

TUGAS 2
PRAKTIKUM ANALISIS ALGORITMA



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PROGRAM STUDI S-1 TEKNIK INFORMATIKA
FAKULTAS MATEMATIKA DAN ILMU PENGETAHUAN ALAM
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1. Mencari nilai max

Algoritma

```
procedure CariMaks(input  $x_1, x_2, \dots, x_n$ : integer, output maks: integer)
{ Mencari elemen terbesar dari sekumpulan elemen larik integer  $x_1, x_2, \dots, x_n$ . Elemen
  terbesar akan disimpan di dalam maks
  Input:  $x_1, x_2, \dots, x_n$ 
  Output: maks (nilai terbesar)
}
```

Deklarasi

i : integer

Algoritma

```
maks  $\leftarrow$   $x_1$ 
 $i \leftarrow 2$ 
while  $i \leq n$  do
  if  $x_i > \text{maks}$  then
    maks  $\leftarrow x_i$ 
  endif
   $i \leftarrow i + 1$ 
endwhile
{ $i > n$ }
```

Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    float arr[100];

    cout << "Masukkan banyak angka : ";
    cin >> n;
    cout << endl;

    for(i = 0; i < n; ++i)
    {
        cout << "masukkan angka ke- " << i + 1 << " : ";
        cin >> arr[i];
    }
}
```

```
}  
  
for(i = 1; i < n; ++i)  
{  
    if(arr[0] < arr[i])  
        arr[0] = arr[i];  
}  
cout << "angka terbesar adalah = " << arr[0];  
  
return 0;  
}
```

Kompleksitas waktu

$\text{maks} \leftarrow x_1$	1 kali
$i \leftarrow 2$	1 kali
$\text{maks} \leftarrow x_i$	n kali
$i \leftarrow i + 1$	n kali

$$T(n) = 1 + 1 + n + n = 2n + 2$$

2. Sequential Search

Algoritma

```
procedure SequentialSearch(input  $x_1, x_2, \dots, x_n$  : integer, y : integer, output idx : integer)
{   Mencari y di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat y ditemukan
    diisi ke dalam idx. Jika y tidak ditemukan, maka idx diisi dengan 0.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output: idx
}
```

Deklarasi

i : integer

found : boolean { bernilai true jika y ditemukan atau false jika y tidak ditemukan }

Algoritma

i \leftarrow 1

found \leftarrow false

while (i \leq n) and (not found) do

if $x_i = y$ then

 found \leftarrow true

else

 i \leftarrow i + 1

endif

endwhile

Code

```
#include <iostream>
using namespace std;

int main() {
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    bool found = false;
    int idx;
    while ((i < n) && (!found)){
        if (x[i] == y)
            found = true;
    }
```

```

else
    i++;
}
if (found)
    idx = i+1;
else
    idx = 0;

cout << "Yang dicari berada di urutan : " << idx << endl;

return 0;
}

```

Kompleksitas waktu

Best Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 = 4$$

Average Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	$\frac{1}{2} n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + \frac{1}{2} n + 1 + 1 = \frac{1}{2} n + 4$$

Worst Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	n kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\max}(n) = 1 + 1 + n + 1 + 1 = n + 4$$

3. Binary Search

Algoritma

```
procedure BinarySearch(input  $x_1, x_2, \dots, x_n$  : integer, x : integer, output : idx : integer)
{ Mencari y di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat y ditemukan diisi ke dalam idx. Jika y tidak ditemukan maka dx diisi dengan 0.
  Input:  $x_1, x_2, \dots, x_n$ 
  Output: idx
}
Deklarasi
  i, j, mid : integer
  found : Boolean
Algoritma
  i  $\leftarrow$  1
  j  $\leftarrow$  n
  found  $\leftarrow$  false
  while (not found) and (i  $\leq$  j) do
    mid  $\leftarrow$  (i + j) div 2
    if  $x_{\text{mid}} = y$  then
      found  $\leftarrow$  true
    else
      if  $x_{\text{mid}} < y$  then {mencari di bagian kanan}
        i  $\leftarrow$  mid + 1
      else {mencari di bagian kiri}
        j  $\leftarrow$  mid - 1
      endif
    endif
  endwhile
  {found or i > j}

  If found then
    Idx  $\leftarrow$  mid
  else
    Idx  $\leftarrow$  0
  endif
  {i < n or found}

  If found then {y ditemukan}
    idx  $\leftarrow$  i
  else
    idx  $\leftarrow$  0 {y tidak ditemukan}
  endif
```

Code

```
#include <iostream>
```

```

using namespace std;

int main() {
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    int j = n-1;
    bool found = false;
    int idx;
    int mid;
    while ((i <= j) && (!found)){
        mid = (i + j)/2;
        if (x[mid] == y)
            found = true;
        else{
            if (x[mid] < y)
                i = mid + 1;
            else
                j = mid - 1;
        }
    }

    if (found)
        idx = mid+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}

```

Kompleksitas waktu

Best Case :

$i \leftarrow 1$

1 kali

$j \leftarrow n$

1 kali

found \leftarrow false	1 kali
mid \leftarrow (i + j) div2	1 kali
found \leftarrow true	1 kali
Idx \leftarrow mid	1 kali

$$T_{min}(n) = 1 + 1 + 1 + 1 + 1 + 1 = 6$$

Average Case :

i \leftarrow 1	1 kali
j \leftarrow n	1 kali
found \leftarrow false	1 kali
mid \leftarrow (i + j) div2	$\frac{1}{2}n + 1$ kali
i \leftarrow mid + 1 or j \leftarrow mid - 1	$\frac{1}{2}n$ kali
found \leftarrow true	1 kali
Idx \leftarrow mid	1 kali

$$T_{avg}(n) = 1 + 1 + 1 + \frac{1}{2}n + 1 + \frac{1}{2}n + 1 + 1 = n + 6$$

Worst Case :

i \leftarrow 1	1 kali
j \leftarrow n	1 kali
found \leftarrow false	1 kali
mid \leftarrow (i + j) div2	$n + 1$ kali
i \leftarrow mid + 1 or j \leftarrow mid - 1	n kali
found \leftarrow true	1 kali
Idx \leftarrow mid	1 kali

$$T_{max}(n) = 1 + 1 + 1 + n + 1 + n + 1 + 1 = 2n + 6$$

4. Insertion Sort

Algoritma

```
procedure InsertionSort(input/output  $x_1, x_2, \dots, x_n$  : integer)
{
    Mengurutkan elemen-elemen  $x_1, x_2, \dots, x_n$  dengan metode insertion sort.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output:  $x_1, x_2, \dots, x_n$  (sudah terurut menaik)
}
Deklarasi
    i, j, insert : integer
Algoritma
    for i  $\leftarrow$  2 to n do
        insert  $\leftarrow$   $x_i$ 
        j  $\leftarrow$  i
        while (j < i) and ( $x[j-i]$  > insert) do
             $x[j] \leftarrow x[j-1]$ 
            j  $\leftarrow$  j-1
        endwhile
         $x[j] =$  insert
    endfor
    {i < n or found}

    If found then {y ditemukan}
```

Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int insert;
    int j;
```

```

for (int i = 1; i < n; i++)
{
    insert = x[i];
    j = i-1;
    while ((j >= 0) && (x[j] > insert))
    {
        x[j+1] = x[j];
        j--;
    }
    x[j+1] = insert;
}

cout << "Data setelah di Sorting : ";
for (int i = 0; i < n; i++)
    cout << x[i] << " ";

return 0;
}

```

Kompleksitas waktu

Best Case :

For i \leftarrow 2 to n do	1 kali
insert \leftarrow x _i	n kali
j \leftarrow i	n kali
x[j] = insert	n kali

$$T_{min}(n) = 1 + n + n + n = 3n + 1$$

Average Case :

For i \leftarrow 2 to n do	1 kali
insert \leftarrow x _i	n kali
j \leftarrow I	n kali
x[j] \leftarrow x[j-1]	n * ½ n kali
j \leftarrow j-1	n * ½ n kali
x[j] = insert	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2} n^2 + \frac{1}{2} n^2 + n = n^2 + 3n + 1$$

Worst Case :

For i \leftarrow 2 to n do	1 kali
insert \leftarrow x _i	n kali
j \leftarrow i	n kali
x[j] \leftarrow x[j-1]	n * n kali
j \leftarrow j-1	n * n kali

$x[j] = \text{insert}$ n kali

$$T_{max}(n) = 1 + n + n + n^2 + n^2 + n = 2n^2 + 3n + 1$$

5. Selection Sort

Algoritma

```
procedure SelectionSort(input/output  $x_1, x_2, \dots, x_n$  : integer)
{ Mengurutkan elemen-elemen  $x_1, x_2, \dots, x_n$  dengan metode selection sort.
  Input:  $x_1, x_2, \dots, x_n$ 
  Output:  $x_1, x_2, \dots, x_n$  (sudah terurut menaik)
}
```

Deklarasi

$i, j, \text{imaks}, \text{temp}$: integer

Algoritma

```
for  $i \leftarrow n$  downto 2 do {pass sebanyak  $n-1$  kali}
  imaks  $\leftarrow 1$ 
  for  $j \leftarrow 2$  to  $i$  do
    if  $x_j > x_{\text{imaks}}$  then
      imaks  $\leftarrow j$ 
    endif
  endfor
  {pertukarkan  $x_{\text{imaks}}$  dengan  $x_i$ }
  temp  $\leftarrow x_i$ 
   $x_i \leftarrow x_{\text{imaks}}$ 
   $x_{\text{imaks}} \leftarrow \text{temp}$ 
endfor
```

Code

```
#include <iostream>
using namespace std;

int main(){
  int n;
  int x[10];
  cout << "Masukkan Jumlah Data : ";
  cin >> n;
  for (int i = 0; i < n; i++){
    cout << "Masukkan Data ke - " << i+1 << " : ";
    cin >> x[i];
  }
  cout << "Data Sebelum di Sorting : ";
  for (int i = 0; i < n; i++)
    cout << x[i] << " ";
  cout << endl;

  int imaks;
  int temp;
  for (int i = n-1; i >= 1; i--){
```

```

    imaks = 0;
    for (int j = 1; j <= i; j++){
        if (x[j] > x[imaks])
            imaks = j;
    }
    temp = x[i];
    x[i] = x[imaks];
    x[imaks] = temp;
}

cout << "Data setelah di Sorting : ";
for (int i = 0; i < n; i++)
    cout << x[i] << " ";

return 0;
}

```

Kompleksitas waktu

Best Case :

for i ← n downto 2 do	1 kali
imaks ← 1	n kali
for j ← 2 to i do	n kali
imaks ← j	n*1 kali
temp ← xi	n kali
xi ← ximaks	n kali
ximaks ← temp	n kali

$$T_{min}(n) = 1 + n + n + n * 1 + n + n + n = 6n + 1$$

Average Case :

for i ← n downto 2 do	1 kali
imaks ← 1	nkali
for j ← 2 to i do	n kali
imaks ← j	n * ½ n kali
temp ← xi	n kali
xi ← ximaks	n kali
ximaks ← temp	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2}n^2 + n + n + n = \frac{1}{2}n^2 + 5n + 1$$

Worst Case :

for i ← n downto 2 do	1 kali
imaks ← 1	n kali
for j ← 2 to i do	n kali
imaks ← j	n * n kali

temp ← xi	n kali
xi ← ximaks	n kali
ximaks ← temp	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n + n + n = n^2 + 5n + 1$$