B'C + A'BC' \\ X_2 &= A'BC + AB'C' + ABC \\ X_3 &= A'BC' + AB'C' + AB'C \\ X_4 &= ABC' + ABC \end{aligned}$$

address

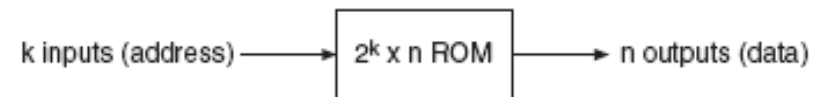
Output

ABC	X_0	X_1	X_2	X_3	X_4
000	1	0	0	0	0
001	1	1	0	0	0
010	0	1	0	1	0
011	0	0	1	0	0
100	0	0	1	1	0
101	1	0	0	1	0
110	0	0	0	0	1
111	0	0	1	0	1

READ ONLY MEMORY(ROM)



Internal Logic of a 32×8 ROM



Block Diagram of ROM

TYPES OF ROM

ROM

- Store information (function) during production
- Mask is used in the production process
- Unalterable
- Low cost for large quantity production --> used in the final products

PROM (Programmable ROM)

- Store info electrically using PROM programmer at the user's site
- Unalterable
- Higher cost than ROM -> used in the system development phase
-> Can be used in small quantity system

EPROM (Erasable PROM)

- Store info electrically using PROM programmer at the user's site
- Stored info is erasable (alterable) using UV light (electrically in some devices) and rewriteable
- Higher cost than PROM but reusable --> used in the system development phase. Not used in the system production due to reusability