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## PA04 - RE-ROOTING

### Time complexity-

For the re-rooting section of this code, the time complexity is  $O[n]$ .

Reason:

The re-rooting occurs on each cutline node. The re-rooting is conducted with first the left child, and then the right child. However, even the leaf nodes have a recursive call in order to conduct re-rooting, but nothing happens in this call as they do not have any children. This means that every cutline node (besides the root node) will have 2 recursive calls, and every leaf node will have one recursive call. The total number of cutline nodes (excepting the root node) is  $((n-1) - 1)$ , which is equal to  $(n - 2)$ . Since each of these cutline nodes have two re-roots, we end up with  $(2n - 4)$  recursive calls. Each leaf node will add one additional recursive call that does nothing, so this adds  $n$  additional calls. Hence the total number of calls is  $(3n - 4)$ . This number is directly proportional to  $n$ . Hence, the overall time complexity is  $O[n]$ .

### Space complexity-

For the re-rooting section of this code, the space complexity is  $O[\log(n)]$ .

Reason:

For the re-rooting of the pre-existing binary tree, no additional memory is required. Only the pointers are moved, and calculations are made. While doing this, the tree is traversed in a post order fashion. Hence, in the worst case scenario, the space complexity will be  $O[h]$ , where  $h$  is the height of the binary tree. This height will be equal to  $\log(n)$ .