

Efficient Symbolic Execution for Software Testing

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Joint work with:

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Symbolic Execution

- Automatically explore program paths
 - *Execute program on “symbolic” input values*
 - *“Fork” execution at each branch*
 - *Record branching conditions*
- Constraint solver
 - *Decides path feasibility*
 - *Generates test cases for paths and bugs*

Symbolic Execution

- (Very brief) history
 - *Test generation by SE in 70s* [King '75] [Boyer et al. '75]
 - *SAT / SMT solvers lead to boom in 2000s* [Godefroid et al. '05][Cadar et al. '06]
- Many successful tools
 - *KLEE, SAGE, PEX, SPF, CREST, Cloud9, S2E, ...*
- Specific advantages
 - *No false positives, useful partial results*
 - *Reduces need for modeling*

Outline

- Symbolic Execution for Testing
- State Merging – Fighting Path Explosion
- Interpreted High-Level Code

$pc = \text{true}$

$x = X$

$r = 0$

```
1 int proc(int x) {  
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3     int r = 0  
5  
6     if (x > 8) {  
7         r = x - 7  
8     }  
9  
10    if (x < 5) {  
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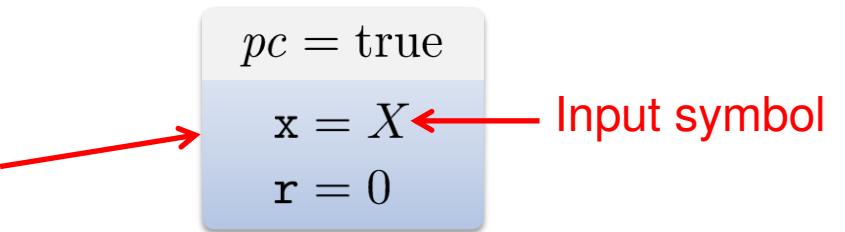
Symbolic
program state



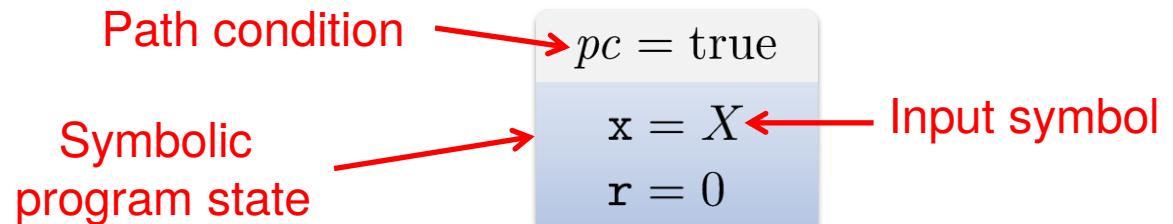
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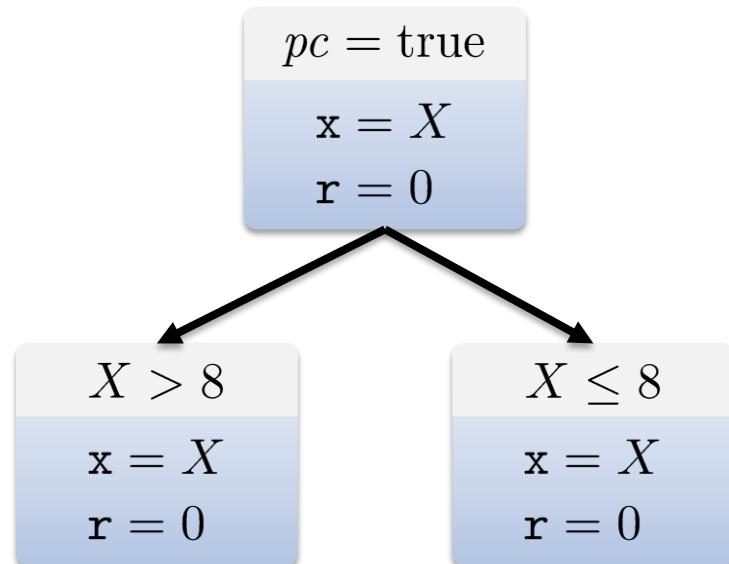
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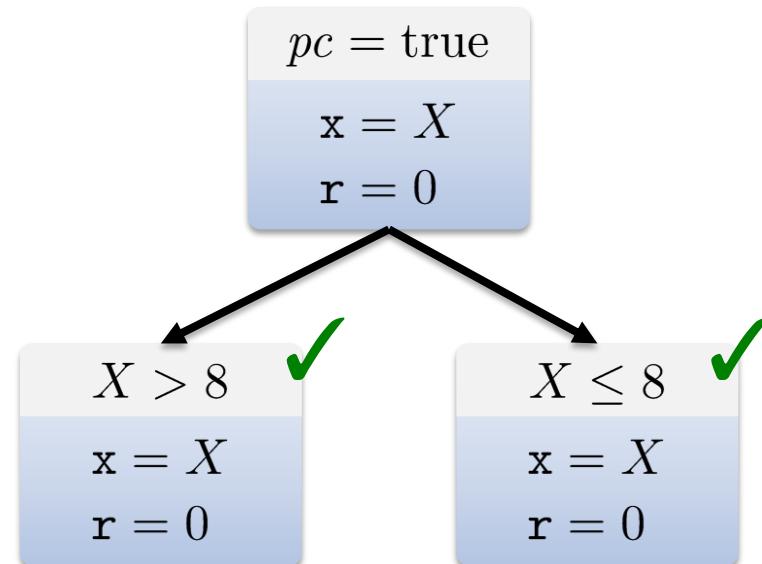


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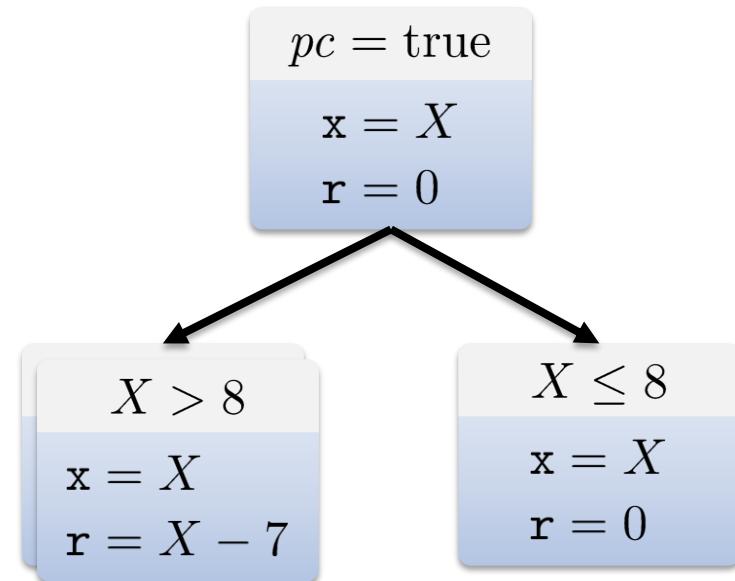
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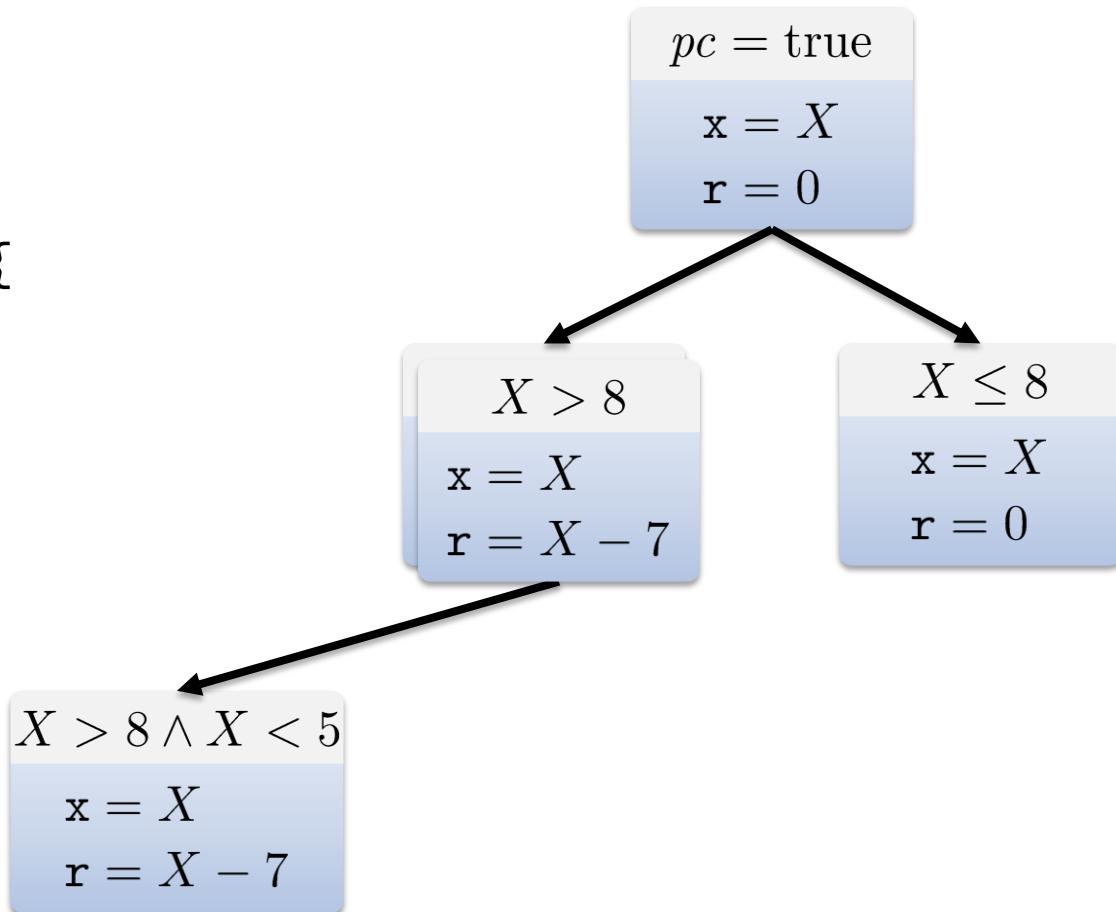
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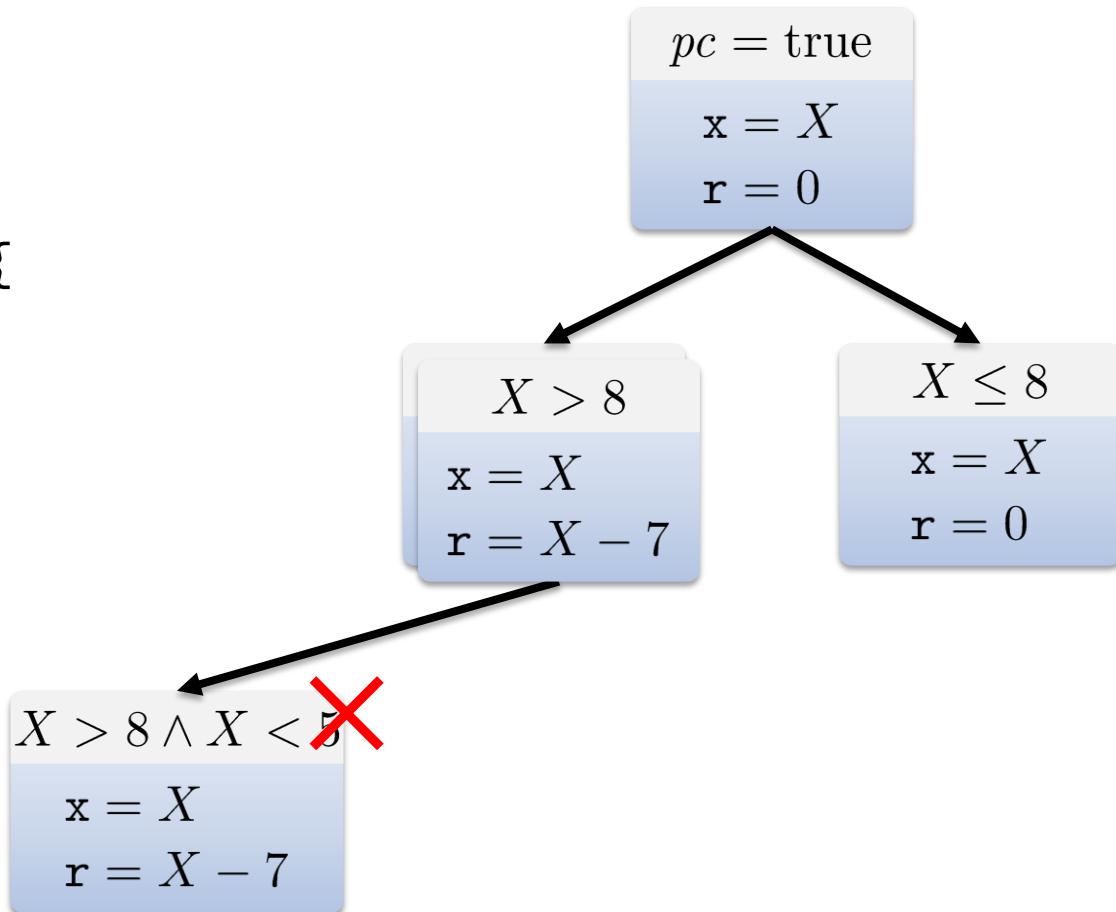
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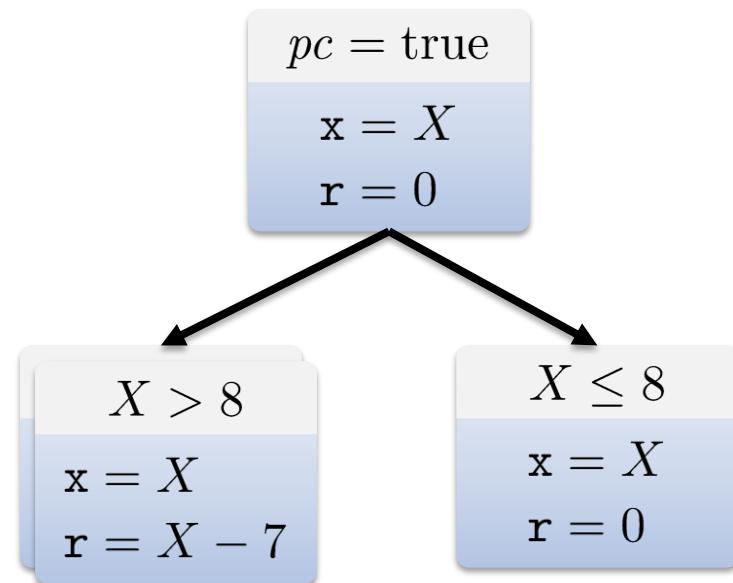
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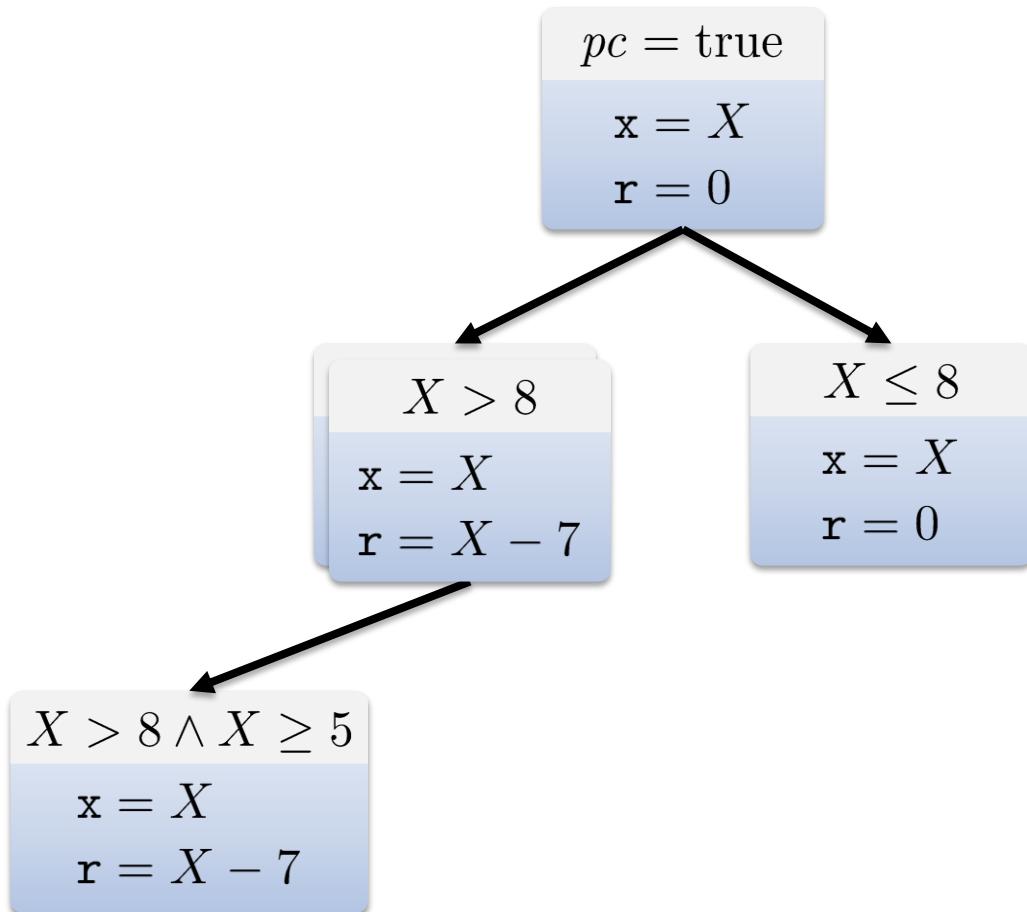
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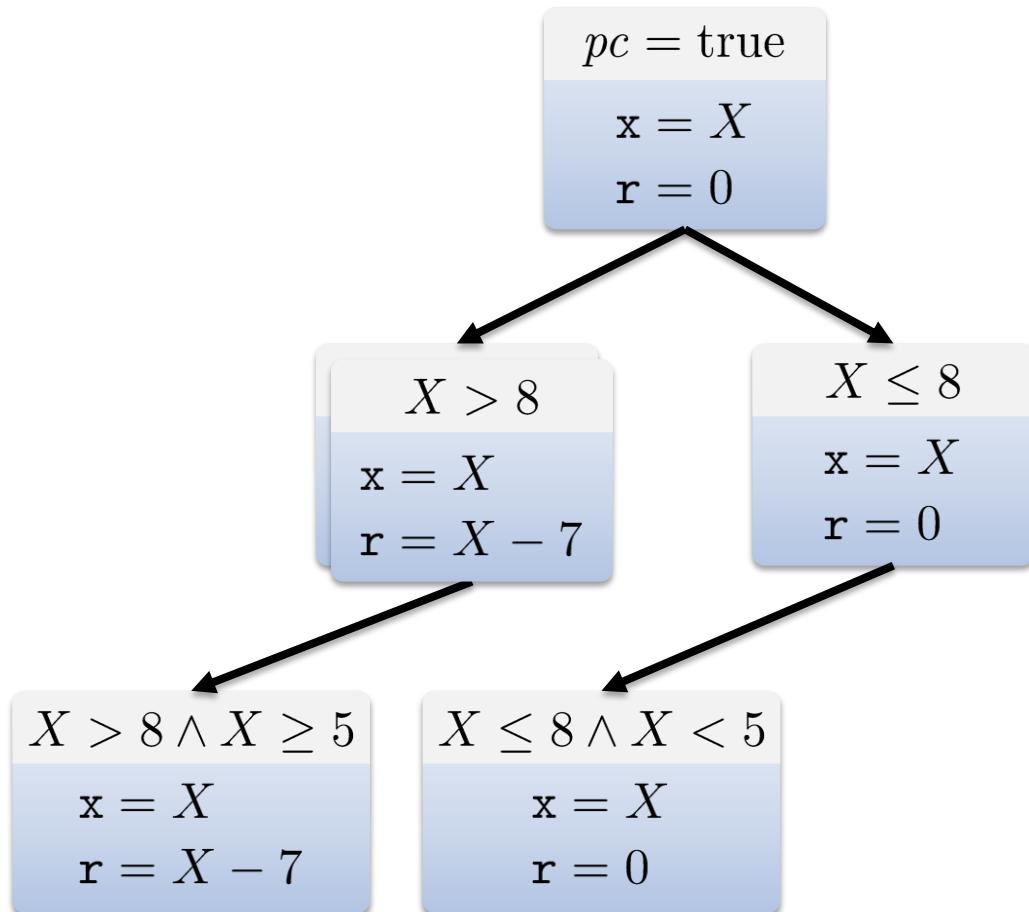
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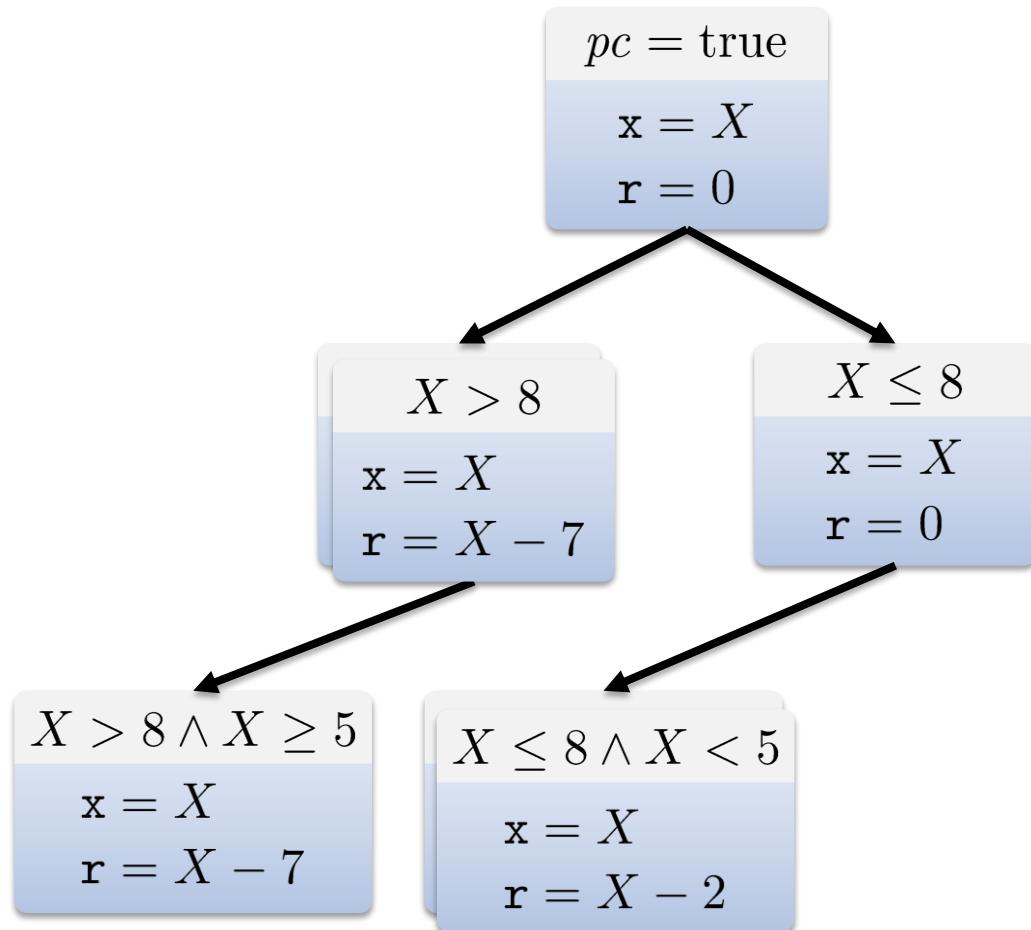
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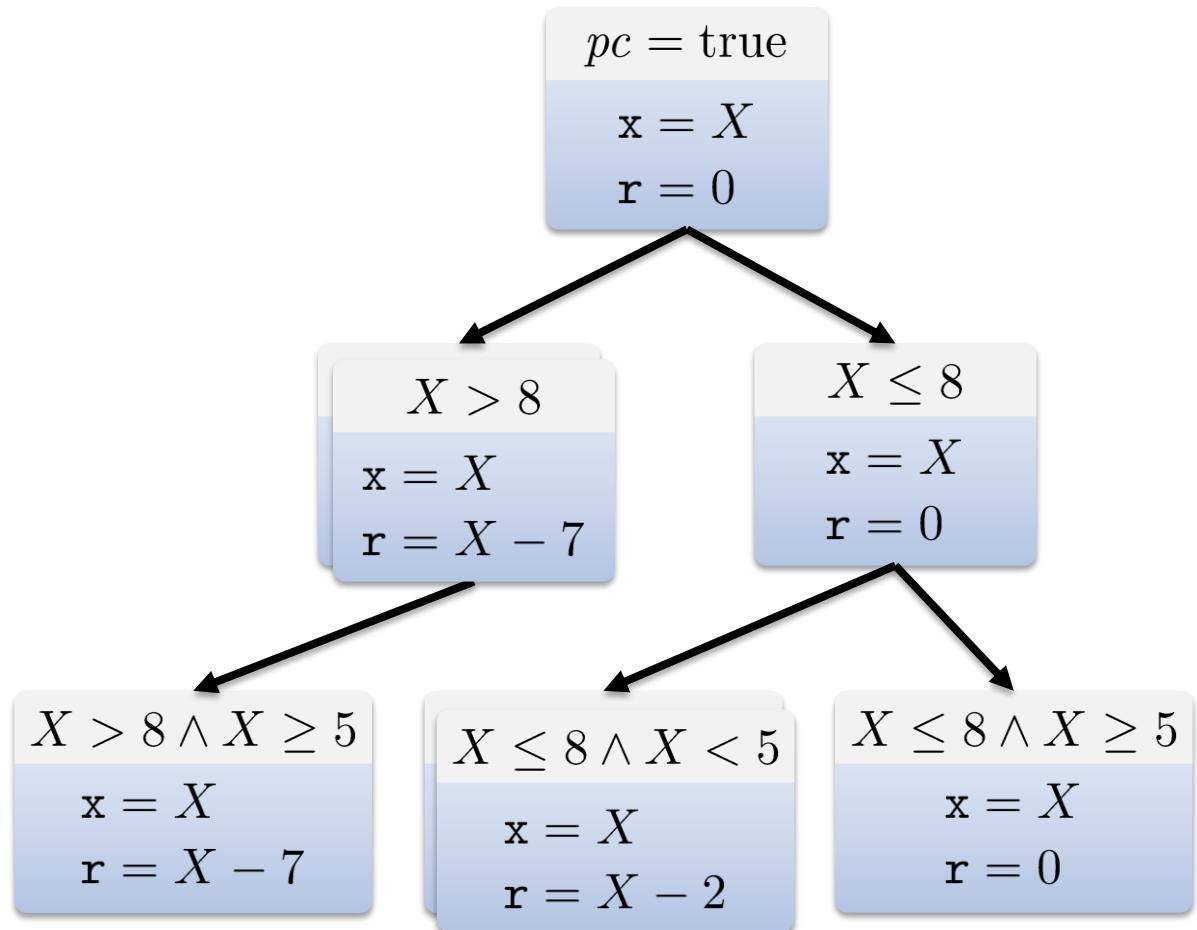
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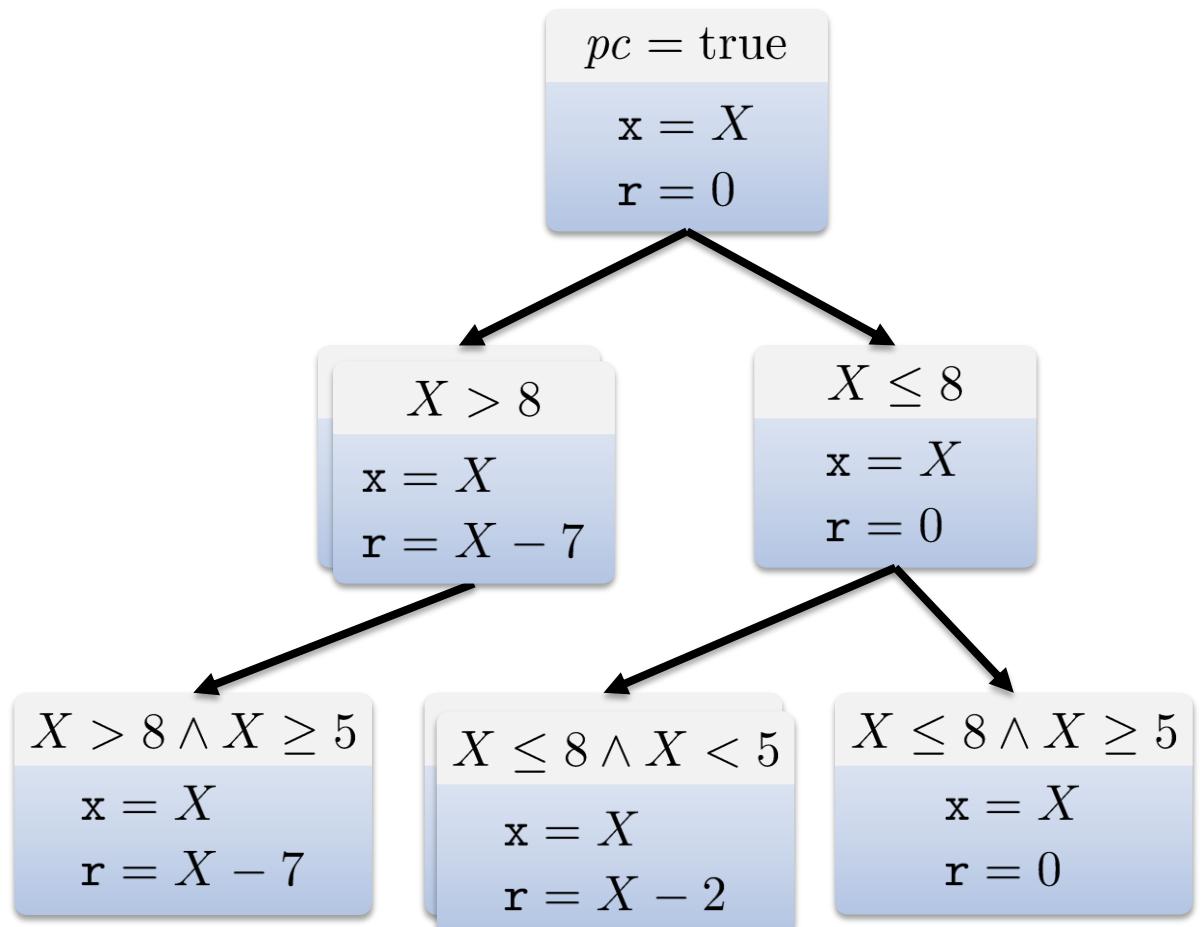
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```



Satisfying assignments:

$$X = 9$$

$$X = 4$$

$$X = 7$$

Test cases:

`proc(9)`

`proc(4)`

`proc(7)`

Finding Bugs

- Symbolic execution enumerates paths
 - *Runs into bugs that trigger whenever path executes*
 - *Assertions, buffer overflows, division by zero, etc., require specific conditions*
- Error conditions
 - *Treat assertions as conditions*

```
assert x != NULL      →      if (x == NULL)
                                abort();
```
 - *Creates explicit error paths*

Finding Bugs

- Instrument program with properties
 - *Translate any safety property to reachability*
- Division by zero

$y = 100 / x$



assert $x \neq 0$
 $y = 100 / x$

Finding Bugs

- Instrument program with properties
 - *Translate any safety property to reachability*
- Division by zero

$y = 100 / x$



```
assert x != 0  
y = 100 / x
```

- Buffer overflows

$a[x] = 10$



```
assert x >= 0 && x < len(a)
```

Finding Bugs

- Instrument program with properties
 - *Translate any safety property to reachability*
- Division by zero

$y = 100 / x$



```
assert x != 0  
y = 100 / x
```

- Buffer overflows

$a[x] = 10$



```
assert x >= 0 && x < len(a)
```

- Implementation is usually implicit

Symbolic Execution Algorithms

- Static symbolic execution
 - *Simulate execution on program source code*
 - *Computes strongest postconditions from entry point*
- Dynamic symbolic execution (DSE)
 - *Run / interpret the program with concrete state*
 - *Symbolic state computed in parallel (“concolic”)*
 - *Solver generates new concrete state*
- DSE-Flavors
 - *EXE-style [Cadar et al. ‘06] vs. DART [Godefroid et al. ‘05]*

EXECUTE

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$$pc = \text{true}$$

$$S(\mathbf{x}) = X \quad C(\mathbf{x}) = 1$$

$$C(\mathbf{r}) = 0$$

EXECUTE

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Symbolic
program state

$pc = \text{true}$
 $S(x) = X \quad C(x) = 1$
 $C(r) = 0$

EXECUTE

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Symbolic
program state

Concrete state	
$pc = \text{true}$	
$S(x) = X$	$C(x) = 1$
	$C(r) = 0$

EXECUTE

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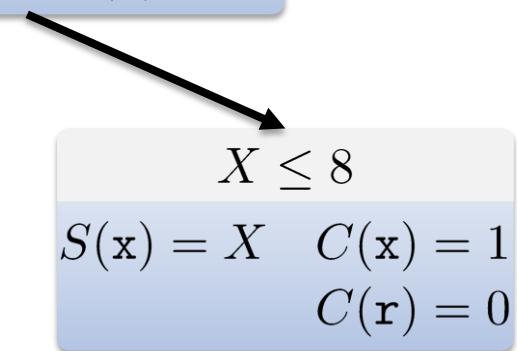
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$$pc = \text{true}$$

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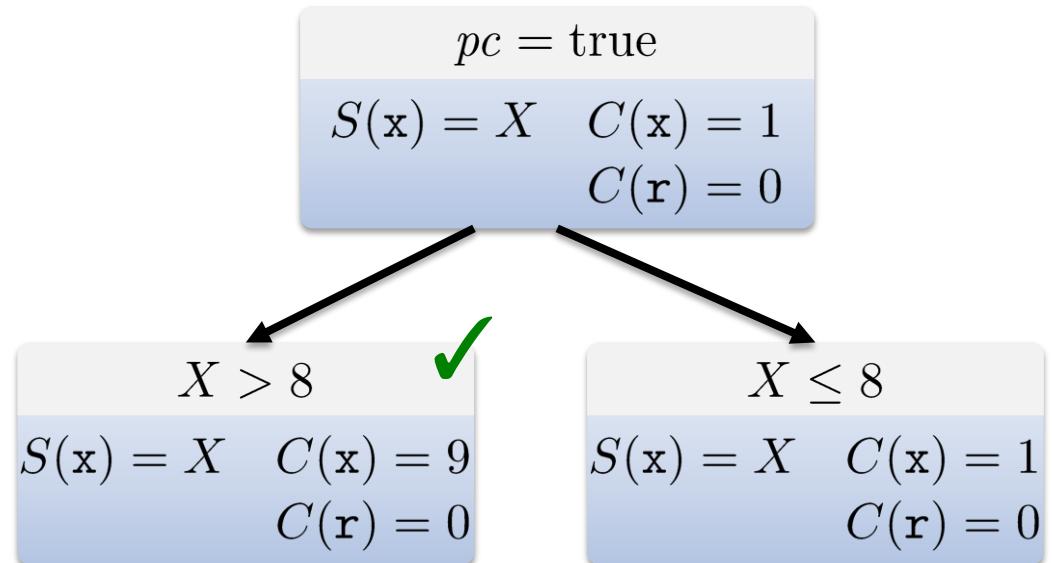
$$X \leq 8$$

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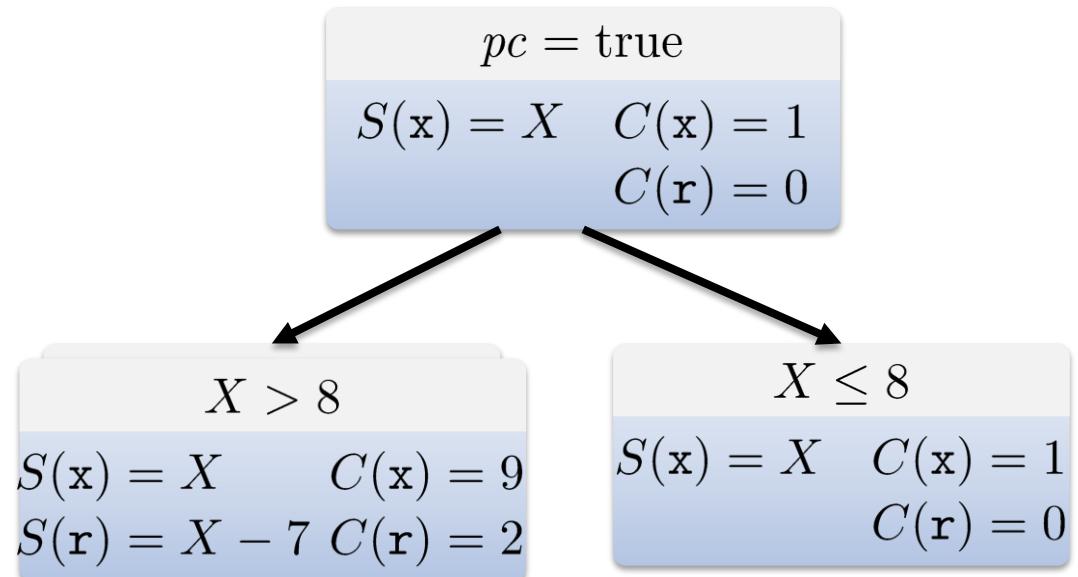
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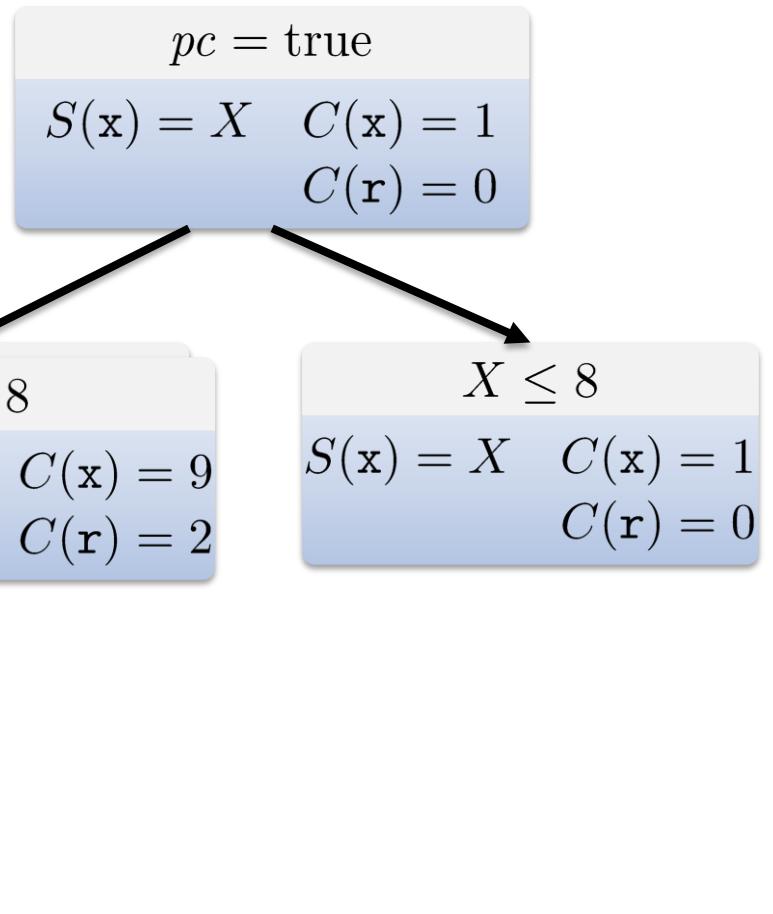
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```

$$pc = \text{true}$$

$$\begin{array}{ll} S(x) = X & C(x) = 1 \\ & C(r) = 0 \end{array}$$

$$X > 8$$

$$\begin{array}{ll} S(x) = X & C(x) = 9 \\ S(r) = X - 7 & C(r) = 2 \end{array}$$

$$X \leq 8$$

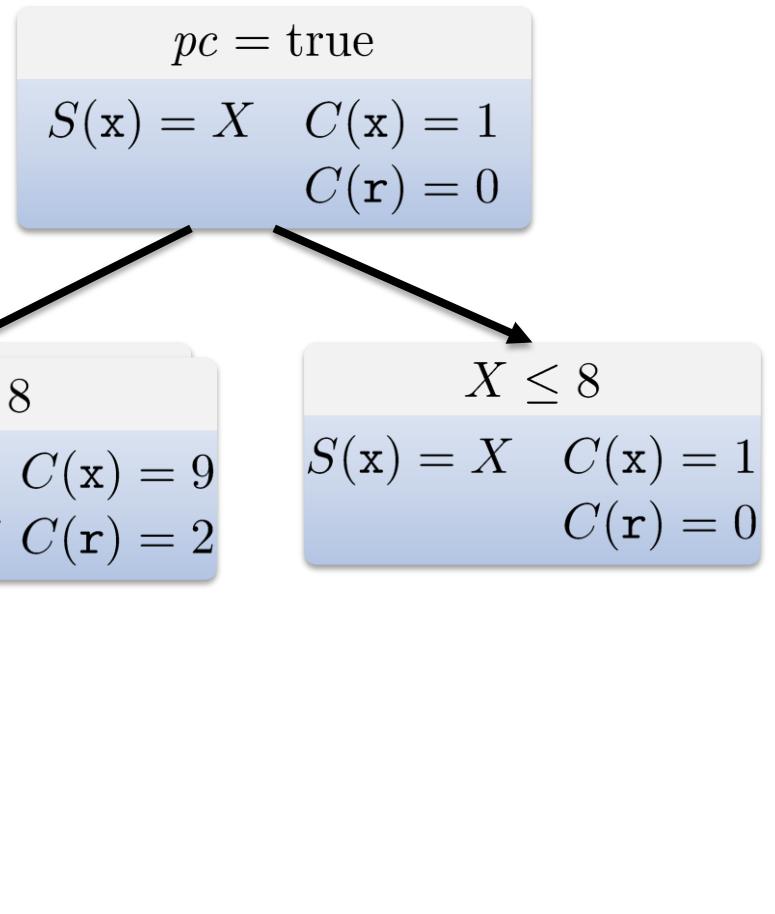
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$$X > 8 \wedge X < 5$$

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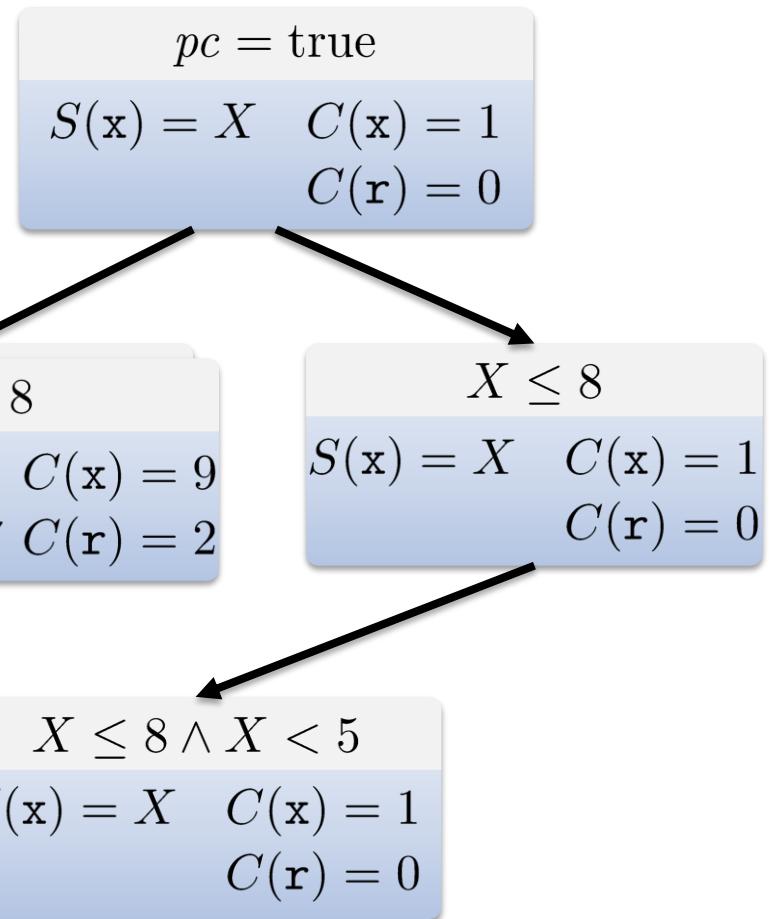
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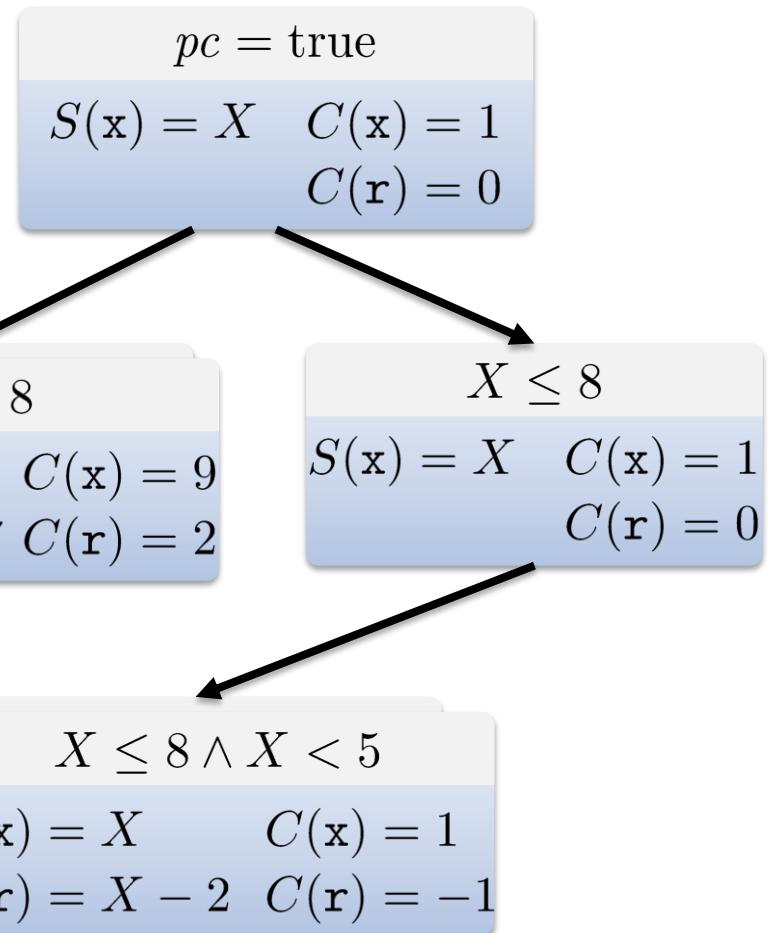
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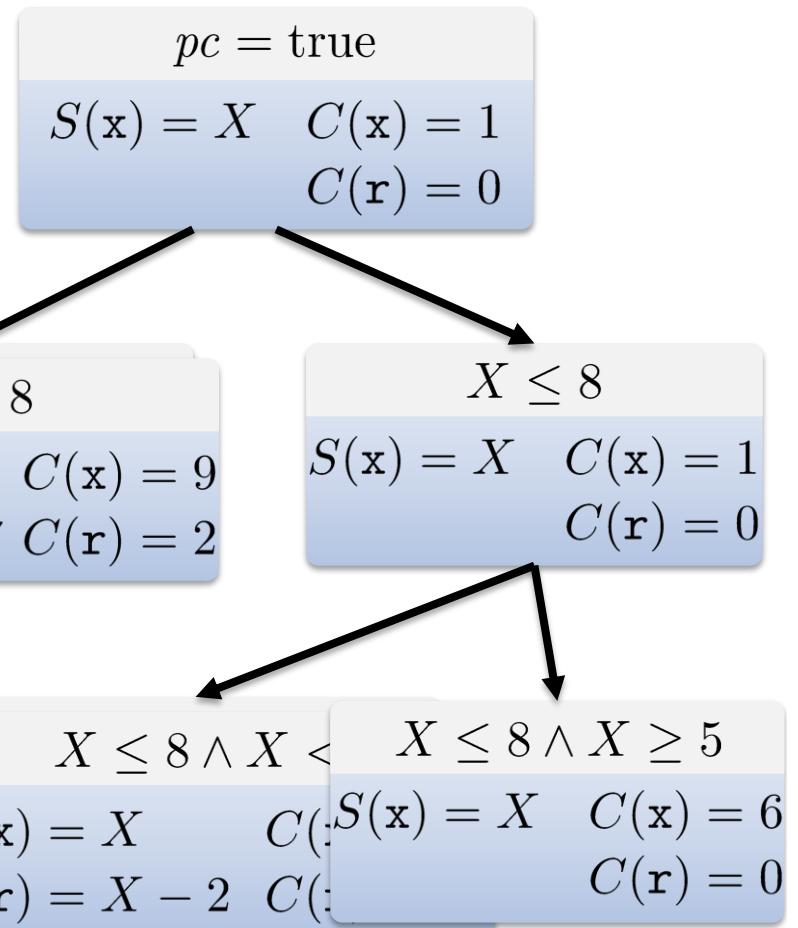
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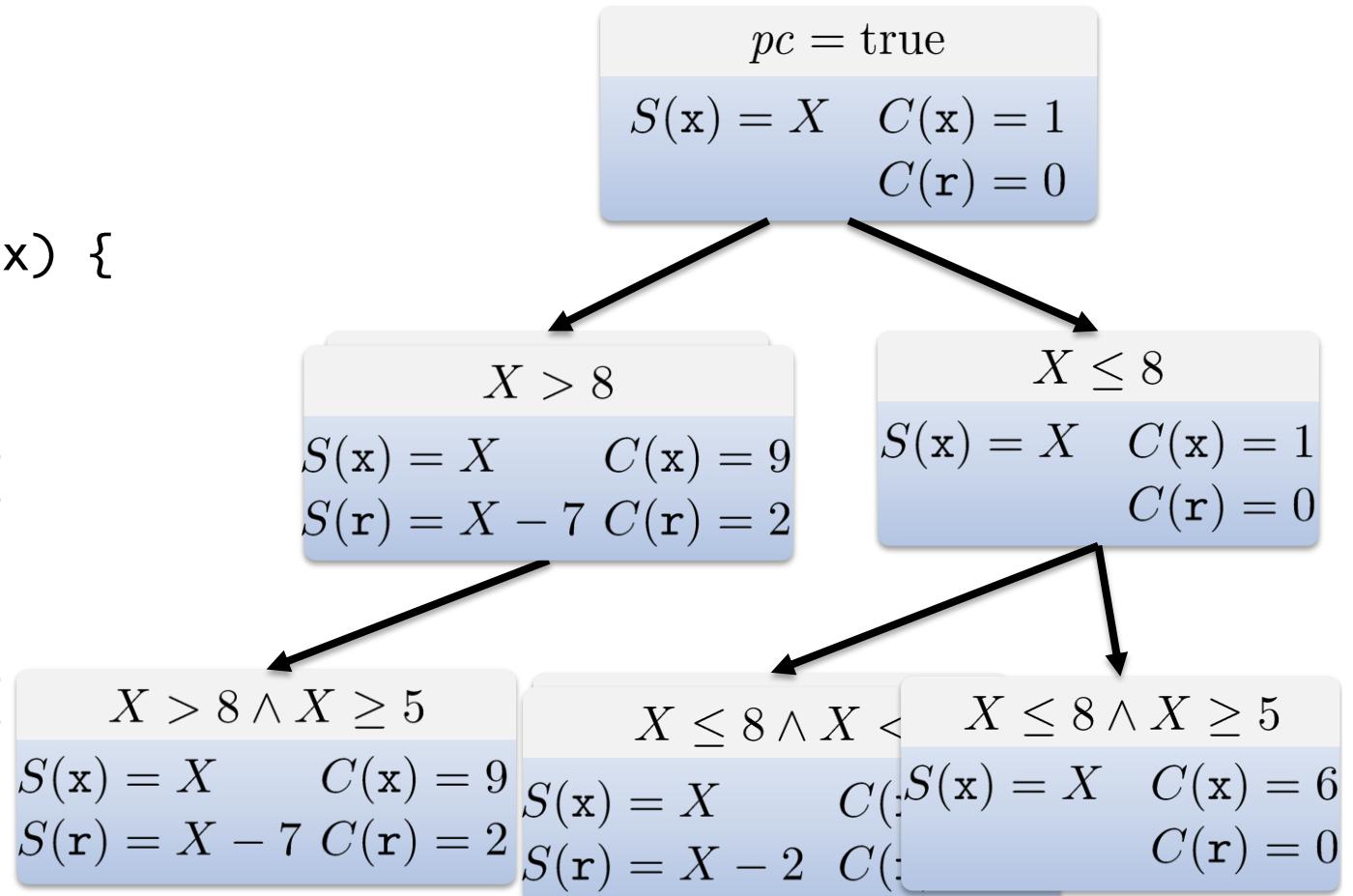
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Satisfying assignments:

$$X = 9$$

$$X = 1$$

$$X = 6$$

Test cases:

`proc(9)`

`proc(1)`

`proc(6)`

DART

```
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Path condition:

Test cases:

proc(1)

DART

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Path condition:

$pc = \text{true}$

$S(x) = X \quad C(x) = 1$
 $C(r) = 0$

$X \leq 8$

$S(x) = X \quad C(x) = 1$
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Test cases:

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Path condition:

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$$X \leq 8$$

$$\begin{aligned}S(\mathbf{x}) &= X & C(\mathbf{x}) &= 1 \\C(\mathbf{r}) &= 0\end{aligned}$$

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$$X \leq 8 \wedge \neg(X < 5)$$

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$$X \leq 8 \wedge X < 5$$

$$\begin{aligned}S(\mathbf{x}) &= X & C(\mathbf{x}) &= 1 \\S(\mathbf{r}) &= X - 2 & C(\mathbf{r}) &= -1\end{aligned}$$

$$X \leq 8 \wedge \neg(X < 5)$$

Test cases:

proc(1)

proc(6) 

DART

```
1 int proc(int x) {  
2  
3     int r = 0  
5  
6     if (x > 8) {  
7         r = x - 7  
8     }  
9  
10    if (x < 5) {  
11        r = x - 2  
12    }  
13  
14    return r;  
15 }
```

Path condition:

$$pc = \text{true}$$

$$\begin{aligned}S(\mathbf{x}) &= X & C(\mathbf{x}) &= 1 \\C(\mathbf{r}) &= 0\end{aligned}$$

$$X \leq 8$$

$$\begin{aligned}S(\mathbf{x}) &= X & C(\mathbf{x}) &= 1 \\C(\mathbf{r}) &= 0\end{aligned}$$

$$X \leq 8 \wedge X < 5 \quad X \leq 8 \wedge X \geq 5$$

$$\begin{aligned}S(\mathbf{x}) &= X & C(\mathbf{x}) &= 6 \\S(\mathbf{r}) &= X - 2 & C(\mathbf{r}) &= 0\end{aligned}$$

$$X \leq 8 \wedge \neg(X < 5)$$

Test cases:

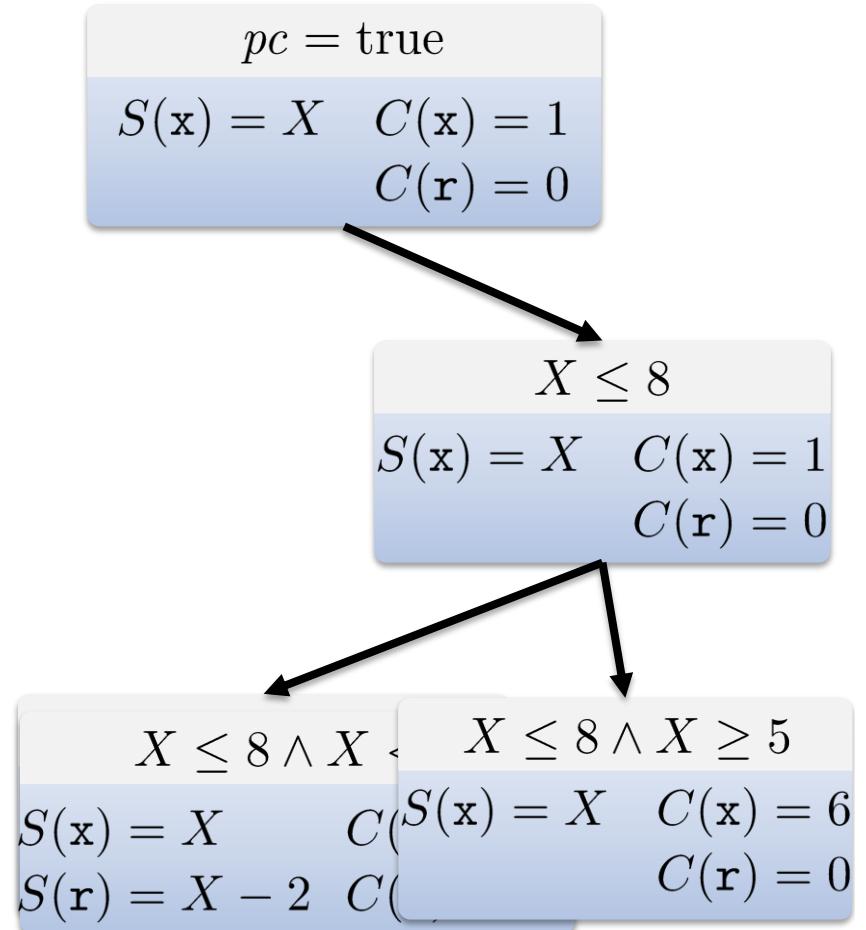
proc(1)

proc(6)

DART

```
1 int proc(int x) {  
2  
3     int r = 0  
5  
6     if (x > 8) {  
7         r = x - 7  
8     }  
9  
10    if (x < 5) {  
11        r = x - 2  
12    }  
13  
14    return r;  
15 }
```

Path condition:



Test cases:

$$\neg(X \leq 8)$$

proc(1)

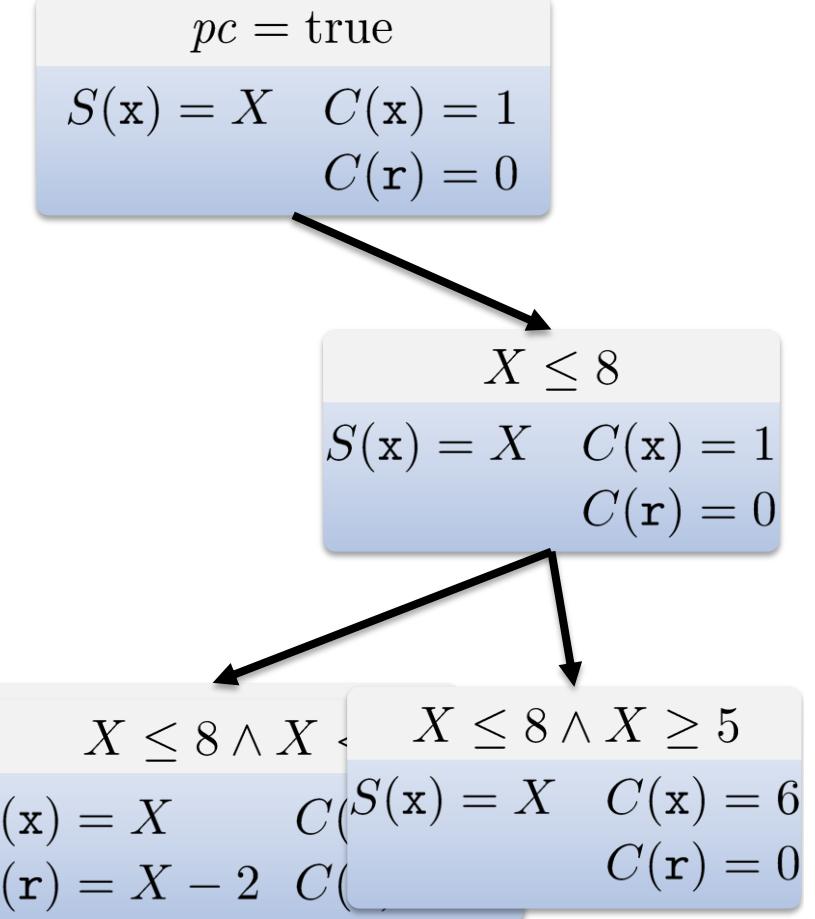
proc(6)

DART

```
1 int proc(int x) {  
2  
3     int r = 0  
5  
6     if (x > 8) {  
7         r = x - 7  
8     }  
9  
10    if (x < 5) {  
11        r = x - 2  
12    }  
13  
14    return r;  
15 }
```

Path condition:

Test cases:



$$\neg(X \leq 8)$$



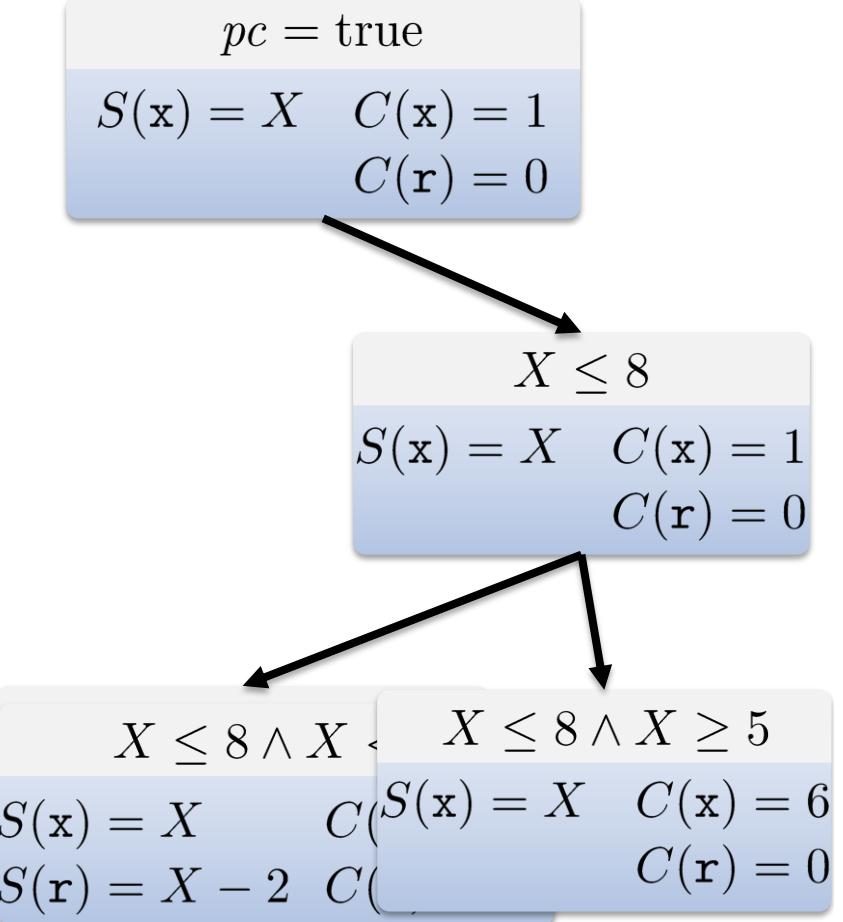
proc(1)

proc(6)

DART

```
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2  
3     int r = 0  
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6     if (x > 8) {  
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8     }  
9  
10    if (x < 5) {  
11        r = x - 2  
12    }  
13  
14    return r;  
15 }
```

Path condition:



Test cases:

proc(9)

proc(1)

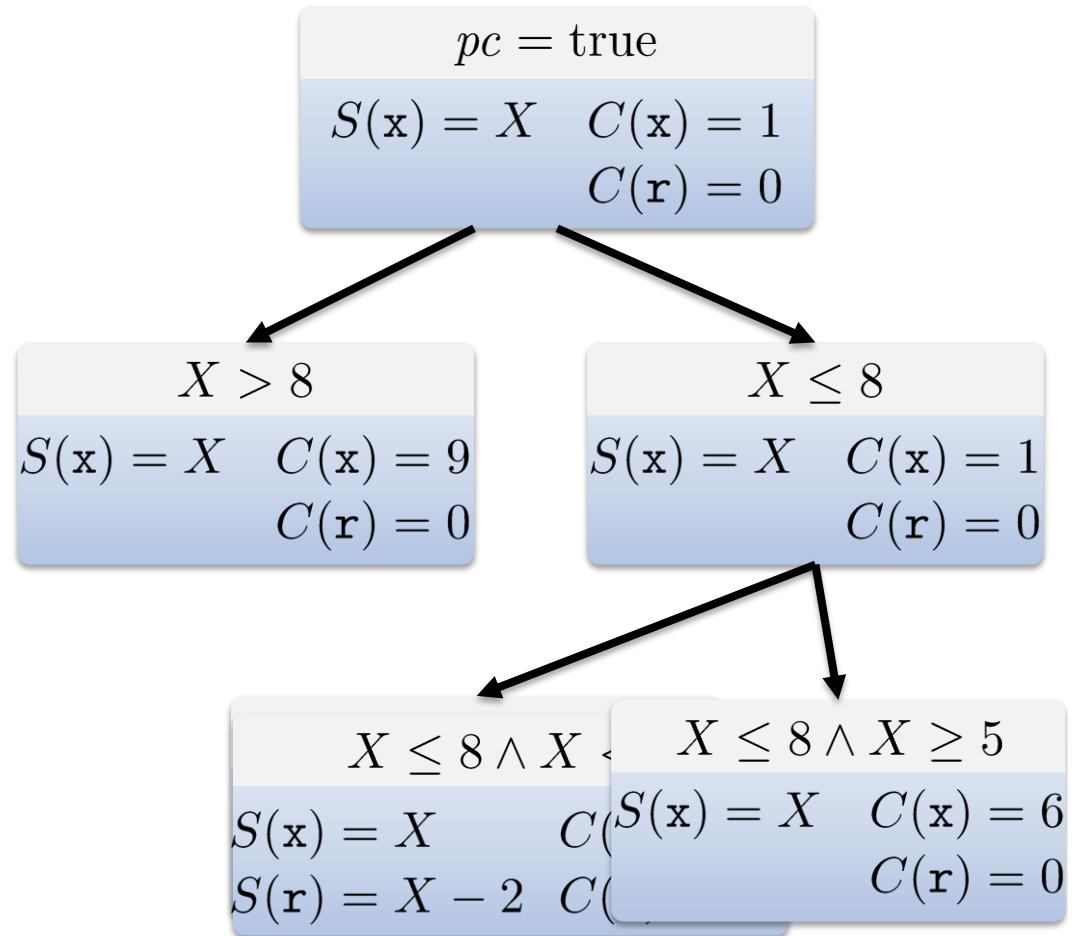
proc(6)

$$\neg(X \leq 8) \quad \checkmark$$

DART

```

1 int proc(int x) {
2
3     int r = 0
4
5     if (x > 8) {
6         r = x - 7
7     }
8
9
10    if (x < 5) {
11        r = x - 2
12    }
13
14    return r;
15 }
```



Path condition:

$$\neg(X \leq 8)$$

Test cases:

`proc(9)`

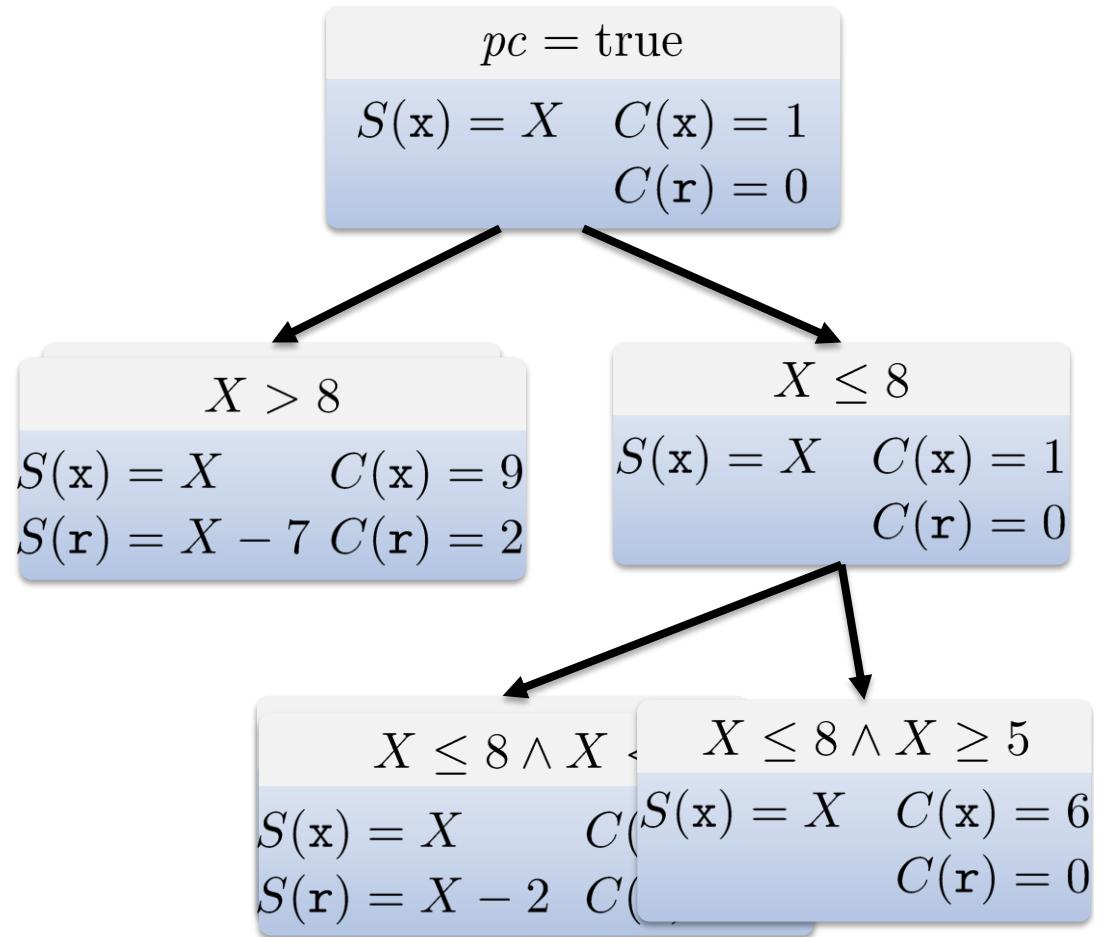
`proc(1)`

`proc(6)`

DART

```

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2
3     int r = 0
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12    }
13
14    return r;
15 }
```



Path condition:

$$\neg(X \leq 8)$$

Test cases:

`proc(9)`

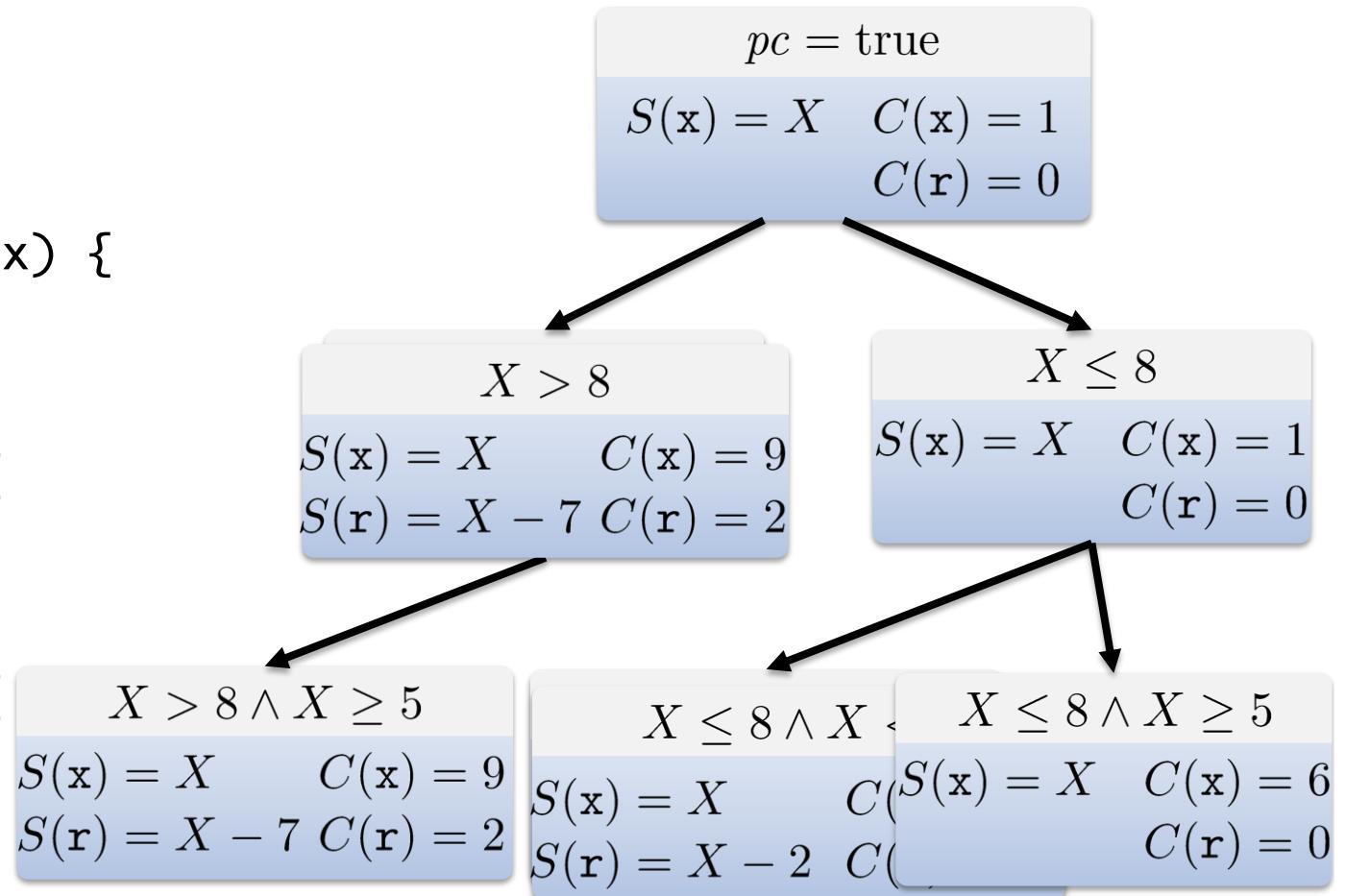
`proc(1)`

`proc(6)`

DART

```

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Path condition:

$$\neg(X \leq 8)$$

Test cases:

`proc(9)`

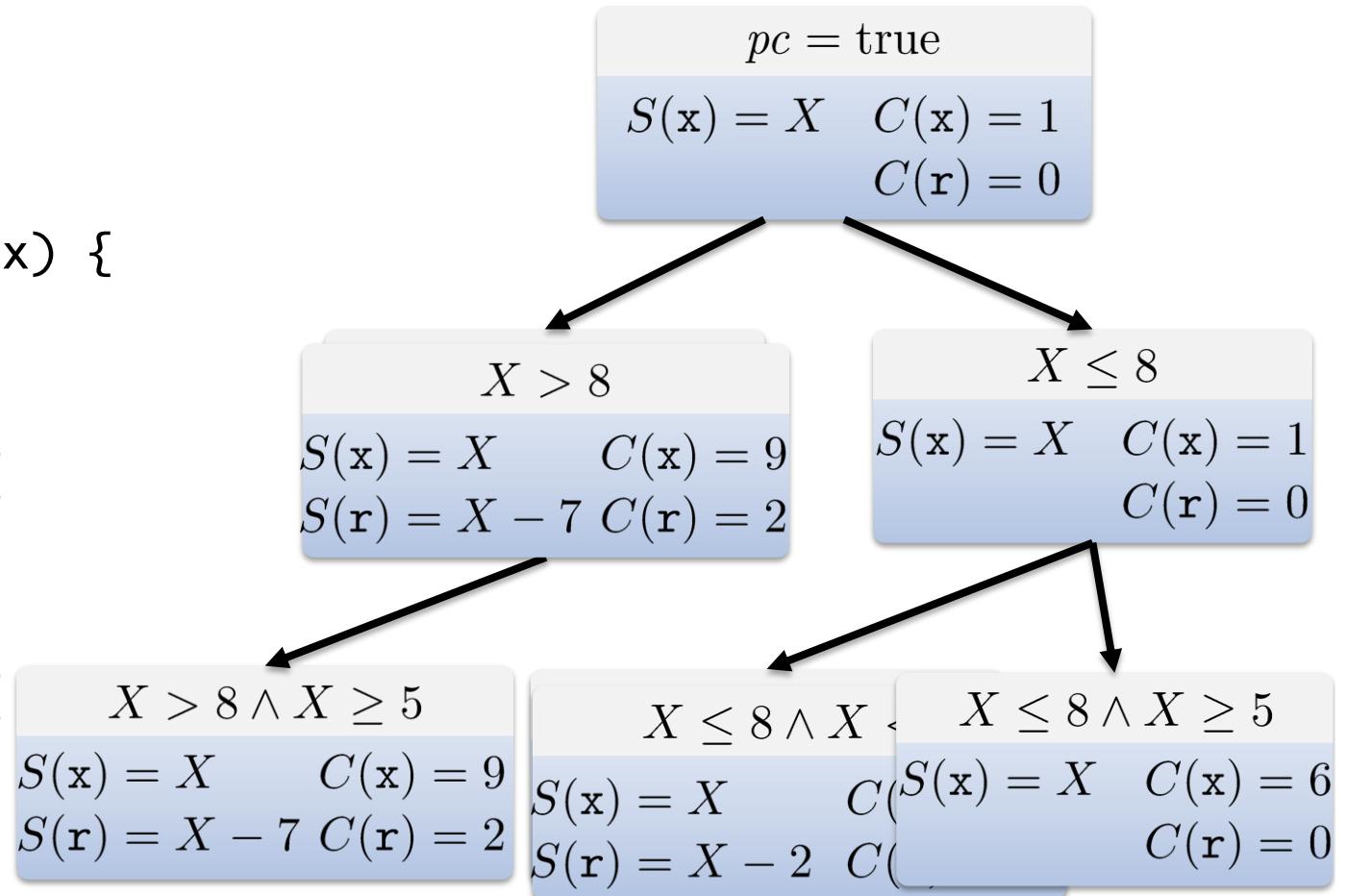
`proc(1)`

`proc(6)`

DART

```

1 int proc(int x) {
2
3     int r = 0
4
5     if (x > 8) {
6         r = x - 7
7     }
8
9
10    if (x < 5) {
11        r = x - 2
12    }
13
14    return r;
15 }
```



Path condition:

$$X > 8 \wedge \neg(X \geq 5)$$

Test cases:

`proc(9)`

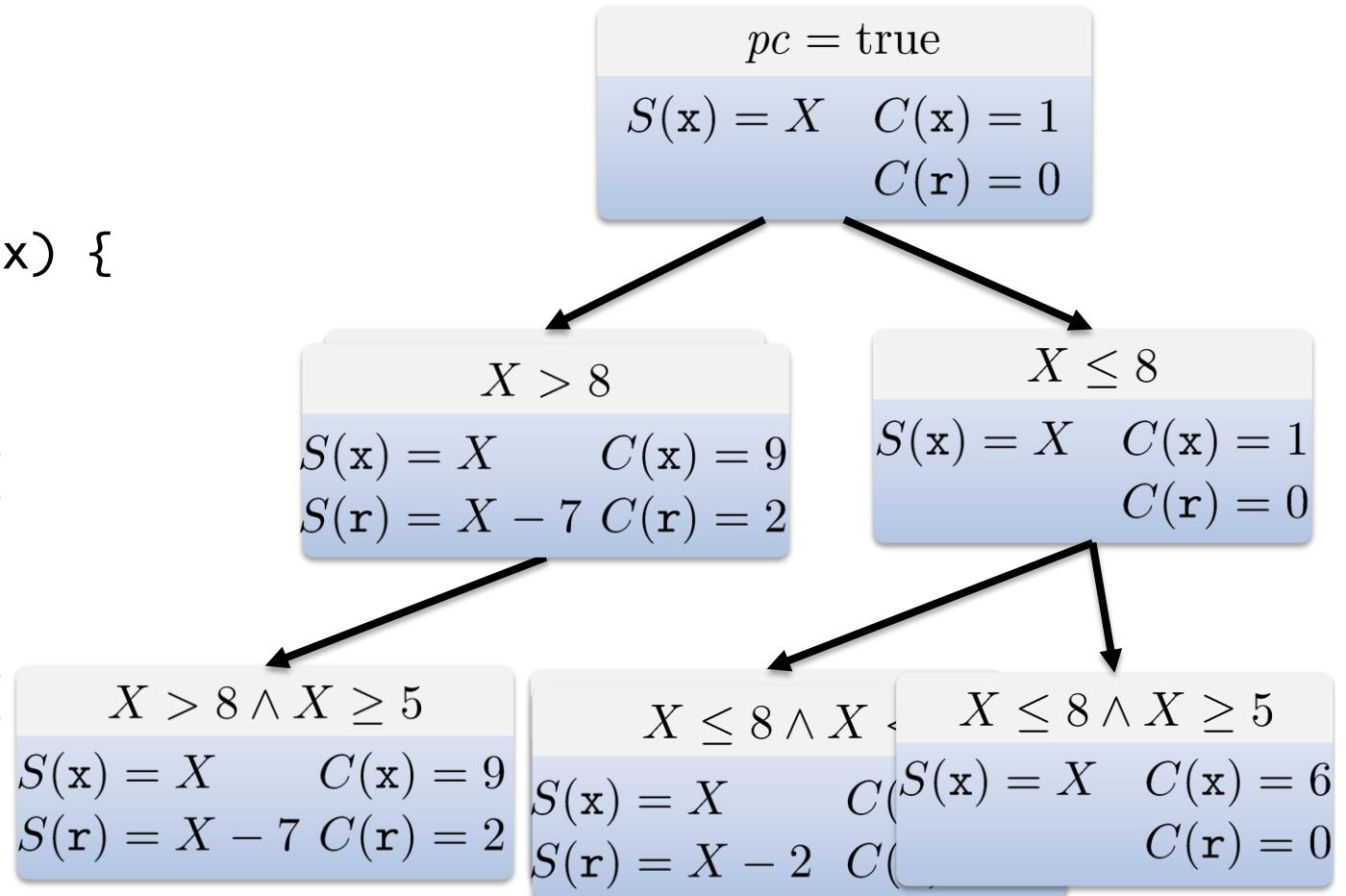
`proc(1)`

`proc(6)`

DART

```

1 int proc(int x) {
2
3     int r = 0
4
5     if (x > 8) {
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7     }
8
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10    if (x < 5) {
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13
14    return r;
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```



Path condition:

$$X > 8 \wedge \neg(X \geq 5)$$

Test cases:

`proc(9)`

`proc(1)`

`proc(6)`

DART vs. EXE

- Complete execution from 1st step
- Deep exploration
 - *One query per run*
- Offline SE possible
 - *Follow recorded trace*
- Fine-grained control of execution
- Shallow exploration
 - *Many queries early on*
- Online SE
 - *SE and interpretation in lockstep*

Concretization

- Parts of input space can be kept concrete
 - *Reduces complexity*
 - *Focuses search*
- Expressions can be concretized at runtime
 - *Avoid expressions outside of SMT solver theories (non-linear etc.)*
- Sound but incomplete

Concretization (Example)

<i>true</i>		if ($m*m > \text{size}$) {
	$C(X) = 5$...
$S(m) = X + 2$	$C(m) = 7$	
$S(\text{size}) = Y$	$C(\text{size}) = 256$	

Concretization (Example)

<i>true</i>		if ($m*m > \text{size}$) {
	$C(X) = 5$...
$S(m) = X + 2$	$C(m) = 7$	
$S(\text{size}) = Y$	$C(\text{size}) = 256$	
		$(X + 2)(X + 2) > Y$

Concretization (Example)

$true$	
	$C(X) = 5$
$S(m) = X + 2$	$C(m) = 7$

$S(\text{size}) = Y$ $C(\text{size}) = 256$

```
if (m*m > size) {
```

...

$$(X + 2)(X + 2) > Y$$

$$(5 + 2)(5 + 2) > Y$$

Concretization (Example)

$true$	
	$C(X) = 5$
$S(m) = X + 2$	$C(m) = 7$

$S(\text{size}) = Y$ $C(\text{size}) = 256$

```
if (m*m > size) {
```

...

$$(X + 2)(X + 2) > Y$$

$$49 > Y$$

Concretization (Example)

true

$$C(X) = 5$$

$$S(m) = X + 2$$

$$C(m) = 7$$

$$S(\text{size}) = Y$$

$$C(\text{size}) = 256$$

if ($m*m > \text{size}$) {

...

$$(X + 2)(X + 2) > Y$$

$$49 > Y$$

$$49 > Y$$

$$C(X) = 5$$

$$S(m) = X + 2$$

$$C(m) = 7$$

$$S(\text{size}) = Y$$

$$C(\text{size}) = 48$$

Concretization (Example)

true

$$\begin{array}{ll} C(X) = 5 & \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 256 \end{array}$$

```
if (m*m > size) {  
    ...  
    if (m < 5) {  
        ...  
    }  
}
```

$49 > Y$

$$\begin{array}{ll} C(X) = 5 & \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 48 \end{array}$$

Concretization (Example)

true

$$\begin{array}{ll} C(X) = 5 \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 256 \end{array}$$

if ($m*m > \text{size}$) {

...
if ($m < 5$) {
...
...

$X + 2 < 5$

$49 > Y$

$$\begin{array}{ll} C(X) = 5 \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 48 \end{array}$$

Concretization (Example)

true

$$\begin{array}{ll} C(X) = 5 & \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 256 \end{array}$$

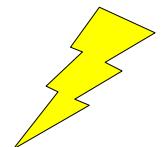
$49 > Y$

$$\begin{array}{ll} C(X) = 5 & \\ S(m) = X + 2 & C(m) = 7 \\ S(\text{size}) = Y & C(\text{size}) = 48 \end{array}$$

if ($m*m > \text{size}$) {

...
if ($m < 5$) {
...
...

$X + 2 < 5$



**Solution diverges from
expected path! (e.g., $X = 2$)**

Concretization (Example)

<i>true</i>		if ($m*m > \text{size}$) {
	$C(X) = 5$...
$S(m) = X + 2$	$C(m) = 7$	if ($m < 5$) {

$S(\text{size}) = Y$ $C(\text{size}) = 256$

...

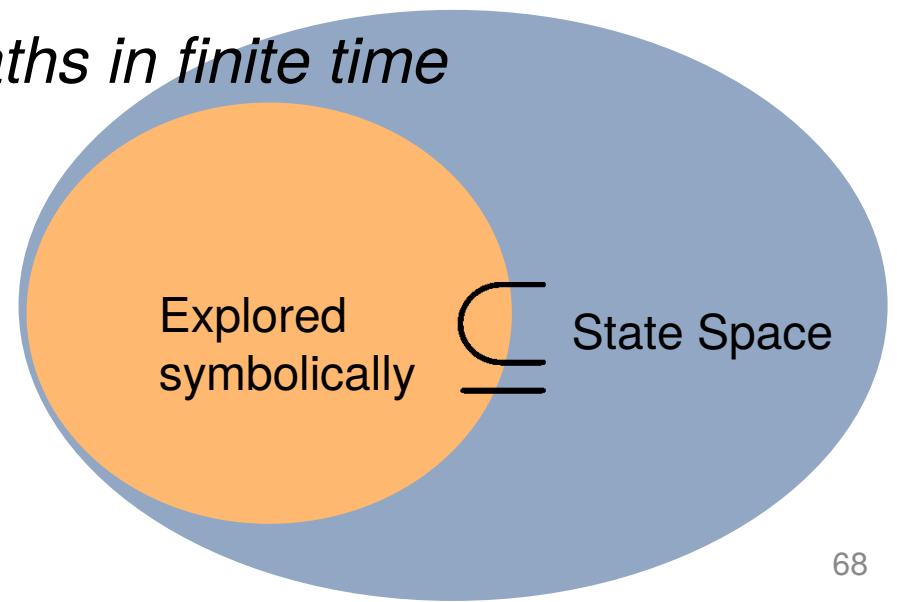
Concretization constraint	
$49 > Y$	$\wedge X = 5$
	$C(X) = 5$
$S(m) = X + 2$	$C(m) = 7$
$S(\text{size}) = Y$	$C(\text{size}) = 48$

Soundness & Completeness

- Conceptually, each path is exact
 - *Strongest postcondition in predicate transformer semantics*
 - *No over-approximation, no under-approximation*
- Globally, SE under-approximates
 - *Explores only subset of paths in finite time*
 - “*Eventual*” completeness

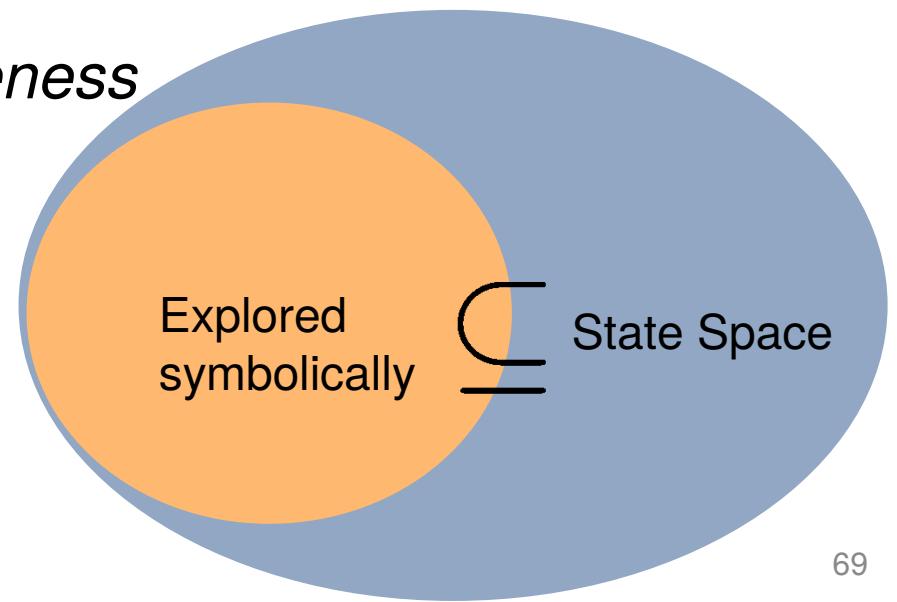
Soundness & Completeness

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 - “*Eventual*” completeness



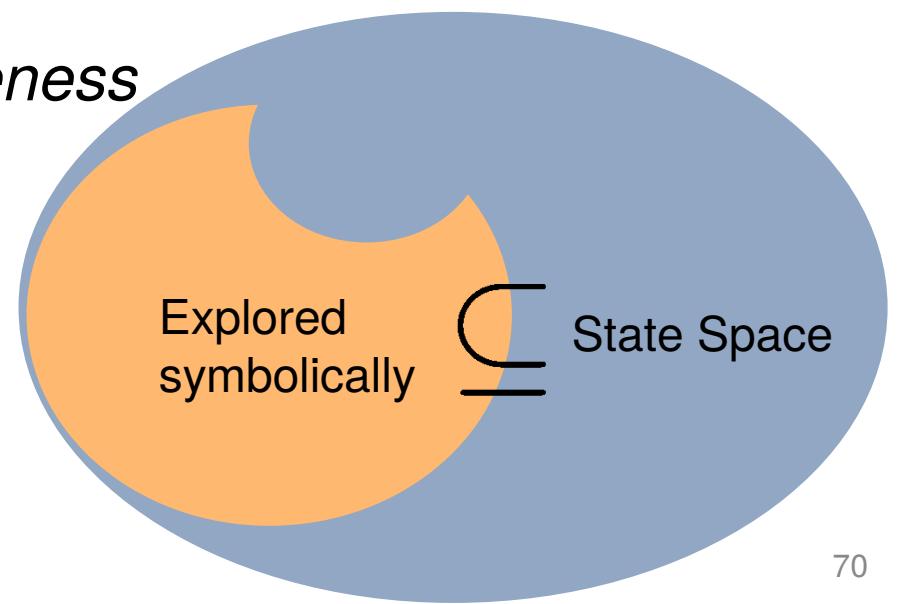
Soundness & Completeness

- Symbolic Execution = Underapproximates
 - *Soundness* = *does not include infeasible behavior*
 - *Completeness* = *explores all behavior*
- Concretization restricts state covered by path
 - *Remains sound*
 - *Loses (eventual) completeness*



Soundness & Completeness

- Symbolic Execution = Underapproximates
 - *Soundness* = *does not include infeasible behavior*
 - *Completeness* = *explores all behavior*
- Concretization restricts state covered by path
 - *Remains sound*
 - *Loses (eventual) completeness*



Concretization

- Key strength of dynamic symbolic execution
- Enables external calls
 - *Concretize call arguments*
 - *Callee executes concretely*
- Concretization constraints can be omitted
 - *Sacrifices soundness (original DART)*
 - *Deal with divergences by random restarts*